

Peut-on factoriser suffisamment rapidement les
nombres en facteurs premiers?

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Introduction et enjeux

La méthode de Dixon

Congruences de carrés

Etapes de la méthode

L'algorithme final

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Les nombres *RSA*

- ▶ Factoriser $N = pq$ où p et q sont premiers et très grands.
- ▶ Dernier nombre non factorisé: RSA-260 (260 chiffres)

$N = 221128255295296664352810852550262309276120895$
0247001539441374831912882294140200198651272972656
9746599085900330031400051170742204560859276357953
7571859542988389587092292384910067030341246205457
8456641366454068421436129301769402084639106587591
4794251435144458199

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Congruence de carrés

$N = pq$, p premier. Supp. $x^2 \equiv y^2 \pmod{N}$ et $x \not\equiv \pm y$.

- ▶ On a $x^2 - y^2 \equiv 0 \pmod{N}$ i.e. $N \mid (x - y)(x + y)$
- ▶ Donc $p \mid (x - y)(x + y)$
- ▶ Lemme d'Euclide: par exemple $p \mid x - y$
- ▶ Alors p divise N et $x - y$: $p \mid N \wedge (x - y)$, ce qui donne $N \wedge (x - y) \neq 1$

Conclusion

$N \wedge (x - y)$ et $N \wedge (x + y)$ sont des facteurs non-triviaux de N

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$$b \in \mathbb{N}$$

2

3

5

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•

•

p_b

$$b \in \mathbb{N}$$

$$(x_1, \quad x_2, \quad x_3, \quad \dots, \quad x_{b+1})$$

2

3

5

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p_b

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$$\cdot \quad \cdot$$

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$$\cdot \qquad \cdot$$

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$$\left. \begin{array}{l} 2 \qquad v_1^{(1)} \\ 3 \qquad v_1^{(2)} \\ 5 \qquad v_1^{(3)} \\ \cdot \qquad \cdot \\ \cdot \qquad \cdot \\ \cdot \qquad \cdot \\ p_b \qquad v_1^{(b)} \end{array} \right\} x_1^2 \pmod{N} = \prod_{i=1}^b p_i^{v_1^{(i)}} = 2^{v_1^{(1)}} 3^{v_1^{(2)}} \dots p_b^{v_1^{(b)}}$$

$$b \in \mathbb{N}$$

$$(x_1, \quad x_2, \quad x_3, \quad \dots, \quad x_{b+1})$$

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2

$$v_1^{(1)}$$

3

$$v_1^{(2)}$$

5

$$v_1^{(3)}$$

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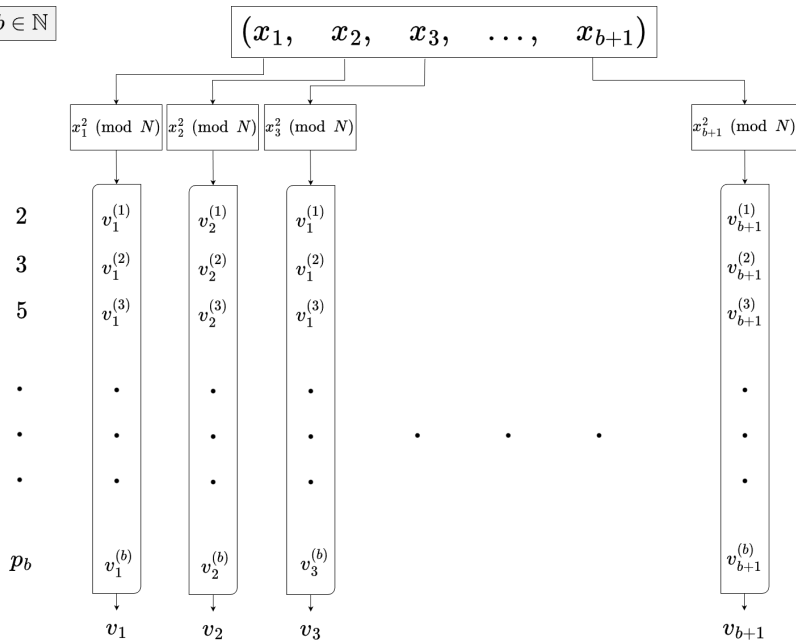
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$$b \in \mathbb{N}$$



Construction de y - Pivot de Gauss

- ▶ On a $b + 1$ vecteurs de \mathbb{F}_2^b et \mathbb{F}_2 corps, cela donne un système lié:

$$\exists (\lambda_i)_{i \in \llbracket 1, b+1 \rrbracket} \in \{0, 1\}^{b+1} \mid \sum_{i=1}^{b+1} \lambda_i v_i = 0_{\mathbb{F}_2^b} = (2\alpha_1, \dots, 2\alpha_b)$$

On pose $y = \prod_{j=1}^b p_j^{\alpha_j}$ et $x = \prod_{j=1}^{b+1} x_j^{\lambda_j}$

- ▶ On peut trouver les λ_i avec un système que l'on résout avec un **pivot de Gauss**.

Résultat admis (calcul)

$$x^2 \equiv y^2 \pmod{N}$$

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Un exemple

- $N = 20382493 = 3467 \times 5879$ et $b = 4, (2, 3, 5, 7)$
- Ces $5 = b + 1$ nombres x_j vérifient x_j^2
 $(\text{mod } N) = 2^{v_j^{(1)}} \dots 7^{v_j^{(4)}} :$

x_j	v_j
16853	(6, 5, 2, 2)
32877	(3, 0, 7, 0)
35261	(3, 2, 1, 0)
48834	(0, 2, 3, 1)

► $N = 20382493 = 3467 \times 5879$ et $b = 4$.

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► On résout dans \mathbb{F}_2^5

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$\lambda = (1, 1, 1, 0, 1)$ solution.

$$\begin{aligned} \text{► } x &= \prod_{j=1}^{b+1} x_j^{\lambda_j} = 7248176 \\ y &= \prod_{j=1}^b p_j^{\alpha_j} = 4837786 \end{aligned}$$

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Ce qu'il faut retenir

Le résultat principal

Étant donné $b \in \mathbb{B}$, trouver $b + 1$ nombres tels que
 $\forall j \in \llbracket 1, b + 1 \rrbracket, x_j^2 \pmod{N}$ a ses facteurs premiers inférieurs à p_b

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Algorithme 1 Recherche de nombres B-friables

Entrée: $N \in \mathbb{N}$ composé, $b \in \mathbb{N}$

Sortie: $(v_i)_{i \in \llbracket 1, b+1 \rrbracket}, (x_i)_{i \in \llbracket 1, b+1 \rrbracket}$

```
1: pour  $i \leftarrow 1 \dots b + 1$  faire
2:    $en\_cours \leftarrow V$ 
3:   tant que  $en\_cours$  faire
4:      $x_i \leftarrow \mathbb{U}(1, N - 1)$ 
5:      $x'_i \leftarrow x_i^2 \bmod N$ 
6:     si  $x'_i$  est B-friable (par algorithme naïf) alors
7:        $en\_cours \leftarrow F$ 
8:        $v_i \leftarrow (v_i^{(1)}, \dots, v_i^{(b)})$ 
   renvoyer  $(v_i)_{i \in \llbracket 1, b+1 \rrbracket}, (x_i)_{i \in \llbracket 1, b+1 \rrbracket}$ 
```

L'algorithme final

Algorithme 2 Factorisation par la méthode de Dixon

Entrée: $N \in \mathbb{N}$ composé, $B \in \mathbb{N}$

Sortie: p et q , tels que $p \mid N$ et $q \mid N$

- 1: $b \leftarrow \pi(B)$
 - 2: $(v_i)_{i \in \llbracket 1, b+1 \rrbracket}, (x_i)_{i \in \llbracket 1, b+1 \rrbracket} \leftarrow \text{RechercheBFriables}(N, b)$
 - 3: $(\lambda_i)_{i \in \llbracket 1, b+1 \rrbracket} \leftarrow \text{PivotdeGauss}((v_i)_{i \in \llbracket 1, b+1 \rrbracket})$
 - 4: $x \leftarrow \prod_{j=1}^{b+1} x_j^{\lambda_j}$
 - 5: $y \leftarrow \prod_{j=1}^b p_j^{\alpha_j}$
- renvoyer** $N \wedge (x - y), N \wedge (x + y)$
-

Etude théorique (Louise Nguyen)

Une minoration de la densité des B -friables

Soit $B : \mathbb{N}^* \rightarrow \mathbb{N}^*$ une fonction telle que $\ln n = o(B(n))$ et $\ln B(n) = o(\ln n)$. Alors on a, pour $n \rightarrow +\infty$,

$$\Psi(B(n), n) \geq n \exp \left(\left(\frac{\ln n}{\ln B(n)} \ln \ln n \right) (-1 + o(1)) \right)$$

Une complexité sous-exponentielle

$$\exp \left((1 + o(1)) 2\sqrt{2} (\ln n \ln \ln n)^{1/2} \right)$$

lorsque $B = \exp \left(\frac{1}{\sqrt{2}} (\ln n \ln \ln n)^{1/2} \right)$

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Principe

- ▶ Utilisation d'un polynôme $Q = (\sqrt{N} + X)^2 - N$ pour générer les x_i
- ▶ Résolution de $Q(x) \equiv 0 \pmod{p}$ grâce à Tonelli-Shanks, 2 solutions x_1 et x_2 dans $\llbracket 1, p \rrbracket$.
- ▶ $p \mid Q(x) \implies \forall k \in \mathbb{N}, p \mid Q(x + kp)$
- ▶ Cribler sur un intervalle $\llbracket 1, S \rrbracket$, puis sur $\llbracket S + 1, 2S \rrbracket$ etc...

