



Cairo University

ANALYSIS AND SIMULATION OF THE TARGET-ATTACKER AND THE TARGET-ATTACKER-DEFENDER PROBLEMS

By

Mostafa Ali Rushdi

A Thesis Submitted to the
Faculty of Engineering at Cairo University
in Partial Fulfillment of the
Requirements for the Degree of
MASTER OF SCIENCE
in
Aerospace Engineering

AEROSPACE ENGINEERING DEPARTMENT
FACULTY OF ENGINEERING, CAIRO UNIVERSITY
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Under the Supervision of

Prof. Ayman H. Kassem

Prof. Gamal El-Bayoumi

.....
Professor

Aerospace Engineering Department
Faculty of Engineering, Cairo University

.....
Professor

Aerospace Engineering Department
Faculty of Engineering, Cairo University

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Approved by the Examining Committee:

Prof. Ayman H. Kassem, Thesis Main Advisor

Prof. Third E. Name, Member

Prof. Second S. Name, Internal Examiner

Prof. First S. Name, External Examiner

**AEROSPACE ENGINEERING DEPARTMENT
FACULTY OF ENGINEERING, CAIRO UNIVERSITY
GIZA, EGYPT
2017**

Engineer's Name: Mostafa Ali Rushdi
Date of Birth: 26/10/1991
Nationality: Egyptian
E-mail: morushdi@gmail.com
Phone: 0111-4000-606
Address: Postal Address
Registration Date: 10/10/2014
Awarding Date: 14/7/2015
Degree: Master of Science
Department: Aerospace Engineering



Supervisors:
Prof. Ayman H. Kassem
Prof. Gamal El-Bayoumi

Examiners:
Prof. First S. Name (External examiner)
Prof. Second S. Name (Internal examiner)
Prof. Ayman H. Kassem (Thesis main advisor)
Prof. Third E. Name (Member)

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Key Words:

Pursuit-Evasion, Target, Attacker, Defender, Apollonius circle, critical speed ratio, Escape region, Voronoi diagram, Optimal strategy.

Summary:

This thesis studies the target-attacker-defender problem, and introduces a unified analytical overview with numerical simulations.

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Abstract

This thesis deals with two scenarios of active target defense. The first scenario is a two-agent pursuit-evasion problem that involves a Target (aircraft) in opposition to an Attacker (missile). The Target tries to evade (avoid being captured) the Attacker. This first problem will be referred to herein as the TA problem as it concerns the Target (T) and Attacker(A). The second scenario is a three-agent pursuit-evasion problem that involves three agents, the Target, the Attacker and the Defender. The Attacker missile pursues a Target aircraft that is being helped by a Defender missile which tries to intercept the Attacker before it reaches the Target. This second problem will be referred to herein as the TAD problem as it concerns the Target (T), the Attacker(A), and the Defender (D). We first consider the TA problem, in which we search for a path that the Target can move on to escape from the Attacker. All the evasion techniques depend on the time of the turn that the Target makes when it detects the Attacker (Missile) and the objective is to maximize the Missile acceleration till the Missile power bleed. We choose the escaping trajectory as a polynomial with unknown coefficients, then decide the values of these coefficients so as to make the Missile exert a maximum acceleration to bleed its power as fast as possible before it reaches the Target. We explain the meaning of proportional navigation, and subsequently simulate two-dimensional proportional-navigation equations using MATLAB and Simulink. After dealing with the TA problem, the thesis shifts to the TAD problem. In this latter problem, a differential game arises in which a team is formed by the Target and the Defender which cooperate to maximize the separation between the Target and the point where the Defender intercepts the Attacker, while the Attacker tries to minimize this separation. This thesis offers a unified analytic treatment of the aforementioned problem based on the construction of two Apollonius circles. The treatment includes all possibilities of the ratio between the speeds of the Attacker and Defender. A criticality condition is derived from which two important entities are obtained, namely: (a) the critical Target speed normalized *w.r.t.* the Attacker speed, and (b) the Voronoi diagram bordering the safe or escape region for the Target Optimal strategies. This Voronoi diagram is shown to obey a complex sixth-degree polynomial when the Defender differs in speed from the Attacker. This polynomial reduces to a real fourth-degree polynomial when the Defender and Attacker are similar, i.e., are of equal speeds. Beside unifying previously published results in a common setting, this thesis simplifies all computations by using intuitionistic plane-geometric arguments rather than the more tedious analytic-geometric manipulations. Moreover, the thesis extends existing results by adding some novel results, thereby giving a complete picture of all cases of interest. The analysis in this thesis is supplemented by extensive computations using MATLAB to solve the complex high-order polynomial equations and to plot the Voronoi diagrams under a variety of pertinent parameters. The numerical results and plots obtained allow useful and insightful interpretation and are in exact agreement with numerical solution of the corresponding two-point boundary value problem.

Part I

INTRODUCTION AND LITERATURE REVIEW

Chapter 1: Introduction and Literature Review

This thesis deals with two scenarios of active target defense involving:

1. Two-agent pursuit-evasion: a Target (aircraft) in opposition to an Attacker (missile). The Target tries to evade (avoid being captured) the Attacker. This problem will be referred to herein as the *TA* problem as it concerns the Target (*T*) and Attacker (*A*).
2. Three-agent pursuit-evasion where each agent has a specific role. A two-agent team consists of a Target (aircraft) and a Defender (missile) cooperating in opposition to an Attacker (missile). The Target tries to evade the Attacker and avoid being captured by him. The Defender cooperates with and assists the Target by trying to intercept(capture and destroy) the Attacker before the latter captures the Target. This problem will be referred to herein as the *TAD* problem as it concerns a triad consisting of the Target (*T*), Attacker (*A*), and Defender (*D*).

The following two sections pertain to a brief introduction and literature survey of the *TA* and *TAD* problem.

1.1 The *TA* problem

1.2 The *TAD* problem

The *TAD* problem constitutes a dynamic differential game [7, 8, 9, 10, 11, 12, 13, 14, 15] and is of interest in aerospace, control, and robotics engineering. We will consider herein recent formulation and treatments of this problem [1, 2, 3, 4, 5, 6], though there exist other formulations and treatments of it that span almost half a century [16, 17, 18, 19, 20, 21, 22, 23, 24, 25]. Table 1.1 shows several variants of this generic problem in a variety of settings, contexts or disciplines. The *TAD* problem is a generalization of the classical problem of a single pursuer and a single evader [26, 27, 28, 29, 30]. In fact, the *TAD* is essentially a duplication of this classical problem, since in the *TAD* problem, the Attacker plays the double role of being a pursuer for the Target, and at the same time an evader for the Defender [19]. The *TAD* problem is also a special case of a more general pursuit-evasion problem in which there are multiple Attackers and multiple Defenders [31, 32, 33, 34, 35]. Many common threads are shared by all these problems, such as the solution of differential games [7, 8, 9, 10, 11, 12, 13, 14, 15] and the construction of Apollonius circles [36, 37, 38, 39] and Voroni diagrams [40, 41, 42, 43, 44, 45, 46, 45].

1.3 Thesis Contribution and Outline

As stated earlier, this thesis has contributions to both *TA* and *TAD* problems.

In the *TA* problem, we search for a path that the Target can move on to escape from the Attacker. All the evasion techniques depend on the time of the turn that the Target makes when it detects the Attacker (Missile) and the objective is to maximize the Missile

Area	Agent 1	Agent 2	Agent 3	References
Aerospace	Target	Attacker(missile)	Defender(missiles)	[1, 2, 3, 4, 5]
Biology	Prey	Predator	Protector	[20, 25]
Society	Lady	Bandits	Bodyguards	[18]
Criminology	Robber	Policemen/Cops	Gangsters	[42]

Table 1.1: A variety of settings or contexts for the same generic problem

acceleration till the Missile power bleed. We choose the escaping trajectory as a polynomial with unknown coefficients, then decide the values of these coefficients so as to make the Missile exert a maximum acceleration to bleed its power as fast as possible before it reaches the Target. We explain the meaning of proportional navigation, and subsequently simulate two-dimensional proportional-navigation equations using MATLAB and Simulink.

This thesis also offers a unified analytic treatment of the *TAD* problem based on the construction of two Apollonius circles. The treatment includes all possibilities of the ratio between the speeds of the Attacker and Defender. Note that the case of a slow Defender appears here for the first time, while the treatment of the cases of fast Defender or similar Defender is extended and augmented with novel results and new insights. A criticality condition is derived from which the two following important entities are obtained

- the critical Target speed normalized w.r.t. the Attacker's speed so as to be a dimensionless quantity,
- the Voronoi diagram bordering the safe or escape region for the Target.

Optimal strategies are also studied, and are shown to obey a complex sixth-degree polynomial when the Defender differs in speed from the Attacker. This polynomial reduces to a real fourth-degree polynomial when the Defender and Attacker are similar in speed. Beside unifying previously published results in a common setting, this thesis simplifies all computations by using intuitionistic plane-geometric arguments rather than the more tedious analytic-geometric manipulations. Moreover, the thesis extends existing results by adding some novel results, thereby giving a complete picture of all cases of interest. The analysis in this thesis is supplemented by extensive computations using MATLAB to solve the complex high-order polynomial equations and to plot the Voronoi diagrams under a variety of pertinent parameters. The numerical results and plots obtained allow useful and insightful interpretations and are in exact agreement with numerical solution of the corresponding two-point boundary value problem.

This thesis consists of four parts. Part I comprises the current chapter of the introduction and literature survey. Part II pertains to the *TA* problem and consists of chapters 2-5. Chapter 2 is an introduction of the *TA* problem, while chapter 3 is devoted to a brief discussion of guidance laws. Chapter 4 constitutes a prominent cornerstone of the thesis as it reports extensive simulation results for the *TA* problem obtained via Matlab and simulink. Chapter 5 discusses a game-theoretic formulation of the *TA* problem. It is followed by Part III of the thesis on the *TAD* problem, which consists of chapters 6-10. Chapter 6 introduces the *TAD* problem, lists the assumptions, notation and nomenclature used therein, discusses the general concepts and properties of an Apollonius circle [36, 37, 38, 39], and then specializes these to the cases of the *AD* Apollonius circle and the *TA* Apollonius circle. Chapter 7

derives a criticality condition which constitutes a quadratic equation that is later solved for the critical dimensionless ratio between the Target speed and Attacker speed. The same criticality condition is then rephrased in Chapter 8 as a relation between the initial coordinates x_T and y_T of the Target. This relation leads to the Voronoi diagram bordering the *safe* or *escape* region R_e of the Target. This Voronoi diagram is published for the first time when the Defender is fast or slow. It includes as a special case the much simpler diagram for a similar Defender that has already appeared in [3]. Chapter 9 develops optimal strategies for the three constituents of the differential game. Chapter 10 explores the simulation of the optimal heading angle. The thesis is concluded with its last part (Part IV) which is a single (chapter 11) that summarizes the thesis contributions and points out possible directions for further work.

Part II

TARGET-ATTACKER PROBLEM

Chapter 2: Introduction for TA problem

2.1 Problem Statement

In this Part, we consider a two-agent pursuit-evasion problem that involves a Target (aircraft) in opposition to an Attacker (missile). The Target tries to evade the Attacker and avoid being captured by him. We will try to find a simple technique for target evasion, so we should have a look at evasion techniques. The missile is on the ground, its position is (0,0). The plane is flying in the air, its position is (10000,40000).

2.2 Evasion techniques

There are some observations we should make before talking about evasion techniques

- The missile is faster than the aircraft, but cannot turn tighter than the aircraft, so it takes a longer path.
- The control mechanism of the Missile is much simpler and is of less capability than that of the aircraft. In order to pull as tight turn as the aircraft, it must exercise an acceleration that is far beyond its capability.
- The missile always attempts to trace the target. Thus if the target changes heading, it will be necessary for the Missile to change heading similarly, but this is too difficult for him to achieve.
- The main problem with evading missiles is their speed, which makes timing somewhat difficult.

Some tactics to evade a missile:

1. If missile is fired head-on at BVR range:

- Turn hard to either left or right so as to fly at roughly 90 degrees angle to attacking aircraft (This forces missile to bleed off the energy and to lead the target).
- Once target aircraft makes a hard turn to reverse a direction, missile with its far larger turn circle will be unable to compensate.



Figure 2.1: Illustration of the first maneuver technique

2. Jinking (useful at short ranges):

- Aircraft must be positioned so that it is at angle (30-60 degrees is optimum) relative to missiles flight path.
- Once missile gets closer, aircraft will make a hard turn in opposite direction.
- As there is a lag between aircraft changing the direction and missile following (for several reasons, most important of which is missiles inertia), this will cause missile to head in wrong direction until it manages to correct, and also to bleed off the energy.
- Missile will fly past the aircraft and miss.

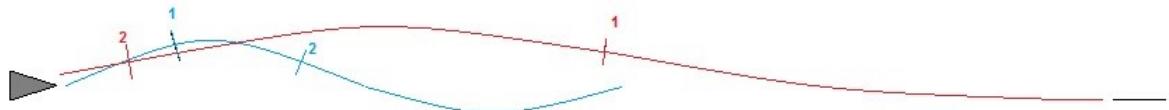


Figure 2.2: Illustration of the second maneuver technique

3. Climb (useful at longer ranges):

- Since at long range missile will have burned out its engine, it will rely on inertia to keep it flying, and climbing will mean that it will bleed off energy rapidly.
- Once missile reaches a close range (maybe around 1,500 meters), dive for the ground, then pull up (This will allow pilot to gain energy and using it to evade the missile).

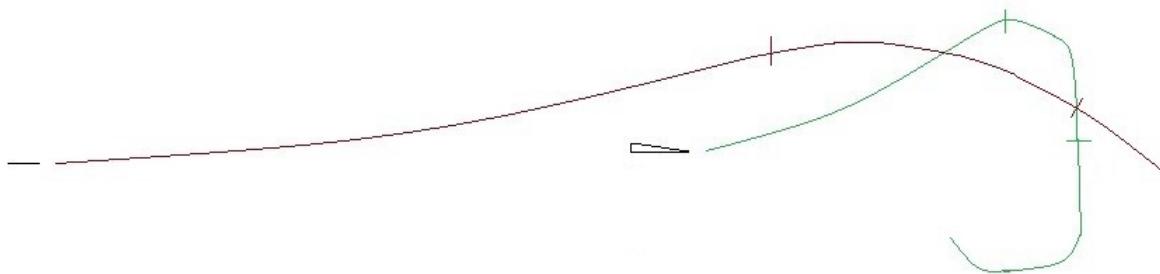


Figure 2.3: Illustration of the third maneuver technique

4. Fourth tactic:

- place the missile at 3 o'clock or 9 o'clock position
- maintain sufficient turn to keep the missile
- This tactic forces the missile to execute a continuous turn, bleeding the energy entire time, making it easier to outturn the missile once it comes close.

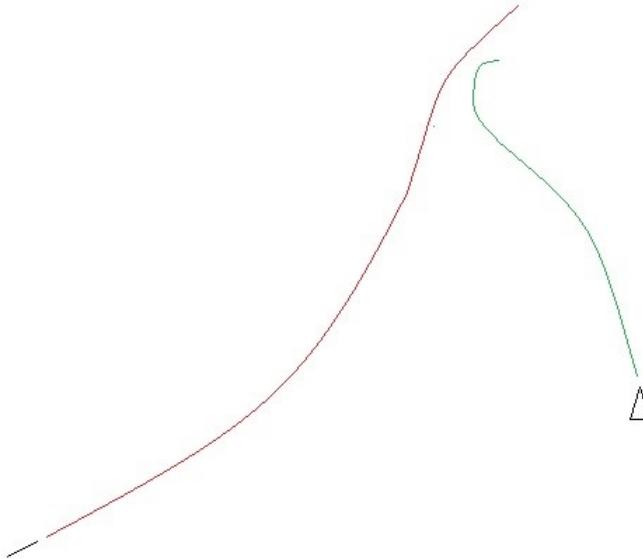


Figure 2.4: Illustration of the forth maneuver technique

Now we want to search for a path that the Target can move on it to escape from the Attacker. All the evasion techniques depend on the time of the turn that the Target makes when it detects the Attacker (Missile) and the objective is to maximize the Missile acceleration till the Missile power bleed. We will choose the escaping trajectory as a polynomial with unknown coefficients, then try to find these coefficients which make the Missile exert a maximum acceleration to bleed its power as fast as we can before it reaches the Target.

2.3 Assumptions, Notation, and Nomenclature

Assumptions

1. Both the Attacker and Target travel at constant speeds.
2. The Attacker's speed is larger than that of the Target. Otherwise, the Target will be unconditionally or trivially capable of escaping.
3. Gravitational and drag effects are neglected for simplicity.
4. The maximum g-force a typical person can handle is 5G

5. The maximum g-force a pilot fighter (with safety equipments) could resist is 10G.
6. The maximum g-force for a sprint missile is 100G.
7. The range for Tactical missile defense, which has short ranges is $18 \rightarrow 20$ Kilometer, approximately $59000 \rightarrow 66000$ feet.

Notation

- n_c : Acceleration command (for the Missile) in m/s^2 .
- N' : Effective navigation ratio, a unit-less designer-chosen gain (usually in the range of $3 \rightarrow 5$).
- V_c : Missile-Target closing velocity.
- λ : line-of-sight angle.
- $\dot{\lambda}$: line of sight rate.
- R_{TM} : length of the line of sight.
- L : Missile lead angle.
- HE : Heading error.
- $\dot{\beta}$: angular velocity of the Target.
- V_{T1}, V_{T2} : Target velocity components in the Earth fixed coordinate system.
- V_{M1}, V_{M2} : Missile velocity components in the Earth fixed coordinate system.

Nomenclature

Inertial coordinate system: fixed to the surface of a flat-Earth model (the 1 axis is down-range and the 2 axis can either be altitude or cross-range).

Missile lead angle: theoretically correct angle for the missile to be on a collision triangle with the Target.

Heading error (HE) : angle representing the initial deviation of the Missile from the collision triangle.

line of sight: The imaginary line connecting the Missile and Target.

length of the line of sight (R_{TM}): Instantaneous separation between Missile and Target.

Miss distance : The point of closest approach of the Missile and Target.

Closing velocity (V_c): the negative rate of change of the distance from the Missile to the Target $V_c = -\dot{R}_{TM} = -\frac{d}{dt}R_{TM}$.

Chapter 3: Guidance Laws

In this chapter, we will illustrate some types of guidance laws. We will be in the side of the Attacker, trying to hit the target. There are three main types of guidance laws:

3.1 Proportional Navigation

In this section, we will illustrate some fundamentals of missile guidance, focusing on the proportional navigation technique, which is one of the simplest guidance laws.

What is proportional navigation?

The proportional navigation guidance law issues acceleration commands, perpendicular to the instantaneous missile-target line-of-sight, which are proportional to the line-of-sight rate and closing velocity. Mathematically, the guidance law can be stated as

$$n_c = N' V_c \dot{\lambda} \quad (3.1)$$

where n_c is the acceleration command (for the missile) in (m/s^2) , N' is the effective navigation ratio, a unit-less designer-chosen gain (usually in the range of $3 \rightarrow 5$), V_c is the missile-target closing velocity in (m/s) and $\dot{\lambda} = \frac{d\lambda}{dt}$ is the rate of the line-of-sight angle and is in (rad/s) . Note that (3.1) is dimensionally homogeneous, since n_c has the dimension

$$\begin{aligned} [n_c] &= [N'][V_c][\dot{\lambda}] \\ &= (1)(LT^{-1})(T^{-1}) \\ &= LT^{-2} \end{aligned} \quad (3.2)$$

which is the appropriate dimension of linear acceleration.

3.1.1 Simulation of proportional navigation equations in 2-D

In this subsection we will introduce the equations of proportional navigation and the sequence to get a simulation for the path of the Target and the Attacker and how the missile acceleration will be affected during this simulation.

The simulation inputs are the initial location of the Missile (R_{M1}, R_{M2}) and the initial location of the Target (R_{T1}, R_{T2}), Target speed V_T , Missile speed V_M , and effective navigation ratio N' .

There are two types of **error sources** that cause the Attacker to miss the Target; they are the heading error (HE) and the acceleration of the Target (n_T) called the Target maneuver.

proportional navigation differential equations

The components of Target velocity

$$V_{T1} = -V_T \cos(\beta) \quad (3.3)$$

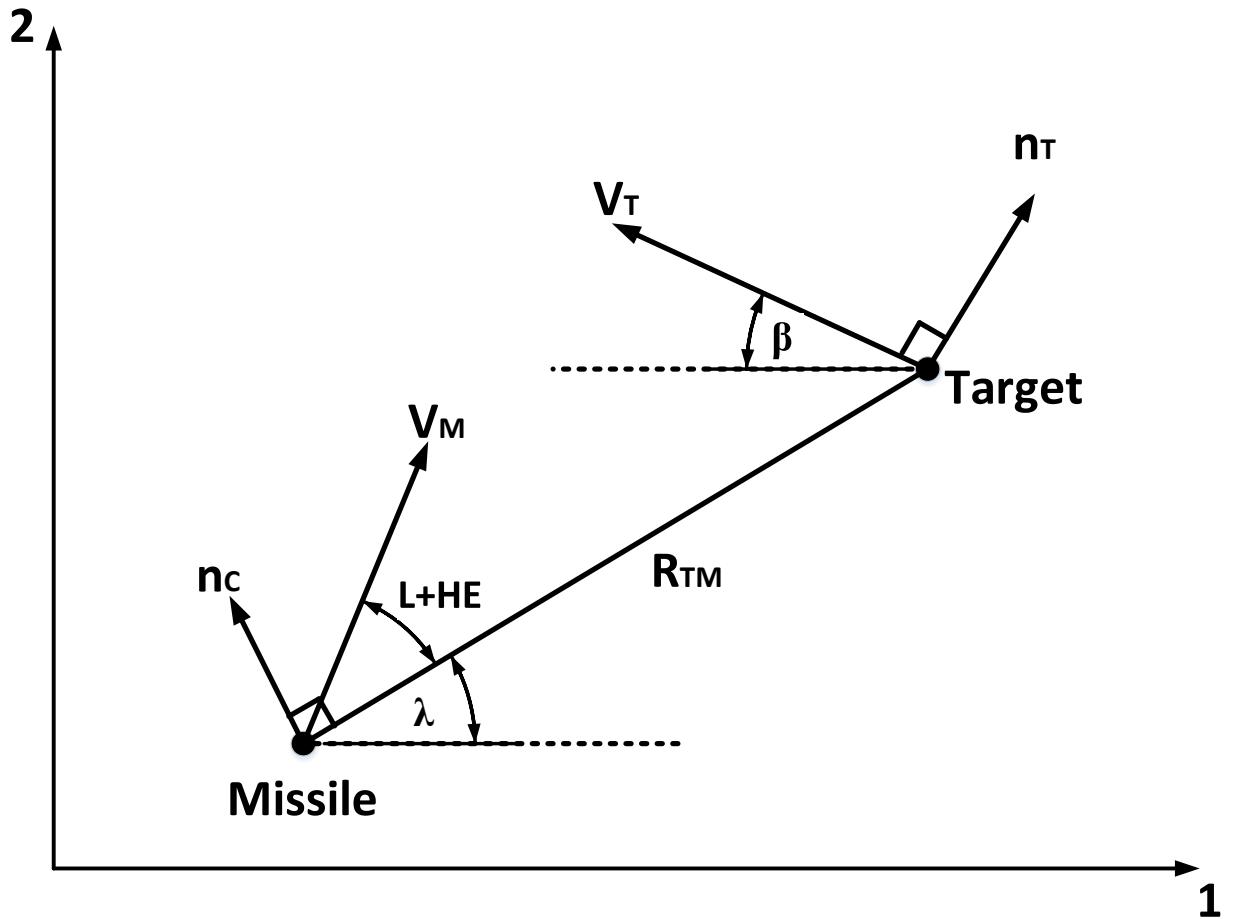


Figure 3.1: Two dimensional Missile-Target engagement geometry.

$$V_{T2} = V_T \sin(\beta) \quad (3.4)$$

Relative missile-target separation

$$R_{TM1} = RT1 - R_{M1} \quad (3.5)$$

$$R_{TM2} = RT2 - R_{M2} \quad (3.6)$$

from the previous 2 equations we get

$$R_{TM} = \sqrt{R_{TM1}^2 + R_{TM2}^2} \quad (3.7)$$

line of sight angle

$$\lambda = \tan^{-1}\left(\frac{R_{TM2}}{R_{TM1}}\right) \quad (3.8)$$

missile lead angle

$$L = \sin^{-1}\left(\frac{V_T \sin(\beta + \lambda)}{V_M}\right) \quad (3.9)$$

the angle between the downrange axis and V_M vector is $\theta = \lambda + L$

Missile velocity components

$$V_{M1} = V_M \cos(\theta + HE) \quad (3.10)$$

$$V_{M2} = V_M \sin(\theta + HE) \quad (3.11)$$

Relative velocity components

$$V_{TM1} = V_{T1} - V_{M1} \quad (3.12)$$

$$V_{TM2} = V_{T2} - V_{M12} \quad (3.13)$$

closing velocity it is the negative rate of change of the distance from the missile to the target $V_c = -R_{TM}$, so we have to differentiate eq(3.7)

$$\dot{R}_{TM} = \frac{1}{2}(R_{TM1}^2 + R_{TM2}^2)^{-\frac{1}{2}} [2R_{TM1}\dot{R}_{TM1} + 2R_{TM2}\dot{R}_{TM2}]$$

we see that so we get

$$V_c = -\dot{R}_{TM} = -\frac{R_{TM1}V_{TM1} + R_{TM2}V_{TM2}}{R_{TM}} \quad (3.14)$$

line of sight rate we have to differentiate eq(3.8) using the rule $\tan^{-1} x = \frac{dx}{1+x^2}$

$$\begin{aligned} \lambda &= \left[\frac{1}{1 + (\frac{R_{TM2}}{R_{TM1}})^2} \right] \left(\frac{\dot{R}_{TM2}}{\dot{R}_{TM1}} \right) \\ &= \frac{R_{TM1}^2}{R_{TM1}^2 + R_{TM2}^2} \left[\frac{R_{TM1}\dot{R}_{TM2} - R_{TM2}\dot{R}_{TM1}}{R_{TM1}^2} \right] \\ &= \frac{R_{TM1}V_{TM2} - R_{TM2}V_{TM1}}{R_{TM1}^2} \end{aligned} \quad (3.15)$$

magnitude of the missile guidance command

$$n_c = N' V_c \lambda \quad (3.16)$$

missile acceleration components

$$a_{M1} = -n_c \sin \lambda \quad (3.17)$$

$$a_{M2} = -n_c \cos \lambda \quad (3.18)$$

angular velocity of the target

$$\dot{\beta} = \frac{n_T}{V_T} \quad (3.19)$$

we will solve all the equations in this section using second-order Runge-Kutta numerical integration procedure. If we have a first order differential equation of the form

$$\dot{x} = f(x, t)$$

where t is time, we seek to find a recursive relationship for x as a function of time. With the second-order RungeKutta numerical technique, the value of x at the next integration interval h is given by

$$x_{k+1} = x_k + \frac{hf(x, t)}{2} + \frac{hf(x, t+h)}{2}$$

3.2 Augmented Proportional Navigation

3.3 Optimal Guidance

Chapter 4: Optimization for escaping maneuver

In this chapter, we will be in the Target side, trying to find an optimized escaping maneuver from the Attacker. We will use two optimization techniques, *Monte Carlo* and *Genetic algorithm*.

4.1 Models & Simulations

In this section we will simulate equations in sec. 3.1.1 for proportional navigation using MATLAB and Simulink. The following figure is a flowchart illustrating the steps of calculations using the equations in sec. 3.1.1 till we plot the trajectories.

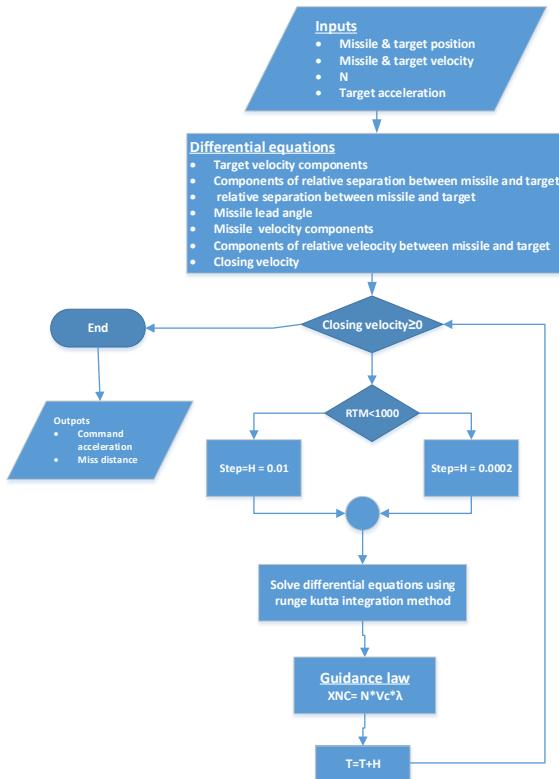


Figure 4.1: Flowchart illustrating the steps of calculations using the equations in sec. 2.3 till we plot the trajectories.

4.1.1 Simulation for some target maneuver cases [using MATLAB]

In this subsection we will use MATLAB to simulate equations in sec. 3.1.1 for proportional navigation. We will solve the differential equations using second-order Runge-Kutta numerical integration technique, then we will draw the trajectories of the pursuit and evader for 4 cases for the Target maneuver error source, and we will deduce the effect of the effective navigation ratio N' and the other type of the error source; Heading error.

4.1.1.1 Zero Target maneuver

In this case the evader (Target - plane) does not do any effort to scape, it just moves in a straight line, as we see in Fig. 4.2. So the pursuer (Attacker - Missile) does not have to bleed much energy to reach the Target.

In the case of **zero Heading error** the effective navigation ratio has no effect on the simulation engagement at all. The missile's acceleration will be zero as in Fig. 4.3.

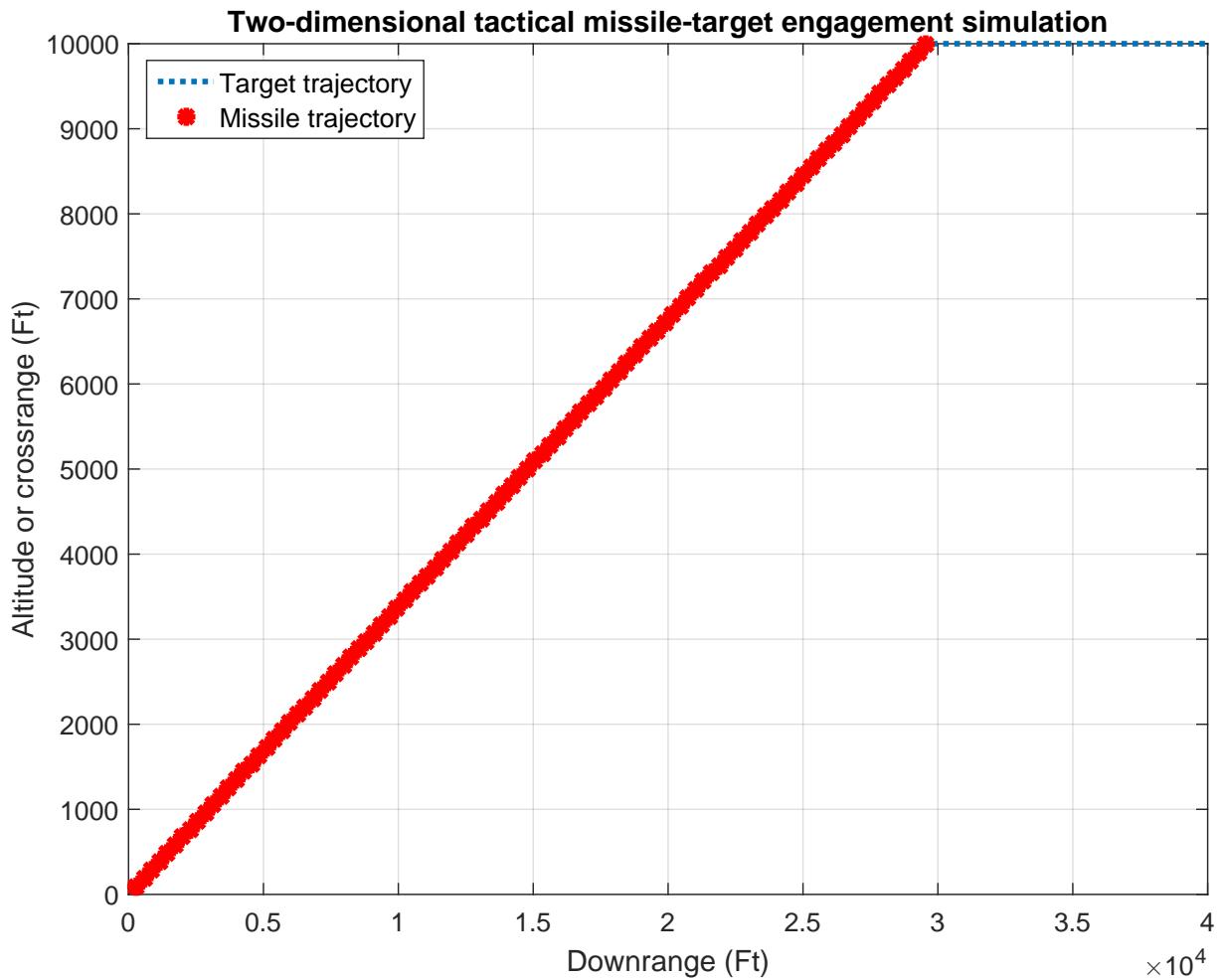


Figure 4.2: Trajectory of the target and attacker in case of zero target maneuver with zero heading error and $N' = 4$.

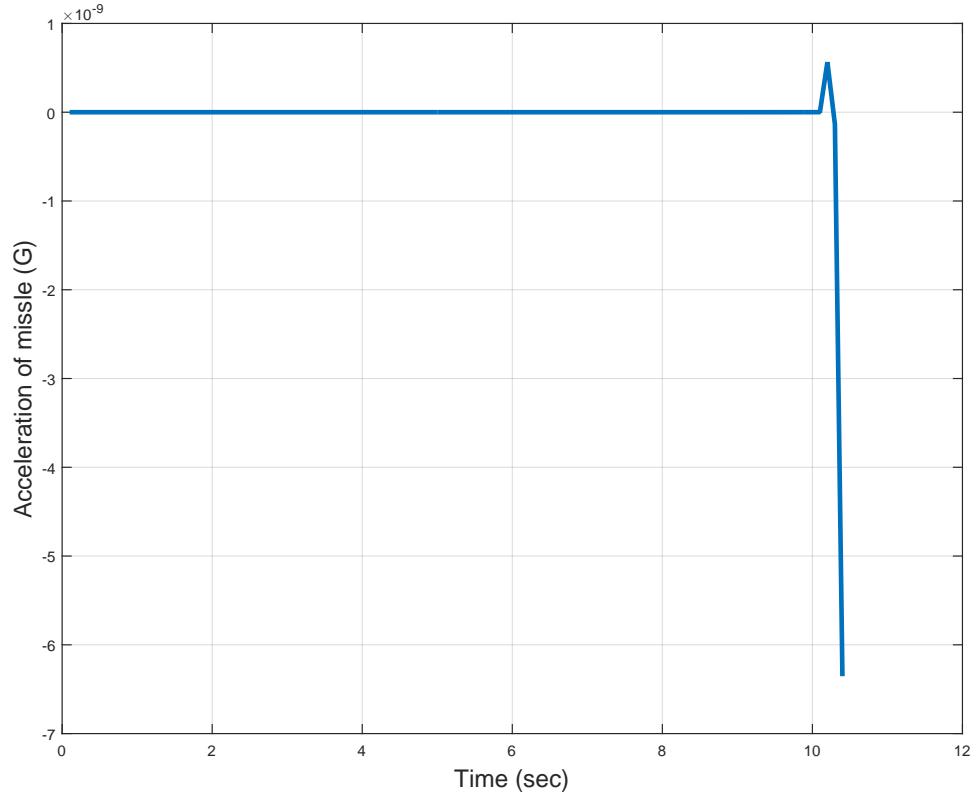


Figure 4.3: Missile acceleration in case of zero target maneuver with zero heading error and $N' = 4$.

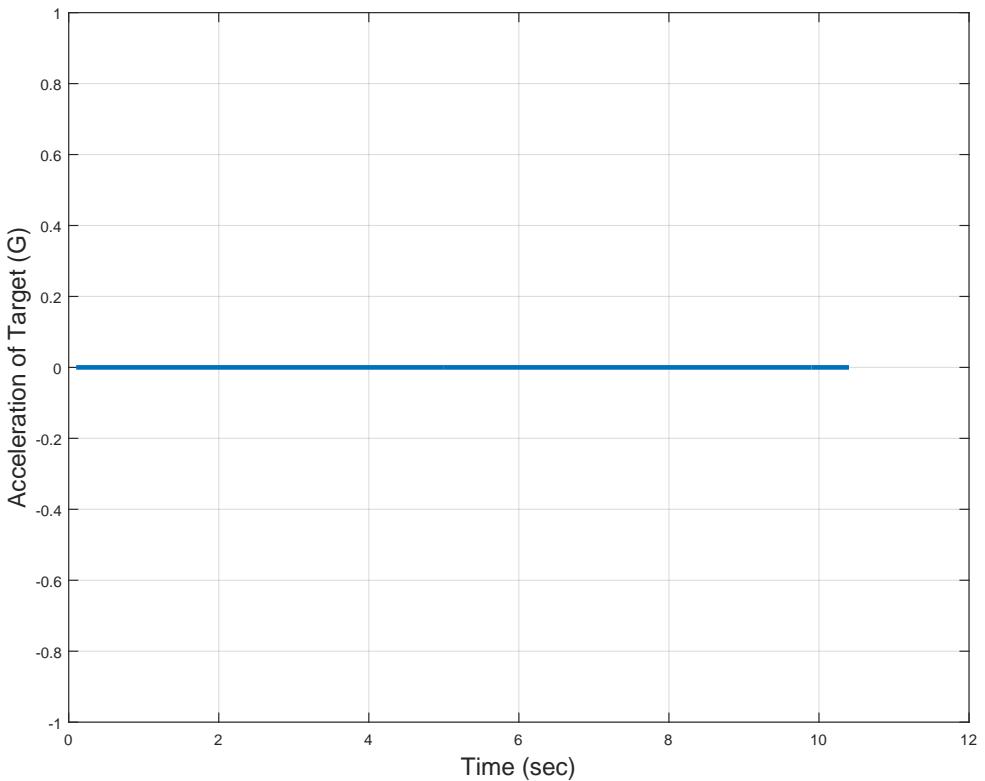


Figure 4.4: Target acceleration in case of zero target maneuver with zero heading error and $N' = 4$.

In the case of **Heading error = -20** increasing the effective navigation ratio causes heading error to be removed rapidly as we see from Fig. 4.5, Fig. 4.7 and Fig. 4.9. The effective navigation ratio has an effect on the acceleration of the missile; the way that the missile will bleed energy as we see from Fig. 4.6 , Fig. 4.8 and Fig. 4.10. The total acceleration (area under the curve) is increasing inversely proportional with the effective navigation ratio N' .

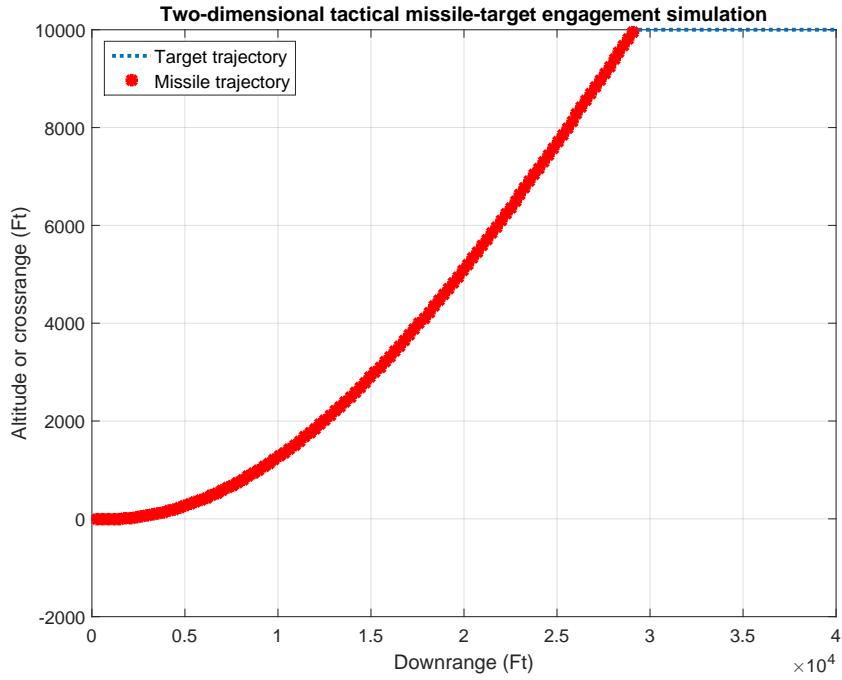


Figure 4.5: Trajectory of the target and attacker in case of zero target maneuver with heading error=-20 and $N' = 3$.

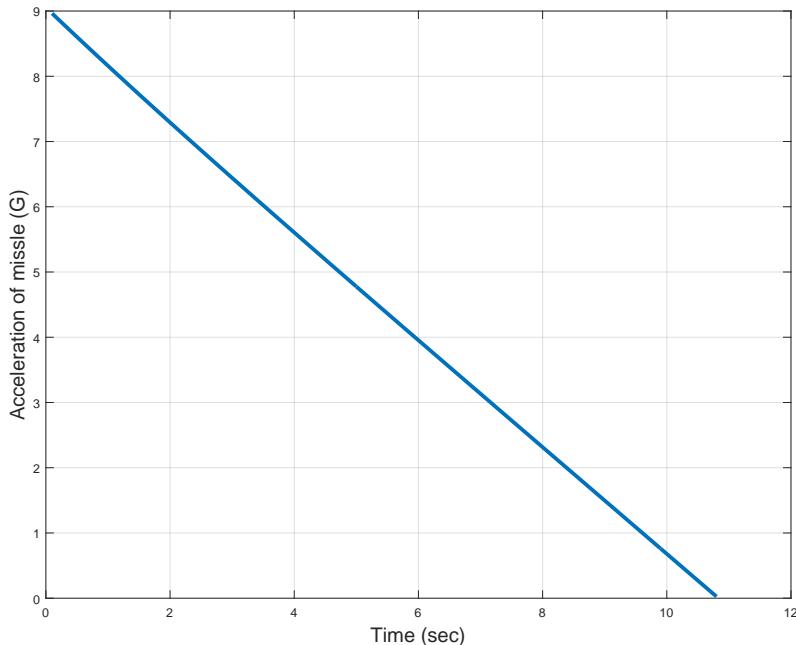


Figure 4.6: Missile acceleration in case of zero target maneuver with heading error=-20 and $N' = 3$.

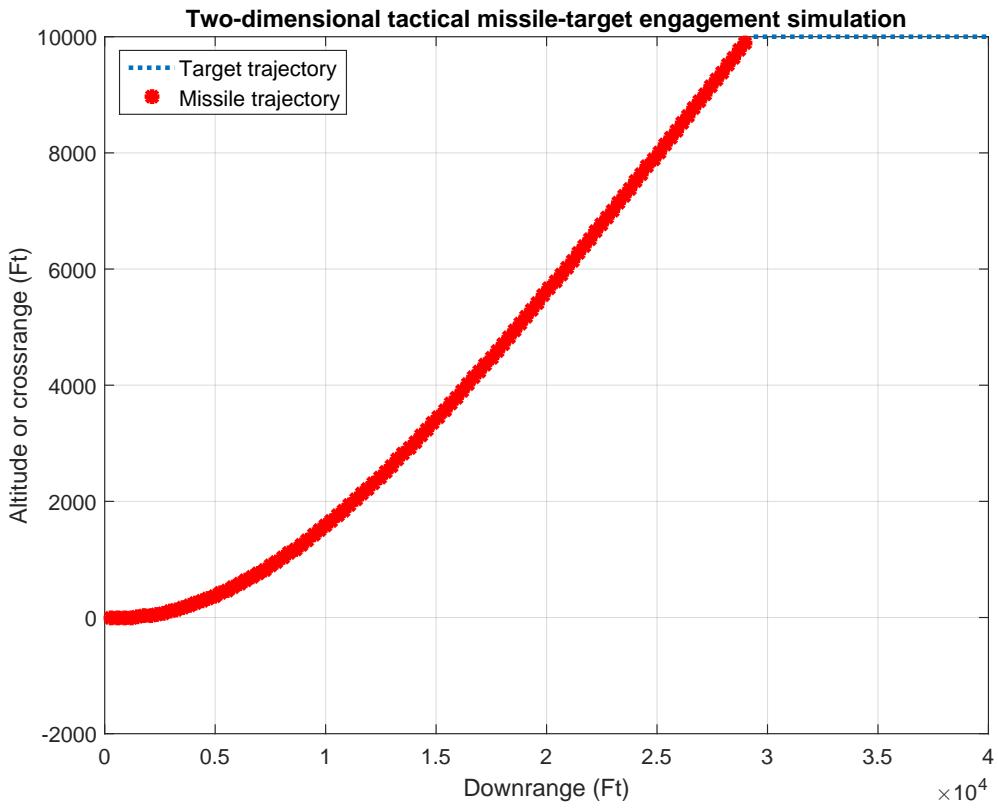


Figure 4.7: Trajectory of the target and attacker in case of zero target maneuver with heading error=-20 and $N' = 4$.

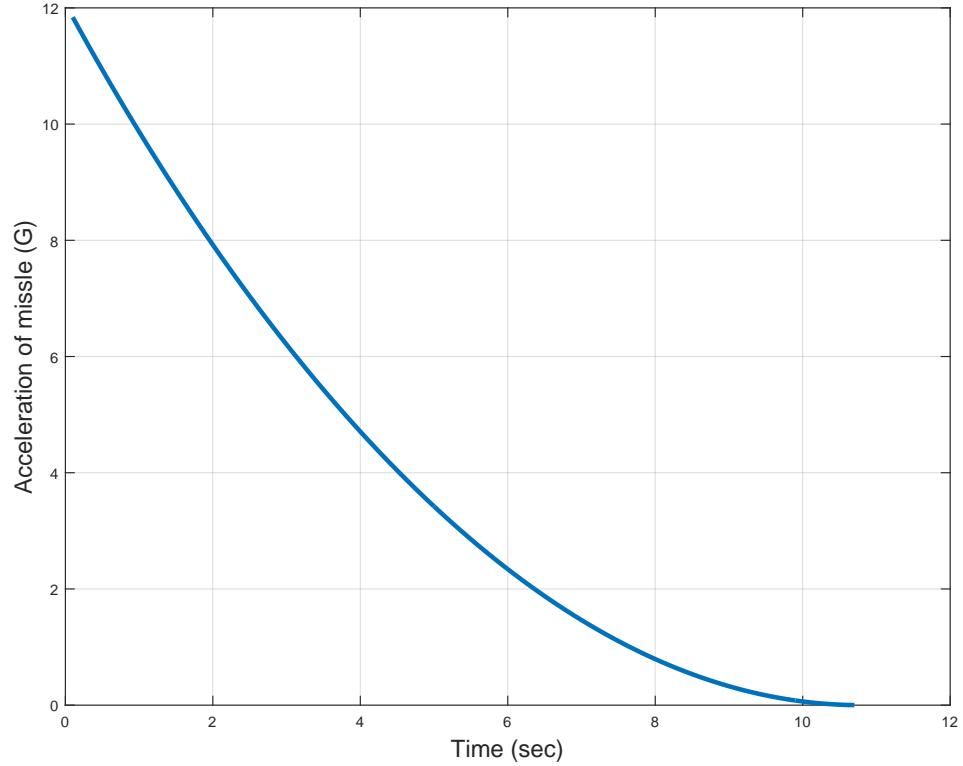


Figure 4.8: Missile acceleration in case of zero target maneuver with heading error=-20 and $N' = 4$.

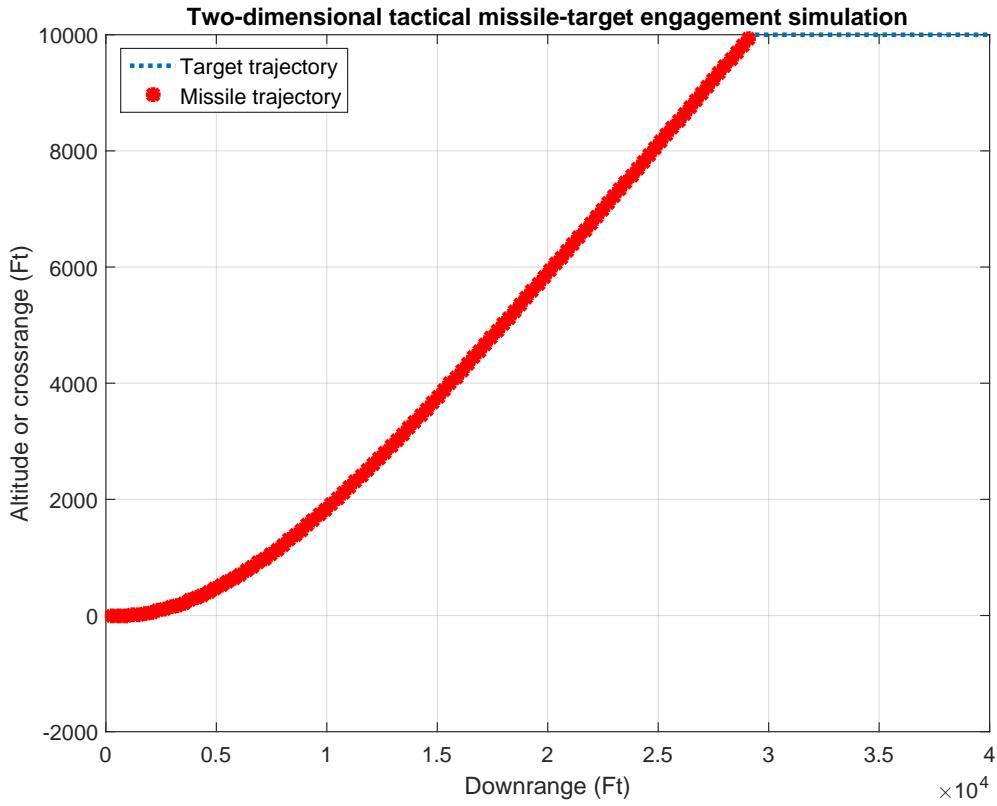


Figure 4.9: Trajectory of the target and attacker in case of zero target maneuver with heading error=-20 and $N' = 5$.

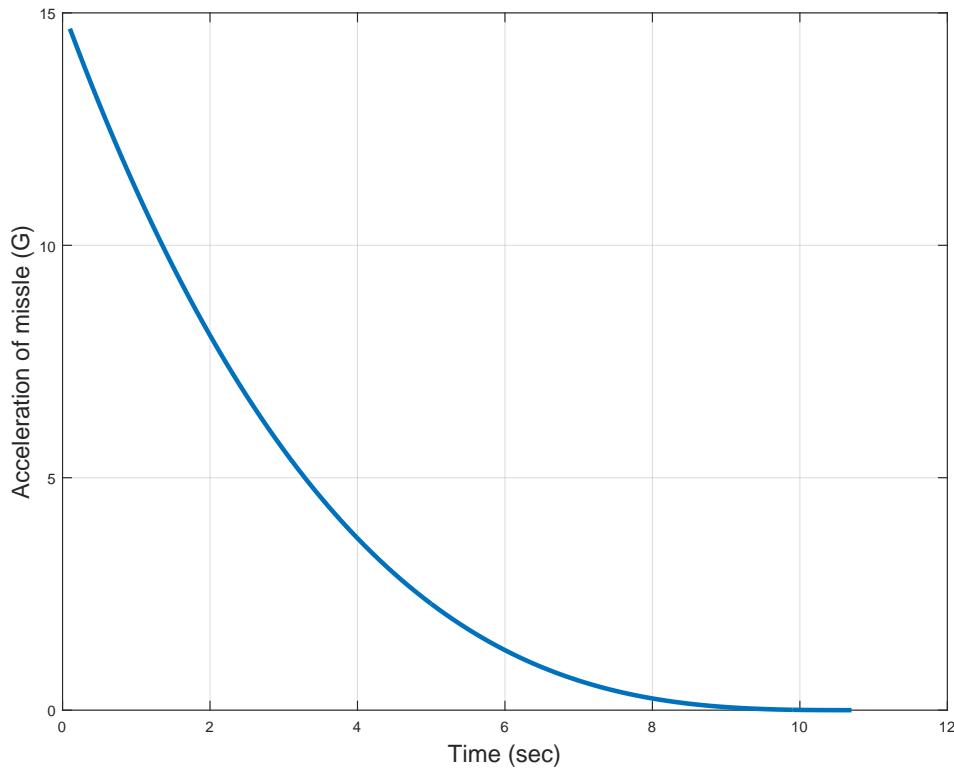


Figure 4.10: Missile acceleration in case of zero target maneuver with heading error=-20 and $N' = 5$.

4.1.1.2 Constant Target maneuver

In this case the target does some effort to escape from the attacker, in the form of constant acceleration (in our example target acceleration = $5G$). We see from Fig. 4.12 and Fig. 4.15 that a higher effective navigation ratio yields less acceleration to hit maneuvering target, and causes the missile to lead the target slightly more than a lower effective navigation ratio does, as we see from Fig. 4.11 and Fig. 4.14.

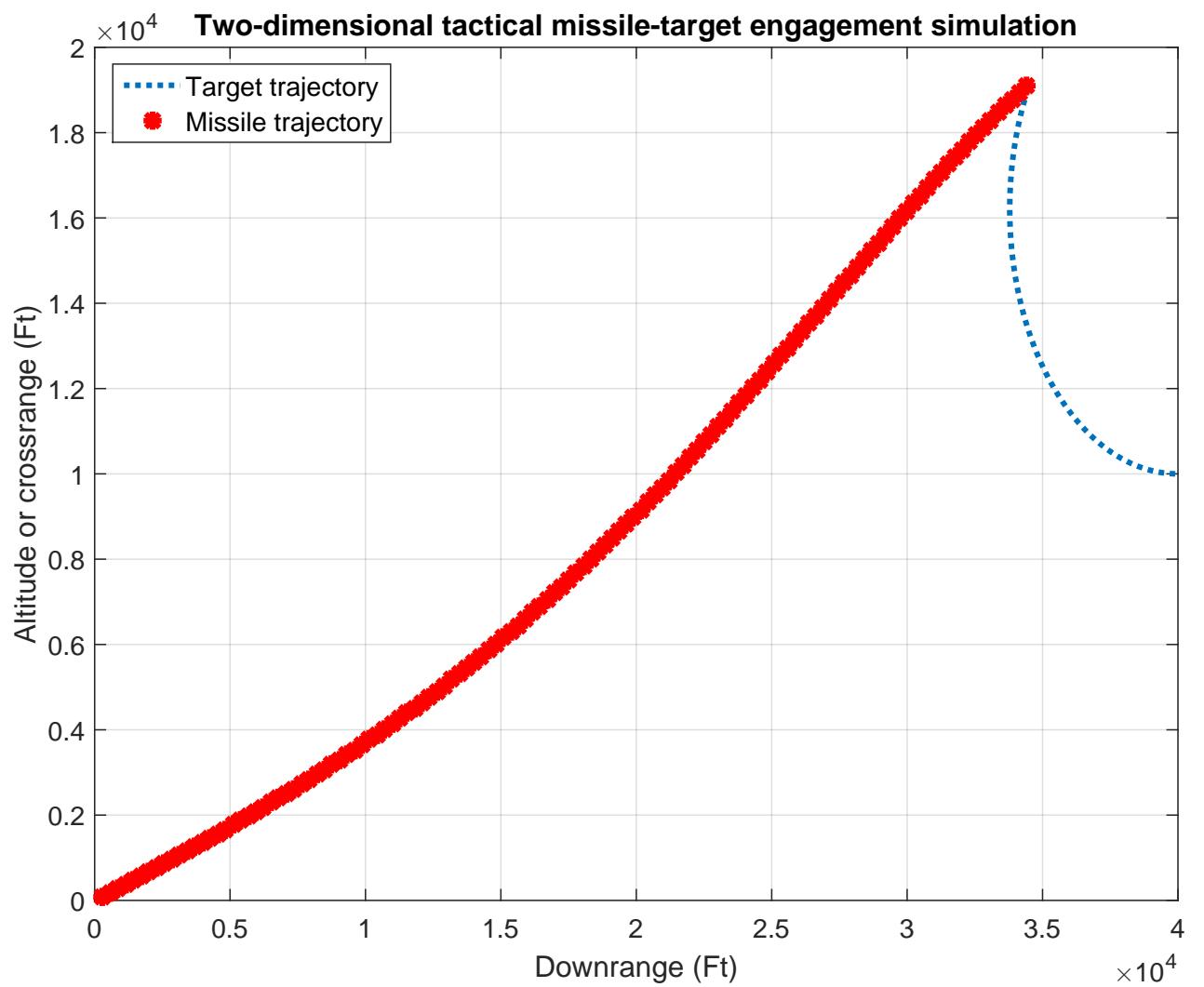


Figure 4.11: Trajectory of the target and attacker in case of constant target maneuver= $5G$ with heading error=0 and $N' = 3$.

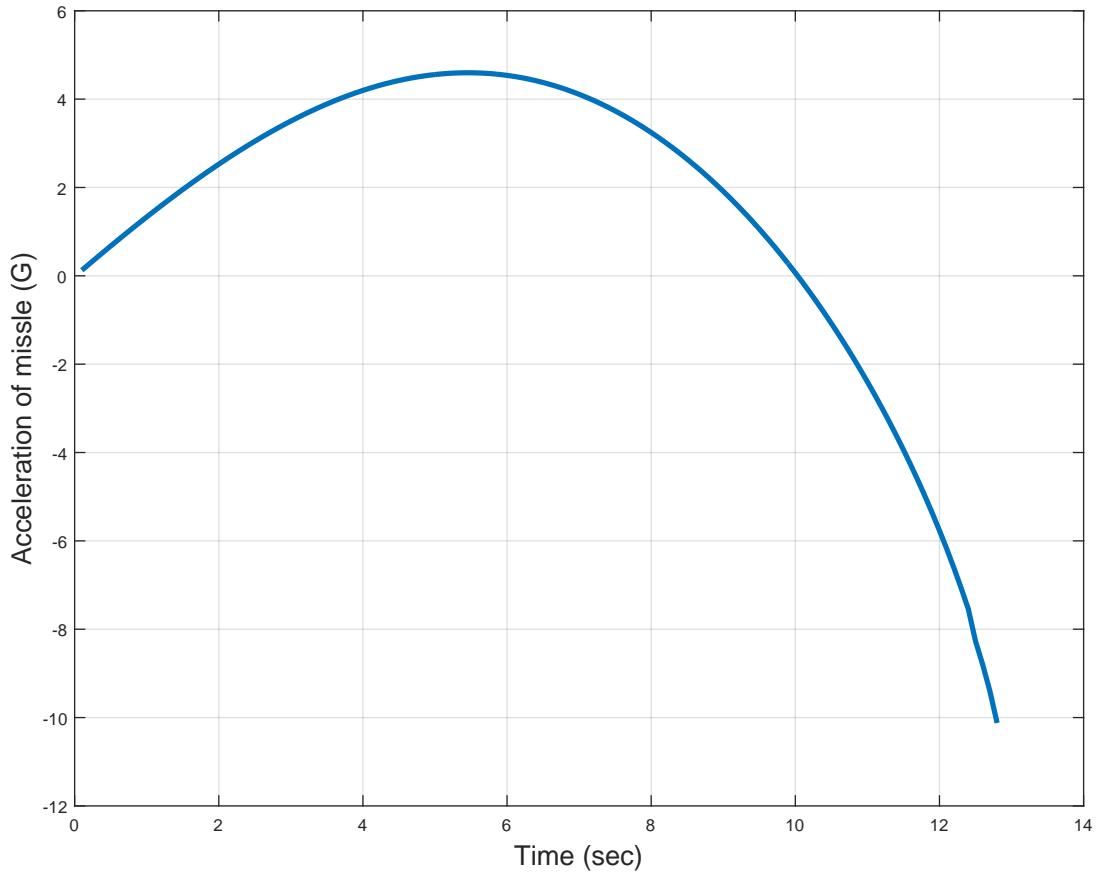


Figure 4.12: Missile acceleration in case of constant target maneuver=5G with heading error=0 and $N' = 3$.

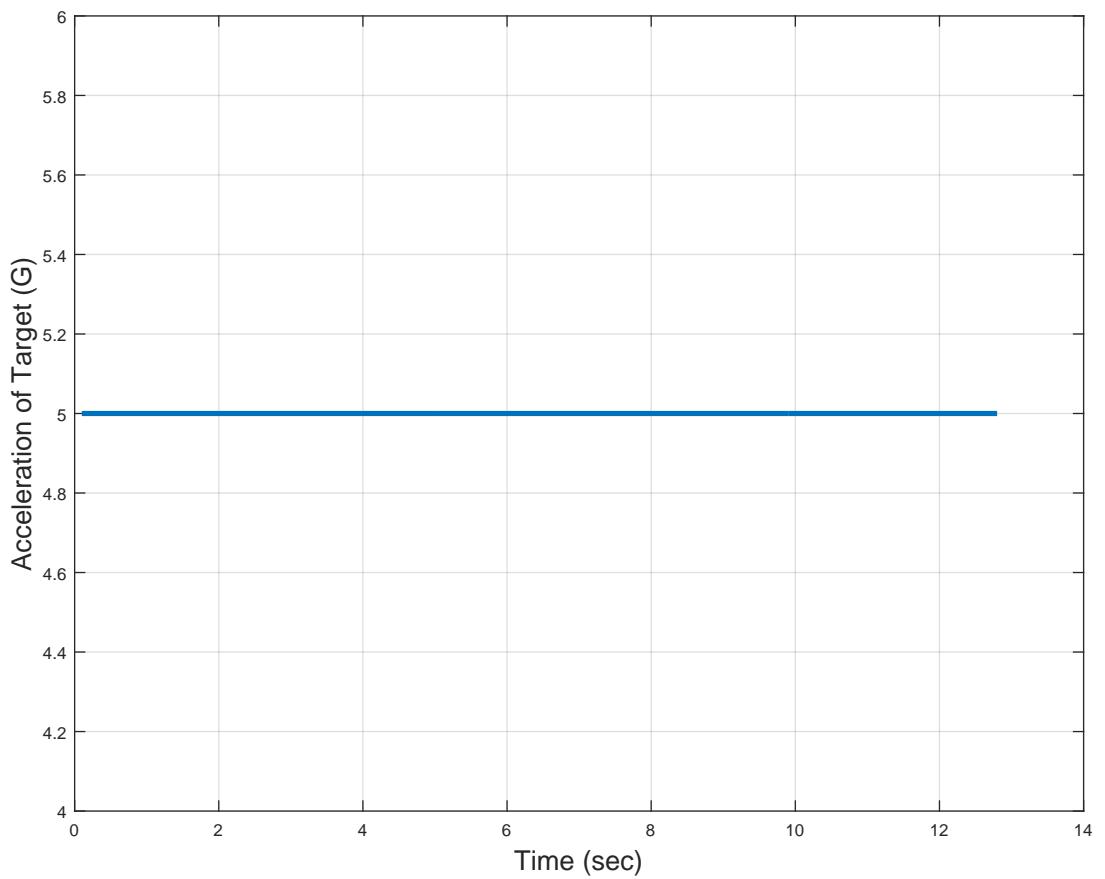


Figure 4.13: Target acceleration in case of zero target maneuver with zero heading error and $N' = 3$. ²⁰

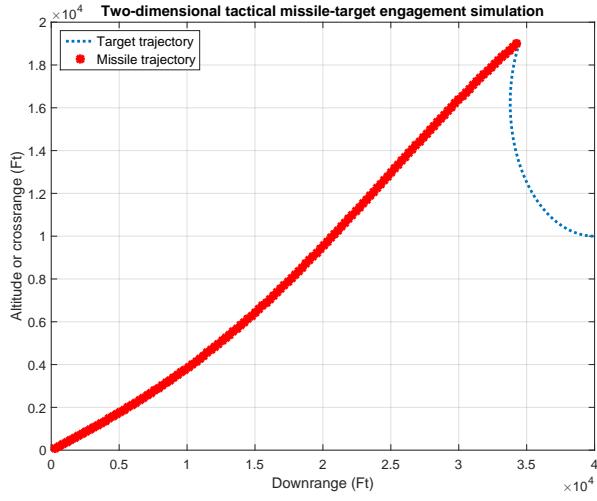


Figure 4.14: Trajectory of the target and attacker in case of constant target maneuver=5G with heading error=0 and $N' = 5$.