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$$S = 10$$

$$N = 20382493$$

$$T = [Q(1), Q(2), Q(3), Q(4), Q(5), Q(6), Q(7), Q(8), Q(9), Q(10)]$$

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$$[2732, \quad 11736, \quad 20796, \quad 29831, \quad 38868, \quad \dots, \quad Q(10)]$$

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A sequence of numbers is shown: $[2732, 11736, 20796, 29831, 38868, \dots, Q(10)]$. Two curved arrows indicate an increment of 2. The first arrow starts above 2732 and points to 20796, with a "+2" label above it. The second arrow starts above 20796 and points to 38868, also with a "+2" label above it.

$$[2732, 11736, 20796, 29831, 38868, \dots, Q(10)]$$

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$$\begin{array}{c}
 \begin{array}{ccccccc}
 & & +2 & & +2 & & \\
 & \swarrow & & \searrow & & \swarrow & \\
 [2732, & 11736, & 20796, & 29831, & 38868, & \dots, & Q(10)] \\
 \div 2^2 & \searrow & & & & & \\
 & [683, & 11736, & 5199, & 29831, & 9717, & \dots, & Q(10)]
 \end{array}
 \end{array}$$

$v_1^{(1)} = 2$

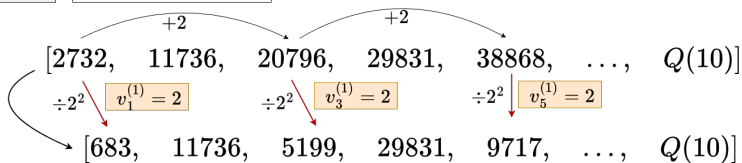
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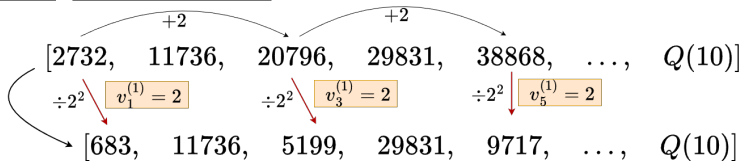
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$$p = 2$$

$$Q(1) \equiv 0 \pmod{2}$$



$$p = 3$$

$$Q(2) \equiv 0 \pmod{3}$$

$$Q(3) \equiv 0 \pmod{3}$$

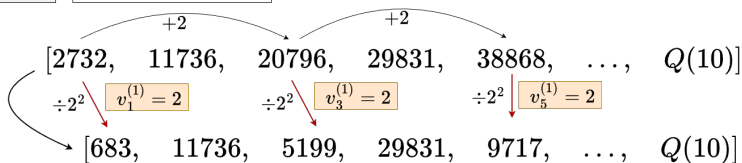
$$S = 10$$

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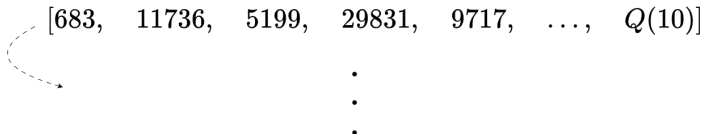
$$Q(1) \equiv 0 \pmod{2}$$



$$p = 3$$

$$Q(2) \equiv 0 \pmod{3}$$

$$Q(3) \equiv 0 \pmod{3}$$



$$p = p_b$$

$$Q(x_1) \equiv 0 \pmod{p_b}$$

$$Q(x_2) \equiv 0 \pmod{p_b}$$

Algorithme 3 Algorithme du crible quadratique

Entrée: $N \in \mathbb{N}^*$, $b \in \mathbb{N}^*$, $S \geq 1$

Sortie: $(v_i)_{i \in \llbracket 1, k \rrbracket}$, $(x_i)_{i \in \llbracket 1, k \rrbracket}$, $k \in \llbracket 0, S \rrbracket$

```

1:  $T \leftarrow$  tableau tel que  $T[i] \leftarrow (i + \lfloor \sqrt{N} \rfloor)^2 - N$  pour  $i \in \llbracket 1, S \rrbracket$ 
2:  $V \leftarrow$  tableau tel que  $V[i] \leftarrow (0, \dots, 0) \in \mathbb{N}^b$  pour  $i \in \llbracket 1, S \rrbracket$ 
3: pour  $p \in \{p_1, \dots, p_b\}$  tel que  $N$  est un carré modulo  $p$  faire
4:    $x_1, x_2 \leftarrow$  les racines de  $(X + \lfloor \sqrt{N} \rfloor)^2 - N$  modulo  $p$ 
5:   pour  $i \in \{1, 2\}$  faire
6:      $q \leftarrow x_i$ 
7:     tant que  $q \leq S$  faire
8:       tant que  $T[q] \bmod p = 0$  faire
9:          $T[q] \leftarrow T[q]/p$ 
10:         $V[q] \leftarrow V[q] + (0, \dots, 1, \dots, 0)$  (en position  $p$ )
11:        $q \leftarrow q + p$ 
renvoyer L'ensemble des  $(i + \lfloor \sqrt{N} \rfloor, V[i])$  tels que  $T[i] = 1$  pour
 $i \in \llbracket 1, S \rrbracket$ 

```

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Justification

- ▶ $O(n)$ au lieu de $O(n^2)$, voire $O(n \log n)$
- ▶ $Q(x) = \prod_{i=1}^k p_i^{\alpha_i}$, soit $\ln(Q(x)) = \sum_{i=1}^k \alpha_i \ln(p_i)$.
Idée: soustraire par $\alpha_i \ln(p_i)$ au lieu de diviser par $p_i^{\alpha_i}$
- ▶ $\log_2(Q(x)) \approx \text{nb_bits}(Q(x))$
- ▶ Problème: on ne connaît pas α_i .
Solution: on soustrait par $\log_2(p_i)$ seulement. Des approximations nécessitent déjà un **seuil**

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$$S = 10$$

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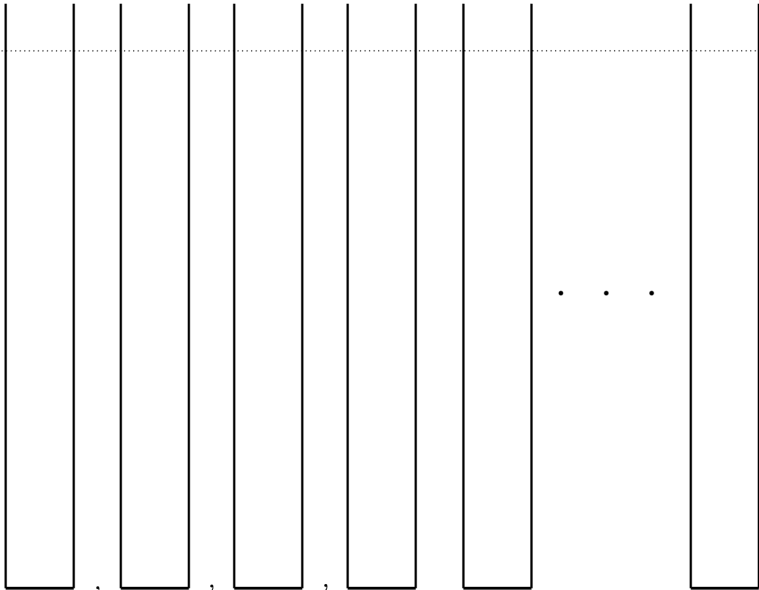
$$T = [\quad , \quad , \quad , \quad , \quad , \quad , \quad]$$

$$S = 10$$

$$N = 20382493$$

Seuil

$T = [$



$$S = 10$$

$$N = 20382493$$

$$Q(1) \equiv 0 \pmod{2}$$

Seuil

$T = [$

$\log_2 2$

,

+2

$\log_2 2$

,

+2

$\log_2 2$

,

...

$\log_2 2$

]

$$S = 10$$

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Seuil

$T = [$

$\log_2 2$

,

$\log_2 3$

,

$\log_2 3$
 $\log_2 2$

,

,

$\log_2 3$
 $\log_2 2$

,

,

$]$

+3

+3

.

.

.

$$S = 10$$

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Seuil

$$T = [\log_2 2, \log_2 3, \log_2 2, \log_2 3, \log_2 2, \dots]$$

$$S = 10$$

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Seuil

$$T = [\log_2 2, \log_2 3, \log_2 2, -, \log_2 3, \log_2 2, \dots]$$

$$\log_2 3$$

$$\log_2 2$$

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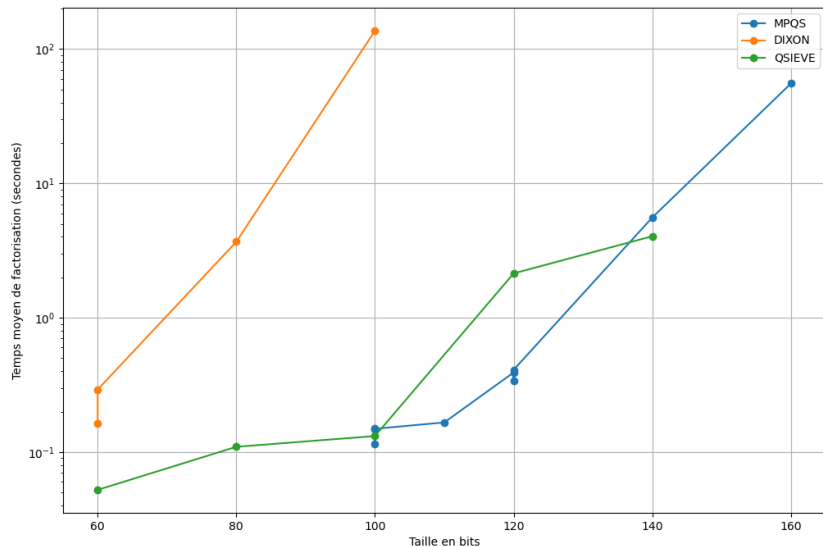
Annexe

Résultats

Après plusieurs centaines de tests, on a les résultats suivants:

Bits	Dixon	QSIEVE	MPQS
60	0.5s	0.05s	-
80	5s	0.1s	-
100	100s	0.1s	0.1s
120	-	2s	0.6s
140	-	5s	5s
160	-	-	80s

Graphique final



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Proposition

Soient $b \in \mathbb{N}$, $(x_i)_{i \in \llbracket 1, b+1 \rrbracket} \in \mathbb{N}^{b+1}$ et $(v_i)_{i \in \llbracket 1, b+1 \rrbracket} \in \mathbb{F}_2^b$ les vecteurs valuations de $x_i^2 \pmod{N}$ pour $i \in \llbracket 1, b+1 \rrbracket$ et finalement $(\lambda_i)_{i \in \llbracket 1, b+1 \rrbracket} \in \{0, 1\}^{b+1}$ tels que,

$$\sum_{i=1}^{b+1} \lambda_i v_i = 0_{\mathbb{F}_2^b} = (2\alpha_1, \dots, 2\alpha_b)$$

On pose $y = \prod_{j=1}^b p_j^{\alpha_j}$ et $x = \prod_{j=1}^{b+1} x_j^{\lambda_j}$, alors $x^2 \equiv y^2 \pmod{N}$

Démonstration

$$\begin{aligned}
 x^2 &= \left(\prod_{i=1}^{b+1} x_i^2 \right)^{\lambda_i} \equiv \prod_{i=1}^{b+1} \prod_{j=1}^b p_j^{\lambda_i v_i^{(j)}} \pmod{N} \\
 &\equiv \prod_{j=1}^b \prod_{i=1}^{b+1} p_j^{\lambda_i v_i^{(j)}} \pmod{N} \\
 &\equiv \prod_{j=1}^b p_j^{\sum_{i=1}^{b+1} \lambda_i v_i^{(j)}} \pmod{N} \\
 &\equiv \left(\prod_{j=1}^b p_j^{\alpha_j} \right)^2 \pmod{N} \quad (\text{déf de } \alpha_j) \\
 &\equiv y^2 \pmod{N}
 \end{aligned}$$

Proposition

Si $Q = (\sqrt{N} + X)^2 - N$, alors $p \mid Q(x) \implies \forall k \in \mathbb{N}, p \mid Q(x + kp)$

Démonstration

En effet, supposons $p \mid Q(x)$, on a:

$$\begin{aligned} Q(x + kp) &= (\sqrt{N} + x + kp)^2 - N \\ &= Q(x) + 2kp(\sqrt{N} + x) + k^2p^2 \end{aligned}$$

d'où $p \mid Q(x + kp)$

../c/vector.h

```
#pragma once
#include <gmp.h>

void mod_vect(int* v, int mod, int n1);
void add_vect(int* sum, int* op, int n1);
```

```
void div_vect(int* v, int d, int n1);
void sub_vect(int** v, int i, int j, int n1);
void prod_vect(mpz_t prod, mpz_t* z, int n1, system_t
s);
```

../c/vector.c

```
#include <gmp.h>
#include <assert.h>
#include <stdlib.h>
#include "system.h"

void mod_vect(int* v, int mod, int n1){
    for(int i = 0; i < n1; i++){
        v[i] = abs(v[i]) % mod;
    }
}

void add_vect(int* sum, int* op, int n1){
    for(int i = 0; i < n1; i++){
        sum[i] += op[i];
    }
}

void div_vect(int* v, int d, int n1){
    for(int i = 0; i < n1; i++){
        assert(v[i] % d == 0);
    }
}
```

```
        v[i] /= d;
    }
}

void sub_vect(int** v, int i, int j, int n1){
    for(int k = 0; k < n1; k++){
        v[i][k] = v[i][k] - v[j][k];
    }
}

void prod_vect(mpz_t prod, mpz_t* z, int n1, system_t
s){
    mpz_set_ui(prod, 1);
    for(int i = 0; i < n1; i++){
        if(s->sol[i]){
            mpz_mul(prod, prod, z[s->perm[i]]);
        }
    }
}
```

../c/tonellishanks.h

#pragma once

#include <gmp.h>

void tonelli_shanks_ui(mpz_t n, **int** p, **int*** x1, **int*** x2);

void tonelli_shanks_mpz(mpz_t a, mpz_t p, mpz_t x1,
mpz_t x2);

../c/tonellishanks.c

```
#include <stdint.h>
#include <gmp.h>
#include <stdio.h>
#include <assert.h>
#include <stdlib.h>

uint64_t modpow(uint64_t a, uint64_t b, uint64_t n) {
    uint64_t x = 1, y = a;
    while (b > 0) {
        if (b % 2 == 1) {
            x = (x * y) % n; // multiplying with base
        }
        y = (y * y) % n; // squaring the base
        b /= 2;
    }
    return x % n;
}

void tonelli_shanks_ui(mpz_t n, unsigned long int p, int
    * x1, int* x2) {
    uint64_t q = p - 1;
    uint64_t ss = 0;
    uint64_t z = 2;
    uint64_t c, r, t, m;

    while ((q & 1) == 0) {
        ss += 1;
        q >>= 1;
    }
```

```
    mpz_t temp, pj;
    mpz_init(temp);
    mpz_init_set_ui(pj, p);
```

```
    if (ss == 1) {
        //uint64_t r1 = modpow(n, (p + 1) / 4, p);
        mpz_powm_ui(temp, n, (p+1)/4, pj);
        uint64_t r1 = mpz_get_ui(temp);

        *x1 = r1;
        *x2 = p - r1;
        mpz_clears(temp, pj, NULL);
        return;
    }

    while (modpow(z, (p - 1) / 2, p) != (unsigned
        long int) p - 1) { // uint_64 only there
        for the compiler to stop complaining
        z++;
    }

    c = modpow(z, q, p);

    //r = modpow(n, (q + 1) / 2, p);
    mpz_powm_ui(temp, n, (q+1)/2, pj);
    r = mpz_get_ui(temp);

    //t = modpow(n, q, p);
    mpz_powm_ui(temp, n, q, pj);
```

```

t = mpz_get_ui(temp);

m = ss;

while(1){
    uint64_t i = 0, zz = t;
    uint64_t b = c, e;
    if (t == 1) {
        *x1 = r;
        *x2 = p - r;
        mpz_clears(temp, pj, NULL);
        return;
    }
    while (zz != 1 && i < (m - 1)) {
        zz = zz * zz % p;
        i++;
    }
    e = m - i - 1;
    while (e > 0) {
        b = b * b % p;
        e--;
    }
    r = r * b % p;
    c = b * b % p;
    t = t * c % p;
    m = i;
}

}

void tonelli_shanks_mpmz(mpz_t n, mpz_t p, mpz_t x1,
    mpz_t x2){
    assert(mpz_legendre(n, p) == 1);

```

```

    mpz_t q, z;
    mpz_init_set(q, p);
    mpz_sub_ui(q, q, 1);
    int ss = 0;
    mpz_init_set_ui(z, 2);

    while(mpz_divisible_ui_p(q, 2) != 0){
        ss += 1;
        mpz_divexact_ui(q, q, 2);
    }

    mpz_t op1;
    mpz_init(op1);

    if (ss == 1) {
        //uint64_t r1 = modpow(n, (p + 1) / 4, p);
        mpz_add_ui(op1, p, 1);
        mpz_divexact_ui(op1, op1, 4);
        mpz_powm(op1, n, op1, p);

        mpz_set(x1, op1);
        mpz_sub(x2, p, x1);

        mpz_clears(q, z, op1, NULL);
        return;
    }

    mpz_t op2, op3;
    mpz_inits(op2, op3, NULL);

    mpz_sub_ui(op1, p, 1);
    mpz_divexact_ui(op1, op1, 2);
    mpz_powm(op2, z, op1, p);

```



```

mpz_sub_ui(op3, p, 1);
while(mpz_cmp(op2, op3) != 0){
    mpz_add_ui(z, z, 1);
    mpz_powm(op2, z, op1, p);
}

mpz_t c, r, t, m, i, zz, b, e;
mpz_inits(c, r, t, m, i, zz, b, e, NULL);
mpz_powm(c, z, q, p);

mpz_add_ui(op1, q, 1);
mpz_divexact_ui(op1, op1, 2);
mpz_powm(r, n, op1, p);

mpz_powm(t, n, q, p);

mpz_set_ui(m, ss);

while(1){
    mpz_set_ui(i, 0);
    mpz_set(zz, t);
    mpz_set(b, c);

    if(mpz_cmp_ui(t, 1) == 0){
        mpz_set(x1, r);
        mpz_sub(x2, p, x1);

        mpz_clears(c, r, t, m, i, zz, b, e, op1, op2,
                    op3, q, z, NULL);
        return;
    }
}

```

```

mpz_sub_ui(op1, m, 1);
while(mpz_cmp_ui(zz, 1) != 0 && mpz_cmp(i,
    op1) < 0){
    mpz_mul(zz, zz, zz);
    mpz_mod(zz, zz, p);
    mpz_add_ui(i, i, 1);
}

mpz_sub(e, m, i);
mpz_sub_ui(e, e, 1);
while(mpz_sgn(e) > 0){
    mpz_mul(b, b, b);
    mpz_mod(b, b, p);
    mpz_sub_ui(e, e, 1);
}

mpz_mul(r, r, b);
mpz_mod(r, r, p);

mpz_mul(c, b, b);
mpz_mod(c, c, p);

mpz_mul(t, t, c);
mpz_mod(t, t, p);

mpz_set(m, i);
}

```

../c/system.h

```
#pragma once  
#include <stdbool.h>
```

```
typedef struct system {  
    int** m;  
    int* perm;  
    int* sol;  
    bool done;  
    int n1, n2, arb;
```

```
} system_s;
```

```
typedef system_s* system_t;
```

```
system_t init_gauss(int** v, int n1, int n2);  
void gaussian_step(system_t s);  
void free_system(system_t s);
```