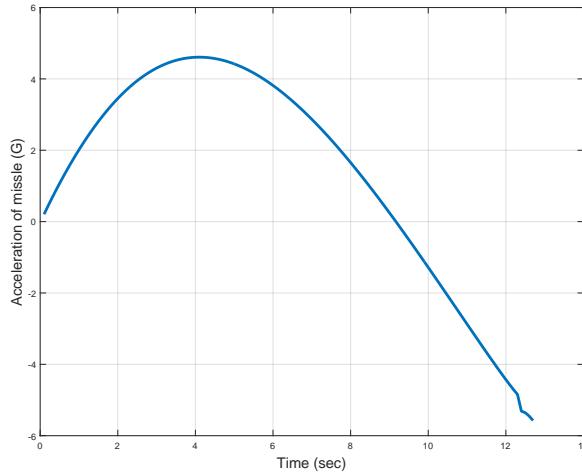


Figure 4.15: Missile acceleration in case of constant target maneuver=5G with heading error=0 and $N' = 5$.

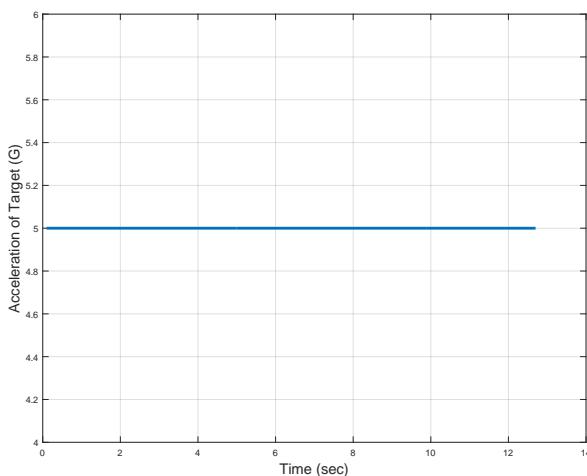


Figure 4.16: Target acceleration in case of zero target maneuver with zero heading error and $N' = 5$.

4.1.1.3 Polynomial Target maneuver

Now we will make the Target maneuver as a polynomial in the form

$$f(t) = c_0 + c_1 T + c_2 T^2 + c_3 T^3 + \dots + c_N T^N \quad (4.1)$$

where T is time , c are unknown coefficients of the polynomial and N is the polynomial degree. In our simulation we choose N reasonably $3 \rightarrow 5$ and the coefficients are randomly generated then we solve the proportional navigation equations to get values for two important parameters, namely the miss distance and the command acceleration for the Missile. Then, we calculate the cost function depends on these two parameters, as our target is to maximize miss distance to make the target safe, and maximize command acceleration for the missile, so it bleeds energy before it catches the target. We will do this process many times till the number of runs we need. The next figures shows the trajectories of the Target and Missile for $N=3,4$ and corresponding missile and target accelerations for each case. In the first case ($N=3$), as we see from Fig. 4.19 the target have to exert $50G$ to escape from the attacker, these amount of G'es is very high and the pilot could not withstand all this amount. Also, form Fig. 4.18 the attacker have to exert $120 G$, this will not happen and the missile's energy will bleed at $100 G$ after 5 sec. Here the simulation did not till us a complete vision about what will happen. We have to take in consideration the amount of G's that the pilot could resist, so we will use genetic algorithm technique to generate the polynomial because it is easy to control the maximum target acceleration in this method. The second case ($N=4$), is hypothetical case because the pilot will die in less than 1 sec. If the technology improved and the scientists invent a pilot suit that could resist these huge amount of G'es, then we can take this solution in consideration.

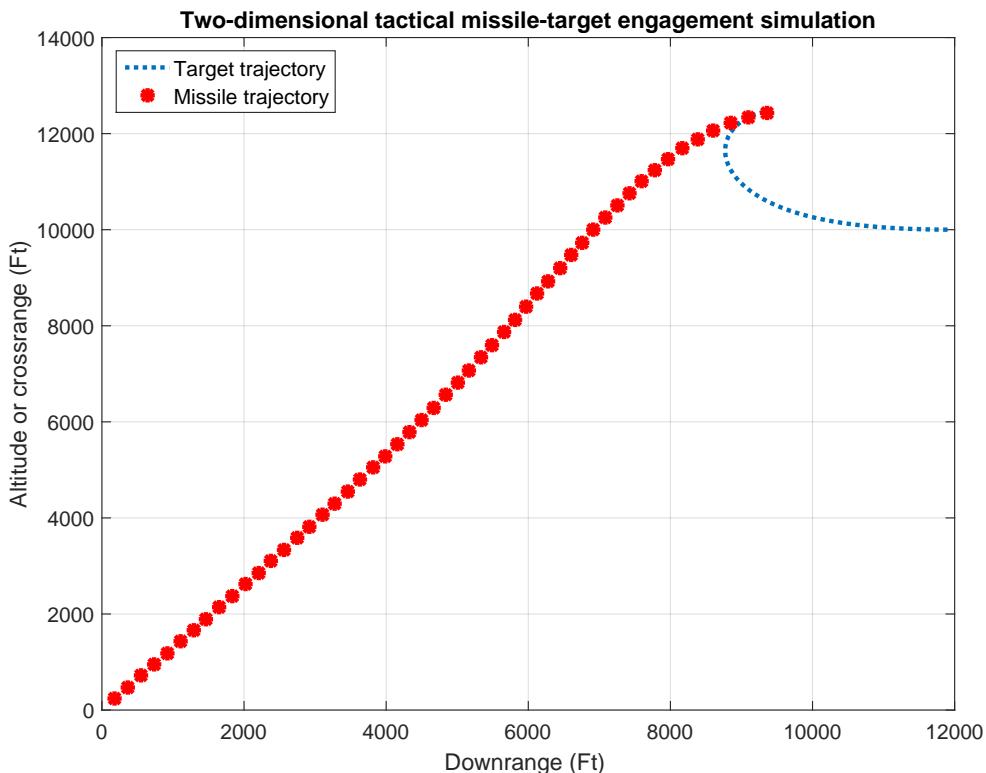


Figure 4.17: Trajectory of the target and attacker in case of polynomial of degree $N=3$ target maneuver with zero heading error and $N' = 3$.

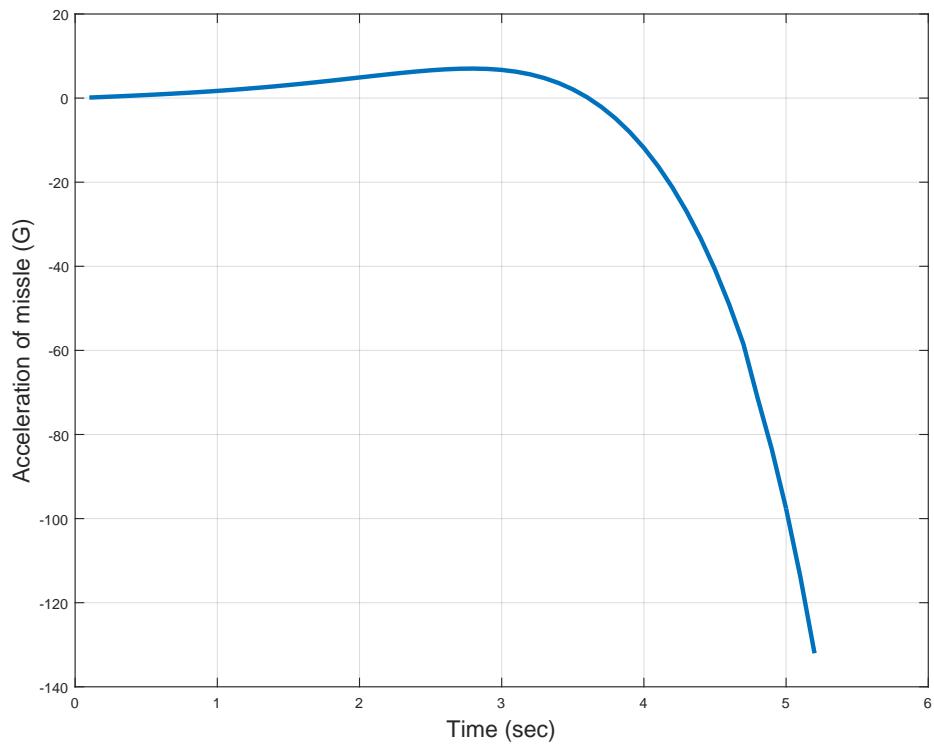


Figure 4.18: Missile acceleration in case of polynomial of degree $N=3$ target maneuver with zero heading error and $N' = 3$.

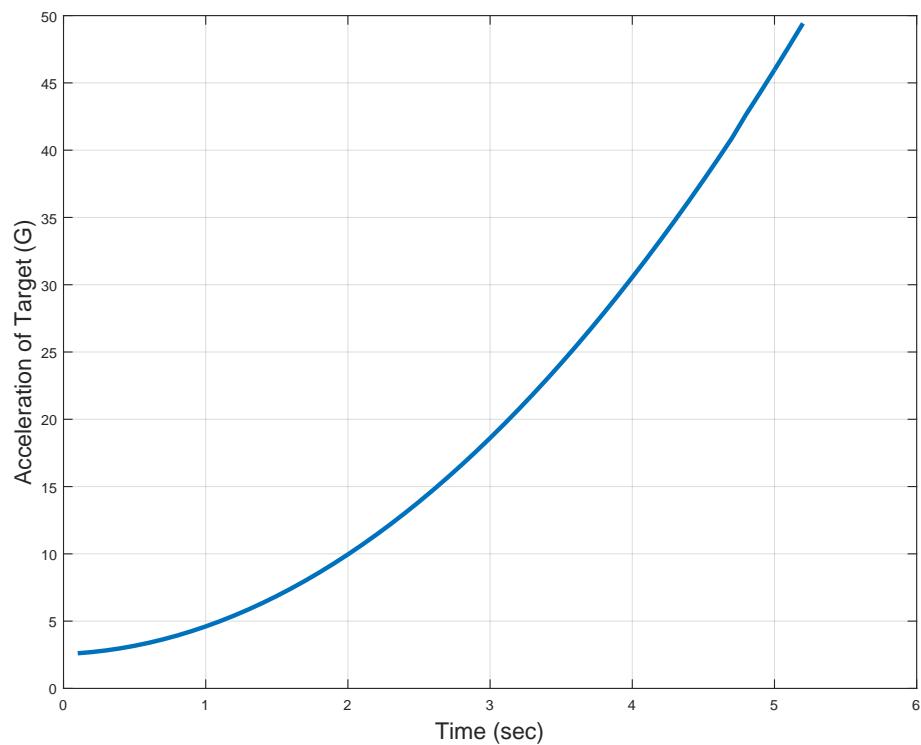


Figure 4.19: Target acceleration in case of polynomial of degree $N=3$ target maneuver with zero heading error and $N' = 3$.

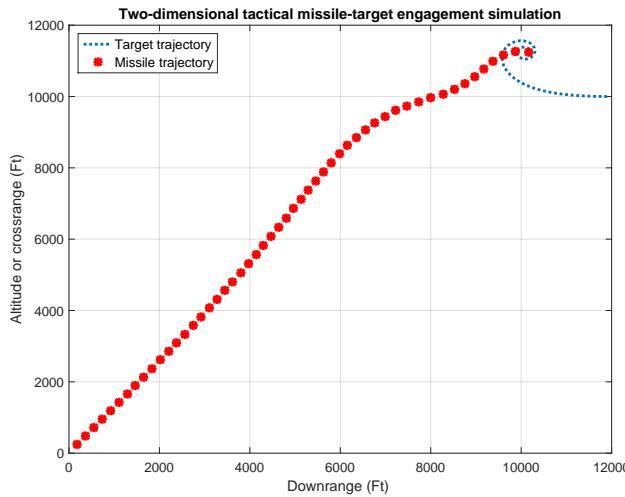


Figure 4.20: Trajectory of the target and attacker in case of polynomial of degree $N=4$ target maneuver with zero heading error and $N' = 3$.

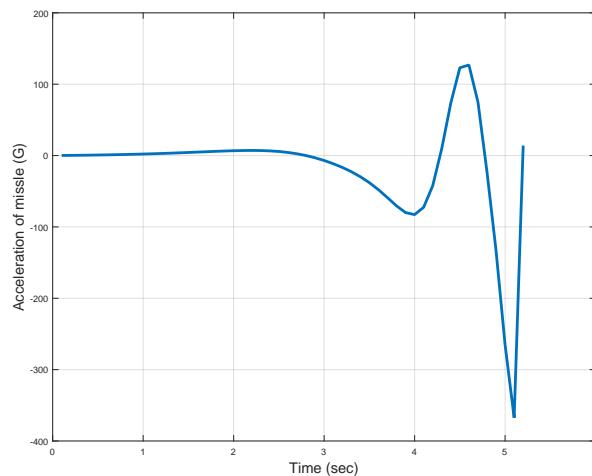


Figure 4.21: Missile acceleration in case of polynomial of degree $N=4$ target maneuver with zero heading error and $N' = 3$.

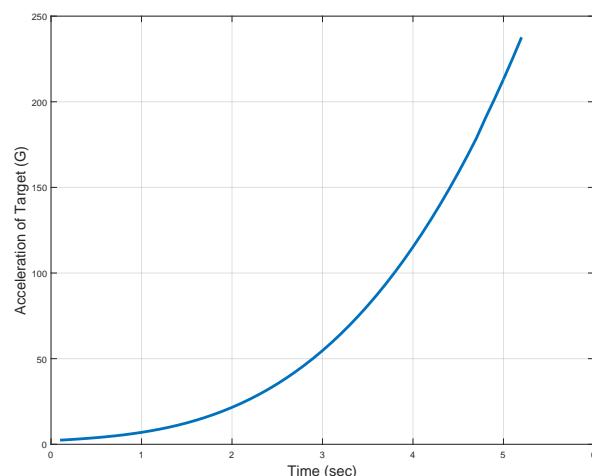


Figure 4.22: Target acceleration in case of polynomial of degree $N=4$ target maneuver with zero heading error and $N' = 3$.

4.1.1.4 Trapezoidal Target maneuver

Now we will make the Target maneuver as a trapezoidal function (as a general case of tooth and square maneuvers) illustrated in Fig. 4.23

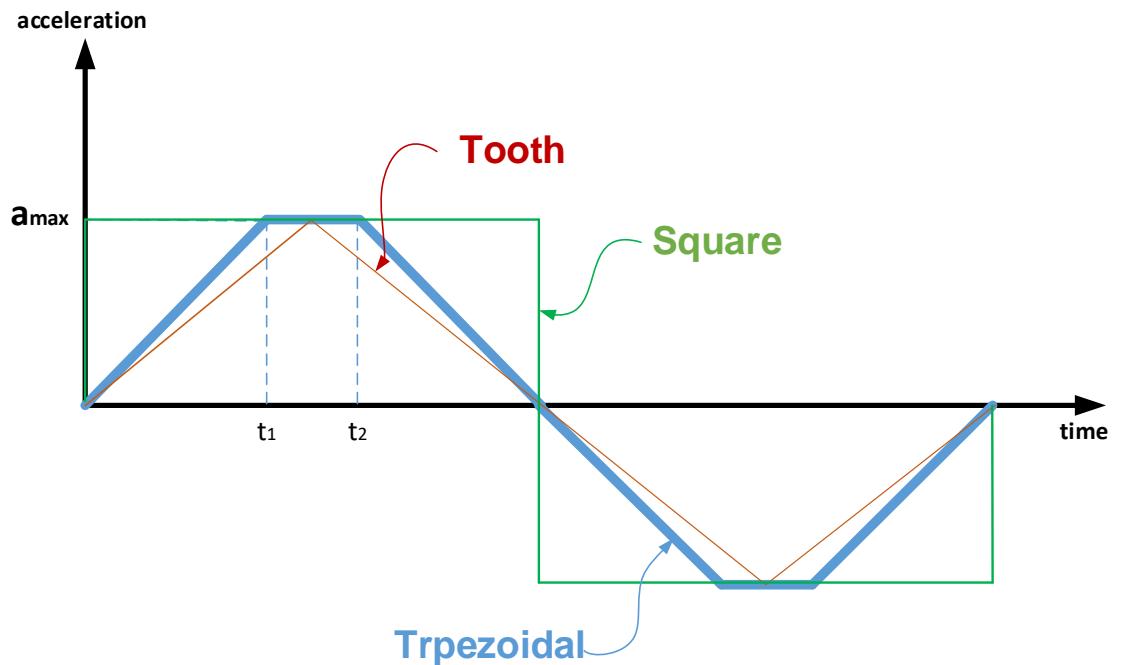


Figure 4.23: Trapezoidal Target maneuver.

In the simulation we choose the value for maximum acceleration. The time of beginning the ramp (which is 0), and the time at the end of the descent, t_1 and t_2 are variables.

We have two cases: **First case:** the maneuver is *symmetric* so we have two extreme cases.

- when t_1 is equal to zero (there is no ascend ramp) so t_2 is at the middle of the maneuver; this leads to *square* target maneuver.
- when t_1 equals t_2 at the quarter of the maneuver; this leads to *tooth* target maneuver.

Varying t_1 and t_2 , starting from the square case till the tooth case, until we find the optimum solution based on or cost function which depends on the values of the miss distance and Missile acceleration.

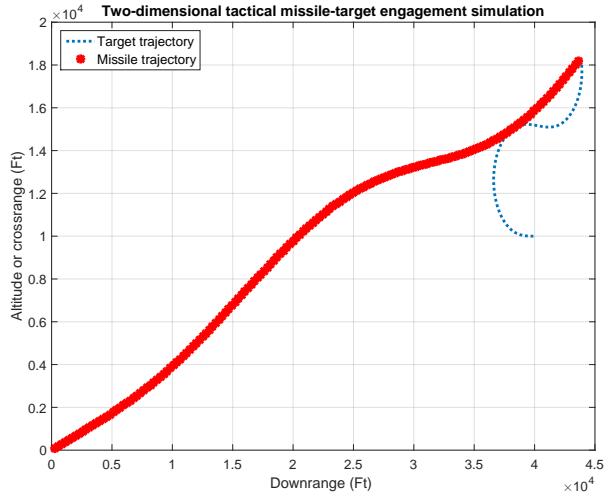


Figure 4.24: Trajectory of the target and attacker in case of symmetric trapezoidal target maneuver with zero heading error and $N' = 5$.

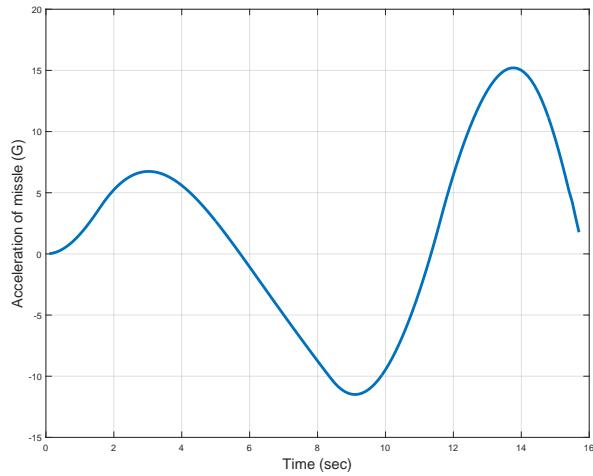


Figure 4.25: Missile acceleration in case of symmetric trapezoidal target maneuver with zero heading error and $N' = 5$.

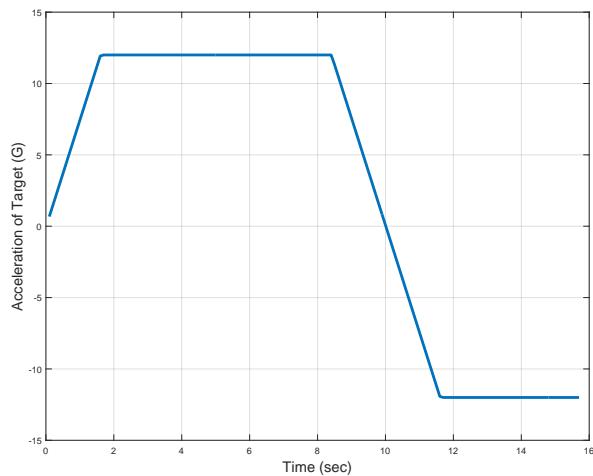


Figure 4.26: Target acceleration in case of symmetric trapezoidal $N=4$ target maneuver with zero heading error and $N' = 5$.

Second case: the maneuver is *not symmetric* so we put a first value for t_1 (starting with zero) and varying t_2 till we scan all the domain, then change t_1 slightly and scan all the domain again with t_2 till we find the optimum solution.

What about if the optimum solution for the target maneuver did not have Identical halves? and all the slopes are not related to each other. We will try to figure that using genetic algorithm in the next section.

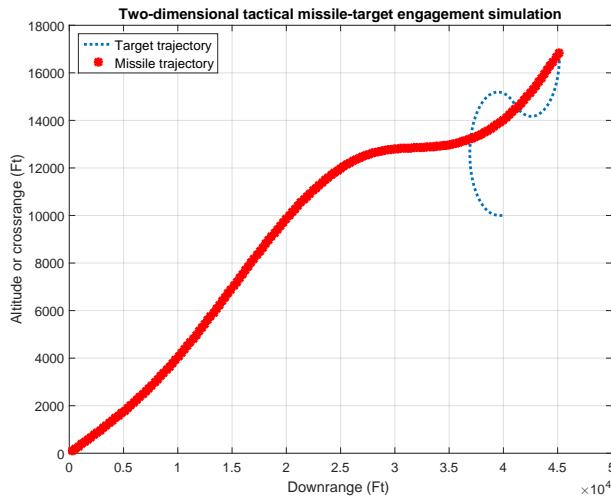


Figure 4.27: Trajectory of the target and attacker in case of trapezoidal target maneuver with zero heading error and $N' = 5$.

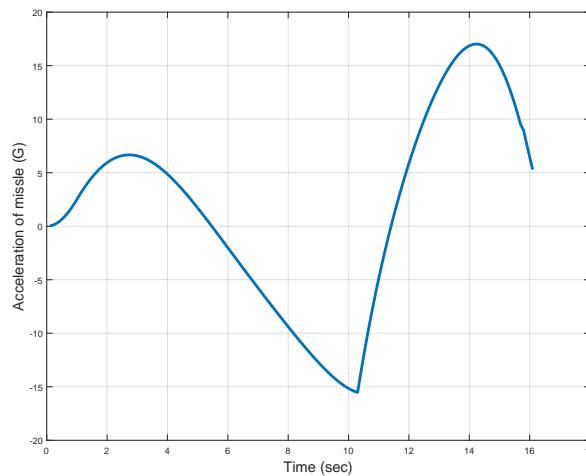


Figure 4.28: Missile acceleration in case of trapezoidal target maneuver with zero heading error and $N' = 5$.

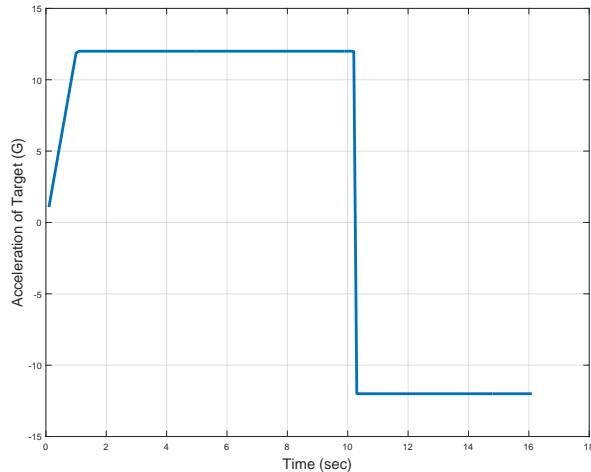


Figure 4.29: Target acceleration in case of trapezoidal $N=4$ target maneuver with zero heading error and $N' = 5$.

4.2 Guidance Toolbox

All the previous results could be generated from a simple GUI (graphic user interface), (Fig.4.30). The inputs for this GUI is the locations and velocities for the Target (plane) and the Attacker (missile). Then the user have to choose the guidance law, which is the way that the Attacker tracking the Target, and there is 3 options :

- Proportional Navigation (PN)
- Augmented Proportional Navigation
- Optimal Guidance

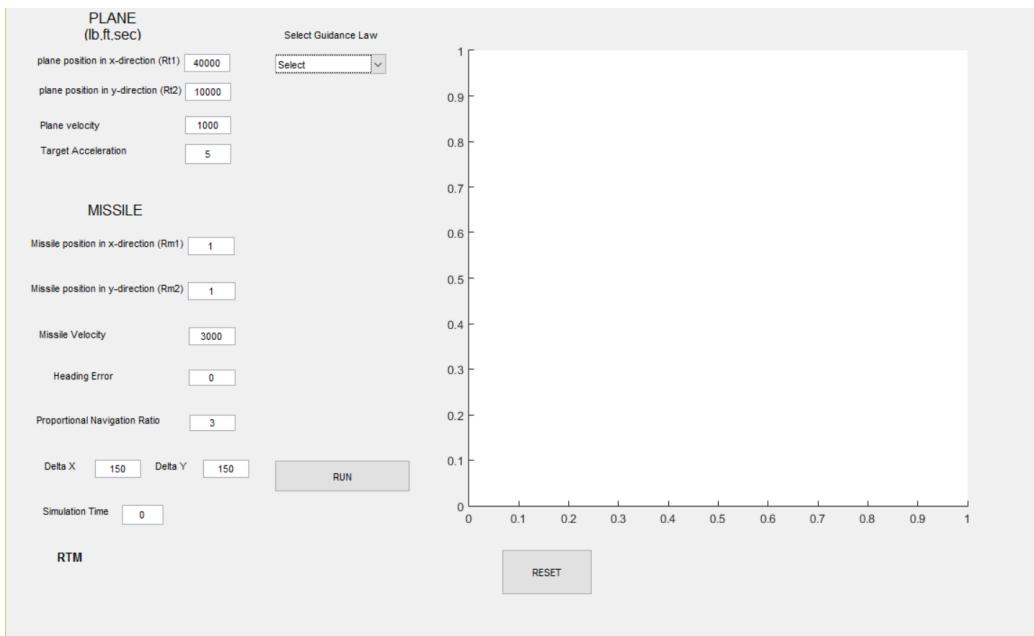


Figure 4.30: Guidance Toolbox

If the user chooses the PN guidance law (Fig.4.31) , there will be more options be available, Now he could select the type of escape maneuver:

- Polynomial
- Trapezoidal
- Symmetric Trapezoidal

and the result will be optimized escape maneuver.

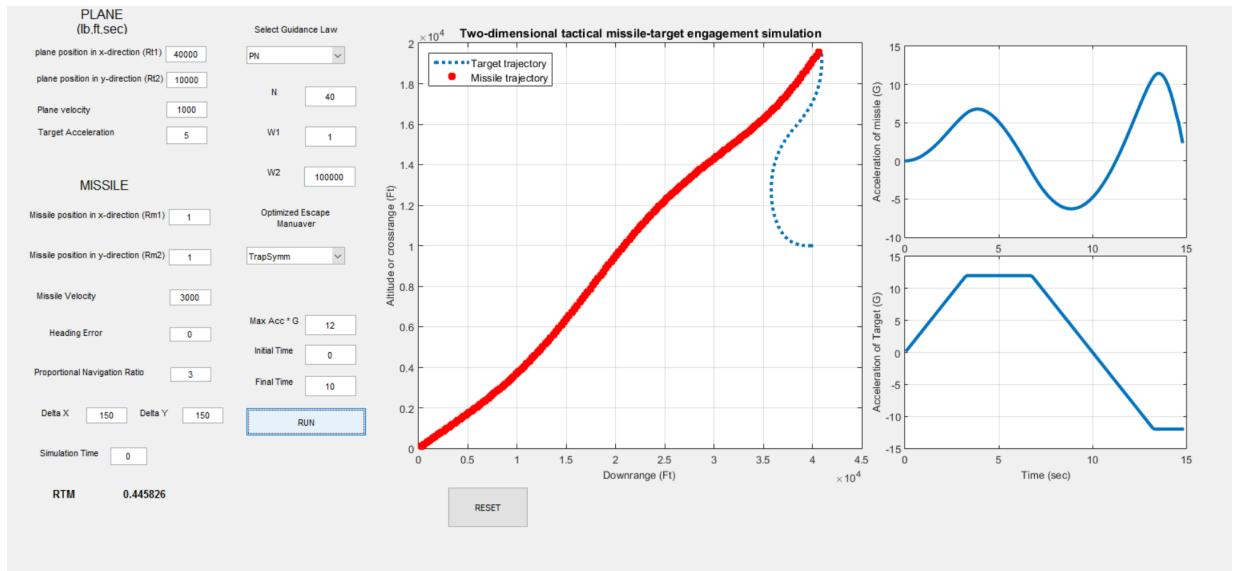


Figure 4.31: Guidance Toolbox options when PN selected

4.3 Genetic Algorithm Solution [using simulink]

In this subsection we will use the power of blocks in Simulink to solve the differential equations in section 3.1.1 in an easy way (Fig. 4.32). We will make these equations as a main block (Fig. 4.33) and as we change the inputs, we get the same results. The advantage of using simulink is to use the main block with an alternative optimization technique (like genetic algorithms) and we will simulate two cases, polynomial and trapezoidal. Now let us have a brief introduction to the genetic algorithm technique.

4.3.1 Introduction to Genetic Algorithm technique

Genetic algorithm is a *stochastic population-based* search method. Genetic Algorithms (GAs) are search algorithms based on the mechanics of the natural selection process (biological evolution). The most basic concept is that the strong tend to adapt and survive while the weak tend to die out. That is, optimization is based on evolution, and the "Survival of the fittest" concept. GAs have the ability to create an initial population of feasible solutions, and then recombine them in a way to guide their search to only the most promising areas of the state space. Each feasible solution is encoded as a chromosome (string) also called a genotype, and each chromosome is given a measure of fitness via a fitness (evaluation or objective) function.

- The fitness of a chromosome determines its ability to survive and produce offspring.
- A finite population of chromosomes is maintained.
- GAs use probabilistic rules to evolve a population from one generation to the next. The generations of the new solutions are developed by genetic recombination operators:
 - Biased Reproduction: selecting the fittest to reproduce.
 - Crossover: combining parent chromosomes to produce children chromosomes.
 - Mutation: altering some genes in a chromosome.

Most Important Parameters in GAs:

- Population Size
- Evaluation Function
- Crossover Method
- Mutation Rate

Use Genetic Algorithms:

- When facing a problem of local minima.
- When a good fitness function is available.
- When a near-optimal, but not optimal solution is acceptable.
- When the state-space is too large for other methods.