../c/system.c

```
#include "system.h"
#include "vector.h"
#include "list_matrix_utils.h"
#include <stdlib.h>
#include <stdio.h>
#include <stdbool.h>

void swap_lines_horz(system_t s, int i, int j){
    int* temp = s->m[i];
    s->m[i] = s->m[j];
    s->m[j] = temp;
}

void swap_lines_vert(system_t s, int i, int j){
    int temp = s->perm[i];
    s->perm[i] = s->perm[j];
    s->perm[j] = temp;

    for(int k = 0; k<s->n1; k++){
        int temp = s->m[k][i];
        s->m[k][i] = s->m[k][j];
        s->m[k][j] = temp;
    }
}

int find_index(system_t s, int from, int look){
    for(int i = from; i < s->n1; i++){
        if(s->m[i][look]){
            return i;
        }
    }
}
```

```
    }
    return -1;
}

system_t transpose(int** v, int n1, int n2){
    system_t s = malloc(sizeof(system_s));

    s->m = malloc(n2*sizeof(int*));
    for(int i = 0; i<n2; i++){
        s->m[i] = malloc(n1*sizeof(int));
        for(int j = 0; j<n1; j++){
            s->m[i][j] = v[j][i];
        }
    }

    s->n1 = n2;
    s->n2 = n1;
    return s;
}

void triangulate(system_t s){
    s->perm = malloc(s->n2*sizeof(int));
    for(int i = 0; i<s->n2; i++){
        s->perm[i] = i;
    }

    int i = 0;
    int j = 0;
    while(i<s->n1 && j<s->n2){
        int k = find_index(s, i, j);
```

```

    if(k != -1){
        if(i != j){
            swap_lines_vert(s, i, j);
        }

        swap_lines_horz(s, i, k);

        for(int l = i + 1; l < s->n1; l++){
            if(s->m[l][i] == 1){
                sub_vect(s->m, l, i, s->n2);
                mod_vect(s->m[l], 2, s->n2);
            }
        }
        i++;
        j = i;
    }
    else{
        j++;
    }
}

}

void get_arbitrary(system_t triangulated){
    for(int i = triangulated->n1-1; i >= 0; i--){
        int j = 0;
        while(j < triangulated->n2 && !triangulated
            ->m[i][j]){
            j++;
        }
        if(j < triangulated->n2){
            triangulated->arb = j+1;
            return;
        }
    }
}

```

```

    }

    fprintf(stderr, "ERROR:-All-vectors-are-zero-in-
        system\n");
    exit(1);
}

void init_sol(system_t s){
    s->sol = malloc(s->n2*sizeof(int));
    for(int i = s->arb; i < s->n2; i++){
        s->sol[i] = 0;
    }
}

void iter_sol(system_t s){
    int i = s->arb;
    while(i < s->n2 && (s->sol[i] == 1)){
        s->sol[i] = 0;
        i++;
    }
    if(i >= s->n2){
        s->done = true;
        return;
    }
    s->sol[i] = 1;
}

system_t init_gauss(int** v, int n1, int n2){
    //printf("Initial vectors\n");
    //print_ll(v, n1, n2);

    system_t s = transpose(v, n1, n2);
    s->done = false;
}

```

```

//printf(" Transposed\n");
//print_ll(s->m, s->n1, s->n2);

for(int i = 0; i<s->n1; i++){
    mod_vect(s->m[i], 2, s->n2);
}

//printf(" Modded\n");
//print_ll(s->m, s->n1, s->n2);

triangulate(s);

//printf(" Triangulated\n");
//print_ll(s->m, s->n1, s->n2);

get_arbitrary(s);
init_sol(s);

return s;
}

void gaussian_step(system_t s){
    iter_sol(s);

    for(int i = s->n1-1; i>=0; i--){
        int j = 0;

```

```

        while(j < s->n2 && !s->m[i][j]){
            j++;
        }

        if(j<s->n2){
            s->sol[j] = 0;

            for(int k = s->n2-1; k>j; k--){
                s->sol[j] -= s->m[i][k] * s->sol[
                    k];
            }
            s->sol[j] = abs(s->sol[j]) % 2;
        }
    }
}

void free_system(system_t s){
    for(int i = 0; i<s->n1; i++){
        free(s->m[i]);
    }
    free(s->m);
    free(s->sol);
    free(s->perm);
    free(s);
}

```

../c/parse_input.h

```
#pragma once
#include <gmp.h>
#include <stdbool.h>

typedef enum {DIXON, QSIEVE, MPQS, PMPQS}
    TYPE;

typedef struct input_s {
    char* output_file;
    int bound, sieving_interval;
```

```
    mpz_t N;
    bool quiet;
    TYPE algorithm;
    int extra;
    int delta;
} input_t;

input_t* parse_input(int argc, char** argv);
void free_input(input_t* input);
```

../c/parse_input.c

```
#include "parse_input.h"
#include <stdlib.h>
#include <string.h>
#include <gmp.h>
#include <stdbool.h>

input_t* init_input(void){
    input_t* input = malloc(sizeof(input_t));
    input->bound = -1;
    input->output_file = NULL;
    input->sieving_interval = -1;
    input->extra = -1;
    input->quiet = false;
    input->algorithm = QSIEVE;
    input->delta = 0;
    mpz_init_set_ui(input->N, 0);
    return input;
}

bool valid_int(char* str){
    int i = 0;
    char c = str[i];
    while(c != '\0'){
        if(c<48 || c>57) return false;
        c = str[++i];
    }
    return true;
}
```

```
void free_input(input_t* input){
    if(input->output_file) free(input->output_file);
    mpz_clear(input->N);
    free(input);
}

input_t* parse_input(int argc, char** argv){
    input_t* input = init_input();

    int i = 1;
    while(i<argc){
        if(strcmp(argv[i], "-b") == 0 || strcmp(argv[
            i], "--bound") == 0){
            i++;
            if(i<argc){
                if(valid_int(argv[i])) input->bound =
                    atoi(argv[i]);
                else return NULL;
            }
            else return NULL;
        }

        else if(strcmp(argv[i], "-s") == 0 || strcmp(
            argv[i], "--sieving_interval") == 0){
            i++;
            if(i<argc){
                if(valid_int(argv[i])) input->
                    sieving_interval = atoi(argv[i]);
                else return NULL;
            }
            else return NULL;
        }
    }
}
```

```

else if(strcmp(argv[i], "--e") == 0 || strcmp(
    argv[i], "--extra") == 0){
    i++;
    if(i < argc){
        if(valid_int(argv[i])) input->extra =
            atoi(argv[i]);
        else return NULL;
    }
    else return NULL;
}

else if(strcmp(argv[i], "-n") == 0 || strcmp(
    argv[i], "--number") == 0){
    i++;
    if(i < argc){
        if(valid_int(argv[i])) mpz_set_str(input
            ->N, argv[i], 10);
        else return NULL;
    }
    else return NULL;
}

else if(strcmp(argv[i], "-d") == 0 || strcmp(
    argv[i], "--delta") == 0){
    i++;
    if(i < argc){
        if(valid_int(argv[i])) input->delta =
            atoi(argv[i]);
        else return NULL;
    }
    else return NULL;
}

else if(strcmp(argv[i], "-o") == 0){
    i++;
}

if(i < argc) input->output_file = argv[i];
else return NULL;

else if(strcmp(argv[i], "-t") == 0 || strcmp(
    argv[i], "--type") == 0){
    i++;
    if(i < argc) {
        if(strcmp(argv[i], "dixon") == 0)
            input->algorithm = DIXON;
        else if(strcmp(argv[i], "qsieve") ==
            0) input->algorithm =
            QSIEVE;
        else if(strcmp(argv[i], "mpqs") == 0)
            input->algorithm = MPQS;
        else if(strcmp(argv[i], "pmpqs") ==
            0) input->algorithm =
            PMPQS;
        else return NULL;
    }
    else return NULL;
}

else if(strcmp(argv[i], "-q") == 0 ||
    strcmp(argv[i], "-stfu") == 0 /*
        easter egg*/ ||
    strcmp(argv[i], "--quiet") == 0){
    input->quiet = true;
}

else return NULL;

i++;
}

```



```
    return input;
}
```

../c/list_matrix_utils.h

#pragma once

void print_list(**int*** l, **int** n);

void print_ll(**int**** ll, **int** n1, **int** n2);

void free_ll(**int**** m, **int** n1);

../c/list_matrix_utils.c

```
#include <stdio.h>
#include <stdlib.h>
```

```
void print_list(int* l, int n){
    for(int i = 0; i<n; i++){
        printf("%d", l[i]);
    }
    printf("\n");
}
```

```
void print_ll(int** ll, int n1, int n2){
    for(int i = 0; i<n1; i++){
```

```
        print_list(ll[i], n2);
    }
    printf("\n");
}
```

```
void free_ll(int** m, int n1){
    for(int i = 0; i<n1; i++){
        free(m[i]);
    }
    free(m);
}
```

../c/factorbase.h

```
#pragma once  
#include <gmp.h>
```

```
// bruh
```

```
bool is_prime(int n);
```

```
// calculates  $\pi(n)$ , the number of prime numbers  $\leq n$ 
```

```
int pi(int n);
```

```
// returns a list of piB first primes
```

```
int* primes(int piB, int B);
```

```
/** Reduces the factor base of the algorithm, refer to:
```

```
 * Quadratic sieve factorisation algorithm
```

```
 * Bc. Ondrej Vladyka
```

```
 * Section 2.3.1 (p.16)
```

```
*/
```

```
int* prime_base(mpz_t n, int* pb_len, int* primes, int  
                piB);
```

../c/factorbase.c

```
#include <stdbool.h>
#include <gmp.h>
#include <stdlib.h>

bool is_prime(int n) {
    // Corner cases
    if (n <= 1)
        return false;
    if (n <= 3)
        return true;

    // This is checked so that we can skip
    // middle five numbers in below loop
    if (n % 2 == 0 || n % 3 == 0)
        return false;

    for (int i = 5; i * i <= n; i = i + 6)
        if (n % i == 0 || n % (i + 2) == 0)
            return false;

    return true;
}

int pi(int n) {
    int k = 0;
    for (int i = 2; i <= n; i++) {
        if (is_prime(i)) k++;
    }
    return k;
}
```

```
int* primes(int piB, int B){
    int* p = malloc(piB*sizeof(int));
    int k = 0;
    for (int i = 2; i <= B; i++) {
        if (is_prime(i)){
            p[k] = i;
            k++;
        }
    }
    return p;
}

/* Used for legendre symbol, exists in gmp already
bool euler_criterion(mpz_t n, int p){
    int e = (p-1)/2;
    mpz_t r, p1;
    mpz_init(r);
    mpz_init_set_ui(p1, p);
    mpz_powm_ui(r, n, e, p1);
    return(mpz_cmp_ui(r, 1) == 0);
}
*/

int* prime_base(mpz_t n, int* pb_len, int* primes, int
piB){
    int* pb = malloc(piB*sizeof(int));
    pb[0] = 2;
```

```

int j = 1;
mpz_t p1;
mpz_init(p1);
for(int i = 1; i < piB; i++){
    mpz_set_ui(p1, primes[i]);
    if(mpz_legendre(n, p1) == 1){
        //printf("%d\n", primes[i]);
        pb[j] = primes[i];
        j++;
    }
}

```

```

    }
}
*pb_len = j;
pb = realloc(pb, (j+1)*sizeof(int)); // +1 used
                                     for mpqs

mpz_clear(p1);
return pb;
}

```

../c/main.c

```
#include <stdbool.h>
#include <gmp.h>
#include <sys/time.h>
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
#include "system.h"
#include "vector.h"
#include "parse_input.h"
#include "factorbase.h"
#include "list_matrix_utils.h"
```

```
// Include algorithms
// Dixon's method
#include "../dixon/dixon.h"
```

```
// The Quadratic Sieve
#include "../qsieve/qsieve.h"
```

```
// Multipolynomial Quadratic Sieve
#include "../mpqs/polynomial.h"
#include "../mpqs/mpqs.h"
#include "../mpqs/parallel_mpqs.h"
```

```
/**
 *
 *
 * START OF ALGORITHM
 *
```

```
*/
```

```
void rebuild_mpqs(mpz_t prod, mpz_t* d, int* v, int*
    primes, int n1, system_t s){
    mpz_set_ui(prod, 1);
    mpz_t temp;
    mpz_init(temp);
    for(int i = 0; i < n1; i++){
        if(s->sol[i]){
            mpz_mul(prod, prod, d[s->perm[i]]);
        }
        mpz_ui_pow_ui(temp, primes[i], v[i]);
        mpz_mul(prod, prod, temp);
    }
    mpz_clear(temp);
}
```

```
void rebuild(mpz_t prod, int* v, int* primes, int n1){
    /** Rebuilds the product of primes to the power
        of half
        * the solution found by the gaussian solve

    * EX:
    *  $v = (1, 2, 3, 1)$ 
    *  $primes = [2, 3, 5, 7]$ 
    *  $prod = 2^{*1} * 3^{*2} * 5^{*3} * 7^{*1}$ 
    * returns prod
    *
    */
}
```

```

    mpz_set_ui(prod, 1);
    mpz_t temp;
    mpz_init(temp);
    for(int i = 0; i < n1; i++){
        mpz_ui_pow_ui(temp, primes[i], v[i]);
        mpz_mul(prod, prod, temp);
    }
    mpz_clear(temp);
}

void sum_lignes(int* sum, int** v, system_t s){
    /** Sums the lines of vectors into 'sum' according
        the solution of the
        * output of the system 's', such that each power
        is even
        */
    for(int i = 0; i < s->n1; i++){
        sum[i] = 0;
    }

    for(int i = 0; i < s->n2; i++){
        if(s->sol[i]){
            add_vect(sum, v[s->perm[i]], s->n1);
        }
    }
}

void factor(input_t* input){
    int piB = pi(input->bound);
    if(!input->quiet) printf(" pi(B)=%d\n", piB);
    int* p = primes(piB, input->bound);

```

```

    int pb_len;
    int* pb;
    switch(input->algorithm){
        case DIXON:
            pb = p;
            pb_len = piB;
            break;
        case QSIEVE:
            pb = prime_base(input->N, &pb_len, p,
                             piB);
            if(!input->quiet) printf(" base-reduction-
                                   %f%%\n", (float)pb_len/piB*100)
                                ;
            free(p);
            break;
        case MPQS:
            pb = prime_base(input->N, &pb_len, p,
                             piB);
            pb[pb_len] = -1;
            if(!input->quiet) printf(" base-reduction-
                                   %f%%\n", (float)pb_len/piB*100)
                                ;
            free(p);
            break;
        case PMPQS:
            pb = prime_base(input->N, &pb_len, p,
                             piB);
            pb[pb_len] = -1;
            if(!input->quiet) printf(" base-reduction-
                                   %f%%\n", (float)pb_len/piB*100)
                                ;
            free(p);
            break;
    }

```



```

}
int target_nb = pb_len + input->extra;

mpz_t* z = malloc((target_nb)*sizeof(mpz_t));
for(int i = 0; i < target_nb; i++){
    mpz_init(z[i]);
}

//Getting zis
int** v;
mpz_t* d;
struct timeval t1, t2;
gettimeofday(&t1, 0);
switch(input->algorithm){
    case DIXON:
        v = dixon(z, input->N, pb_len, pb, input
            ->extra, input->quiet);
        break;
    case QSIEVE:
        v = qsieve(z, input->N, pb_len, pb,
            input->extra, input->
            sieving_interval, input->quiet);
        break;
    case MPQS:
        d = malloc(target_nb*sizeof(mpz_t));
        for(int i = 0; i < target_nb; i++){
            mpz_init(d[i]);
        }
        v = mpqs(z, d, input->N, pb_len, pb,
            input->extra, input->
            sieving_interval, input->delta,
            input->quiet);
        break;

```

```

    case PMPQS:
        d = malloc(target_nb*sizeof(mpz_t));
        for(int i = 0; i < target_nb; i++){
            mpz_init(d[i]);
        }
        v = parallel_mpqs(z, d, input->N, pb_len
            , pb, input->extra, input->
            sieving_interval, input->delta,
            input->quiet);
        break;
}

gettimeofday(&t2, 0);
long seconds = t2.tv_sec - t1.tv_sec;
long microseconds = t2.tv_usec - t1.tv_usec;
double time_spent = seconds + microseconds*1e
    -6;
if(!input->quiet) printf("Time-to-get-zi:-%fs\n",
    time_spent);

mpz_t f, Z1, Z2, test1, test2;
mpz_inits(f, Z1, Z2, test1, test2, NULL);

//gaussian init
system_t s;
int* sum;
switch(input->algorithm){
    case DIXON:
        s = init_gauss(v, target_nb, pb_len);
        sum = malloc(pb_len*sizeof(int));
        break;
    case QSIEVE:
        s = init_gauss(v, target_nb, pb_len);

```

```

        sum = malloc(pb_len*sizeof(int));
        break;
    case MPQS:
        // for -1
        s = init_gauss(v, target_nb, pb_len+1);
        sum = malloc((pb_len+1)*sizeof(int));
        break;
    case PMPQS:
        // for -1
        s = init_gauss(v, target_nb, pb_len+1);
        sum = malloc((pb_len+1)*sizeof(int));
        break;
}
if(!input->quiet) printf("2^%d-solutions-to-iterate
    \n", s->n2 - s->arb);

```

```

bool done = false;
while(!done){
    gaussian_step(s);

```

```

    prod_vect(Z1, z, target_nb, s);
    sum_lignes(sum, v, s);
    div_vect(sum, 2, pb_len);

    switch(input->algorithm){
        case DIXON:
            rebuild(Z2, sum, pb, pb_len);
            break;
        case QSIEVE:
            rebuild(Z2, sum, pb, pb_len);
            break;
        case MPQS:
            rebuild_mpqs(Z2, d, sum, pb, pb_len,

```

```

                s);
            break;
        case PMPQS:
            rebuild_mpqs(Z2, d, sum, pb, pb_len,
                s);
            break;
    }

    // TEST
    mpz_set(test1, Z1);
    mpz_mul(test1, test1, test1);
    mpz_set(test2, Z2);
    mpz_mul(test2, test2, test2);
    assert(mpz_congruent_p(test1, test2, input->
        N) != 0);
    // END TEST

    mpz_sub(f, Z1, Z2);
    mpz_gcd(f, f, input->N);

    if(mpz_cmp_ui(f, 1) != 0 && mpz_cmp(f,
        input->N) != 0){
        assert(mpz_divisible_p(input->N, f));
        if(!input->quiet) gmp_printf("%Zd==0-
            [%Zd]\n", input->N, f);
        done = true;
    }

    mpz_add(f, Z1, Z2);
    mpz_gcd(f, f, input->N);

    if(mpz_cmp_ui(f, 1) != 0 && mpz_cmp(f,
        input->N) != 0){

```

```

    assert(mpz_divisible_p(input->N, f));
    if(!input->quiet) gmp_printf("%Zd==0-
        [%Zd]\n", input->N, f);
    done = true;
}

if(s->done){
    if(!input->quiet) fprintf(stderr, "ERROR
        :-no-solution-for-this-set-of-zi\n");
    exit(1);
}

free(sum);
free(pb);
free_system(s);
free_ll(v, target_nb);
for(int i = 0; i < target_nb; i++){
    mpz_clear(z[i]);
}
free(z);
switch(input->algorithm){
    case DIXON:
        break;
    case QSIEVE:
        break;
    case MPQS:
        for(int i = 0; i < target_nb; i++){
            mpz_clear(d[i]);
            free(d);
            break;
        }
    case PMPQS:
        for(int i = 0; i < target_nb; i++){

```

```

            mpz_clear(d[i]);
            free(d);
            break;
        }

    mpz_clears(f, Z1, Z2, test1, test2, NULL);
}

int main(int argc, char** argv){
    input_t* input = parse_input(argc, argv);
    if(input==NULL){
        fprintf(stderr, "ERROR:-Invalid-input\n");
        return 1;
    }

    if(mpz_cmp_ui(input->N, 0) == 0){
        fprintf(stderr, "ERROR:-No-input-number,-use-
            -n-%%number%%\n");
        return 1;
    }

    if(input->bound == -1) input->bound =
        10000;
    if(input->sieving_interval == -1) input->
        sieving_interval = 100000;
    if(input->extra == -1) input->extra = 1;

    struct timeval t1, t2;
    gettimeofday(&t1, 0);
    factor(input);
    gettimeofday(&t2, 0);
    long seconds = t2.tv_sec - t1.tv_sec;