4.3.2 COMPONENTS OF GENETIC ALGORITHMS

Representation (definition of individuals):

Objects forming possible solution within original problem context are called phenotypes, their encoding, the individuals within the GA, are called genotypes. The representation step specifies the mapping from the phenotypes onto a set of genotypes. Candidate solution, phenotype and individual are used to denote points of the space of possible solutions. This space is called phenotype space. Chromosome, and individual can be used for points in the genotype space. Elements of a chromosome are called genes. A value of a gene is called an allele. The GA works with a coding of the parameter rather than the actual parameter. Depending on the particular problem, the coding may be Boolean-valued, integer-valued, real-valued, complex-valued, vector-valued, symbolic valued, or multiple-valued.

Evaluation function:

This is just the cost or the objective function which represent what we want to maximize. For minimization problem use the reciprocal (i.e. $1/f$).

Population:

The role of the population is to hold possible solutions. A population is a multiset of genotypes. In GA, the population size is (almost always) constant. The bigger the size, the better the final solution, and the larger the execution time.

Parent Selection Mechanism:

The role of parent selection (mating selection) is to distinguish among individuals based on their quality to allow the better individuals to become parents of the next generation. Parent selection is probabilistic. Thus, high quality individuals get a higher chance to become parents than those with low quality. Nevertheless, low quality individuals are often given a small, but positive chance; otherwise the whole search could become too greedy and get stuck in a local optimum.

Variation Operators:

The role of variation operators is to create new individuals from old ones. Variation operators form the implementation of the elementary steps with the search space.

Mutation Operator:

A unary variation operator is called *mutation*. It is applied to one genotype and delivers a modified mutant, the child or *offspring* of it. In general, mutation is supposed to cause a random unbiased change.

Crossover Operator:

A binary variation operator is called recombination or crossover. This operator merges information from two parent genotypes into one or two offspring genotypes. Similarly to mutation, crossover is a stochastic operator: the choice of what parts of each parent are combined, and the way these parts are combined, depend on random drawings. The principle behind crossover is simple: by mating two individuals with different but desirable features, we can produce an offspring which combines both of those features.

Survivor Selection Mechanism:

The role of survivor selection is to distinguish among individuals based on their quality. In GA, the population size is (almost always) constant, thus a choice has to be made on which individuals will be allowed in the next generation. This decision is based on their fitness values, favoring those with higher quality. As opposed to parent selection which is stochastic, survivor selection is often deterministic, for instance, ranking the unified multiset of parents and offspring and selecting the top segment (fitness biased), or selection only from the offspring (age-biased).

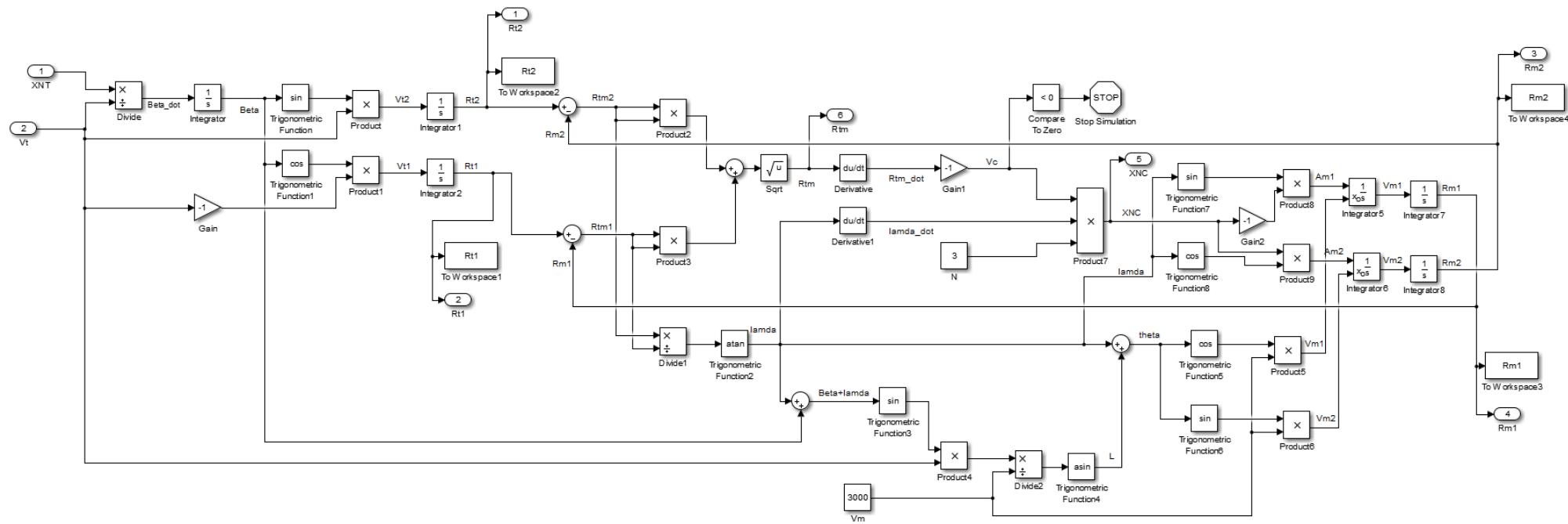


Figure 4.32: The Simulink model for proportional navigation equations in sec 3.1.1

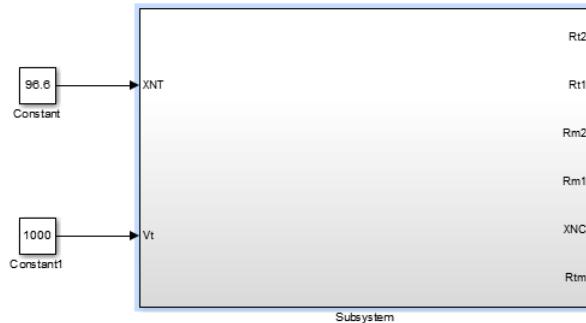


Figure 4.33: The Simulink main block for solving proportional navigation equations in sec 3.1.1

4.3.3 Results

We will use the main block we have created with the genetic algorithm toolbox in MATLAB in the same manner we have done, let the target maneuver unknown, then it is required to find it using genetic algorithm to maximize miss distance and missile acceleration. We have two cases of the target maneuver:

4.3.3.1 Polynomial Target Maneuver

let the target maneuver ba a polynomial with unknown coefficients, then it is required to find the values of these coefficients using genetic algorithm toolbox in Matlab (Fig. 4.34) to maximize miss distance and missile acceleration. The result for the best fitness values are in Fig.4.35. The optimized target acceleration and trajectory simulation are in figures 4.36 and 4.37.

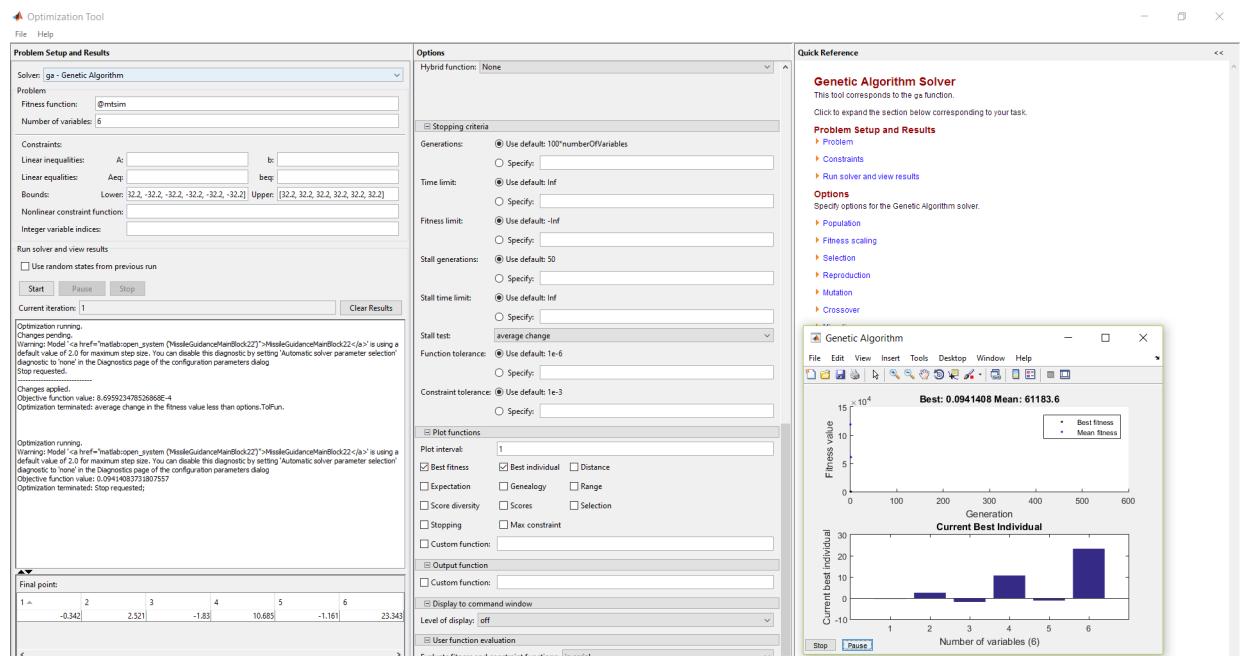


Figure 4.34: Genetic Algorithm Toolbox in Matlab.

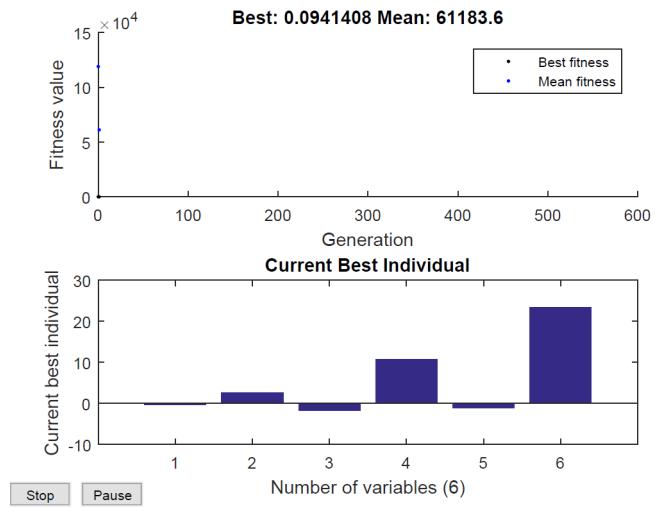


Figure 4.35: Best fitness values

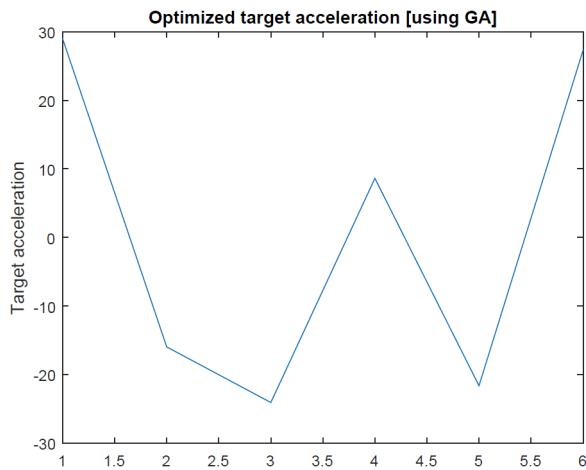


Figure 4.36: Optimized target acceleration using GA.

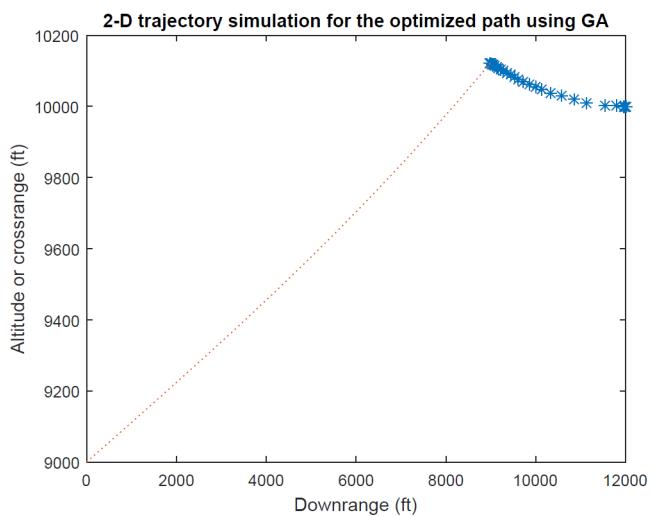


Figure 4.37: Optimized trajectory using GA.

4.3.3.2 Trapezoidal Target Maneuver

Chapter 5: Game

In this chapter we want to have the contribution of the human intelligence in making decisions, so we made a simulation game using *Unity*.

5.1 Introduction to Unity

Unity is a cross-platform game engine (Fig. 5.1) developed by Unity Technologies and used to develop video games for PC, consoles, mobile devices and websites. It is a free software (<https://unity3d.com/>), working with C# programming language.

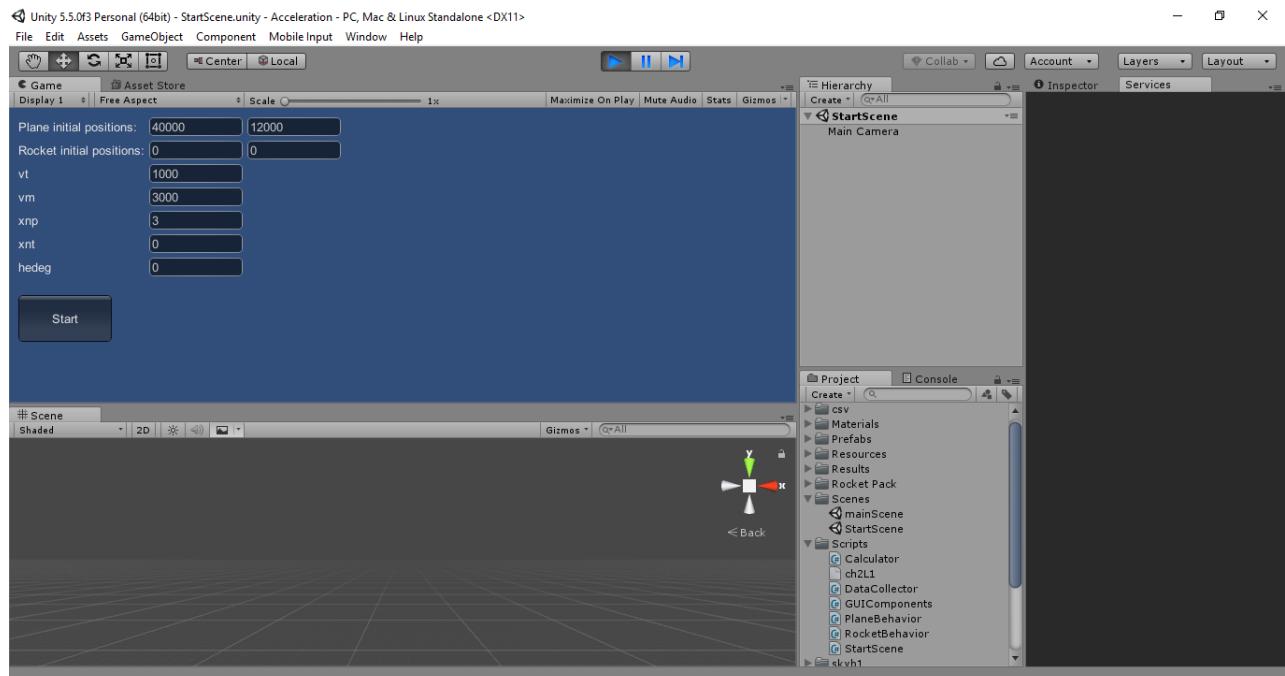


Figure 5.1: Unity game engine interface

5.2 Methodology of the game

The game (Fig. 5.2) consists of two objects: plane (Target) and missile (Attacker). A human player controls the increasing and decreasing of the target acceleration (XNT) by two arrows on the keyboard. The missile object is moving on according to the proportional navigation guidance law (in sec. 3.1).



Figure 5.2: Unity game for simulating Target-Attacker engagement

5.3 Results

Part III

TARGET-ATTACKER-DEFENDER PROBLEM

Chapter 6:

TARGET-ATTACKER-DEFENDER PROBLEM

6.1 Problem Statement

6.2 Assumptions, Notation, and Nomenclature

Assumptions

1. The speeds V_T , V_A , and V_D of the Target, Attacker, and Defender are constant.
2. The Attacker missile is faster than the Target aircraft, i.e., $\alpha \equiv \frac{V_T}{V_A} < 1$ (In the case $\alpha \geq 1$, the Target is guaranteed to survive by following an optimal strategy, even without assistance from the Defender).
3. There are three distinct cases for the ratio of the speed of the Attacker to that of the Defender $\gamma = \frac{1}{\beta} = \frac{V_A}{V_D}$
 - $\gamma < 1$ (fast Defender) discussed earlier in Garcia et el. [2], and extended, expounded, simplified and exposed herein.
 - $\gamma = 1$ (same speed Defender) discussed in Garcia et el. [1, 3], and extended, expounded, simplified and exposed herein.
 - $\gamma > 1$ (slow Defender), which is a novel case, discussed herein for the first time, and unified with the two previous cases.
4. The Defender intercepts the Attacker if their separation becomes zero (point capture).
5. The optimal trajectories of the three agents are straight lines.
6. A Cartesian frame is attached to the initial positions A and D of the Attacker and Defender in such a way the X -axis is the infinite extension of the straight segment \overline{AD} and the Y -axis is the perpendicular bisector of \overline{AD} .

Notation

- A : Initial position of the Attacker.
- D : Initial position of the Defender.
- T : Initial position of the Target.
- T' Terminal position of the Target, i.e., its position at the time the Defender intercepts the Attacker.

- $\alpha = \frac{V_T}{V_A}$ (assumed < 1 , otherwise the target trivially survives).
- $\gamma = \frac{1}{\beta} = \frac{V_A}{V_D}$ (studied for a fast Defender ($\gamma < 1$), a similar Defender ($\gamma = 1$), and a slow Defender ($\gamma > 1$)).
- u, v, w : Aim points on the AD Apollonius circle by the Attacker, Target and Defender, respectively.

Nomenclature

The AD Apollonius circle. The locus of a point such that the ratio of its distance from the initial positions $A = (x_A, 0)$ and $D = (-x_A, 0)$ of the Attacker and Defender, respectively, is a fixed ratio $\gamma = \frac{V_A}{V_D}$. This circle degenerates into the perpendicular bisector of \overline{AD} if $\gamma = 1$. Points on the circumference of this circle are reached simultaneously by the Attacker and Defender. For $\gamma < 1$, A belongs to the interior of this circle and points within the circle are reached by the Attacker before the Defender. For $\gamma > 1$, D belongs to the interior of this circle and points within the circle are reached by the Defender before the Attacker. For $\gamma = 1$, the AD circle becomes of infinite radius and degenerates into a straight line, namely the perpendicular bisector of \overline{AD} . In this case, A belongs to the R.H.S. of the XY -plane where points are reached by the Attacker before the Defender, while D belongs to the L.H.S. of the XY -plane where points are reached by the Defender before the Attacker.

The TA Apollonius circle. The locus of a point such that the ratio of its distances from the initial positions $T = (x_T, y_T)$ and $A = (x_A, 0)$, of the Target and Attacker, respectively, is a fixed ratio $\alpha = \frac{V_T}{V_A}$ strictly less than 1. Naming of the AD and TA Apollonius circles herein follows a common practice in mathematical circles [36, 37, 38] and is opposite to the style used by Garcia et al. [1, 2, 3].

Region reachable by the Defender before the Attacker (Reachability region R_r):

For $\gamma < 1$, R_r is the exterior of the AD Apollonius circle.

For $\gamma = 1$, R_r is the L.H.S. of the XY -plane.

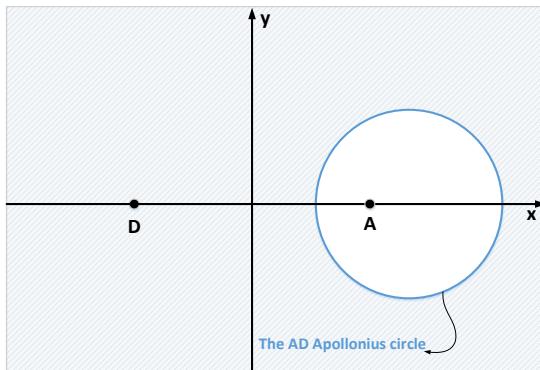
For $\gamma > 1$, R_r is the interior of the AD Apollonius circle.

Figure 6.1 illustrates this region as a shaded region to which point D belongs.

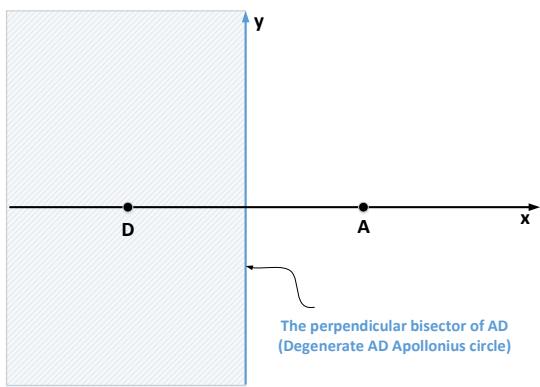
Region of guaranteed Target's escape (Escape region R_e). The escape region R_e is the set of all coordinate pairs (x, y) such that if the Target initial position $T = (x_T, y_T)$ is inside this region, then it is guaranteed to escape the Attacker if both the Target and Defender implement their corresponding optimal strategies. This set is a strict superset of the set of points reachable by the Defender before the Attacker ($R_e \supseteq R_r$). The boundary of this set is called a Voronoi Diagram.

The critical speed ratio $\bar{\alpha}$. A lower limit on the speed ratio $\gamma = \frac{V_T}{V_A}$, attained when the TA Apollonius circle is tangent to the AD Apollonius circle (or to the perpendicular bisector of \overline{AD} in the degenerate case $\gamma = 1$).

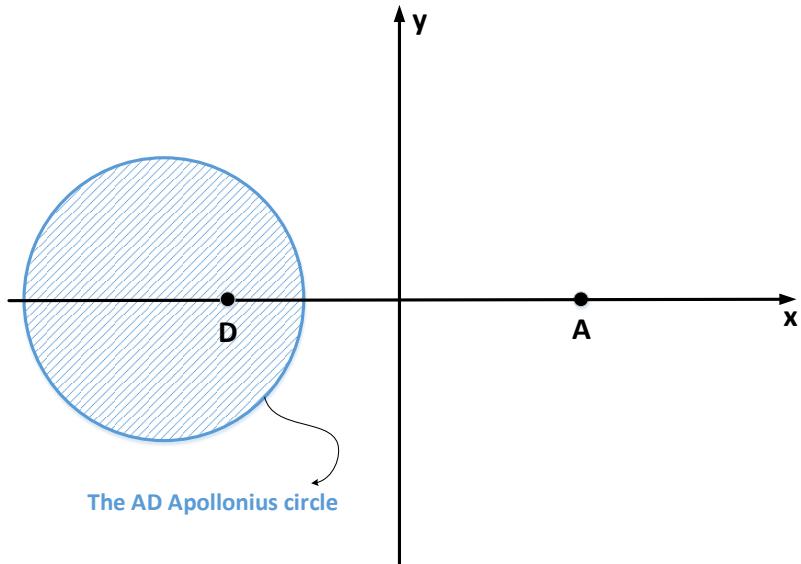
Escape Condition. For a given speed ratio $\alpha = \frac{V_T}{V_A}$ and a given Attacker's initial position $A = (x_A, 0)$, the escape condition is that the TA Apollonius circle (of center O_2 and radius r_2)



(a) $\gamma < 1$.



(b) $\gamma = 1$.



(c) $\gamma > 1$.

Figure 6.1: The reachability region R_r (one including D whose points are reached by the Defender before the Attacker) is shown shaded.

intersects the AD Apollonius circle (of center O_1 and radius r_1). This happens for $\gamma < 1$ when

$$|\mathbf{O}_1 - \mathbf{O}_2| + r_2 > r_1 \quad (6.1)$$

It happens for $\gamma > 1$ when

$$|\mathbf{O}_1 - \mathbf{O}_2| - r_2 < r_1 \quad (6.2)$$

In the degenerate case of $\gamma = 1$, it is required that the TA Apollonius circle intersects the Y -axis, i.e.,

$$r_2 > \text{the abscissa of } \mathbf{O}_2 \quad (6.3)$$

6.3 APOLLONIUS CIRCLES PERTAINING TO THE PROBLEM

A circle is the locus moving at a constant distance (called the circle's radius r) from a fixed point (called the circle's centre O). In the limit of an infinite radius ($r \rightarrow \infty$), the circle degenerates into a straight line. Another definition of the circle is that it is the locus of a point moving such that the ratio of its distances from two fixed points A and B is a constant k . With this definition, the circle is called a circle of Apollonius (in honour of Apollonius of Perga (ca. 262-190 BC), the Great Geometer of Antiquity). In the limit ($k \rightarrow 1$) this circle degenerates into the perpendicular bisector of \overline{AD} , while in the two limits ($k \rightarrow 0$) and ($k \rightarrow \infty$), this circle collapses to the two points A and B , respectively.

Two Apollonius circles are studied herein, namely the AD Apollonius circle, and the TA Apollonius circle. The study will be based on simple and intuitive plane-geometric arguments and will avoid the more involved treatment of analytic geometry. Figures 6.2 and 6.3 shows the Apollonius circles for a moving point P such that $\frac{AP}{PB} = k \neq 1$. The case $k > 1$ is shown in Fig. 6.2, while the case $k < 1$ is shwon in Fig. 6.3. The points I and E are the two special cases of P that lie on the straight line extension of the straight segment \overline{AB} . These points divide the straight segment \overline{AB} internally and externally in the ratio k ($k \neq 1$), i.e.,

$$\boxed{\frac{AI}{IB} = \frac{AE}{EB} = k, \{k \neq 1\}.} \quad (6.4)$$

If we denote the position vector of a point by a bold version of its name, we can rewrite (6.4) in vectorial form as

$$(\mathbf{I} - \mathbf{A}) = k(\mathbf{B} - \mathbf{I}), \quad (6.5)$$

$$(\mathbf{E} - \mathbf{A}) = k(\mathbf{B} - \mathbf{E}), \quad (6.6)$$

which can be used to express \mathbf{I} and \mathbf{E} in terms of \mathbf{A} and \mathbf{B} as

$$\mathbf{I} = \frac{1}{k+1}(\mathbf{A} + k\mathbf{B}), \quad (6.7)$$

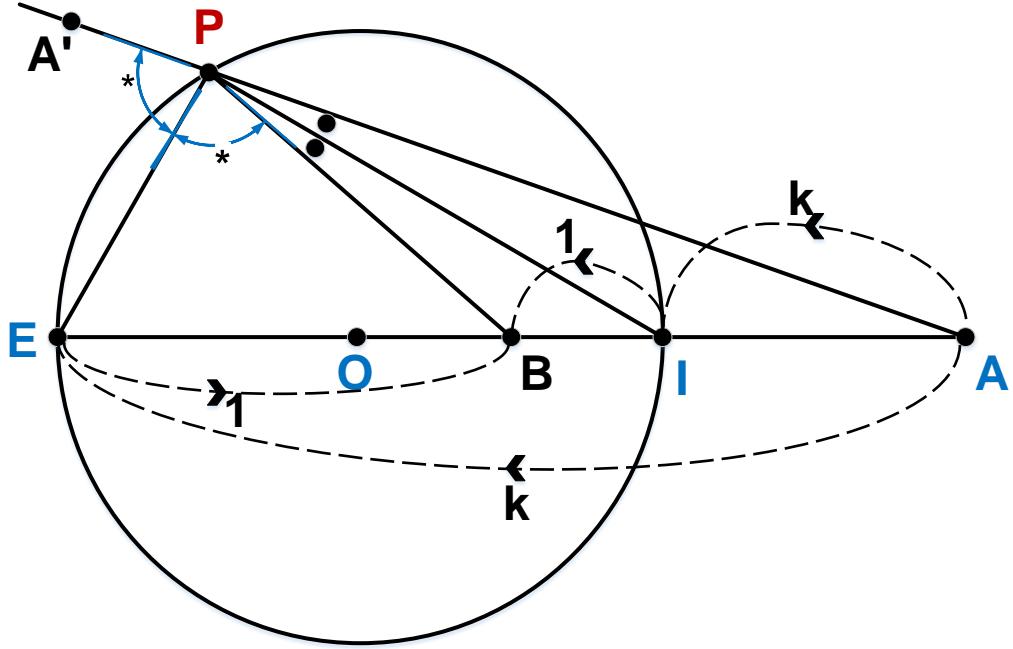


Figure 6.2: Apollonius circle for a moving point P such that $\frac{AP}{PB} = k > 1$. Here $m\angle API = m\angle BPI$ and $m\angle A'PE = m\angle BPE$.

$$\mathbf{E} = \frac{1}{k-1}(-\mathbf{A} + k\mathbf{B}). \quad (6.8)$$

The Apollonius circle is easily characterized by the points I and E , since they are the two end points of one of its diameters. The center of the circle is the midpoint of points I and E , namely:

$$\begin{aligned} \mathbf{O} &= \frac{1}{2}(\mathbf{I} + \mathbf{E}) \\ &= \frac{1}{2}\left[\left(\frac{1}{k+1} - \frac{1}{k-1}\right)\mathbf{A} + k\left(\frac{1}{k+1} + \frac{1}{k-1}\right)\mathbf{B}\right] \\ &= -\frac{1}{k^2-1}\mathbf{A} + \frac{k^2}{k^2-1}\mathbf{B}, \end{aligned} \quad (6.9)$$

while the radius of the circle is half the length of the displacement from I to E ,

$$\begin{aligned} r &= \frac{1}{2}|\mathbf{I} - \mathbf{E}| \\ &= \frac{1}{2}\left|\left(\frac{1}{k+1} + \frac{1}{k-1}\right)\mathbf{A} + k\left(\frac{1}{k+1} - \frac{1}{k-1}\right)\mathbf{B}\right| \\ &= \left|\frac{k}{k^2-1}\mathbf{A} - \frac{k}{k^2-1}\mathbf{B}\right| \\ &= \left|\frac{k}{k^2-1}\|\mathbf{A} - \mathbf{B}\| = \frac{k}{k^2-1}(AB)\right|. \end{aligned} \quad (6.10)$$

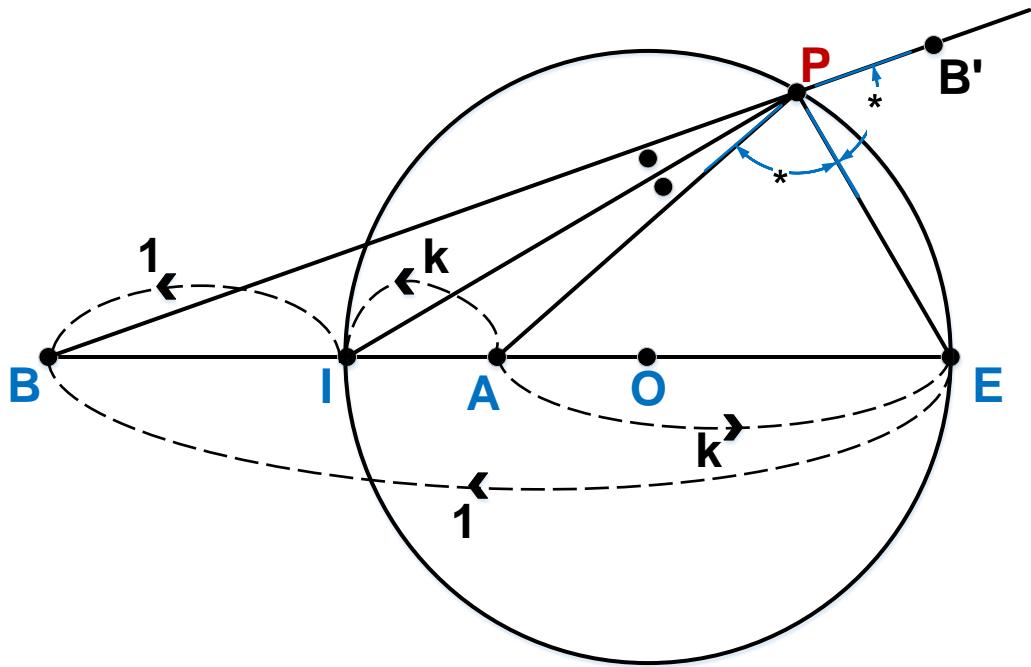


Figure 6.3: Apollonius circle for a moving point P such that $\frac{AP}{PB} = k < 1$. Here $m\angle API = m\angle BPI$ and $m\angle APE = m\angle B'PE$.

It is clear from (6.9) and (6.10) that $\lim_{k \rightarrow 1} |O| \rightarrow \infty$ and $\lim_{k \rightarrow 1} r \rightarrow \infty$, and hence for $k = 1$, the Apollonius circle degenerates into a straight line, namely the perpendicular bisector of the straight segment \overline{AB} (Fig. 6.4).

6.3.1 The AD Apollonius circle

For the *AD* Apollonius circle, the two fixed points are the initial positions of the Attacker $A = (x_A, 0)$ and the initial position of the defender $D = (-x_A, 0)$. The fixed ratio of the circle k is replaced by the following dimensionless ratio which is the Attacker's speed normalized w.r.t the Defender speed:

$$\gamma = \frac{V_A}{V_D}. \quad (6.11)$$

We will consider the three cases of

1. $\gamma < 1$ (fast Defender) discussed in Garcia et al. [2].
2. $\gamma = 1$ (same-speed Defender) discussed in Garcia et al. [1, 3].
3. $\gamma > 1$ (slow Defender), which is a novel case.

Substituting the values of A and D above for A and B in (6.7),(6.8),(6.9) and (6.10), respectively, and replacing k therein by γ we obtain for $\gamma \neq 1$,

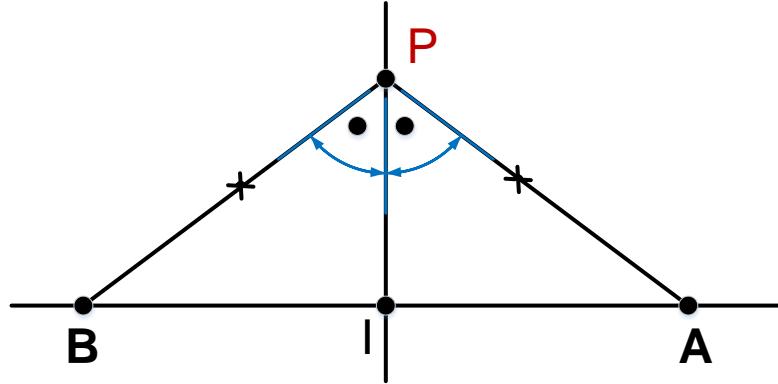


Figure 6.4: For $k = 1$, the Apollonius circle in figures 6.2 or 6.3 degenerates into the perpendicular bisector of the straight segment \overline{AB} . The point E disappears in this figure as it goes to ∞

$$\mathbf{I}_1 = \left(\frac{1-\gamma}{1+\gamma} x_A, 0 \right), \quad (6.12)$$

$$\mathbf{E}_1 = \left(\frac{1+\gamma}{1-\gamma} x_A, 0 \right), \quad (6.13)$$

$$\mathbf{O}_1 = \left(\frac{1+\gamma^2}{1-\gamma^2} x_A, 0 \right), \quad (6.14)$$

$$r_1 = \frac{2\gamma}{|1-\gamma^2|} x_A. \quad (6.15)$$

Three interesting sets of limits are now considered. The first set of limits are those when ($\gamma \rightarrow 1$), namely

$$\lim_{\gamma \rightarrow 1} \mathbf{I}_1 = (0, 0), \quad (6.16)$$

$$\lim_{\gamma \rightarrow 1} \mathbf{E}_1 = (\pm\infty, 0), \quad (6.17)$$

$$\lim_{\gamma \rightarrow 1} \mathbf{O}_1 = (\pm\infty, 0), \quad (6.18)$$

$$\lim_{\gamma \rightarrow 1} r_1 = \infty, \quad (6.19)$$

which means that the AD Apollonius circle degenerates in the limit ($\gamma \rightarrow 1$) to a circle of an infinite radius with center still on the x -axis but infinitely distant from the origin. More precisely, this limit is identified as a straight line, namely the perpendicular bisector of the straight-line segment \overline{AD} .

The second set of limits are those when ($\gamma \rightarrow 0$), namely

$$\lim_{\gamma \rightarrow 0} \mathbf{I}_1 = (x_A, 0), \quad (6.20)$$

$$\lim_{\gamma \rightarrow 0} \mathbf{E}_1 = (x_A, 0), \quad (6.21)$$

$$\lim_{\gamma \rightarrow 0} \mathbf{O}_1 = (x_A, 0), \quad (6.22)$$

$$\lim_{\gamma \rightarrow 0} r_1 = 0, \quad (6.23)$$

which means that the *AD* Apollonius circle collapses in the limit ($\gamma \rightarrow 0$) to a single point, namely the initial position of the Attacker $\mathbf{A} = (x_A, 0)$.

The third set of limits are those when ($\gamma \rightarrow \infty$), namely

$$\lim_{\gamma \rightarrow \infty} \mathbf{I}_1 = (-x_A, 0), \quad (6.24)$$

$$\lim_{\gamma \rightarrow \infty} \mathbf{E}_1 = (-x_A, 0), \quad (6.25)$$

$$\lim_{\gamma \rightarrow \infty} \mathbf{O}_1 = (-x_A, 0), \quad (6.26)$$

$$\lim_{\gamma \rightarrow \infty} r_1 = 0, \quad (6.27)$$

which means that the *AD* Apollonius circle collapses in the limit ($\gamma \rightarrow \infty$) to a single point, namely the initial position of the Defender $\mathbf{D} = (-x_A, 0)$.

6.3.2 The *TA* Apollonius circle

For the *TA* Apollonius circle, the two fixed points are the initial position of the Target $\mathbf{T} = (x_T, y_T)$ and the initial position of the Attacker $\mathbf{A} = (x_A, 0)$. Again, the fixed ratio of the circle k is replaced by a dimensionless quantity, namely the Target's speed normalized w.r.t. the Attacker's speed:

$$\alpha = \frac{V_T}{V_A}. \quad (6.28)$$

Here, we will consider only the case $\alpha < 1$, since for $\alpha \geq 1$, the Target always survives, even without any assistance from the Defender.

Now, we substitute the values of \mathbf{T} and \mathbf{A} above for \mathbf{A} and \mathbf{B} in (6.7),(6.8),(6.9) and (6.10), respectively, and replace k therein by α to obtain:

$$\mathbf{I}_2 = \frac{1}{1+\alpha}(\mathbf{T} + \alpha\mathbf{A}) = \frac{1}{1+\alpha}(x_T + \alpha x_A, y_T), \quad (6.29)$$

$$\mathbf{E}_2 = \frac{1}{\alpha-1}(-\mathbf{T} + \alpha\mathbf{A}) = \frac{1}{1-\alpha}(x_T - \alpha x_A, y_T), \quad (6.30)$$

$$\mathbf{O}_2 = \frac{1}{1-\alpha^2}(\mathbf{T} - \alpha^2\mathbf{A}) = \frac{1}{1-\alpha^2}(x_T - \alpha^2 x_A, y_T), \quad (6.31)$$

$$r_2 = \frac{\alpha}{1-\alpha^2}(TA) = \frac{\alpha d}{1-\alpha^2}, \quad (6.32)$$

where $d = TA$ = the initial distance between the target and Attacker, namely

$$d = \sqrt{(x_T - x_A)^2 + y_T^2}. \quad (6.33)$$

Chapter 7: CRITICAL SPEED RATIO

Target survival is guaranteed if there is an overlapping of the region reachable by the Target before the Attacker (the interior of the TA Apollonius circle) and the region R_r reachable by the Defender before the Attacker (Fig. 6.1), since within this overlapping, the Defender can perform its intended role of intercepting the Attacker before the Attacker captures the Target. Target survival is critical when the aforementioned overlapping diminishes into a single point at which the aforementioned two regions barely touch, or are tangent to one another. For this situation, the normalized Target speed (the speed ratio) $\alpha = \frac{V_T}{V_A}$ attains its minimal or critical value $\bar{\alpha}$. Now, we consider three cases, in which the TA -Apollonius circle is tangent from outside to the boundary of the shaded region R_r . An implicit assumption throughout the forthcoming analysis is that $\mathbf{T} = (x_T, Y_T)$ is outside R_r .

7.1 The case $\gamma < 1$ (fast defender)

The critical speed ratio $\bar{\alpha}$, occurs when the TA Apollonius circle is *internally* tangent to the AD Apollonius circle, i.e., when the centers \mathbf{O}_2 and \mathbf{O}_1 of these two circles and their tangency point \mathbf{C} are collinear (Fig. ??). This happens when

$$r_1 - r_2 = |\mathbf{O}_1 - \mathbf{O}_2|. \quad (7.1)$$

Substituting for r_1, r_2, \mathbf{O}_1 and \mathbf{O}_2 from (6.15), (6.32), (6.15) and (6.31) respectively, and noting that $\gamma < 1$, one obtains

$$\begin{aligned} \frac{2\gamma}{1-\gamma^2}x_A - \frac{\alpha}{1-\alpha^2}d &= |(\frac{1+\gamma^2}{1-\gamma^2}x_A, 0) - \frac{1}{1-\alpha^2}(x_T - \alpha^2 x_A, y_T)| \\ &= [(\frac{1+\gamma^2}{1-\gamma^2}x_A - \frac{1}{1-\alpha^2}(x_T - \alpha^2 x_A))^2 + \frac{1}{(1-\alpha^2)^2}y_T^2]^{\frac{1}{2}}. \end{aligned} \quad (7.2)$$

In (7.2), we deliberately replaced $|1 - \gamma^2|$ by $(1 - \gamma^2)$ because $\gamma < 1$. We now square both sides of (7.2) to obtain:

$$\begin{aligned} \frac{4\gamma^2}{(1-\gamma^2)^2}x_A^2 + \frac{\alpha^2}{(1-\alpha^2)^2}d^2 - \frac{4\gamma ad}{(1-\gamma^2)(1-\alpha^2)}x_A &= \\ = \frac{(1+\gamma^2)^2}{(1-\gamma^2)^2}x_A^2 + \frac{x_T^2 - 2\alpha^2 x_T x_A + \alpha^4 x_A^2}{(1-\alpha^2)^2} - \frac{2(1+\gamma^2)(x_T - \alpha^2 x_A)x_A}{(1-\gamma^2)(1-\alpha^2)} + \frac{1}{(1-\alpha^2)^2}y_T^2. \end{aligned} \quad (7.3)$$

By squaring both sides of equation (7.2), the cardinality of the solution set for equation (7.2) is doubled. This means that when we solve the resulting equation, half of the resulting solutions will be extraneous or irrelevant and have to be rejected. Fortunately, we will have genuine reasons that enable us to identify such solutions, and compel us to reject them. The solutions retained after such a rejection are the correct solutions of (7.1) We now use (6.33) to replace y_T^2 by $(d^2 - x_A^2 - x_T^2 - 2x_A x_T)$ and rearrange terms in (7.3) to obtain

$$\begin{aligned}
& \left[\frac{4\gamma^2}{(1-\gamma^2)^2} - \frac{(1+\gamma^2)^2}{(1-\gamma^2)^2} - \frac{\alpha^4}{(1-\alpha^2)^2} - \frac{2\alpha^2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} + \frac{1}{(1-\alpha^2)^2} \right] x_A^2 \\
& \quad + \left[\frac{\alpha^2}{(1-\alpha^2)^2} - \frac{1}{(1-\alpha^2)^2} \right] d^2 \\
& \quad - \frac{4\gamma\alpha x_A d}{(1-\gamma^2)(1-\alpha^2)} + \left[\frac{-1}{(1-\alpha^2)^2} + \frac{1}{(1-\alpha^2)^2} \right] x_T^2 \\
& \quad + \left[\frac{2\alpha^2}{(1-\alpha^2)^2} + \frac{2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} - \frac{2}{(1-\alpha^2)^2} \right] x_T x_A = 0
\end{aligned} \tag{7.4}$$

The coefficient of x_A^2 in (7.4) is

$$\begin{aligned}
& \frac{4\gamma^2}{(1-\gamma^2)^2} - \frac{(1+\gamma^2)^2}{(1-\gamma^2)^2} - \frac{\alpha^4}{(1-\alpha^2)^2} - \frac{2\alpha^2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} + \frac{1}{(1-\alpha^2)^2} \\
& = \frac{4\gamma^2 - 1 - \gamma^4 - 2\gamma^2}{(1-\gamma^2)^2} + \frac{(1-\alpha^4)}{(1-\alpha^2)^2} - \frac{2\alpha^2(1-\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} \\
& = -1 + \frac{1+\alpha^2}{(1-\alpha^2)} - \frac{2\alpha^2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} \\
& = \frac{2\alpha^2}{(1-\alpha^2)} - \frac{2\alpha^2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} \\
& = \frac{2\alpha^2}{(1-\alpha^2)} \left[1 - \frac{(1+\gamma^2)}{(1-\gamma^2)} \right] = \frac{2\alpha^2}{(1-\alpha^2)} \left(\frac{-2\gamma^2}{1-\gamma^2} \right) \\
& = \frac{-4\alpha^2\gamma^2}{(1-\alpha^2)(1-\gamma^2)},
\end{aligned} \tag{7.5}$$

and the coefficient of d^2 in (7.4) is

$$\frac{\alpha^2}{(1-\alpha^2)^2} - \frac{1}{(1-\alpha^2)^2} = \frac{-(1-\alpha^2)}{(1-\alpha^2)^2} = -\frac{1}{(1-\alpha^2)}, \tag{7.6}$$

while the coefficient of $x_T x_A$ in (7.4) is

$$\begin{aligned}
& \frac{2\alpha^2}{(1-\alpha^2)^2} + \frac{2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} - \frac{2}{(1-\alpha^2)^2} \\
& = \frac{2(\alpha^2 - 1)}{(1-\alpha^2)^2} + \frac{2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} \\
& = -\frac{2}{(1-\alpha^2)} + \frac{2(1+\gamma^2)}{(1-\gamma^2)(1-\alpha^2)} \\
& = \frac{2}{(1-\alpha^2)} \left[-1 + \left(\frac{1+\gamma^2}{1-\gamma^2} \right) \right] = \frac{2}{(1-\alpha^2)} \left(\frac{2\gamma^2}{1-\gamma^2} \right) \\
& = \frac{4\gamma^2}{(1-\alpha^2)(1-\gamma^2)}
\end{aligned} \tag{7.7}$$

This means that equation (7.4) can be simplified considerably to the equivalent form

$$\frac{-4\alpha^2\gamma^2}{(1-\alpha^2)(1-\gamma^2)}x_A^2 - \frac{1}{1-\alpha^2}d^2 - \frac{4\gamma\alpha x_A d}{(1-\gamma^2)(1-\alpha^2)} + \frac{4\gamma^2}{(1-\alpha^2)(1-\gamma^2)}x_T x_A = 0 \quad (7.8)$$

Now, we multiply (7.8) by $[-(1-\alpha^2)(1-\gamma^2)]$ noting that $\alpha \neq 1$ and $\gamma \neq 1$, to obtain

$$(4\gamma^2 x_A^2)\alpha^2 + (4\gamma x_A d)\alpha + (1-\gamma^2)d^2 - 4\gamma^2 x_T x_A = 0 \quad (7.9)$$

or equivalently

$$a^2 + \frac{d}{\gamma x_A}\alpha + \left[\left(\frac{1-\gamma^2}{4\gamma^2} \right) \left(\frac{d}{x_A} \right)^2 - \frac{x_T}{x_A} \right] = 0 \quad (7.10)$$

Equation (7.10) is an improvement of equation (36) in [3], while (7.10) is a quadratic equation in α , equation (36) in [3] is a quartic in α that has the same two roots of (7.10) in addition to the two irrelevant roots $\alpha \mp 1$. As a bonus, equation (7.10) is written in a self-verifying dimensionless form.

7.2 The case $\gamma = 1$ (similar Defender)

The critical speed ratio $\bar{\alpha}$ occurs when the TA Apollonius circle touches the *L.H.S.* of the XY -plane (Fig. ??), i.e., when

$$r_2 = \text{Abscissa of } \mathbf{O}_2. \quad (7.11)$$

i.e., thanks to (6.31) and (6.32) when:

$$\frac{\alpha d}{1-\alpha^2} = \frac{x_T - \alpha^2 x_A}{1-\alpha^2}. \quad (7.12)$$

Multiplying (7.12) by $(1-\alpha^2)$ (thanks to the fact that $\alpha \neq 1$), and rearranging, one obtains the following quadratic equation in α

$$x_A \alpha^2 + d\alpha - x_T = 0, \quad (7.13)$$

which is equation (12) in [3]. It is interesting to note that equation (7.12) is the special case ($\gamma = 1$) of (7.10), though (7.10) was derived under the assumption $\gamma \neq 1$.

7.3 The case $\gamma > 1$ (Slow Defender)

The critical speed ratio $\bar{\alpha}$ occurs when the TA Apollonius circle is *externally tangent* to the AD Apollonius circle, i.e., when the centre \mathbf{O}_2 , the tangency point \mathbf{C} and the centre \mathbf{O}_1 are collinear (Fig. ??), i.e., when

$$r_1 + r_2 = |\mathbf{O}_1 - \mathbf{O}_2| \quad (7.14)$$

Since $\gamma > 1$, we write

$$r_1 + r_2 = \frac{2\gamma}{|1-\gamma^2|} + \frac{\alpha}{1-\alpha^2}d = \frac{-2\gamma}{1-\gamma^2} + \frac{\alpha}{1-\alpha^2}d. \quad (7.15)$$

The above expression for $(r_1 + r_2)$ is exactly the negative of $(r_1 - r_2)$ in (7.2). Hence, squaring of (7.14) results in (7.3), and consequently lead to (7.4) and (7.8), and finally to (7.9) and (7.10). Equation (7.10) is thereby shown to hold for a slow Defender besides being true for a fast one.

Therefore, we will use the quadratic formula (7.10) for all values of γ , and will solve it to obtain $\bar{\alpha}$ for any γ as

$$\begin{aligned}\bar{\alpha} &= \frac{1}{2} \left[-\frac{d}{\gamma x_A} \pm \left[\left(\frac{d}{\gamma x_A} \right)^2 - 4 \left(\frac{1-\gamma^2}{4\gamma^2} \right) \left(\frac{d}{x_A} \right)^2 + \frac{4x_T}{x_A} \right]^{1/2} \right] \\ &= \frac{1}{2\gamma x_A} [-d \mp \gamma \sqrt{d^2 + 4x_T x_A}] \\ &= \frac{1}{2\gamma x_A} [-\sqrt{(x_A - x_T)^2 + y^2} + \gamma \sqrt{(x_A + x_T)^2 + y_T^2}]\end{aligned}\quad (7.16)$$

Equation (7.16) was derived in [3] for $\gamma < 1$ and in [1] for $\gamma = 1$, provided $x_T > 0$. It is shown herein to hold for all values in γ , including $\gamma > 1$. Note that $\bar{\alpha}$ is definitely nonnegative and hence the minus sign in the solution for $\bar{\alpha}$ is rejected. The question remains whether (7.16) yields a non-negative value for $\bar{\alpha}$ when the positive sign is chosen in (7.16). It could be argued that (7.16) should be used (with the positive sign, of course) as long as it yields a nonnegative value for $\bar{\alpha}$, and otherwise $\bar{\alpha}$ should be set to zero. For example, if $\gamma = 1$ and $x_T < 0$, equation (7.16) yields a negative value for $\bar{\alpha}$ which need to be rejected and replaced by zero. Physically, the situation $\gamma = 1, x_T < 0$ is one in which the Target is initially in the region reachable by the Defender before the Attacker, i.e., it is within the effective protection of the Defender, even if it does not move itself at all. Therefor, we rewrite (7.16) in the following dimensionless form:

$$\bar{\alpha} = \frac{1}{2\gamma} \left[-\sqrt{\left(1 - \frac{x_T}{x_A}\right)^2 + \left(\frac{y_T}{x_A}\right)^2} + \gamma \sqrt{\left(1 + \frac{x_T}{x_A}\right)^2 + \left(\frac{y_T}{x_A}\right)^2} \right] \quad (7.17)$$

provided that

$$\gamma \sqrt{\left(1 + \frac{x_T}{x_A}\right)^2 + \left(\frac{y_T}{x_A}\right)^2} > \sqrt{\left(1 - \frac{x_T}{x_A}\right)^2 + \left(\frac{y_T}{x_A}\right)^2} \quad (7.18)$$

and whenever (7.18) is not valid, we simply set

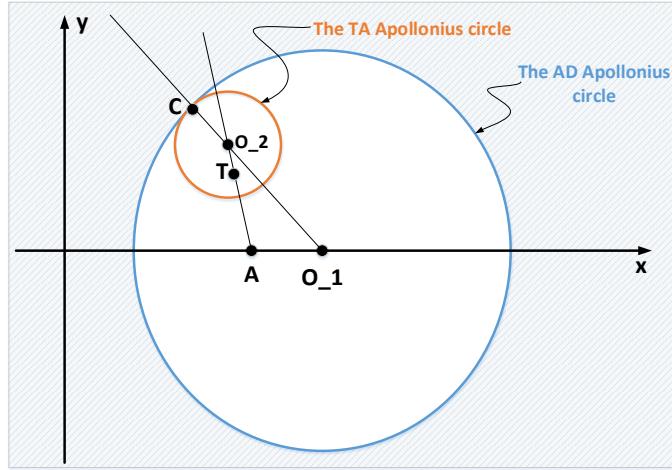
$$\bar{\alpha} = 0. \quad (7.19)$$

Figures 7.2 to 7.4 plot α as a function of $(\frac{x_T}{x_A})$ with $(\frac{y_T}{x_A})$ as a parameter for the specific γ values of 0.8, 1.0, and 1.25 respectively.

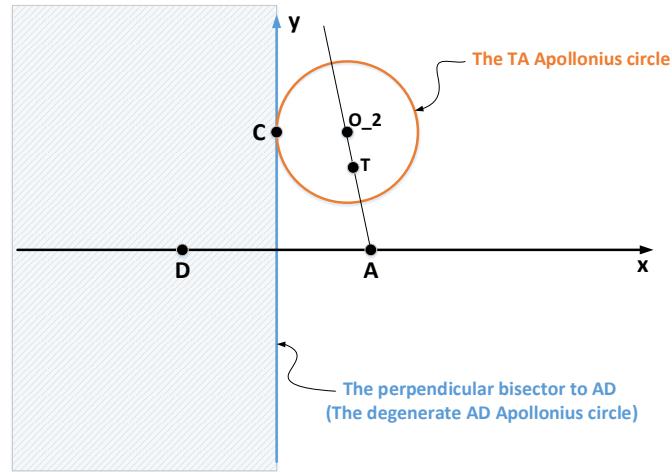
We pay a special attention to the special case when the Target is initially on the x -axis (i.e., when the Target, Attacker, and Defender are initially collinear). In this case, equation (7.17) reduces to

$$\bar{\alpha} = \frac{1}{2\gamma} \left[-\sqrt{\left(1 - \frac{x_T}{x_A}\right)^2} + \gamma \sqrt{\left(1 + \frac{x_T}{x_A}\right)^2} \right], \quad (7.20)$$

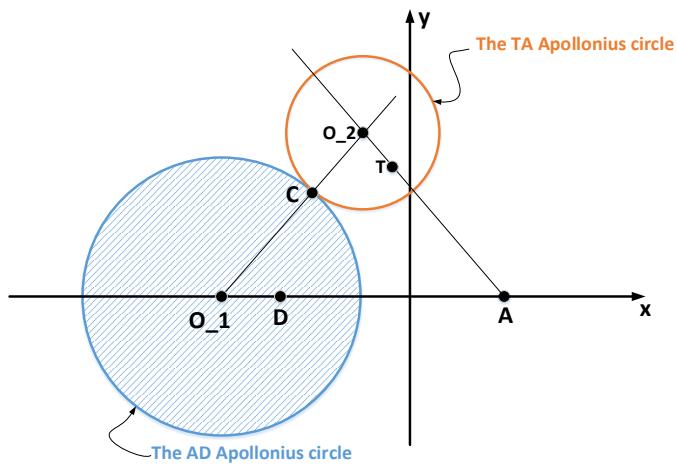
care must be taken that the $\sqrt{}$ sign indicates the principal branch of the inverse square



(a) $\gamma < 1$



(b) $\gamma = 1$



(c) $\gamma > 1$

Figure 7.1: The critical speed ratio $\bar{\alpha}$ is obtained when the TA Apollonius circle is tangent to the boundary of the shaded region R_r which is the region reachable by the Defender before the Attacker.

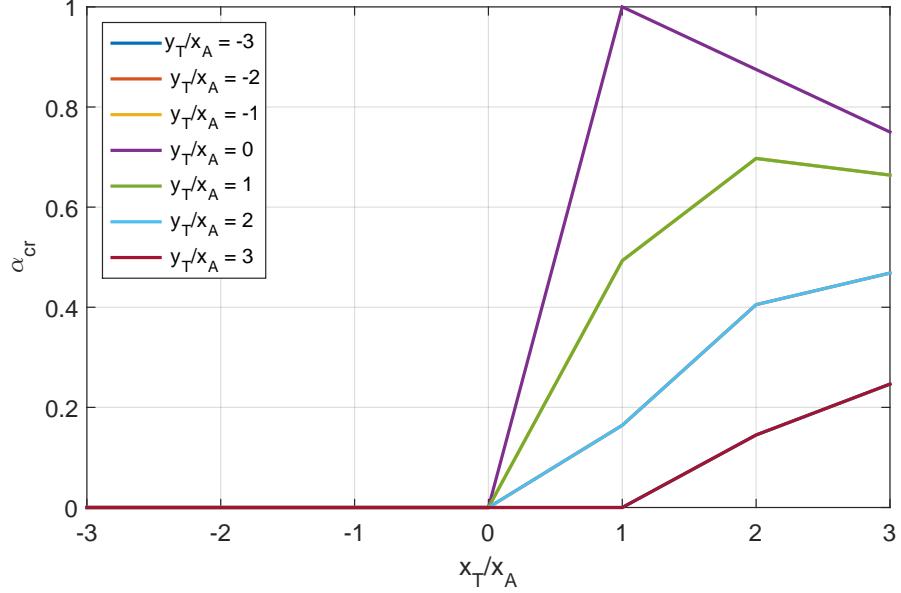


Figure 7.2: The critical normalized speed $\bar{\alpha}$ or α_{cr} as a function of $(\frac{x_T}{x_A})$ with $(\frac{y_T}{x_A})$ as a parameter when $\gamma = 0.8$ (fast defender).

function. i.e.,

$$\sqrt{y^2} = |y| \quad (7.21)$$

Therefore, we write

$$\sqrt{(1 - \frac{x_T}{x_A})^2} = \begin{cases} 1 - \frac{x_T}{x_A}, & \frac{x_T}{x_A} \leq 1 \\ \frac{x_T}{x_A} - 1, & \frac{x_T}{x_A} > 1, \end{cases} \quad (7.22)$$

$$\sqrt{(1 + \frac{x_T}{x_A})^2} = \begin{cases} 1 + \frac{x_T}{x_A}, & \frac{x_T}{x_A} \geq -1 \\ -1 - \frac{x_T}{x_A}, & \frac{x_T}{x_A} < -1, \end{cases} \quad (7.23)$$

We now consider the three specific values of γ considered in Figs. 7.2 to 7.4 , namely $\gamma = 0.8$, $\gamma = 1$, and $\gamma = 1.25$.

For $\gamma = 0.8$

$$\bar{\alpha} = \frac{1}{1.6} \left[-\sqrt{(1 - \frac{x_T}{x_A})^2} + 0.8 \sqrt{(1 + \frac{x_T}{x_A})^2} \right]. \quad (7.24)$$

For $(\frac{x_T}{x_A}) < -1$

$$\bar{\alpha} = \frac{1}{1.6} \left[-(1 - \frac{x_T}{x_A}) + 0.8(-1 - \frac{x_T}{x_A}) \right] = \frac{1}{1.6} \left[-1.8 + 0.2 \frac{x_T}{x_A} \right] = \frac{1}{8} \left(-9 + \frac{x_T}{x_A} \right)$$

which is definitely negative, and hence

$$\bar{\alpha} = 0, \quad (\frac{x_T}{x_A}) < -1, \quad \gamma = 0.8 \quad (7.25)$$

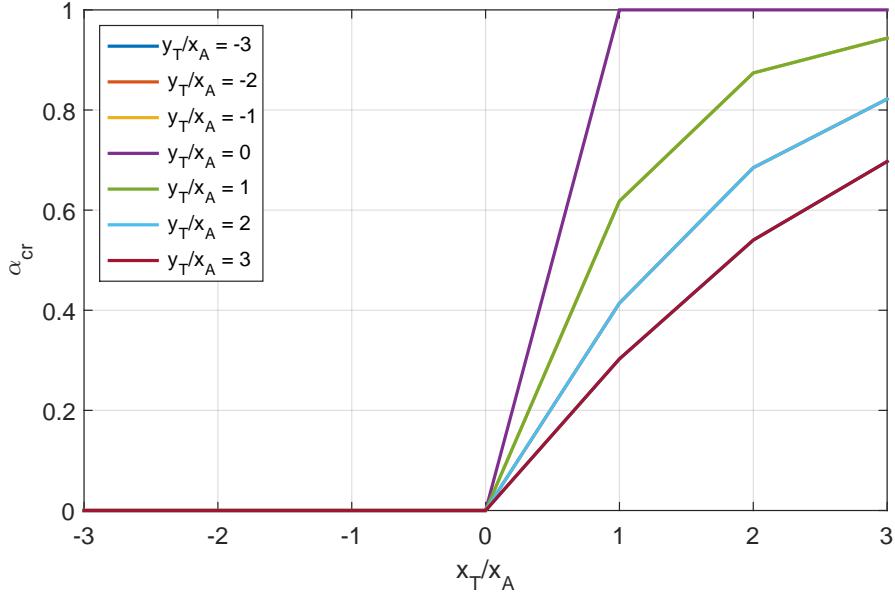


Figure 7.3: The critical normalized speed $\bar{\alpha}$ or α_{cr} as a function of $(\frac{x_T}{x_A})$ with $(\frac{y_T}{x_A})$ as a parameter when $\gamma = 1$ (similar defender).