```

```
long microseconds = t2.tv_usec - t1.tv_usec;
double time_spent = seconds + microseconds*1e
    -6;
if(!input->quiet) printf("Total time: %fs\n",
    time_spent);

free_input(input);

return 0;
}
```

../c/dixon/dixon.h

`#pragma once`

`extra, bool tests);`

`int** dixon(mpz_t* z, mpz_t N, int pb_len, int* pb, int`

../c/dixon/dixon.c

```
#include <gmp.h>
#include <stdbool.h>
#include <stdio.h>
#include <stdlib.h>

bool vectorize_dixon(mpz_t n, int* v, int pb_len, int*
    pb){
    /** Attempts naive factorisation to 'n' with the
        primes in
        * the prime base 'pb' and putting the result into '
            v', vector of powers of
        * the primes in the prime base
        * If it succeeds, returns true, otherwise, returns
            false
    */
    for(int i = 0; i < pb_len; i++){
        v[i] = 0;
    }

    for(int i = 0; i < pb_len && (mpz_cmp_ui(n, 1) !=
        0); i++){
        while (mpz_divisible_ui_p(n, pb[i])){
            v[i]++;
            mpz_divexact_ui(n, n, pb[i]);
        }
    }

    if(mpz_cmp_ui(n, 1) == 0)
        return true;
    return false;
}
```

```
}

int** dixon(mpz_t* z, mpz_t N, int pb_len, int* pb, int
    extra, bool tests){
    /** Gets pb_len+extra b—smooth realtions
        defined at:
        * Quadratic sieve factorisation algorithm
        * Bc. Ondrej Vladyka
        * Definition 1.11 (p.5)
    */

    //ceil(sqrt(n))
    mpz_t sqrt_N;
    mpz_init(sqrt_N);
    mpz_sqrt(sqrt_N, N);
    mpz_add_ui(sqrt_N, sqrt_N, 1);

    mpz_t zi;
    mpz_t zi_cpy;
    mpz_init_set(zi, sqrt_N);
    mpz_init(zi_cpy);

    int** v = malloc((pb_len+extra)*sizeof(int*));

    for(int i = 0; i < pb_len+extra; i++){
        bool found = false;
        int* vi = malloc(pb_len*sizeof(int));

        while(!found){
            mpz_add_ui(zi, zi, 1);
        }
    }
}
```

```

    mpz_mul(zi_cpy, zi, zi);
    mpz_mod(zi_cpy, zi_cpy, N);

    found = vectorize_dixon(zi_cpy, vi, pb_len,
                           pb);
}
if(!tests){
    printf("\r");
    printf("%.1f%%", (float)i/(pb_len+extra
    -1)*100);
    fflush(stdout);
}

```

```

    v[i] = vi;
    mpz_set(z[i], zi);
}
if(!tests) printf("\n");

mpz_clears(sqrt_N, zi, zi_cpy, NULL);

return v;
}

```

../c/qsieve/qsieve.h

```
#pragma once  
#include <gmp.h>  
#include <stdbool.h>
```

```
bool vectorize_qsieve(mpz_t n, int* v, int pb_len, int*
```

```
    pb);  
int** qsieve(mpz_t* z, mpz_t N, int pb_len, int* pb,  
    int extra, int s, bool tests);
```


../c/qsieve/qsieve.c

```
#include <gmp.h>
#include <stdbool.h>
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
#include <math.h>
```

```
#include "../system.h"
#include "../tonellishanks.h"
```

```
bool vectorize_qsieve(mpz_t n, int* v, int pb_len, int*
pb){
    /** Attempts naive factorisation to 'n' with the
    primes in
    * the prime base 'pb' and putting the result into '
    v', vector of powers of
    * the primes in the prime base
    * If it succeeds, returns true, otherwise, returns
    false
    */
    for(int i = 0; i < pb_len; i++){
        v[i] = 0;
    }

    for(int i = 0; i < pb_len && (mpz_cmp_ui(n, 1) !=
0); i++){
        while (mpz_divisible_ui_p(n, pb[i])){
            v[i]++;
            mpz_divexact_ui(n, n, pb[i]);
        }
    }
}
```

```
    }

    if(mpz_cmp_ui(n, 1) == 0)
        return true;
    return false;
}

float* prime_logs(int* pb, int pb_len){
    float* plogs = malloc(pb_len*sizeof(float));

    for(int i = 0; i < pb_len; i++){
        plogs[i] = log2(pb[i]);
    }

    return plogs;
}

int calculate_threshold(mpz_t N, mpz_t sqrt_N, int s,
int loop_number, int* pb, int pb_len){

    mpz_t qstart;
    mpz_init_set_ui(qstart, s);
    mpz_mul_ui(qstart, qstart, loop_number);
    mpz_add(qstart, qstart, sqrt_N);
    mpz_mul(qstart, qstart, qstart);
    mpz_sub(qstart, qstart, N);

    int t = mpz_sizeinbase(qstart, 2) - (int) log2(pb[
pb_len-1]);
    mpz_clear(qstart);
}
```

```

    return t;
}

int** qsieve(mpz_t* z, mpz_t N, int pb_len, int* pb,
    int extra, int s, bool quiet){
    /** Gets pb_len+extra zis that are b-smooth,
        defined at:
        * Quadratic sieve factorisation algorithm
        * Bc. Ondrej Vladyka
        * Definition 1.11 (p.5)
        */

    //ceil(sqrt(n))
    mpz_t sqrt_N;
    mpz_init(sqrt_N);
    mpz_sqrt(sqrt_N, N);
    mpz_add_ui(sqrt_N, sqrt_N, 1);

    mpz_t zi;
    mpz_init_set(zi, sqrt_N);
    mpz_t qx;
    mpz_init(qx);

    int** v = malloc((pb_len+extra)*sizeof(int*));
    for(int i = 0; i < pb_len+extra; i++){
        v[i] = malloc(pb_len*sizeof(int*));
    }
    float* sinterval = malloc(s*sizeof(float));
    float* plogs = prime_logs(pb, pb_len);

    // TESTS
    mpz_t temp;

```

```

    mpz_init(temp);
    // END TESTS

```

```

    int* x1 = malloc(pb_len*sizeof(int));
    int* x2 = malloc(pb_len*sizeof(int));

    // find solution for 2
    mpz_set(temp, sqrt_N);
    mpz_mul(temp, temp, temp);
    mpz_sub(temp, temp, N);
    x1[0] = 0;
    if(mpz_divisible_ui_p(temp, 2) == 0) x1[0] = 1;

```

```

    int sol1, sol2;
    for(int i = 1; i < pb_len; i++){

        tonelli_shanks_ui(N, pb[i], &sol1, &sol2);
        x1[i] = sol1;
        x2[i] = sol2;

        // change solution from  $x = n[p]$  to  $(\text{sqrt}(N) + x) = n[p]$ 
        mpz_set_ui(temp, x1[i]);
        mpz_sub(temp, temp, sqrt_N);
        mpz_mod_ui(temp, temp, pb[i]);

        x1[i] = mpz_get_ui(temp);

        mpz_set_ui(temp, x2[i]);
        mpz_sub(temp, temp, sqrt_N);
        mpz_mod_ui(temp, temp, pb[i]);
    }

```

```

        x2[i] = mpz_get_ui(temp);
    }
    mpz_clear(temp);

    int loop_number = 0;
    int relations_found = 0;
    int tries = 0;
    while(relations_found < pb_len + extra){

        for(int i = 0; i < s; i++){
            sinterval[i] = 0;
        }

        // sieve for 2
        while(x1[0] < s){
            sinterval[x1[0]] += plogs[0];
            x1[0] += pb[0];
        }
        x1[0] = x1[0] - s;

        // sieve other primes
        for(int i = 1; i < pb_len; i++){

            while(x1[i] < s){
                sinterval[x1[i]] += plogs[i];
                x1[i] += pb[i];
            }

            while(x2[i] < s){

                sinterval[x2[i]] += plogs[i];
                x2[i] += pb[i];
            }
        }
    }

```

```

        //next interval
        x1[i] = x1[i] - s;
        x2[i] = x2[i] - s;
    }

    int t = calculate_threshold(N, sqrt_N, s,
                                loop_number, pb, pb_len);
    //printf("t = %d\n", t);

    bool found;
    for(int i = 0; i < s && relations_found <
        pb_len + extra; i++){
        if(sinterval[i] > t){
            tries++;

            // zi = sqrt(n) + x where x = s *
            // loopnumber + i
            mpz_set_ui(zi, s);
            mpz_mul_ui(zi, zi, loop_number);
            mpz_add_ui(zi, zi, i);
            mpz_add(zi, zi, sqrt_N);

            // qx = zi**2 - N
            mpz_mul(qx, zi, zi);
            mpz_sub(qx, qx, N);

            found = vectorize_qsieve(qx, v[
                relations_found], pb_len, pb);

            if(found){
                mpz_set(z[relations_found], zi);
                relations_found++;
            }
        }
    }

```

```

        found = false;
        if(!quiet){
            printf("\r");
            printf(" %.1f%%%-|-%.1f%%%",
                (float)relations_found
                /(pb_len+extra)*100,
                (float)relations_found
                /tries*100);
            fflush(stdout);
        }
    }
}
loop_number++;

```

```

    }

    if(!quiet) printf("\n");

    mpz_clears(sqrt_N, zi, qx, NULL);
    free(x1);
    free(x2);
    free(sinterval);
    free(plogs);

    return v;
}

```

../c/mpqs/common_mpqs.h

```
#pragma once  
#include <gmp.h>  
#include <stdbool.h>
```

```
int calculate_threshold_mpqs(mpz_t sqrt_N, int s, int*  
    pb, int pb_len, int delta);
```

```
float* prime_logs_mpqs(int* pb, int pb_len);  
bool vectorize_mpqs(mpz_t n, int* v, int pb_len, int*  
    pb);  
bool already_added(mpz_t zi, mpz_t* z, int  
    relations_found);
```

../c/mpqs/common_mpqs.c

```
#include <gmp.h>
#include <stdbool.h>
#include <math.h>
#include <stdlib.h>
#include <stdio.h>

int calculate_threshold_mpqs(mpz_t sqrt_N, int s, int*
    pb, int pb_len, int delta){

    mpz_t qstart;
    mpz_init_set_ui(qstart, s);
    mpz_mul(qstart, qstart, sqrt_N);

    int t = mpz_sizeinbase(qstart, 2) - (int) log2(pb[
        pb_len-1]) - delta;
    mpz_clear(qstart);
    return t;
}

float* prime_logs_mpqs(int* pb, int pb_len){
    float* plogs = malloc(pb_len*sizeof(float));

    for(int i = 0; i < pb_len; i++){
        plogs[i] = log2(pb[i]);
    }

    return plogs;
}

bool vectorize_mpqs(mpz_t n, int* v, int pb_len, int*
```

```
    pb){
    /** Attempts naive factorisation to 'n' with the
        primes in
        * the prime base 'pb' and putting the result into '
        v', vector of powers of
        * the primes in the prime base
        * If it succeeds, returns true, otherwise, returns
        false
    */
    for(int i = 0; i < pb_len; i++){
        v[i] = 0;
    }
    if(mpz_sgn(n) < 0){
        v[pb_len] = 1;
        mpz_neg(n, n);
    }
    else{
        v[pb_len] = 0;
    }

    for(int i = 0; i < pb_len && (mpz_cmp_ui(n, 1) !=
        0); i++){
        while (mpz_divisible_ui_p(n, pb[i])){
            v[i]++;
            mpz_divexact_ui(n, n, pb[i]);
        }
    }

    if(mpz_cmp_ui(n, 1) == 0)
        return true;
```

```
    return false;
}

bool already_added(mpz_t zi, mpz_t* z, int
relations_found){
    for(int i = 0; i<relations_found; i++){
        if(mpz_cmp(zi, z[i]) == 0){
```

```
            return true;
        }
    }
    return false;
}
```

../c/mpqs/polynomial.h

```
#pragma once
#include <gmp.h>
#include <stdbool.h>
```

```
struct poly_s {
    mpz_t d;
    mpz_t N;

    mpz_t a;
    mpz_t b;
    mpz_t c;

    mpz_t zi;
    mpz_t qx;
```

```
// used to make operations without declaring and  
freeing everytime
```

```
    mpz_t op1, op2, op3;
```

```
};
```

```
typedef struct poly_s* poly_t;
```

```
void get_next_poly(poly_t p);
```

```
poly_t init_poly(mpz_t N, int M);
```

```
void calc_poly(poly_t p, mpz_t x);
```

```
poly_t copy_poly(poly_t p);
```

```
void free_poly(poly_t p);
```


../c/mpqs/polynomial.c

```
#include "polynomial.h"
#include <gmp.h>
#include <stdlib.h>
#include <assert.h>
#include <stdio.h>

#include "../tonellishanks.h"

void calc_coefficients(poly_t p){
    mpz_mul(p->a, p->d, p->d);

    mpz_t x1, x2;
    mpz_inits(x1, x2, NULL);
    tonelli_shanks_mpz(p->N, p->d, x1, x2);

    // getting ready for congruence solve for raising
    // solution
    mpz_mul_ui(p->op1, x1, 2);

    mpz_mul(p->op2, x1, x1);
    mpz_sub(p->op2, p->op2, p->N);
    mpz_divexact(p->op2, p->op2, p->d);
    mpz_neg(p->op2, p->op2);
    mpz_mod(p->op2, p->op2, p->d);

    mpz_t g, n, m;
    mpz_inits(g, n, m, NULL);
    mpz_gcdext(g, n, m, p->d, p->op1);
    assert(mpz_cmp_ui(g, 1) == 0);
    mpz_mul(p->op1, p->op1, m); // t
```

```
    mpz_clears(g, n, m, NULL);

    mpz_set(p->b, p->d);
    mpz_mul(p->b, p->b, p->op1);
    mpz_add(p->b, p->b, x1);

    mpz_mul(p->op1, p->b, p->b);
    assert(mpz_congruent_p(p->op1, p->N, p->a)
           != 0);

    mpz_sub(p->c, p->op1, p->N);
    mpz_divexact(p->c, p->c, p->a);

    mpz_clears(x1, x2, NULL);
}

void get_next_poly(poly_t p){
    mpz_nextprime(p->d, p->d);
    while(mpz_legendre(p->N, p->d) != 1){
        mpz_nextprime(p->d, p->d);
    }
    calc_coefficients(p);
}

poly_t init_poly(mpz_t N, int M){
    poly_t p = malloc(sizeof(struct poly_s));

    mpz_inits(p->d, p->N, p->a, p->b, p->c,
              p->op1, p->op2, p->op3, p->zi, p
              ->qx, NULL);
```

```

mpz_set(p->N, N);

// choose value of d according to 2.4.2
// sqrt( (sqrt(2N))/M )
mpz_mul_ui(p->op1, N, 2);
mpz_sqrt(p->op1, p->op1);
mpz_div_ui(p->op1, p->op1, M);
mpz_sqrt(p->op1, p->op1);
mpz_prevprime(p->d, p->op1);

// get next prime such that  $(n/p) = 1$ 
while(mpz_legendre(N, p->d) != 1){
    mpz_nextprime(p->d, p->d);
}

calc_coefficients(p);
return p;
}

void calc_poly(poly_t p, mpz_t x){
    mpz_mul(p->zi, p->a, x);
    mpz_add(p->zi, p->zi, p->b);

    mpz_mul(p->qx, x, x);
    mpz_mul(p->qx, p->qx, p->a);

    mpz_mul(p->op1, p->b, x);
    mpz_mul_ui(p->op1, p->op1, 2);
    mpz_add(p->qx, p->qx, p->op1);

```

```

    mpz_add(p->qx, p->qx, p->c);
}

void free_poly(poly_t p){
    mpz_clears(p->d, p->N, p->a, p->b, p->c,
        p->op1, p->op2, p->op3, p->zi, p->qx, NULL);
    free(p);
}

poly_t copy_poly(poly_t p){
    poly_t cpy = malloc(sizeof(struct poly_s));

    mpz_inits(cpy->d, cpy->N, cpy->a, cpy->b,
        cpy->c, cpy->op1, cpy->op2, cpy->op3, cpy->zi, cpy->qx, NULL);

    mpz_set(cpy->d, p->d);
    mpz_set(cpy->N, p->N);

    mpz_set(cpy->a, p->a);
    mpz_set(cpy->b, p->b);
    mpz_set(cpy->c, p->c);

    return cpy;
}

```

../c/mpqs/mpqs.h

```
#pragma once
```

```
#include <gmp.h>
```

```
#include <stdbool.h>
```

```
int** mpqs(mpz_t* z, mpz_t* d, mpz_t N, int pb_len,  
           int* pb, int extra, int s, int delta, bool quiet);
```

../c/mpqs/mpqs.c

```
#include <gmp.h>
#include <stdbool.h>
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
#include <math.h>
#include <time.h>

#include "polynomial.h"
#include "common_mpqs.h"
#include "../system.h"
#include "../tonellishanks.h"

int** mpqs(mpz_t* z, mpz_t* d, mpz_t N, int pb_len,
    int* pb, int extra, int s, int delta, bool quiet){
    /** Gets pb_len+extra zis that are b-smooth,
        defined at:
        * Quadratic sieve factorisation algorithm
        * Bc. Ondrej Vladyka
        * Definition 1.11 (p.5)
        */

    //ceil(sqrt(n))
    mpz_t sqrt_N;
    mpz_init(sqrt_N);
    mpz_sqrt(sqrt_N, N);
    mpz_add_ui(sqrt_N, sqrt_N, 1);

    mpz_t x;
    mpz_init(x);
```

```
poly_t Q = init_poly(N, s);

int** v = malloc((pb_len+extra)*sizeof(int*));
for(int i = 0; i < pb_len+extra; i++){
    v[i] = malloc((pb_len+1)*sizeof(int*)); // +1
    for -1
}
float* sinterval = malloc(2*s*sizeof(float));
float* plogs = prime_logs_mpqs(pb, pb_len);
int t = calculate_threshold_mpqs(sqrt_N, s, pb,
    pb_len, delta);

// TESTS
mpz_t temp;
mpz_init(temp);
// END TESTS

int* r = malloc(pb_len*sizeof(int));
int* x1 = malloc(pb_len*sizeof(int));
int* x2 = malloc(pb_len*sizeof(int));

int sol1, sol2;
for(int i = 1; i < pb_len; i++){
    tonelli_shanks_ui(N, pb[i], &sol1, &sol2);
    r[i] = sol1;
}

mpz_t g, m, n, pi;
```

```

mpz_inits(g, m, n, pi, NULL);

int relations_found = 0;
clock_t start;
start = clock();
int tries = 0;
while(relations_found < pb_len + extra){

    // for 2
    mpz_set_ui(temp, 0);
    calc_poly(Q, temp);
    x1[0] = 0;
    if(mpz_divisible_ui_p(Q->qx, 2) == 0) x1[0]
        = 1;

    //others
    for(int i = 1; i < pb_len; i++){
        mpz_set_ui(pi, pb[i]);
        mpz_gcdext(g, m, n, Q->a, pi);
        if(mpz_cmp_ui(g, 1) != 0){
            fprintf(stderr, "ERROR:-Number-is-too
                ~small-for-the-current-
                implementation-of-MPQS\n");
            exit(1);
        }

        mpz_set_ui(temp, r[i]);
        mpz_sub(temp, temp, Q->b);
        mpz_mul(temp, temp, m);
        mpz_mod(temp, temp, pi);

        x1[i] = mpz_get_ui(temp);

```

```

//calc_poly(Q, temp);
//assert(mpz_divisible_ui_p(Q->qx, pb[i])
    != 0);

```

```

mpz_set_ui(temp, pb[i]);
mpz_sub_ui(temp, temp, r[i]);
mpz_sub(temp, temp, Q->b);
mpz_mul(temp, temp, m);
mpz_mod(temp, temp, pi);