For $-1 \leq (\frac{x_T}{x_A}) \leq 1$

$$\bar{\alpha} = \frac{1}{1.6} [-(1 - \frac{x_T}{x_A}) + 0.8(1 + \frac{x_T}{x_A})] = \frac{1}{1.6} [-0.2 + 1.8 \frac{x_T}{x_A}] = \frac{1}{8} (-1 + 9 \frac{x_T}{x_A})$$

The above $\bar{\alpha}$ is negative when $(-1 + 9 \frac{x_T}{x_A})$ is negative and hence we again set

$$\bar{\alpha} = 0, \quad -1 \leq (\frac{x_T}{x_A}) < \frac{1}{9}, \quad \gamma = 0.8. \quad (7.26)$$

while $\bar{\alpha}$ increases linearly from 0 to 1 in $\frac{x_T}{x_A} \in [\frac{1}{9}, 1]$

$$\bar{\alpha} = \frac{1}{8} (-1 + 9 \frac{x_T}{x_A}), \quad \frac{1}{9} \leq (\frac{x_T}{x_A}) \leq 1, \quad \gamma = 0.8 \quad (7.27)$$

Finally, when $(\frac{x_T}{x_A}) > 1$

$$\bar{\alpha} = \frac{1}{1.6} [-(\frac{x_T}{x_A} - 1) + 0.8(\frac{x_T}{x_A} + 1)] = \frac{1}{1.6} [1.8 - 0.2(\frac{x_T}{x_A})] = \frac{1}{8} (9 - \frac{x_T}{x_A}). \quad (7.28)$$

This means that after peaking at 1 when $(\frac{x_T}{x_A}) = 1$, the value of $\bar{\alpha}$ starts to decrease linearly to be 0 at $(\frac{x_T}{x_A}) = 9$ and then it remains 0 (rather than becoming negative), i.e.,

$$\bar{\alpha} = \frac{1}{8} (9 - \frac{x_T}{x_A}), \quad 1 < (\frac{x_T}{x_A}) \leq 9, \quad \gamma = 0.8, \quad (7.29)$$

$$\bar{\alpha} = 0, \quad (\frac{x_T}{x_A}) > 9, \quad \gamma = 0.8. \quad (7.30)$$

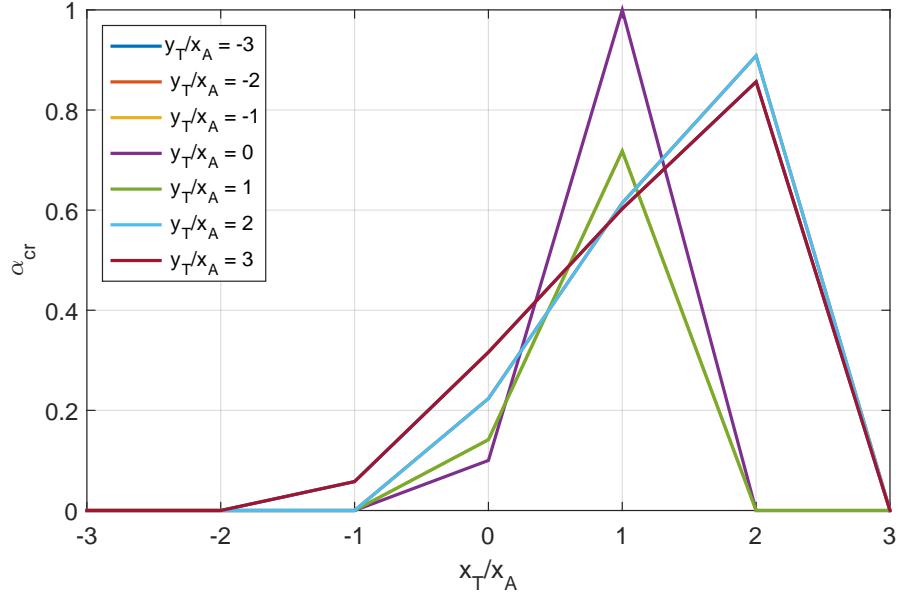


Figure 7.4: The critical normalized speed $\bar{\alpha}$ or α_{cr} as a function of $(\frac{x_T}{x_A})$ with $(\frac{y_T}{x_A})$ as a parameter when $\gamma = 1.25$ (slow defender).

For $\gamma = 1$

$$\bar{\alpha} = \frac{1}{2} \left[-\sqrt{(1 - \frac{x_T}{x_A})^2} + \sqrt{(1 + \frac{x_T}{x_A})^2} \right]. \quad (7.31)$$

For $(\frac{x_T}{x_A}) < -1$

$$\bar{\alpha} = \frac{1}{2} \left[-(1 - \frac{x_T}{x_A}) + (-1 - \frac{x_T}{x_A}) \right] = -1.$$

Set

$$\bar{\alpha} = 0, \quad (\frac{x_T}{x_A}) < -1, \quad \gamma = 1 \quad (7.32)$$

For $-1 \leq (\frac{x_T}{x_A}) \leq 1$

$$\begin{aligned} \bar{\alpha} &= \frac{1}{2} \left[-(1 - \frac{x_T}{x_A}) + (1 + \frac{x_T}{x_A}) \right] \\ \bar{\alpha} &= (\frac{x_T}{x_A}), \quad -1 \leq (\frac{x_T}{x_A}) \leq 1, \quad \gamma = 1 \end{aligned} \quad (7.33)$$

For $-1 \leq (\frac{x_T}{x_A}) < 0$, $\bar{\alpha}$ is again negative and is to be set to 0

$$\bar{\alpha} = (\frac{x_T}{x_A}), \quad 0 \leq (\frac{x_T}{x_A}) \leq 1, \quad \gamma = 1 \quad (7.34)$$

and hence $\bar{\alpha}$ increases linearly from 0 at $(\frac{x_T}{x_A}) = 0$ to 1 at $(\frac{x_T}{x_A}) = 1$

For $(\frac{x_T}{x_A}) > 1$

$$\bar{\alpha} = \frac{1}{2} \left[-\left(\frac{x_T}{x_A} - 1 \right) + \left(1 + \frac{x_T}{x_A} \right) \right] = 1, \quad \left(\frac{x_T}{x_A} \right) > 1, \quad \gamma = 1 \quad (7.35)$$

Note that the situation of ($\gamma = 1$) differs significantly from that of ($\gamma = 0.8$). For ($\gamma = 1$), $\bar{\alpha}$ increases from 0 to 1 and remains so indefinitely, but for ($\gamma = 0.8$), $\bar{\alpha}$ increases from 0 to 1 and immediately decreases again to 0.

For $\gamma = 1.25$

$$\bar{\alpha} = \frac{1}{2.5} \left[-\sqrt{\left(1 - \frac{x_T}{x_A} \right)^2} + 1.25 \sqrt{\left(1 + \frac{x_T}{x_A} \right)^2} \right] \quad (7.36)$$

For $\left(\frac{x_T}{x_A} \right) < -1$

$$\bar{\alpha} = \frac{1}{2.5} \left[-\left(1 - \frac{x_T}{x_A} \right) + 1.25 \left(-1 - \frac{x_T}{x_A} \right) \right] = \frac{1}{2.5} \left[-2.25 - 0.25 \frac{x_T}{x_A} \right] = \frac{1}{10} \left(-9 - \frac{x_T}{x_A} \right),$$

which results in setting

$$\bar{\alpha} = 0, \quad -9 \leq \left(\frac{x_T}{x_A} \right) < -1, \quad \gamma = 1.25, \quad (7.37)$$

$$\bar{\alpha} = \frac{1}{10} \left(-9 - \frac{x_T}{x_A} \right), \quad \left(\frac{x_T}{x_A} \right) < -9, \quad \gamma = 1.25. \quad (7.38)$$

For $-1 \leq \left(\frac{x_T}{x_A} \right) \leq 1$

$$\bar{\alpha} = \frac{1}{2.5} \left[-\left(1 - \frac{x_T}{x_A} \right) + 1.25 \left(1 + \frac{x_T}{x_A} \right) \right] = \frac{1}{2.5} \left[0.25 + 2.25 \left(\frac{x_T}{x_A} \right) \right] = \frac{1}{10} \left[1 + 9 \frac{x_T}{x_A} \right],$$

which reduces to

$$\bar{\alpha} = 0, \quad -1 \leq \frac{x_T}{x_A} < -\frac{1}{9}, \quad \gamma = 1.25, \quad (7.39)$$

$$\bar{\alpha} = \frac{1}{10} \left(1 + 9 \frac{x_T}{x_A} \right), \quad -\frac{1}{9} \leq \frac{x_T}{x_A} \leq 1, \quad \gamma = 1.25. \quad (7.40)$$

While $\bar{\alpha}$ remains 0 for $\left(\frac{x_T}{x_A} \right) \in [-9, -\frac{1}{9}]$ it starts to increase linearly to 1 for $\left(\frac{x_T}{x_A} \right) \in [-\frac{1}{9}, 1]$.

Finally, when $\left(\frac{x_T}{x_A} \right) > 1$

$$\bar{\alpha} = \frac{1}{2.5} \left[-\left(\frac{x_T}{x_A} - 1 \right) + 1.25 \left(\frac{x_T}{x_A} + 1 \right) \right] = \frac{1}{2.5} \left[2.25 + 0.25 \frac{x_T}{x_A} \right] = \frac{1}{10} \left[9 + \frac{x_T}{x_A} \right], \quad (7.41)$$

which means that $\bar{\alpha}$ increases linearly. Note that equations (7.38) and (7.41) predict values of $\bar{\alpha}$ exceeding 1, a phenomenon taking place in the current analysis only when $\gamma > 1$, i.e., for a slow defender.

Chapter 8: ESCAPE REGION AND VORONOI DIAGRAM

In this chapter, we develop a novel analytic expression for the pursuit-evasion Voronoi diagram [1] that marks or borders the "safe" or "escape" region R_e for the Target, i.e., the region defined by the set of all coordinate points (x, y) such that if the target's initial position (x_T, y_T) is inside this region, then the target is guaranteed to survive (to escape the Attacker) provided both the Target and Defender implement their optimal strategies, and regardless of the policy adopted by the Attacker, whether optimal or not. The Voronoi diagram obtained divides the XY -plane, into two regions

- a safe or escape region R_e to which the Defender's initial position \mathbf{D} belongs.
- a *potentially unsafe* region to which the Attacker's initial position \mathbf{A} belongs.

Garcia et al. [1] treated this problem for $\gamma = 1$ only using some involved arguments with inequalities developed from scratch, and without due reference to their earlier analysis concerning the critical speed ratio.

A very important *insight* that markedly simplifies our work is to note that required Voronoi diagram is simply a rephrasing of the criticality equation (7.10) for any value of γ including the cases $\gamma > 1$ or $\gamma < 1$ and $\gamma = 1$, viewed as a relation between y_T and x_T rather than a polynomial in α . Working with equalities is definitely simpler than working with inequalities, and it specifies a Voronoi diagram that divides the whole plane into two distinct parts, namely: the safe region and the (potentially) unsafe region. The safe region is then identified simply as the region to which \mathbf{D} belongs. Note that the set of points of the safe regions R_e is a superset of the points reachable by the Defender before the Attacker R_r , i.e., $R_e \supseteq R_r$. The two regions R_e and R_r are identical if the Target does not move at all ($\alpha = 0$), a situation in which the Target does not exert any own effort to secure its survival and relies solely on the Defender's helping efforts. The family of Voronoi diagrams for various values of the normalized speed α will be bounded by the one for $\alpha = 0$, which is the border of R_r (Fig. 6.1).

Now, we obtain the general Voronoi diagram by rewriting the critically equation (7.10) or (7.9) as a relation between y_T and x_T . This is achieved via (6.33) to substitute $(x_A^2 + x_T^2 - 2x_A x_T + y_T^2)$ for d^2 . To avoid mistakes in the manipulations, one must make sure that the pertinent equation at every step is dimensionally homogeneous. From (7.9), we write

$$(4\gamma\alpha x_A)d = -(1 - \gamma^2)d^2 + 4\gamma^2 x_A x_T - 4\gamma^2 \alpha^2 x_A^2 \quad (8.1)$$

Now, we square both sides of (8.1) to obtain

$$\begin{aligned} 16\gamma^2 \alpha^2 x_A^2 (x_A^2 + x_T^2 - 2x_A x_T + y_T^2) &= [-(1 - \gamma^2)(x_A^2 + x_T^2 - 2x_A x_T + y_T^2) + 4\gamma^2 x_A x_T - 4\gamma^2 \alpha^2 x_A^2]^2 \\ &= [-(1 - \gamma^2)(x_A^2 + x_T^2 + y_T^2) + 2(1 + \gamma^2)x_A x_T - 4\gamma^2 \alpha^2 x_A^2]^2 \end{aligned} \quad (8.2)$$

Now we rearrange (8.2) as a quartic equation (forth-degree polynomial) in x_T and y_T , namely

$$\begin{aligned}
p(x_T, y_T) = & 16\gamma^2 \alpha^2 x_A^4 \\
& + 16\gamma^2 \alpha^2 x_A^2 x_T^2 \\
& - 32\gamma^2 \alpha^2 x_A^3 x_T \\
& + 16\gamma^2 \alpha^2 x_A^2 y_T^2 \\
& - (1 - \gamma^2)^2 [x_A^4 + x_T^4 + y_T^4 + 2x_A^2 x_T^2 + 2x_A^2 y_T^2 + 2x_T^2 y_T^2] \\
& - (1 + \gamma^2)^2 * 4x_A^2 x_T^2 \\
& - 16\gamma^4 \alpha^4 x_A^4 \\
& + 2(1 - \gamma^4)(2x_A^3 x_T + 2x_A x_T^3 + 2x_A x_T y_T^2) \\
& - 8\gamma^2(1 - \gamma^2)\alpha^2(x_A^4 + x_A^2 x_T^2 + x_A^2 y_T^2) \\
& + 16\gamma^2(1 + \gamma^2)\alpha^2 x_A^3 x_T
\end{aligned} \tag{8.3}$$

Equation (8.3) is satisfied by the required Voronoi diagram. The quartic polynomial $p(x_T, y_T)$ satisfies

$$p(x_T, -y_T) = p(x_T, y_T), \tag{8.4}$$

since it has even powers only for y_T . This means $p(x_T, y_T)$ is an even function in y_T and its graph in the $x_T - y_T$ plane is *symmetric* w.r.t. the x_T -axis.

Though (8.3) is a quartic equation, the actual Voronoi diagram is expected to be a quadratic curve and not a quartic curve. Note that we performed a squaring operation in obtaining (7.9) and hence for obtaining (8.3), which doubled the pertinent degree from two to four. Hence, the Voronoi diagram is just a second-degree factor, part or branch of the fourth-degree curve depicted by (8.3). We will be able to identify the correct Voronoi diagram by rejecting any part of the curve that lies in R_r , i.e., the correct Voronoi diagram should define the border of R_e such that $R_e \supseteq R_r$.

Note that (8.3) is the general equation for the Voronoi diagram in our current *TAD* problem and it appears here for the first time. As a check on its correctness, we use it to produce the correct curve for $\gamma = 1$, earlier obtained by Garcia et al. [3], namely

$$16x_A^2[-(1 - \alpha^2)x_T^2 + \alpha^2 y_T^2] = -16x_A^4 \alpha^2(1 - \alpha^2), \tag{8.5}$$

or equivalently

$$\frac{x_T^2}{\alpha^2 x_A^2} - \frac{y_T^2}{(1 - \alpha^2)x_A^2} = 1. \tag{8.6}$$

Equation (8.5) represents a hyperbola which is an open curve of two branches. Interestingly enough the term hyperbola is believed to have been coined by our old friend Apollonius of Perga in his work on conic sections. This hyperbola is represented in Fig. 8.1 and has the following properties:

- The hyperbola has two focal points $\mathbf{F}_1 = (c, 0)$ and $\mathbf{F}_2 = (-c, 0)$, where

$$c^2 = \alpha^2 x_A^2 + (1 - \alpha^2)x_A^2 = x_A^2, \tag{8.7}$$

i.e. the focal points coincide with the initial positions of the Attacker $\mathbf{A} = (x_A, 0)$ and

the Defender $\mathbf{D} = (-x_A, 0)$. Note that these two focal points are shared by an infinite set of curves for various values of α ($0 < \alpha < 1$).

- The line joining the focal points, i.e., the *transverse axis* coincides with the x_T -axis. The midpoint of the two focal points, i.e., the *center* of the hyperbola, is the origin (0,0). The conjugate axis (the line perpendicular to the hyperbola's transverse axis through its center) is the y_T -axis.
- The eccentricity of the hyperbola e is

$$e = \frac{c}{\alpha x_A} = \frac{1}{\alpha} > 1, \quad (8.8)$$

where the fact that $e > 1$ distinguishes the hyperbola from other conic sections, namely the ellipse (of $e < 1$) and the parabola (of $e = 1$). Note that the eccentricity e equals the ratio of the distance PF_i from a point P on the hyperbola to one focus F_i to the distance of the point to the corresponding directrix PQ_i , i.e.

$$PF_i = e PQ_i, \quad i = 1, 2, \quad (8.9)$$

- The vertices of the hyperbola (its intersections with the x_T -axis) are at $(\alpha x_A, 0)$ and $(-\alpha x_A, 0)$, and hence approach the two focal points as α tends to 1 (from below).
- Perpendicular to the transverse axis are the two directrices D_1 and D_2 given by

$$x_T = \mp \frac{\alpha x_A}{e} = \mp \alpha^2 x_A \quad (8.10)$$

These directrices also approach the focal points when $\alpha \rightarrow 1$.

- The hyperbola has two asymptotes that intersect at its center, and are given by

$$y_T = \mp \left(\frac{\sqrt{1 - \alpha^2}}{\alpha} \right) x_T. \quad (8.11)$$

For $0 < \alpha < 1$, the magnitude of the slope for an asymptote increases as α decreases and vice versa. Unlike the vertices and directrices, the asymptotes are independent of x_A .

- As a special case of the polynomial $p(x_T, y_T)$, in (8.3) the hyperbola (8.6) is an even function of y_T and hence is symmetric about the x_T -axis. Incidentally the hyperbola is also even in x_T and hence its two branches are symmetric about the y_T -axis but this observation is irrelevant since only one branch is retained in further study.
- In the limit ($\alpha \rightarrow 0$), the hyperbola (8.6) degenerates into the y_T -axis.
- In the limit ($\alpha \rightarrow 1$), the right branch of the hyperbola (8.6) degenerates into a ray emanating from \mathbf{A} in the positive direction of the x_T -axis, while its left branch degenerates into a ray emanating from \mathbf{D} in the opposite direction.

In passing, we note that though equation (8.6) was obtained herein as the special case ($\gamma = 1$) of equation (8.3), we have first obtained it by a rephrasing of the criticality condition

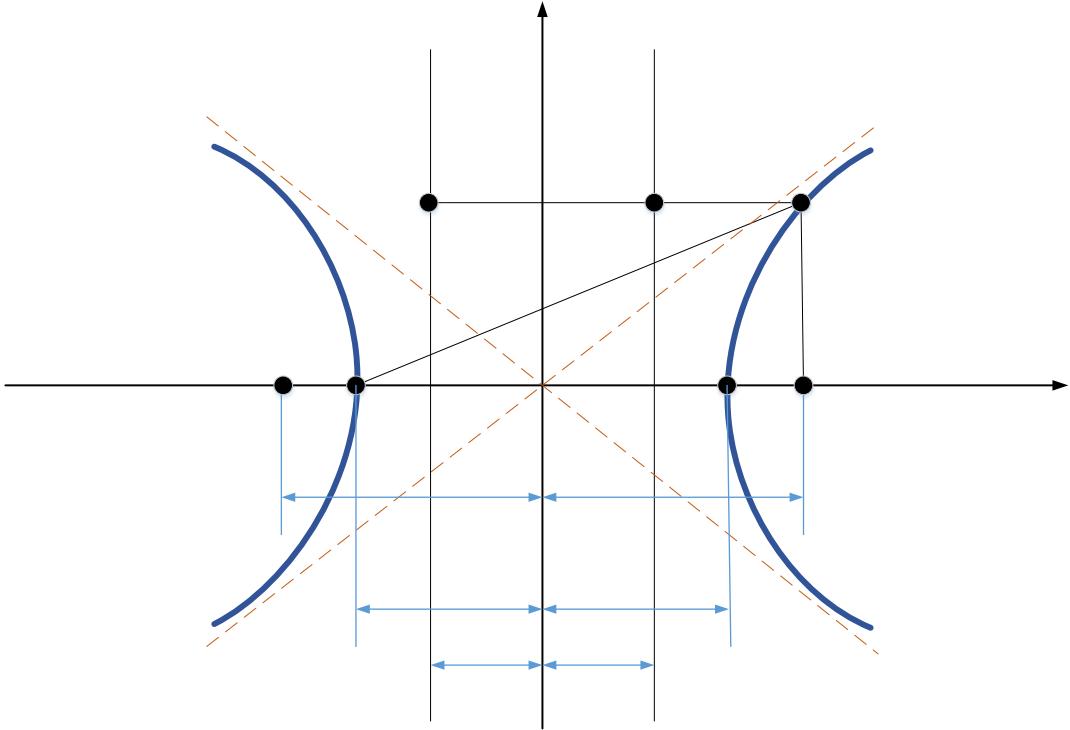


Figure 8.1: Graph and features of the hyperbola $\frac{x_T^2}{\alpha^2 x_A^2} - \frac{y_T^2}{(1-\alpha^2)x_A^2} = 1$

(7.13) which holds for $\gamma = 1$.

Now we note that the left branch of the hyperbola lies entirely in R_r , and hence it is rejected. The right branch of the hyperbola is the required Voronoi diagram, and it is bordering the escape region R_e which extends to its left. Clearly, it is an enlargement of the reachability region R_r . Likewise, we expect the escape region to be a subset of the AD Apollonius circle in Fig. ?? where $\gamma < 1$, and to be a superset of this circle in Fig. ?? where $\gamma > 1$.

Figure 8.2 is a sketch of the hyperbola (8.6) for $\alpha = 0.25$. the figure stresses that the right branch of the hyperbola is accepted as a Voronoi diagram while the left branch is rejected as such. Figure 5.3 is computer generated Voronoi diagrams for $\gamma = 1$ and α varying from 0 (the y_T -axis) to slightly less than 1 (the straight-line ray emanating from A and coinciding with the x_T -axis).

Now, we come to a novel and really important contribution of this thesis, where we visualize the escape regions in the case of a fast Defender and a slow Defender. Figure 8.4 shows the AD-Apollonius circle for $\gamma = 0.8$ and $x_A = 4$ for which $I_1 = (0.444, 0)$, $E_1 = (36, 0)$, $O_1 = (18.22, 0)$, and $r_1 = 17.78$. Imposed on this circle is the quartic curve given by (8.3) which consists of two closed curves. There is one closed curve entirely outside the AD-Apollonius circle, i.e., entirely inside the Reachability region R_r and hence is rejected. The other closed curve is entirely inside the AD-Apollonius circle, i.e., the reachability region R_r and hence it is accepted as the Voronoi diagram. The escape region (safe region) is the *unshaded* area in this figure. Figure 5.5 generalizes Fig. 8.4 by demonstrating various accepted branches of the Voronoi diagram for $\gamma = 0.8$ and α as s parameter ranging from 0 (the AD-Apollonius circle) to a value approaching 1 from below (a closed curve collapsing to

one barely encircling point A). This means that as α increases from 0 to 1, the safe or escape region expands from the exterior of the AD-Apollonius circle to almost the whole $x_T - y_T$ plane (excluding point A).

Now, we discuss the strikingly similar or dual case of a slow Defender. Figure 8.6 shows the AD-Apollonius circle for $\gamma = 1.25$ and $x_A = 4$ for which $\mathbf{I}_1 = (-0.444, 0)$, $\mathbf{E}_1 = (-36, 0)$, $\mathbf{O}_1 = (-18.22, 0)$, $r_1 = 17.78$. Imposed on this circle is the quartic curve given by (8.3). Note that the whole graph in Fig. 8.6 is a *mirror image* of that in Fig. 8.4. The quartic curve again consists of two closed curves. There is one closed curve entirely inside the AD-Apollonius circle. i.e., entirely inside the Reachability region R_r and hence is rejected. The other closed curve is entirely outside the AD-Apollonius circle, i.e., entirely outside the Reachability region R_r , and hence is accepted as the Voronoi diagram. The escape region (safe region) is the *shaded* area in this figure. Figure 5.7 generalizes Fig. 8.6 by demonstrating various accepted branches of the Voronoi diagram for $\gamma = 1.25$ and α as a parameter ranging from 0 (the AD-Apollonius circle) to a value approaching 1 from below.

In passing, we discuss the situation of $\gamma = 0$ (when the Target does not move at all). In this case the TA circle collapse into a single point T and the critically condition (7.9) reduces to

$$(1 - \alpha^2)d^2 = 4\gamma^2 x_T x_A \quad (8.12)$$

which can readily be written as

$$(x_T - (\frac{1 + \gamma^2}{1 - \gamma^2})x_A)^2 + (y_T - 0)^2 = (\frac{2\gamma x_A}{1 - \gamma^2})^2 \quad (8.13)$$

an equation that can be identified as the equation of the AD Apollonius circle in the $Y_T - X_T$ coordinates. This circle serves as the Voronoi diagram for $\alpha = 0$, a fact confirming our assertion that the escape region R_e reduces to the reachability region R_r when $\alpha = 0$. Likewise, when $\gamma = 1$, imposing the condition $\alpha = 0$ on the condition (7.13) results in

$$x_T = 0 \quad (8.14)$$

which means that the right branch of the hyperbolic Voronoi diagram (8.6) degenerates into the y_T -axis (the perpendicular bisector of \overline{AD}) and again the region R_e concides with the region R_r in this case.

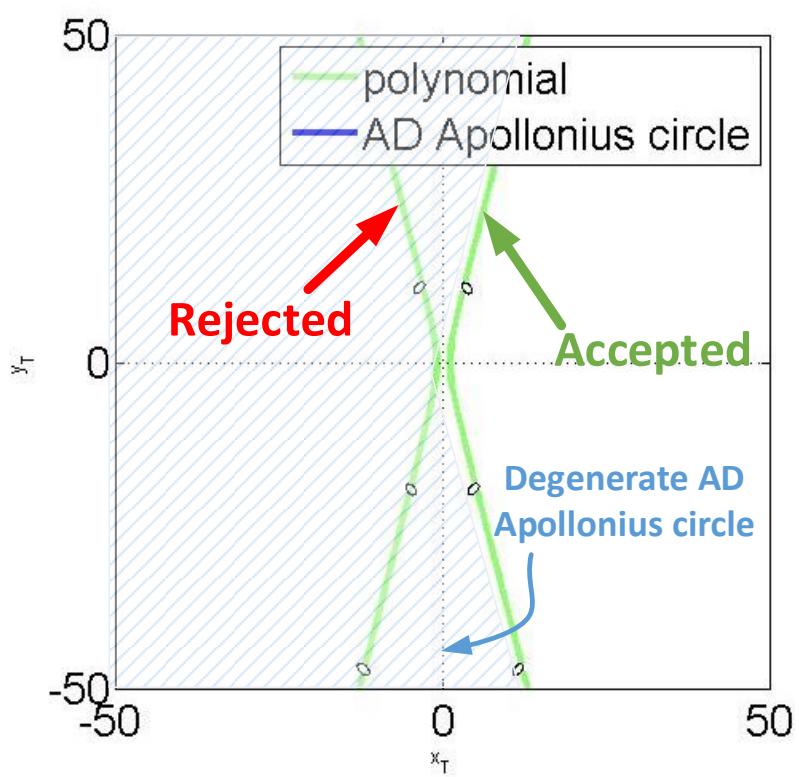


Figure 8.2: Generated computer output for the Voronoi diagram bordering the safe region for $x_A = 4$, $\alpha = 0.25$, $\gamma = 1$ (the safe region is the shaded area)

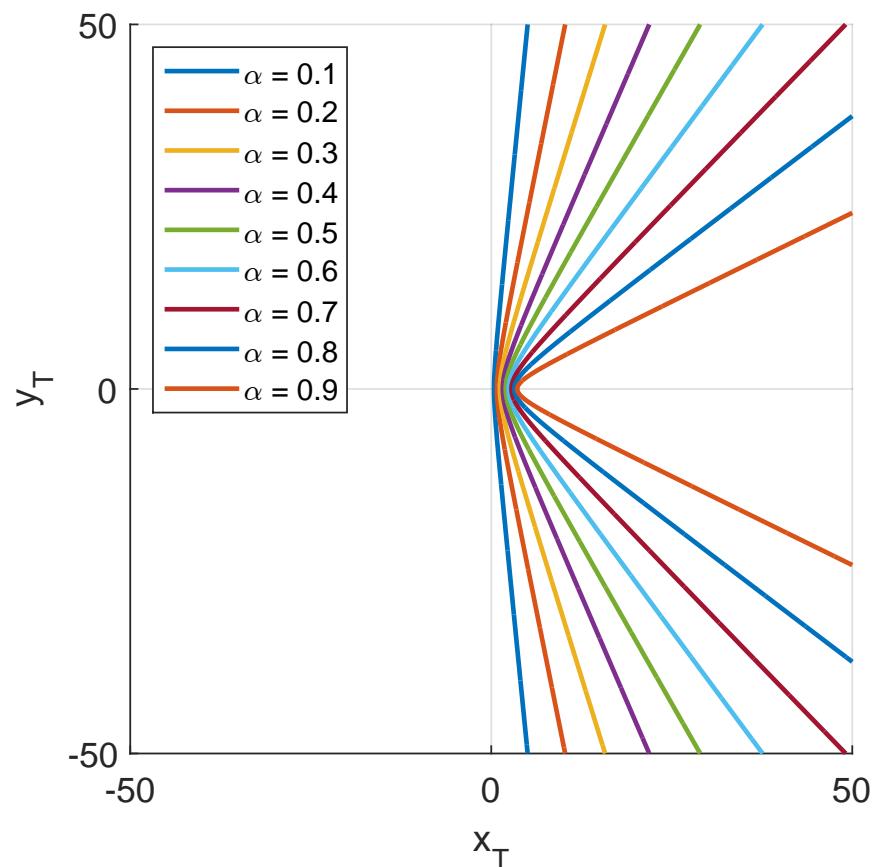


Figure 8.3: Various accepted branches of the voronoi diagram for $\gamma = 1$ and α as a parameter ranging from 0 to 1. These curves are computer generated from (8.3) and (8.6)

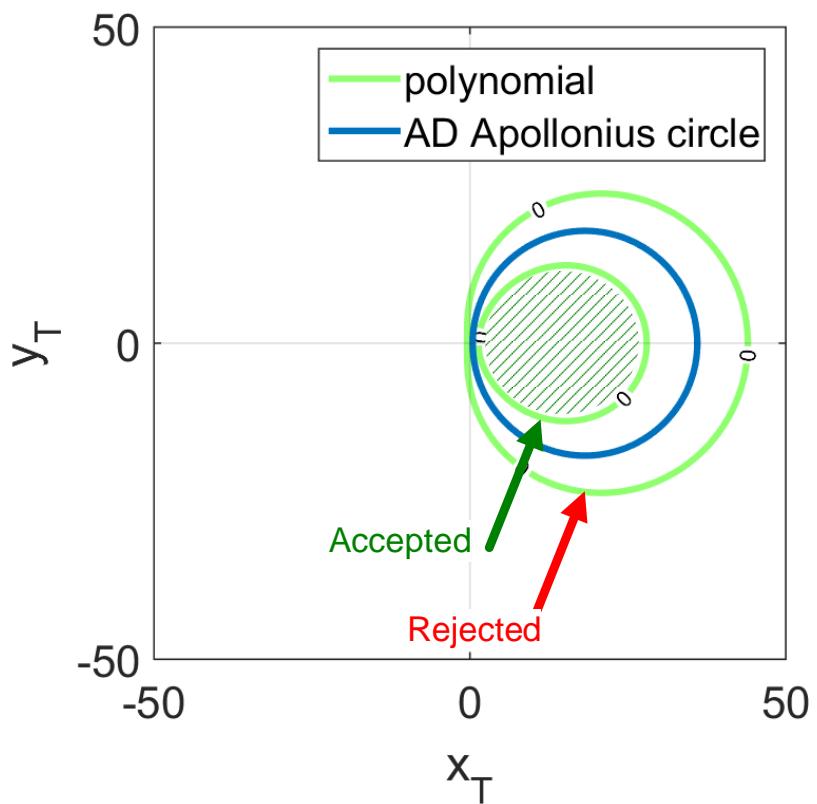


Figure 8.4: generated computer output for the Voronoi diagram bordering the safe region for $x_A = 4$, $\alpha = 0.25$, $\gamma = 0.8$ (the safe region is the unshaded area) the quartic in (8.2) or (8.3) produces two closed curves: one outside the AD-Apollonius circle (rejected) and the other inside the circle (accepted as the Voronoi diagram)

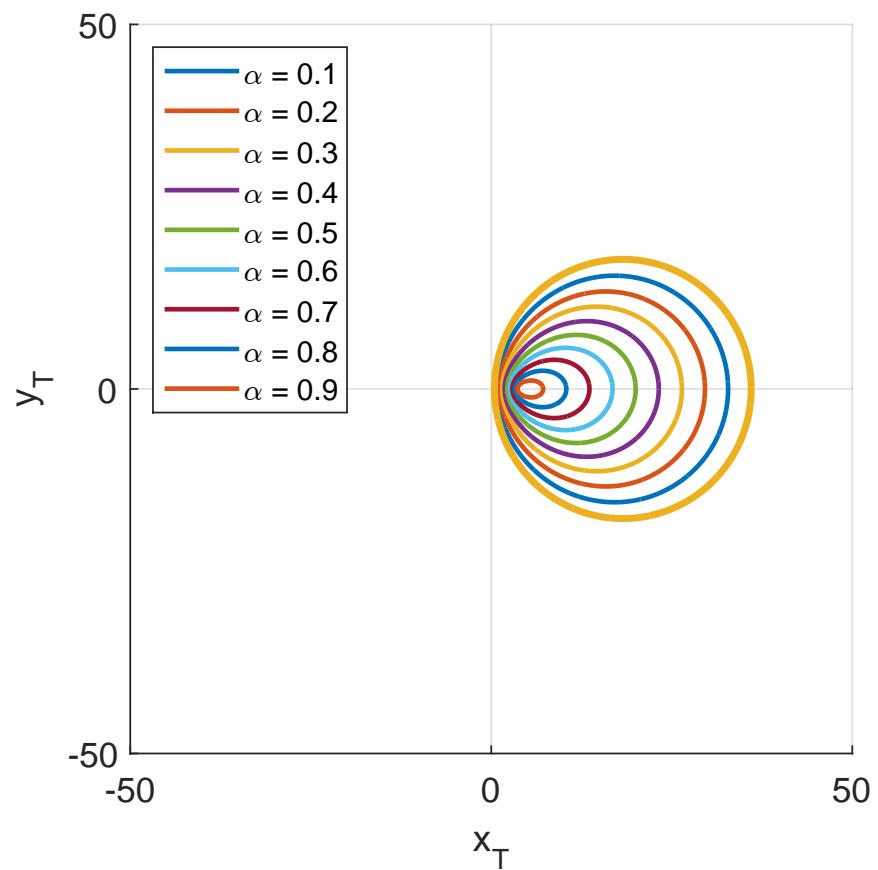


Figure 8.5: Various accepted branches of the voronoi diagram for $\gamma = 0.8$ and α as a parameter ranging from 0 to 1. These curves are computer generated from (8.3) and (8.6)

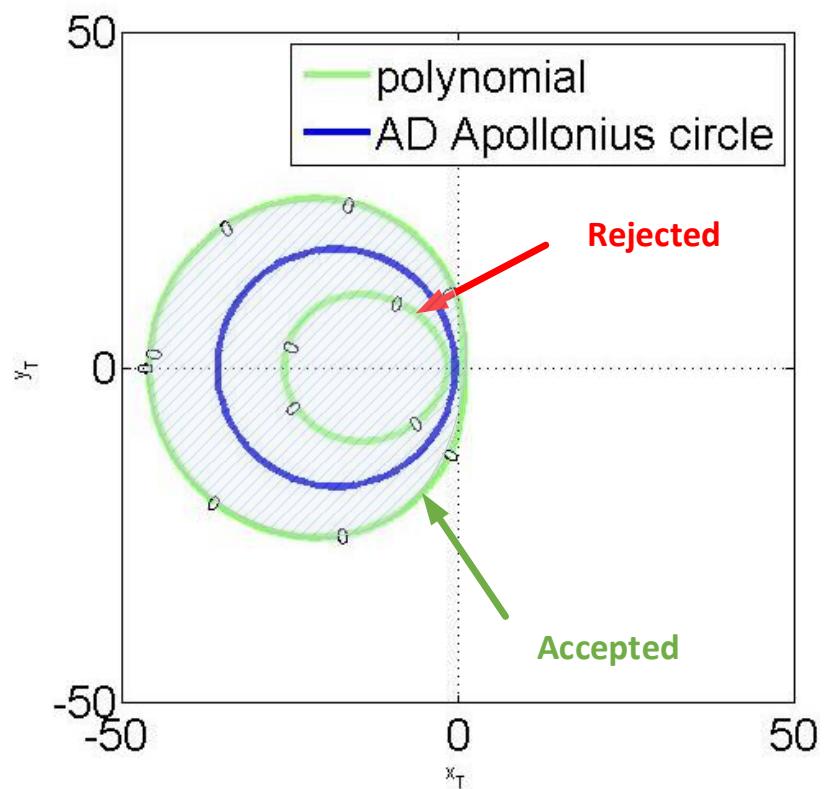


Figure 8.6: generated computer output for the Voronoi diagram bordering the safe region for $x_A = 4$, $\alpha = 0.25$, $\gamma = 1.25$ (the safe region is the shaded area) the quartic in (8.2) or (8.3) produces two closed curves: one inside the AD-Apollonius circle (rejected) and the other outside the circle (accepted as the Voronoi diagram)

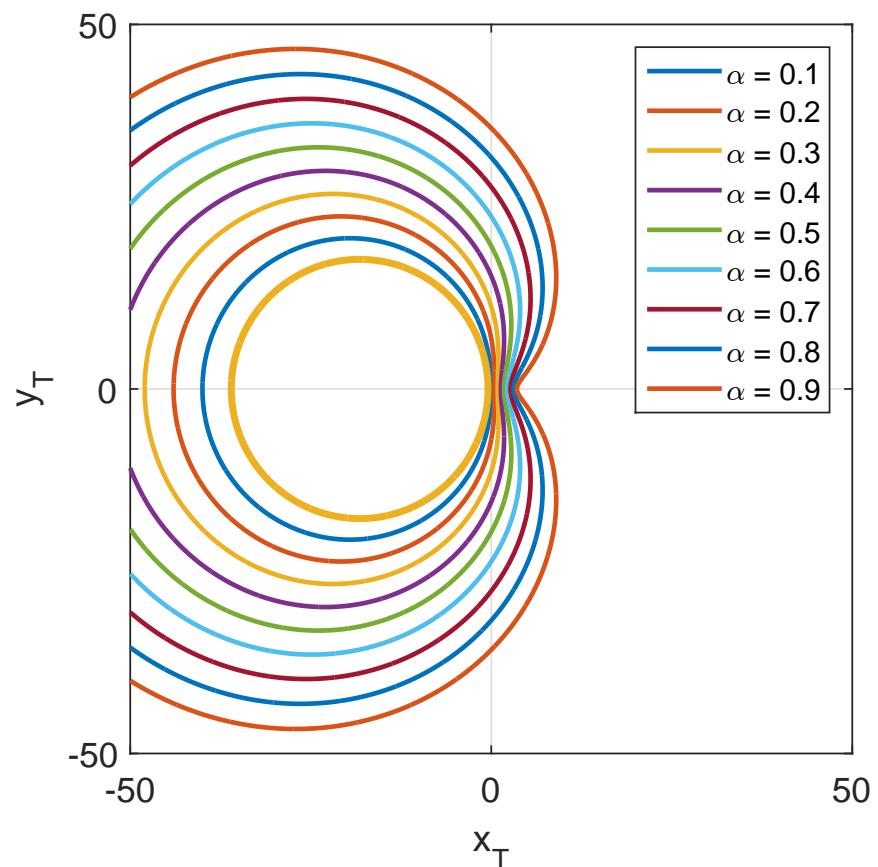


Figure 8.7: Various accepted branches of the voronoi diagram for $\gamma = 1.25$ and α as a parameter ranging from 0 to 1. These curves are computer generated from (8.3) and (8.6)

Chapter 9: OPTIMAL STRATEGIES

The *TAD* pursuit-evasion differential game is discussed and solved herein to obtain the optimal heading angels for the Target and the Defender team to maximize the terminal separation between the Target and Attacker, and also to obtain the optimal heading angle for the Attacker to minimize this distance. In the following sections, we discuss a variety of cases differing according to the initial position of the Target $\mathbf{T} = (x_T, y_T)$ relative to the *AD*-Apollonius circle, and also according to whether the Defender is fast ($\gamma < 1$), similar ($\gamma = 1$) or slow ($\gamma > 1$).

9.1 Target initial position is outside the *AD* Apollonius circle, and $\gamma < 1$

Assume that the Target, Attacker, and Defender aim at points v, u and w on the *AD* Apollonius circle. While the Attacker and Defender try to go to points u and w , the Target tries to run away from the point u . These three agents try collectively to solve the minmax optimization problem.

$$\min_u \max_{v,w} J(u, v, w) \quad (9.1)$$

where the cost/payoff function $J(u, v, w)$ is the final separation, i.e., the distance between the Target terminal position T' and the point C on the *AD* Apollonius circle where the Defender intercepts the Attacker, i.e.,

$$\begin{aligned} J(u, v, w) &= CT' = CT + TT' \\ &= CT + \alpha AC \\ &= [(x_C - x_T)^2 + (y_C - y_T)^2]^{\frac{1}{2}} + \alpha[(x_A - x_C)^2 + y_C^2]^{\frac{1}{2}} \\ &= J(x_C, y_C) \end{aligned} \quad (9.2)$$

Note that the Target is initially outside the *AD*-Apollonius circle, i.e., it is within the Reachability Region R_r whose points are reachable by the Defender before the Attacker. Hence, the Defender can help the Target regardless of the effort contributed by the Target, i.e., regardless of the speed ratio α ($0 \leq \alpha < 1$). This means that the critical speed ratio in this case $\bar{\alpha} = 0$. The Defender's optimal policy is to choose its aim point w as w^* such that:

$$w^*(u, v) = u, \quad (9.3)$$

in order to guarantee interception of the Attacker (which is aiming at u). The cost/payoff function $J(u, v, w)$ can now be viewed as $J(u, v)$, i.e., a function of u and v only. Gaecia, et al. [2] assert also the solution of the optimization problem (9.1) is such that

$$u^* = v^*. \quad (9.4)$$

The point C chosen in (9.2) is in fact the common optimal aim points (according to (9.3) and (9.4)), i.e.,

$$C = u^* = v^* = w^*, \quad (9.5)$$

which means that the statement made in (9.2) is indeed correct. Hence the optimization problem (9.1) is replaced by

$$\min_{(x_C, y_C)} J(x_C, y_C) \text{ subject to } (x_C, y_C) \text{ being on the AD circle} \quad (9.6)$$

With angle ϕ defined as the complement to the polar angle of position vector of C w.r.t. the centre of the AD Apollonius circle and the angle λ defined as the complement to the polar angle of the position vector of T w.r.t. that center (Fig. 9.1), the *constrained* optimization in (9.6) is replaced by the *unconstrained* optimization [2] with a single independent variable ϕ ,

$$\min_{\phi} J(\phi) \quad (9.7)$$

where

$$J(\phi) = [r_1^2 + N^2 - 2Nr_1 \cos(\phi - \lambda)]^{\frac{1}{2}} + \alpha[r_1^2 + M^2 - 2Mr_1 \cos \phi]^{\frac{1}{2}} \quad (9.8)$$

where by view of (6.15) and (6.15), we have

$$r_1 = \frac{2\gamma}{1-\gamma^2} x_A, \quad (9.9)$$

$$M = AO_1 = \left(\frac{1+\gamma^2}{1-\gamma^2}\right)x_A - x_A = \frac{2\gamma^2}{1-\gamma^2} x_A = \gamma r_1, \quad (9.10)$$

$$N = O_1 T = \left[\left(\frac{1+\gamma^2}{1-\gamma^2}x_A - x_T\right)^2 + y_T^2\right]^{\frac{1}{2}}, \quad (9.11)$$

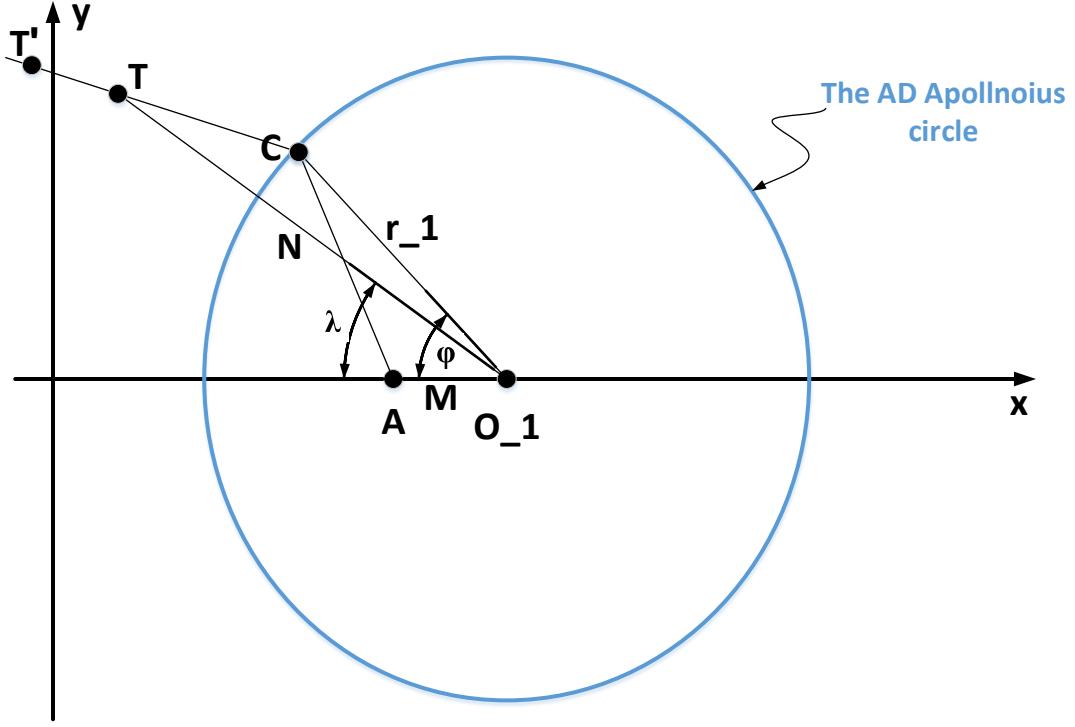


Figure 9.1: Pertaining to expressing the final separation in terms of shown distances and angles

We now solve (9.7) by setting $\frac{dJ(\phi)}{d\phi}$ to zero, thereby obtaining

$$Nr_1 \sin(\phi - \lambda)[r_1^2 + N^2 - 2Nr \cos(\phi - \lambda)]^{\frac{-1}{2}} = -\alpha Mr_1 \sin \phi [r_1^2 + M^2 - 2Mr \cos \phi]^{\frac{-1}{2}} \quad (9.12)$$

which is squared and simplified via (9.10) to obtain:

$$N^2 r_1^2 \sin^2(\phi - \lambda)[1 + \gamma^2 - 2\gamma \cos \phi] = \alpha^2 M^2 \sin^2 \phi [r_1^2 + N^2 - 2Nr_1 \cos(\phi - \lambda)] \quad (9.13)$$

Note that both sides of (9.13) now have the dimension of $(length)^4$. Now, we substitute for $\sin \phi, \cos \phi, \sin(\phi - \lambda), \cos(\phi - \lambda)$ by their exponential expressions

$$\sin \phi = \frac{1}{2j}[e^{j\phi} - e^{-j\phi}],$$

$$\cos \phi = \frac{1}{2}[e^{j\phi} + e^{-j\phi}],$$

$$\sin(\phi - \lambda) = \frac{1}{2j}[e^{j(\phi-\lambda)} - e^{-j(\phi-\lambda)}]$$

,

$$\cos(\phi - \lambda) = \frac{1}{2}[e^{j(\phi-\lambda)} + e^{-j(\phi-\lambda)}] = \frac{1}{2}[e^{j\phi}l^{-1} + e^{-j\phi}l],$$

to obtain

$$\sin^2 \phi = -\frac{1}{4}[e^{j(2\phi)} + e^{-j(2\phi)} - 2]$$

$$\sin^2(\phi - \lambda) = -\frac{1}{4}[e^{j(2\phi)}e^{-j(2\lambda)} + e^{-j(2\phi)}e^{j(2\lambda)} - 2] = -\frac{1}{4}[e^{j(2\phi)}l^{-2} + e^{-j(2\phi)}l^2 - 2]$$

$$\begin{aligned} & N^2 r_1^2 [e^{j(2\phi)}l^{-2} + e^{-j(2\phi)}l^2 - 2][1 + \gamma^2 - \gamma(e^{j\phi} + e^{-j\phi})] \\ &= \alpha^2 M^2 [e^{j(2\phi)} + e^{-j(2\phi)} - 2][r_1^2 + N^2 - Nr_1(e^{j\phi}l^{-1} + e^{-j\phi}l)] \end{aligned} \quad (9.14)$$

where $l = e^{j\lambda}$

$$\begin{aligned} & [-N^2 r_1^2 l^{-2} \gamma + Nr_1 l^{-1} \alpha^2 M^2] e^{j(6\phi)} \\ &+ [N^2 r_1^2 l^{-2} (1 + \gamma^2) - \alpha^2 M^2 (r_1^2 + N^2)] e^{j(5\phi)} \\ &+ [-N^2 r_1^2 l^{-2} \gamma + 2\gamma N^2 r_1^2 + \alpha^2 M^2 Nr_1 l - 2\alpha^2 M^2 Nr_1 l^{-1}] e^{j(4\phi)} \\ &+ [-2N^2 r_1^2 (1 + \gamma^2) + 2\alpha^2 M^2 (r_1^2 + N^2)] e^{j(3\phi)} \\ &+ [-N^2 r_1^2 l^2 \gamma + 2N^2 r_1^2 \gamma + \alpha^2 M^2 Nr_1 l^{-1} - 2\alpha^2 M^2 Nr_1 l] e^{j(2\phi)} \\ &+ [N^2 r_1^2 l^2 (1 + \gamma^2) - \alpha^2 M^2 (r_1^2 + N^2)] e^{j\phi} \\ &+ [-N^2 r_1^2 l^2 \gamma + \alpha^2 M^2 Nr_1 l] = 0 \end{aligned} \quad (9.15)$$

Now, we rewrite (9.15) in the following form (in an attempt to make it resemble equation (41) in [2]) by invoking (9.10) and dividing by $\alpha^2 M^2 = \alpha M \gamma r_1$.

$$\begin{aligned} & \frac{Nr_1}{l} \left(1 - \frac{N}{\alpha^2 M l}\right) e^{j(6\phi)} \\ &+ \left[\left(\frac{N}{\alpha M l}\right)^2 (r_1^2 + M^2) - (r_1^2 + N^2)\right] e^{j(5\phi)} \\ &+ Nr_1 \left[\frac{N}{\alpha^2 M l^2} (2l^2 - 1) + \left(l - \frac{2}{l}\right)\right] e^{j(4\phi)} \\ &+ 2[r_1^2 + N^2 - \left(\frac{N}{\alpha M}\right)^2 (r_1^2 + M^2)] e^{j(3\phi)} \\ &+ Nr_1 \left[\left(\frac{N}{\alpha^2 M}\right) (2 - l^2) - 2l + \frac{1}{l}\right] e^{j(2\phi)} \\ &+ \left[\left(\frac{Nl}{\alpha M}\right)^2 (r_1^2 + M^2) - (r_1^2 + N^2)\right] e^{j\phi} \\ &+ Nr_1 l \left(1 - \frac{Nl}{\alpha^2 M}\right) = 0 \end{aligned} \quad (9.16)$$

Equation (9.16) is a sixth-degree polynomial in $e^{j\phi}$ that has six roots. Each of these roots should be substituted in the expression (9.8) for $J(\phi)$ to determine which root yields the minimal value of $J(\phi)$.

9.2 Target initial position is inside the AD-Apollonius circle, and $\gamma < 1$

Again, the Target, Attacker, and Defender aim at points u, v and W on the *AD*-Apollonius circle, but now each of them is going toward its respective aim. The three agents now try to collectively solve the *maxmin* optimization problem

$$\max_{v,w} \min_u J(u, v, w) \quad (9.17)$$

Garcia et al. [2] use arguments similar to those of Sec. 6.1 to deduce conditions similar to (9.3) - (9.5), and then reduce to (9.17) to a form similar to that of (9.6) [albeit with minimization replaced by maximization]

$$\max_{(x_c, y_c)} J(x_c, y_c) \quad \text{subject to } (x_c, y_c) \text{ being on the } AD \text{ circle}, \quad (9.18)$$

where

$$\begin{aligned} J(x_c, y_c) &= CT' = TT' - CT \\ &= \alpha AC - CT \\ &= \alpha[(x_A - x_c)^2 + y_c^2]^{\frac{1}{2}} - [(x_c - x_T)^2 + (y_c - y_T)^2]^{\frac{1}{2}} \end{aligned} \quad (9.19)$$

Again, the constrained optimization (9.18) can be replaced by constrained optimization of single-variable cost/payoff function, $J(\phi)$, namely

$$\max_{\phi} J(\phi), \quad (9.20)$$

where

$$J(\phi) = \alpha[r_1^2 + M^2 - 2r_1M \cos \phi]^{\frac{1}{2}} - [r_1^2 + N^2 - 2r_1N \cos(\lambda - \phi)]^{\frac{1}{2}}. \quad (9.21)$$

where, again, $M = AO_1$ and $N = TO_1$. Equation (9.21) differs from (9.8) in the sign of one of its two terms. Setting $\frac{dJ(\phi)}{d(\phi)}$ to zero results in an equation that has a sign difference from (9.12) which is lost upon squaring to obtain (9.13).

Therefore, (9.14)-(9.16) are obtained in this case also. Again, one has to obtain each of the six roots of (9.16), but now he has to substitute it in the expression (9.21) for $J(\phi)$ to determine which root yields the maximal (rather than the minimal) value for $J(\phi)$.

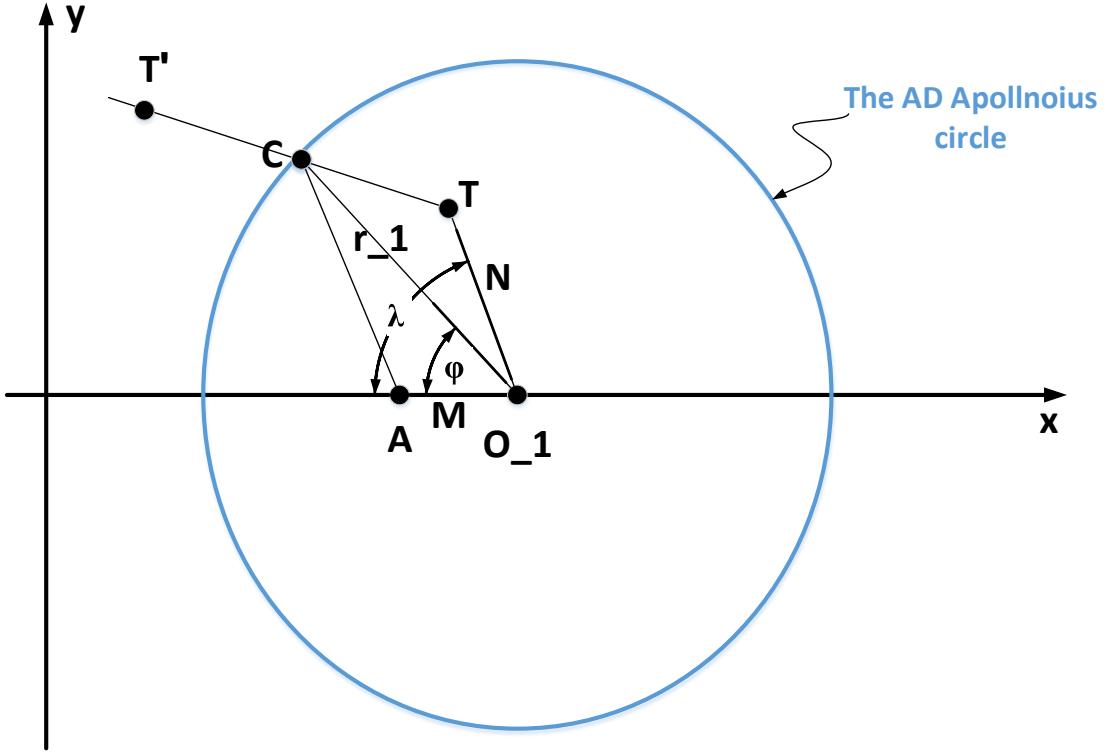


Figure 9.2: Pertaining to expressing the final separation in terms of the shown distance and angles for $\gamma < 1$ and T inside the AD Apollonius circle, i.e., outside R_r .

9.3 Target initial position is in the R.H.S. of the XY-plane, and $\gamma = 1$

This case, which appears here for the first time, resembles that of Sec. 6.1, with the exception that the condition $\gamma = 1$ makes the AD-Apollonius circle degenerate into the perpendicular bisector or the Y -axis (Fig. 9.4). Now the cost/payoff function is replaced by

$$\begin{aligned}
 J(y_c) &= CT' = CT + TT' \\
 &= CT + \alpha AC \\
 &= [x_T^2 + (y_c - y_T)^2]^{\frac{1}{2}} + \alpha[x_A^2 + y_c^2]^{\frac{1}{2}}.
 \end{aligned} \tag{9.22}$$

and one has to solve a single-variable unconstrained minimization problem.

$$\min_{y_c} J(y_c) \tag{9.23}$$

by setting $\frac{dJ(y_c)}{dy_c}$ to zero, i.e.,

$$\frac{dJ(y_c)}{dy_c} = \frac{(y_c - y_T)}{[x_T^2 + (y_c - y_T)^2]^{\frac{1}{2}}} + \frac{\alpha y_c}{[x_A^2 + y_c^2]^{\frac{1}{2}}} = 0. \tag{9.24}$$

which can be arranged as

$$(y_c - y_T)(x_A^2 + y_c^2)^{\frac{1}{2}} = -\alpha y_c(x_T^2 + (y_c - y_T)^2)^{\frac{1}{2}} \quad (9.25)$$

Equation (9.25) can be squared to yield the following quartic (fourth-degree) real equation in y_c .

$$(y_c^2 - 2y_c y_T + y_T^2)(x_A^2 + y_c^2) = \alpha^2 y_c^2 [x_T^2 + y_c^2 + y_T^2 - 2y_c y_T] \quad (9.26)$$

or

$$(1 - \alpha^2)y_c^4 - 2(1 - \alpha^2)y_T y_c^3 + [(1 - \alpha^2)y_T^2 + x_A^2 - \alpha^2 x_T^2]y_c^2 - 2x_A^2 y_T y_c + x_A^2 y_T^2 = 0 \quad (9.27)$$

Equation (9.27) yields four roots (expected to be all real, or to have two real and two complex conjugate roots). Each of the real roots is substituted into $J(y_c)$ given by (9.22) to select the one minimizing it.

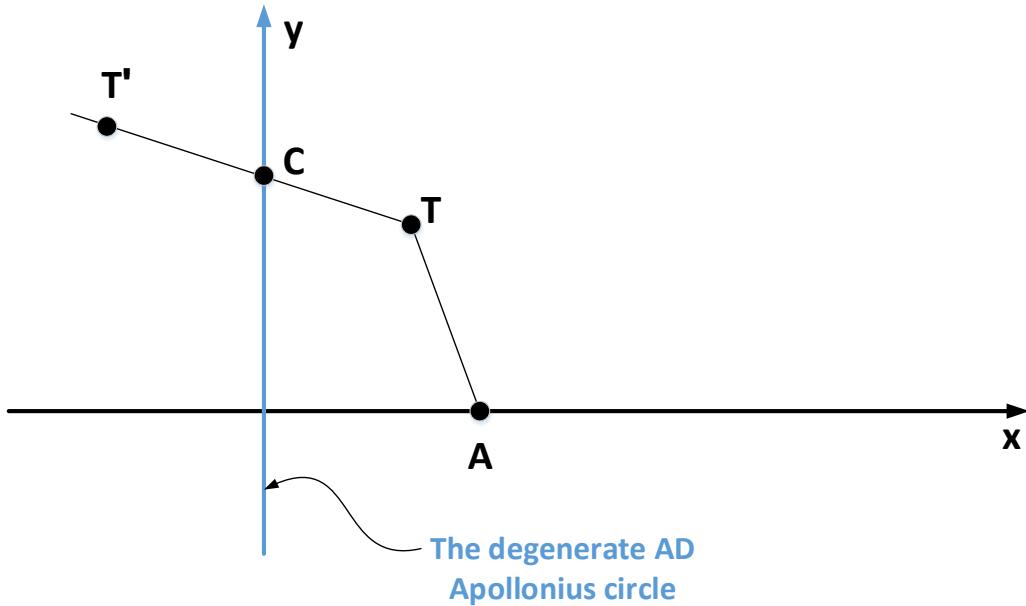


Figure 9.3: Pertaining to expressing the final separation CT' when $\gamma = 1$ and T is in the R.H.S. of the Y-plane, i.e., outside R_r .

9.4 Target initial position is in the L.H.S. of the XY -plane, and $\gamma = 1$

This case was discussed previously by Garcia et al. [], and resemble that of Sec. 6.1, with the exception that the condition $\gamma = 1$ makes the AD -Apollonius circle degenerate into the perpendicular bisector or the Y -axis (Fig. 9.3). The current case is also similar to the case in Sec. 6.3, with the exception that T is in the L.H.S. rather than the R.H.S. of the Y -plane. Now, the cost/payoff function is replaced by

$$\begin{aligned} J(y_c) &= CT' = TT' - CT \\ &= \alpha AC - CT \\ &= \alpha[x_A^2 + y_c^2]^{\frac{1}{2}} - [x_T^2 + (y_c - y_T)^2]^{\frac{1}{2}}, \end{aligned} \quad (9.28)$$

and one has to solve a single-variable unconstrained maximization problem

$$\max_{y_c} J(y_c) \quad (9.29)$$

by setting $\frac{dJ(y_c)}{dy_c}$ to zero, i.e.,

$$\frac{dJ(y_c)}{dy_c} = \frac{\alpha y_c}{[x_A^2 + y_c^2]^{\frac{1}{2}}} - \frac{(y_c - y_T)}{[x_T^2 + (y_c - y_T)^2]^{\frac{1}{2}}} \quad (9.30)$$

Equation (9.30) can be rearranged and then squared to yield (9.26) and consequently (9.27). Each of the real roots or the quartic equation (9.27) should be substituted into $J(y_c)$ given by (9.28) to decide which of them maximize $J(y_c)$

9.5 Target initial position is inside the AD Apollonius circle, and $\gamma > 1$

Again, the case discussed in this section has not appear earlier in the open literature. This case resembles the cases in sections 6.1 and 6.3 in the fact that it deals with a Target initially within R_2 . The current situation is depicted in Fig. 9.5, in which the distances M , N and angels ϕ and λ are defined. Note that the angles ϕ and λ are the polar angles (rather than complements thereof) for the position vector of C and T w.r.t. the center O_1 of the AD Apollonius circle. The distance $N = O_1T$ is still given by (9.11), while $M = AO_1$ is given by

$$\begin{aligned} M + AO_1 &= x_A - \frac{1 + \gamma^2}{1 - \gamma^2} x_A \\ &= \frac{-2\gamma^2}{1 - \gamma^2} x_A = \frac{2\gamma^2}{\gamma^2 - 1} x_A \\ &= \gamma r_1, \end{aligned} \quad (9.31)$$

a result that can be unified with (9.10) in the form

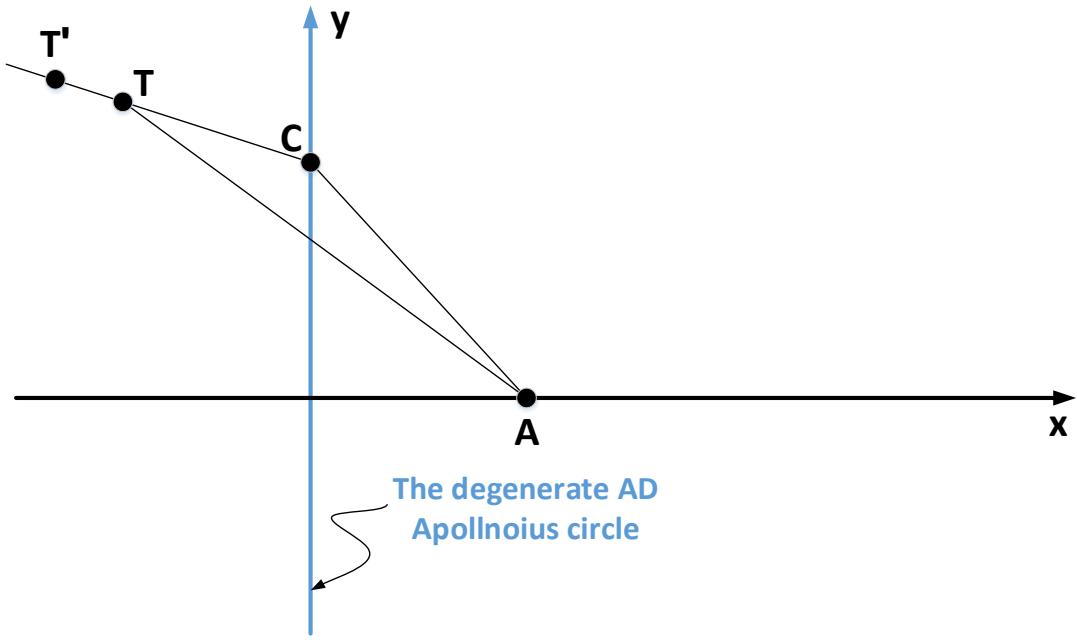


Figure 9.4: Pertaining to expressing the final separation CT' when $\gamma = 1$ and T is in the L.H.S. of the Y-plane, i.e., within R_r .

$$M = \frac{2\gamma^2}{|1 - \gamma^2|} x_A = \gamma r_1. \quad (9.32)$$

with these minor observations, we can easily verify that equations (9.1)-(9.16) are still valid in the current case.