```

```

x2[i] = mpz_get_ui(temp);

```

```

//calc_poly(Q, temp);
//assert(mpz_divisible_ui_p(Q->qx, pb[i])
    != 0);

```

```

//realign sieving interval to [-s, s]

```

```

int k = (x1[i] + s)/pb[i];
x1[i] -= k * pb[i];
x1[i] += s;

```

```

k = (x2[i] + s)/pb[i];
x2[i] -= k * pb[i];
x2[i] += s;

```

```

//mpz_set_si(temp, -s);
//mpz_add_ui(temp, temp, x1[i]);
//calc_poly(Q, temp);
//assert(mpz_divisible_ui_p(Q->qx, pb[i])
    != 0);

```

```

}

```

```

for(int i = 0; i < 2*s; i++){
    interval[i] = 0;
}

/*
// sieve for 2
while(x1[0] < 2*s){
    interval[x1[0]] += plogs[0];
    x1[0] += pb[0];
}
*/

// sieve other primes
for(int i = 30; i < pb_len; i++){
    while(x1[i] < 2*s){
        interval[x1[i]] += plogs[i];
        x1[i] += pb[i];
    }

    while(x2[i] < 2*s){
        interval[x2[i]] += plogs[i];
        x2[i] += pb[i];
    }
}

bool found;
bool update_time = false;
for(int i = 0; i < 2*s && relations_found < pb_len + extra; i++){
    if(interval[i] > t){
        tries++;
    }
}

```

```

mpz_set_si(x, -s);
mpz_add_ui(x, x, i);
calc_poly(Q, x);

if(!already_added(Q->zi, z,
    relations_found)){
    found = vectorize_mpqqs(Q->qx,
        v[relations_found], pb_len,
        pb);
    if(found){
        mpz_set(z[relations_found], Q
            ->zi);
        mpz_set(d[relations_found], Q
            ->d);
        relations_found++;
        update_time = true;
        found = false;
    }
    if(!quiet){
        printf("\r");
        printf("%.1f%%-|-%.1f
            %%%", (float)
            relations_found/(
            pb_len+extra)
            *100, (float)
            relations_found/
            tries*100);
        fflush(stdout);
    }
}

```

```

    if(update_time && !quiet) printf("(~%.0fs-left
        )-----", (double)(clock() - start)/
        CLOCKS_PER_SEC/relations_found*((
        pb_len+extra - relations_found)));
    get_next_poly(Q);
}

if(!quiet) printf("\n");
mpz_clears(sqrt_N, temp, g, m, n, pi, x, NULL);
free(x1);

    free(x2);
    free(r);
    free(sinterval);
    free(plogs);
    free_poly(Q);

    return v;
}

```

../c/mpqs/parallel_mpqs.h

```
#pragma once
#include <gmp.h>
#include "polynomial.h"
#include <sys/time.h>
#include <stdint.h>
```

```
struct sieve_arg_s {
    // used for sieving
    int* pb;
    int pb_len;
    int extra;
    int* r;
    float* plogs;
    int s;
    int t;
    int* relations_found;
    int** v;
    bool quiet;
    mpz_t* z;
    mpz_t* d;
};
```

```
poly_t Qinit;
```

```
// used to print progress and predicted time left
struct timeval begin;
uint_fast64_t* tries;
```

```
// used to constantly have a certain number of
    threads running
```

```
int thread_id;
bool* threads_running;
```

```
};
```

```
typedef struct sieve_arg_s sieve_arg_t;
```

```
bool already_added(mpz_t zi, mpz_t* z, int
    relations_found);
```

```
void* sieve_100_polys (void* args);
```

```
int** parallel_mpqs(mpz_t* z, mpz_t* d, mpz_t N, int
    pb_len, int* pb, int extra, int s, int delta, bool
    quiet);
```


../c/mpqs/parallel_mpqs.c

```
#include <gmp.h>
#include <stdbool.h>
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
#include <math.h>
#include <time.h>
#include <pthread.h>
#include <sys/time.h>
```

```
#include "polynomial.h"
#include "common_mpqs.h"
#include "parallel_mpqs.h"
#include "../system.h"
#include "../tonellishanks.h"
```

```
pthread_mutex_t mutex;
```

```
void* sieve_100_polys (void* args){
    sieve_arg_t* arg = (sieve_arg_t*) args;

    poly_t Q = copy_poly(arg->Qinit);

    mpz_t temp, g, m, n, pi, x;
    mpz_inits(temp, g, m, n, pi, x, NULL);
    float* sinterval = malloc(2*arg->s*sizeof(float))
    ;
    int* x1 = malloc(arg->pb.len*sizeof(int));
    int* x2 = malloc(arg->pb.len*sizeof(int));
```

```
for(int i = 0; i<100 && *(arg->relations_found)
    < arg->pb.len + arg->extra; i++){
    get_next_poly(Q);
```

```
//get sol for 2
mpz_set_ui(temp, 0);
calc_poly(Q, temp);
x1[0] = 0;
if(mpz_divisible_ui_p(Q->qx, 2) == 0) x1[0]
    = 1;
```

```
//get sol for others
for(int i = 1; i<arg->pb.len; i++){
    mpz_set_ui(pi, arg->pb[i]);
    mpz_gcdext(g, m, n, Q->a, pi);
    if(mpz_cmp_ui(g, 1) != 0){
        fprintf(stderr, "ERROR:-Number is too
            ~small for the current
            implementation of MPQS\n");
        exit(1);
    }
```

```
    mpz_set_ui(temp, arg->r[i]);
    mpz_sub(temp, temp, Q->b);
    mpz_mul(temp, temp, m);
    mpz_mod(temp, temp, pi);
```

```
x1[i] = mpz_get_ui(temp);
```

```

//calc_poly(Q, temp);
//assert(mpz_divisible_ui_p(Q->qx, arg
    ->pb[i]) != 0);

mpz_set_ui(temp, arg->pb[i]);
mpz_sub_ui(temp, temp, arg->r[i]);
mpz_sub(temp, temp, Q->b);
mpz_mul(temp, temp, m);
mpz_mod(temp, temp, pi);

x2[i] = mpz_get_ui(temp);

//calc_poly(Q, temp);
//assert(mpz_divisible_ui_p(Q->qx, arg
    ->pb[i]) != 0);

//realign sieving interval to [-s, s]
int k = (x1[i] + arg->s)/arg->pb[i];
x1[i] -= k * arg->pb[i];
x1[i] += arg->s;

k = (x2[i] + arg->s)/arg->pb[i];
x2[i] -= k * arg->pb[i];
x2[i] += arg->s;

//mpz_set_si(temp, -arg->s);
//mpz_add_ui(temp, temp, x1[i]);
//calc_poly(Q, temp);
//assert(mpz_divisible_ui_p(Q->qx, arg
    ->pb[i]) != 0);
}

//reset sieving_interval

```

```

for(int i = 0; i < 2*arg->s; i++){
    sinterval[i] = 0;
}

/*
// sieve for 2
while(x1[0] < 2*arg->s){
    sinterval[x1[0]] += arg->plogs[0];
    x1[0] += arg->pb[0];
}
*/

// sieve other primes
for(int i = 30; i < arg->pb.len; i++){
    while(x1[i] < 2*arg->s){
        sinterval[x1[i]] += arg->plogs[i];
        x1[i] += arg->pb[i];
    }
    while(x2[i] < 2*arg->s){
        sinterval[x2[i]] += arg->plogs[i];
        x2[i] += arg->pb[i];
    }
}

bool found;
bool update_time = false;
pthread_mutex_lock(&mutex);
for(int i = 0; i < 2*arg->s && *(arg->
    relations_found) < arg->pb.len + arg
    ->extra; i++){
    if(sinterval[i] > arg->t){
        *(arg->tries) += 1;
        mpz_set_si(x, -arg->s);
    }
}

```

```

mpz_add_ui(x, x, i);
calc_poly(Q, x);

if(!already_added(Q->zi, arg->z,
    *(arg->relations_found))) {
    found = vectorize_mpps(Q->qx,
        arg->v[*(arg->
            relations_found)], arg->
            pb_len, arg->pb);
    if(found) {
        mpz_set(arg->z[*(arg->
            relations_found)], Q
            ->zi);
        mpz_set(arg->d[*(arg->
            relations_found)], Q
            ->d);
        *(arg->relations_found) +=
            1;
        found = false;
        update_time = true;
        if(!arg->quiet) {
            printf("\r");
            printf("%.1f%%-|-%.1f
                %%%", (float)(*(
                    arg->
                    relations_found))
                /(arg->pb_len+
                    arg->extra)
                *100, (float)(*(
                    arg->
                    relations_found))
                /(arg->tries))
                *100);
        }

        fflush(stdout);
    }
}

struct timeval current;
gettimeofday(&current, 0);
long seconds = current.tv_sec - arg->begin.
    tv_sec;
long microseconds = current.tv_usec - arg
    ->begin.tv_usec;
double elapsed = seconds + microseconds*1e
    -6;
if(update_time && !arg->quiet) printf("-
    (~%.0fs-left)~~~~~", elapsed/(arg->
    relations_found)*(arg->pb_len+arg
    ->extra - (arg->relations_found))
    );
pthread_mutex_unlock(&mutex);
}

mpz_clears(temp, g, m, n, pi, x, NULL);
free(x1);
free(x2);
free(sinterval);
free_poly(Q);

arg->threads_running[arg->thread_id] = false;
return NULL;

```



```

        threads_running
    };
    threads_running[i] = true;
    pthread_create(&threads[i], NULL,
        sieve_100_polys, args+i);
}
for(int i = 0; i < 100; i++){
    get_next_poly(Q);
}
}
if(!quiet) printf("\n");
for(int i = 0; i < 8; i++){

```

```

        pthread_join(threads[i], NULL);
    }

    free(threads);
    free(args);
    free(r);
    free(plogs);
    free(threads_running);
    free_poly(Q);
    mpz_clear(sqrt_N);

    return v;
}

```