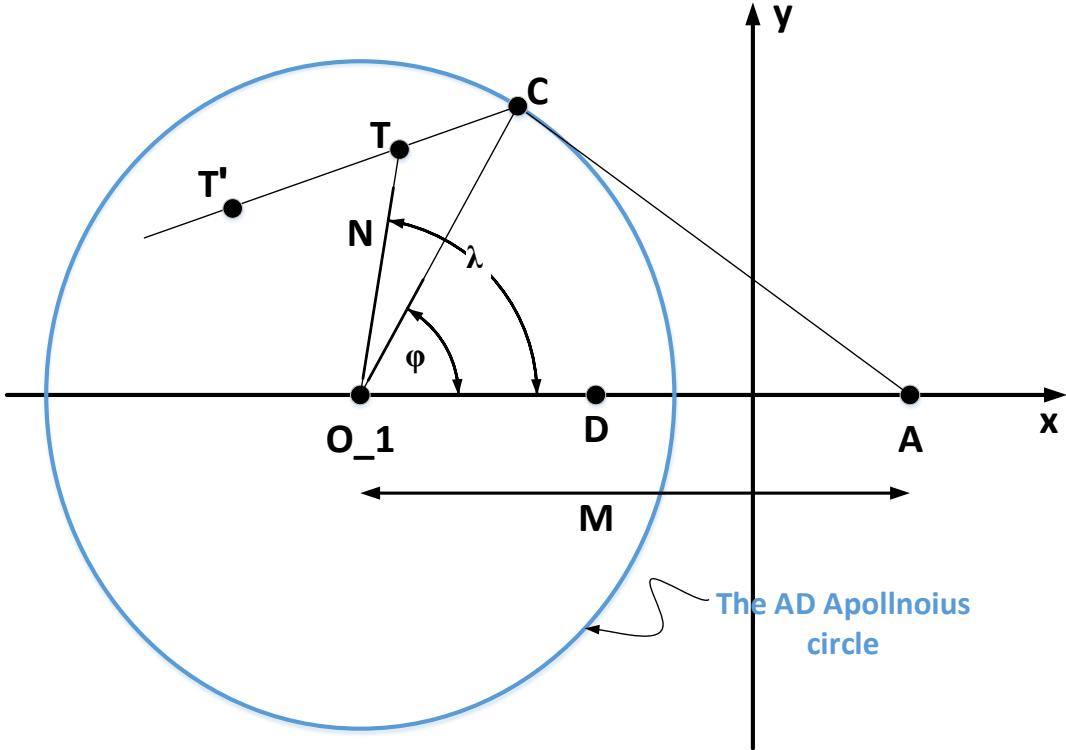


Figure 9.5: Pertaining to expressing the final separation in terms of the shown distance and angles for $\gamma > 1$ and T inside the AD Apollonius circle, i.e., within R_r .

9.6 Target initial position is outside the AD Apollonius circle, and $\gamma > 1$

This case is again a novel case. It resembles the cases in the Sections 6.2 and 6.4 in the fact that it deals with a Target initially outside R_r . It is similar to the case in the previous section (Sec. 6.5) as both sections deals with a slow Defender. The situation is demonstrated by Fig. 9.6, which has similarities and differences with Fig. 9.2 and Fig. 9.5. We can easily verify that equations (9.18)-(9.21), and consequently equations (9.14)-(9.16) are still valid in the current case.

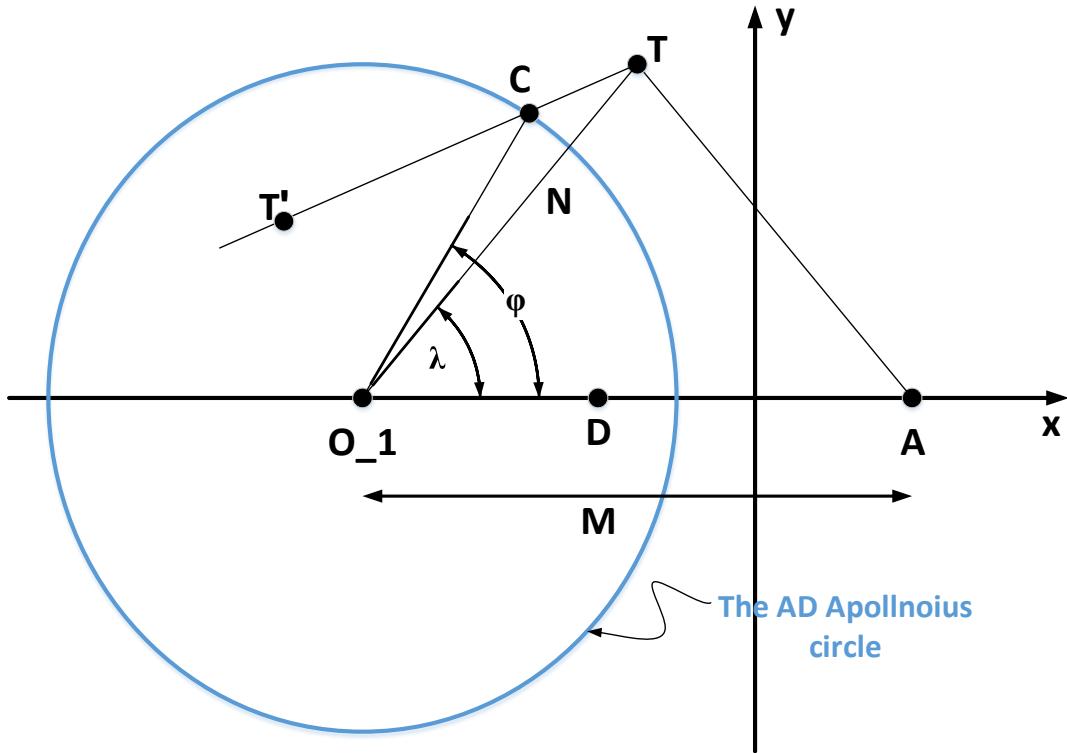


Figure 9.6: Pertaining to expressing the final separation in terms of the shown distance and angles for $\gamma > 1$ and T outside the AD Apollonius circle, i.e., outside R_r .

9.7 Numerical Results

We study the optimal values of ϕ or y_c for $x_A = 4$, $\alpha = 0.25$ and the following choices of γ and (x_T, y_T)

9.7.1 $\gamma = 0.8$, $x_T = 1.0$, $y_T = 20.0$

This is the case in Sec. 6.1, and is handled by obtaining six roots via (9.16) and selecting the one which minimizes $J(\phi)$ in (9.8). Results are shown in Table 6.1 and the optimal value of ϕ is $\phi_{optimal} =$

i	$\phi_i = Root \neq i$ of (9.16)	$J(\phi_i)$ in (9.8) [choose minimum]
1		
2		
3		
4		
5		
6		

9.7.2 $\gamma = 0.8, x_T = 20.0, y_T = 1.0$

This is the case in Sec. 6.2, and is handled by obtaining six roots via (9.16) and selects the one that maximizes $J(\phi)$ in (9.21). Results are shown in Table 6.2, and the optimal value of ϕ is $\phi_{optimal} =$

i	$\phi_i = Root \neq i$ of (9.16)	$J(\phi_i)$ in (9.21) [choose minimum]
1		
2		
3		
4		
5		
6		

9.7.3 $\gamma = 1.0, x_T = -1.0, y_T = 1.0$

This is the case in Sec. 6.3, and is handled by obtaining four roots via (9.27) and selecting the one that minimizes $J(\phi)$ in (9.24). Results are shown in Table 6.3, and the optimal value of $y_c =$

i	$\phi_i = Root \neq i$ of (9.21)	$J(\phi_i)$ in (9.24) [choose minimum]
1		
2		
3		
4		

9.7.4 $\gamma = 1.0, x_T = 1.0, y_T = 1.0$

This is the case in Sec. 6.4, and is handled by obtaining four roots via (9.27) and selecting the one that maximizes $J(y_c)$ in (9.28). Results are shown in Table 6.4, and the optimal value of $y_c =$

i	$y_i = Root \neq i$ of (9.27)	$J(\phi_i)$ in (9.28) [choose minimum]
1		
2		
3		
4		

9.7.5 $\gamma = 1.25, x_T = -20.0, y_T = 1.0$

This is the case in Sec. 6.5, and is handled by obtaining six roots via (9.16) and selecting the one that minimizes $J(\phi)$ in (9.8). Results are shown in Table 6.5 and the optimal value of ϕ is $\phi_{optimal} =$

i	$\phi_i = Root \neq i$ of (9.16)	$J(\phi_i)$ in (9.8) [choose minimum]
1		
2		
3		
4		
5		
6		

9.7.6 $\gamma = 1.25$, $x_T = -1.0$, $y_T = 20.0$

This is the case in Sec. 6.6, and is handled by obtaining six roots via (9.16) and selecting the one that maximizes $J(\phi)$ in (9.21). Results are shown in Table 6.6, and the optimal value of ϕ is $\phi_{optimal} =$

i	$\phi_i = Root \neq i$ of (9.16)	$J(\phi_i)$ in (9.21) [choose minimum]
1		
2		
3		
4		
5		
6		

Chapter 10: OPTIMAL HEADING ANGLE SIMULATION

In this chapter, we address the active target defense differential game where an Attacker missile pursues a Target aircraft. A Defender missile is fired by the Targets wingman in order to intercept the Attacker before it reaches the aircraft. Thus, a team is formed by the Target and the Defender which cooperate to maximize the distance between the Target aircraft and the point where the Attacker missile is intercepted by the Defender missile, while the Attacker tries to minimize said distance. The results shown here extend previous work. We consider here the case where the Defender is faster than the Attacker. The solution to this differential game provides optimal heading angles for the Target and the Defender team to maximize the terminal separation between Target and Attacker and it also provides the optimal heading angle for the Attacker to minimize the said distance.

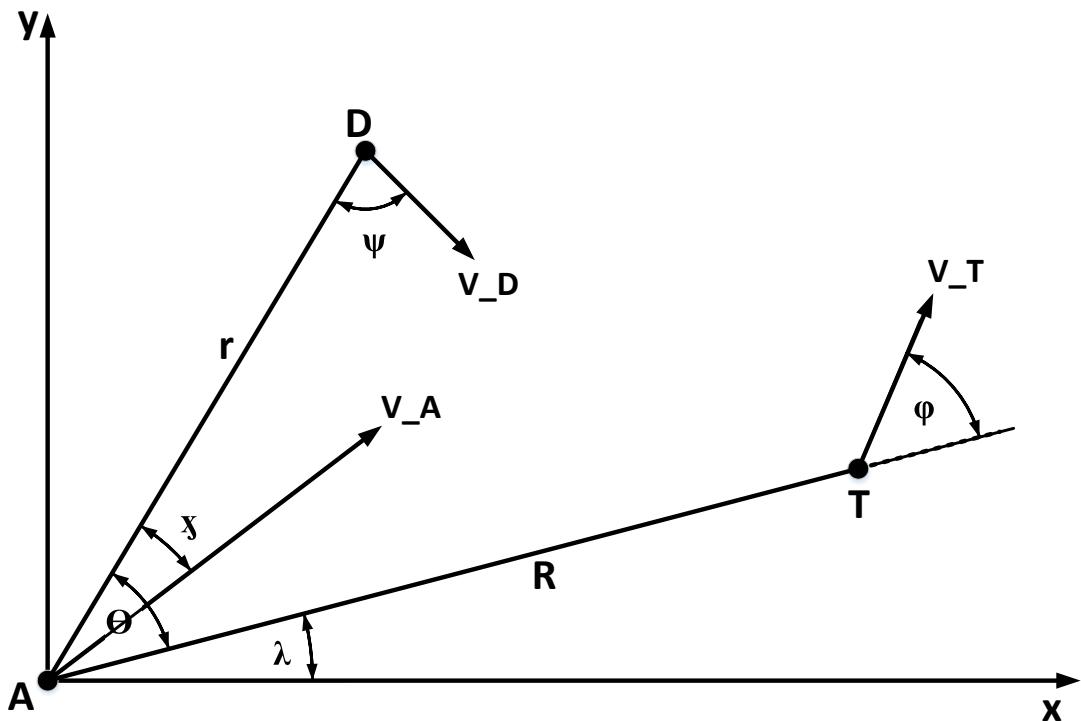


Figure 10.1: State space for the Attacker (A), Target (T), and Defender (D). The origin is arbitrary situated at the Attacker position (A). In polar coordinates with respect to arbitrary coordinates (x, y) , the various positions and speeds are

Figure 10.1 illustrates the state space for the ATD game. The variables R and r are the separations between the Target and Attacker, respectively, and hence denote the radial polar coordinates of the Target and Defender when the origin of coordinates is arbitrarily situated

at the position of the Attacker, The speeds of the Attacker, Target and Defender are denoted by \mathbf{V}_A , \mathbf{V}_T , \mathbf{V}_D with corresponding magnitudes V_A , V_T and V_D , and polar angles

$$\text{Ang } \mathbf{V}_A = \hat{\chi} = \lambda + \theta - \chi \quad (10.1)$$

$$\text{Ang } \mathbf{V}_T = \hat{\phi} = \lambda + \phi \quad (10.2)$$

$$\text{Ang } \mathbf{V}_D = \hat{\psi} = \lambda + \theta - (\pi - \psi) \quad (10.3)$$

The derivatives \dot{R} and \dot{r} can be obtained as differences in radial speeds, while $\dot{\theta}$ can be obtained in terms of the difference in circumferential speeds. The end points of the distance R are moving at a speed V_T at an angle ϕ with respect to the radial directions and a speed V_A at an angle $(\theta - \chi)$ with respect to the radial directions. Hence;

$$\begin{aligned} \dot{R} &= V_t \cos \phi - V_A \cos(\theta - \chi) \\ &= V_A(\alpha \cos \phi - \cos(\theta - \chi)) \end{aligned} \quad (10.4)$$

Similarly, we obtain

$$\begin{aligned} \dot{r} &= -V_A \cos \chi - V_D \cos \psi \\ &= V_A(-\cos \chi - \gamma \cos \psi) \end{aligned} \quad (10.5)$$

In the following, we will deliberately differ from Garcia et.al. [] We will not use reduced equations in which $V_A = V_D = 1$. Our equations will look dimensionally homogeneous to any reader, and we will allow $\gamma = \frac{V_D}{V_A}$ to differ from 1.

The circumferential speeds are

$$\begin{aligned} R \dot{\lambda} &= V_T \sin \phi - V_A \sin(\theta - \chi) \\ \dot{\lambda} &= V_A \left[\frac{\alpha}{R} \sin \phi - \frac{1}{R} \sin(\theta - \chi) \right] \\ r(\theta + \lambda) \cdot &= -V_D \sin \psi + V_A \sin \chi \end{aligned}$$

Hence, one obtains

$$\dot{\lambda} = V_A \left[\frac{\alpha}{R} \sin \phi - \frac{1}{R} \sin(\theta - \chi) \right] \quad (10.6)$$

$$\dot{\theta} + \dot{\lambda} = V_A \left[-\frac{\gamma}{r} \sin \psi + \frac{1}{r} \sin \chi \right] \quad (10.7)$$

Now, subtract (10.7) minus (10.6) to obtain

$$\dot{\theta} = V_A \left[-\frac{\alpha}{R} \sin \phi + \frac{1}{R} \sin(\theta - \chi) - \frac{\gamma}{r} \sin \psi + \frac{1}{r} \sin \chi \right] \quad (10.8)$$

Equations (10.4), (10.5), (10.8) reduce to equations (10.6) of Garcia et.al [] when we set both V_A and $\gamma = \frac{V_D}{V_A}$ equal to 1.

The system dynamics are given by (10.4), (10.5), and (10.8) for $0 \leq t \leq t_f$, together with the initial conditions

$$\begin{aligned} R(t_0) &= R_0, \\ r(t_0) &= r_0, \\ \theta(t_0) &= \theta_0, \end{aligned}$$

The objective of the Target-Defender team is to maximize the separation between the Target and the Attacker at the interception time $R(t_f)$, where the terminal time t_f is free, such that $r(t_f) = r_c$. The objective of the Attacker is to minimize the same distance $R(t_f)$. This can be expressed as

$$\max_{\phi, \psi} \min_{\chi} J = \int_{t_0}^{t_f} \dot{R} dt$$

The Hamiltonian is

$$\begin{aligned} H &= \cos(\theta - \chi) - \alpha \cos \phi \\ &\quad + [\alpha \cos \phi - \cos(\theta - \chi)] \lambda_R \\ &\quad - [\cos \chi + \cos \psi] \lambda_r \\ &\quad + [-\frac{\alpha}{R} \sin \phi + \frac{1}{R} \sin(\theta - \chi) - \frac{1}{r} \sin \psi + \frac{1}{r} \sin \chi] \lambda_\theta \end{aligned} \tag{10.9}$$

$$\begin{aligned} H &= -(1 - \lambda_R)[\alpha \cos \phi - \cos(\theta - \chi)] \\ &\quad - [\cos \chi + \beta \cos \psi] \lambda_R \\ &\quad + [-\frac{\alpha}{R} \sin \phi + \frac{1}{R} \sin(\theta - \chi) - \frac{\beta}{r} \sin \psi + \frac{1}{r} \sin \chi] \lambda_\theta \end{aligned}$$

We now find the optimal heading angles ψ^* , ϕ^* and χ^* . Since these angles are known to be positive or negative angles (ranging from $-\frac{\pi}{2}$ to $\frac{\pi}{2}$), it suffices to determine the tangent $\tan a$ of each angle a and then determine the cosine and sine from

$$\cos a = [1 + \tan^2 a]^{\frac{-1}{2}} \tag{10.10}$$

$$\sin a = \tan a [1 + \tan^2 a]^{\frac{-1}{2}} \tag{10.11}$$

We obtain ψ^* by partially differentiating the Hamiltonian (10.9) w.r.t. ψ and equating the derivative $\frac{\partial H}{\partial \psi}$ to zero, namely

$$\frac{\partial H}{\partial \psi} = \beta \lambda_r \sin \psi - \frac{\beta}{r} \lambda_\theta \cos \psi = 0 \tag{10.12}$$

Hence the optimal heading ψ^* is given by:

$$\tan \psi^* = \frac{\lambda_\theta / r}{\lambda_r} \tag{10.13}$$

Figure 10.5 shows that ψ^* is a (positive or negative) acute angle in a right-angled triangle

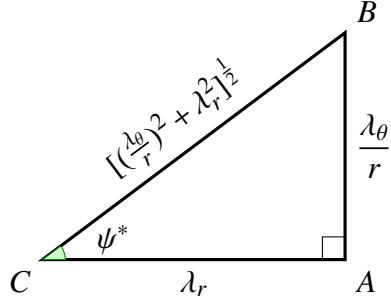


Figure 10.2: Defender

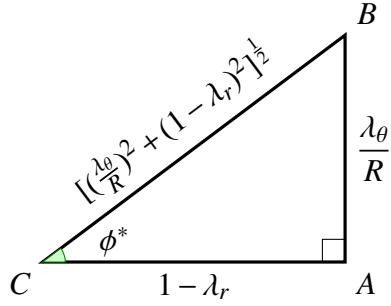


Figure 10.3: Target

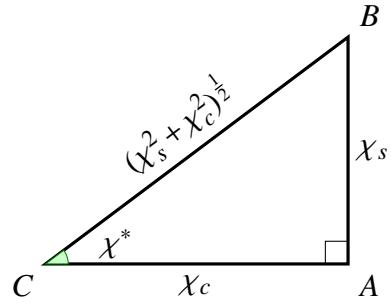


Figure 10.4: Attacker

Figure 10.5: Right-angled triangles that define the optimal headings ψ^* , ϕ^* and χ^* for the Defender, Target, and Attacker, respectively

with an opposite side equal to (λ_θ/r) , an adjacent side equal to λ_r and a hypotenuse equal to $[(\frac{\lambda_\theta}{r})^2 + \lambda_r^2]^{1/2}$. Using Fig.10.5 or relations (10.10) and (10.11) one obtains

$$\cos \psi^* = \lambda_r [(\frac{\lambda_\theta}{r})^2 + \lambda_r^2]^{-1/2} = \frac{r \lambda_r}{\sqrt{\lambda_\theta^2 + r^2 \lambda_r^2}} \quad (10.14)$$

$$\sin \psi^* = (\frac{\lambda_\theta}{r}) [(\frac{\lambda_\theta}{r})^2 + \lambda_r^2]^{-1/2} = \frac{\lambda_\theta}{\sqrt{\lambda_\theta^2 + r^2 \lambda_r^2}} \quad (10.15)$$

We further compute the second partial derivative of the Hamiltonian w.r.t. ψ , that is

$$\begin{aligned}\frac{\partial^2 H}{\partial \psi^2} &= \beta \lambda_r \cos \psi + \beta \frac{\lambda_\theta}{r} \sin \psi \\ &= [\beta r \lambda_r^2 + \beta \frac{\lambda_\theta^2}{r}] [\lambda_\theta^2 + r^2 \lambda_r^2]^{\frac{-1}{2}} > 0\end{aligned}\tag{10.16}$$

The fact that $\frac{\partial^2 H}{\partial \psi^2} > 0$ means that the extremum obtained by (10.12) is a minimum, i.e., the optimal value ψ^* minimizes the cost $(-J)$.

We now obtain ϕ^* by partially differentiating Hamiltonian (10.9) w.r.t. ϕ and equating the partial derivative $\frac{\partial H}{\partial \phi}$ to zero, namely

$$\frac{\partial H}{\partial \phi} = \alpha(1 - \lambda_R) \sin \phi - \frac{\alpha}{R} \lambda_\theta \cos \phi = 0\tag{10.17}$$

Hence the optimal heading ϕ^* is given by

$$\tan \phi^* = \frac{(\lambda_\theta/R)}{1 - \lambda_R}\tag{10.18}$$

Figure 10.5 shows that ϕ^* is a positive or negative acute angle in a right-angled triangle with an opposing side of length (λ_θ/R) , an adjacent side of length $(1 - \lambda_R)$, and a hypotenuse of length $[(\lambda_\theta/R)^2 + (1 - \lambda_R)^2]^{\frac{1}{2}}$.

Using Fig.10.5 or relations (10.10) and (10.11), we obtain

$$\cos \phi^* = (1 - \lambda_R)[(\lambda_\theta/R)^2 + (1 - \lambda_R)^2]^{\frac{-1}{2}} = \frac{R(1 - \lambda_R)}{\sqrt{\lambda_\theta^2 + R^2(1 - \lambda_R)^2}}\tag{10.19}$$

$$\sin \phi^* = (\lambda_\theta/R)[(\lambda_\theta/R)^2 + (1 - \lambda_R)^2]^{\frac{-1}{2}} = \frac{\lambda_\theta}{\sqrt{\lambda_\theta^2 + R^2(1 - \lambda_R)^2}}\tag{10.20}$$

The second partial derivative of the Hamiltonian w.r.t. ϕ is given by

$$\begin{aligned}\frac{\partial^2 H}{\partial \phi^2} &= \alpha(1 - \lambda_R) \cos \phi + \frac{\alpha}{R} \lambda_\theta \sin \phi \\ &= [\alpha R(1 - \lambda_R)^2 + \frac{\alpha}{R} \lambda_\theta^2][\lambda_\theta^2 + R^2(1 - \lambda_R)^2]^{\frac{-1}{2}} > 0\end{aligned}\tag{10.21}$$

The fact that $\frac{\partial^2 H}{\partial \phi^2} > 0$ means that the extrema obtained by (10.17) is minimum, i.e., the optimal heading ϕ^* minimizes the cost $(-J)$.

Finally, we obtain the optimal heading χ^* by partially differentiating the Hamiltonian (10.9) w.r.t. χ and equating the partial derivative $\frac{\partial H}{\partial \chi}$ to zero, namely

$$\frac{\partial H}{\partial \chi} = (1 - \lambda_R) \sin(\theta - \chi) + \lambda_r \sin \chi - \frac{\lambda_\theta}{R} \cos(\theta - \chi) + \frac{\lambda_\theta}{r} \cos \chi = 0\tag{10.22}$$

Using the trigonometric expansions:

$$\sin(\theta - \chi) = \sin \theta \cos \chi - \cos \theta \sin \chi\tag{10.23}$$

$$\cos(\theta - \chi) = \cos \theta \cos \chi + \sin \theta \sin \chi \quad (10.24)$$

We can rewrite (10.22) in the form

$$\frac{\partial H}{\partial \chi} = \chi_s \cos \chi - \chi_c \sin \chi = 0 \quad (10.25)$$

where

$$\chi_s = (1 - \lambda_R) \sin \theta - \frac{\lambda_\theta}{R} \cos \theta + \frac{\lambda_\theta}{r} \quad (10.26)$$

$$\chi_c = (1 - \lambda_R) \cos \theta + \frac{\lambda_\theta}{R} \sin \theta - \lambda_r \quad (10.27)$$

The optimal heading χ^* is obtained from (10.25) as

$$\tan \chi^* = \frac{\chi_s}{\chi_c} \quad (10.28)$$

Figure 10.5 shows that χ^* is a positive or negative acute angle in a right-angled triangle with an opposing leg of length χ_s , an adjacent leg of length χ_c , and a hypotenuse of length $[\chi_s^2 + \chi_c^2]^{\frac{1}{2}}$. Using Fig. 10.5 or relations (10.10) and (10.11), we obtain

$$\cos \chi^* = \chi_c [\chi_s^2 + \chi_c^2]^{\frac{-1}{2}} = \frac{\chi_c}{\sqrt{\chi_s^2 + \chi_c^2}} \quad (10.29)$$

$$\sin \chi^* = \chi_s [\chi_s^2 + \chi_c^2]^{\frac{-1}{2}} = \frac{\chi_s}{\sqrt{\chi_s^2 + \chi_c^2}} \quad (10.30)$$

The second derivative of the Hamiltonian w.r.t. χ is obtained from (10.25) as

$$\begin{aligned} \frac{\partial^2 H}{\partial \chi^2} &= -\chi_s \sin \chi - \chi_c \cos \chi \\ &= -[\chi_s^2 + \chi_c^2] (\chi_s^2 + \chi_c^2)^{\frac{-1}{2}} \\ &= -(\chi_s^2 + \chi_c^2)^{\frac{1}{2}} > 0 \end{aligned} \quad (10.31)$$

Equation (10.31) demonstrates definitely that $\frac{\partial^2 H}{\partial \chi^2}$ is negative. The corresponding claim via Eq.(28) of Garcia et. al. [] is not complete. The fact that $\frac{\partial^2 H}{\partial \chi^2} < 0$ means that the extremum obtained via (10.25) is a maximum, i.e., the optimal heading χ^* maximizes the cost $(-J)$.

In passing, we note that (10.22) can be simplified via (10.14), (10.15), (10.19) and (10.20) to give:

$$\begin{aligned} \frac{\partial H}{\partial \chi} &= [(\frac{\lambda_\theta}{R})^2 + (1 - \lambda_R)^2]^{\frac{1}{2}} [\cos \phi \sin(\theta - \chi) - \sin \phi \cos(\theta - \chi)] + [(\frac{\lambda_\theta}{r})^2 + \lambda_r^2]^{\frac{1}{2}} [\cos \psi \sin \chi + \sin \psi \cos \chi] \\ &= [(\frac{\lambda_\theta}{R})^2 + (1 - \lambda_R)^2]^{\frac{1}{2}} \sin(\theta - \chi - \phi) + [(\frac{\lambda_\theta}{r})^2 + \lambda_r^2]^{\frac{1}{2}} \sin(\chi + \psi) = 0 \end{aligned} \quad (10.32)$$

The expressions for the optimal heading angles (10.13), (10.18), and (10.28) are to be used in the numerical solution of the Two-Point Boundary Value Problem (TPBVP). That solution is found by substituting the optimal headings into the state equations (), and the co-state rquations () with the terminal conditions ().

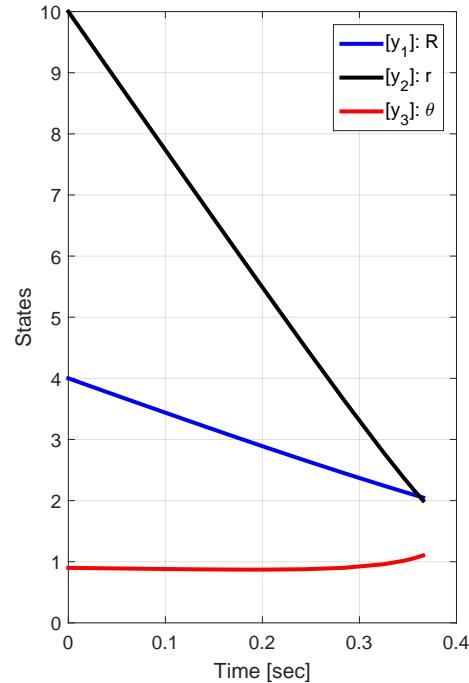


Figure 10.6: States

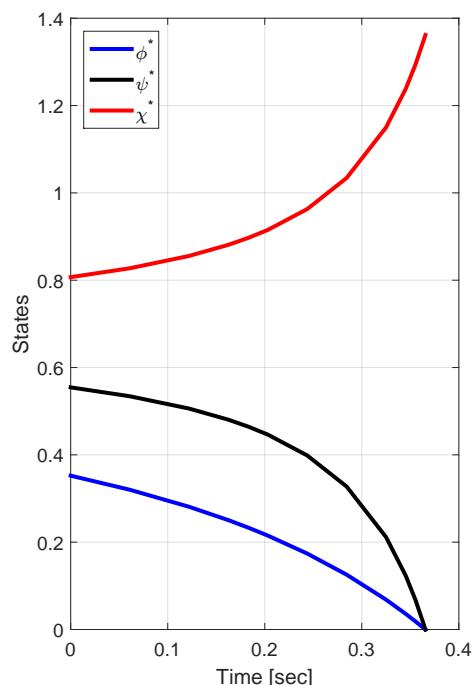


Figure 10.7: Optimal heading angles

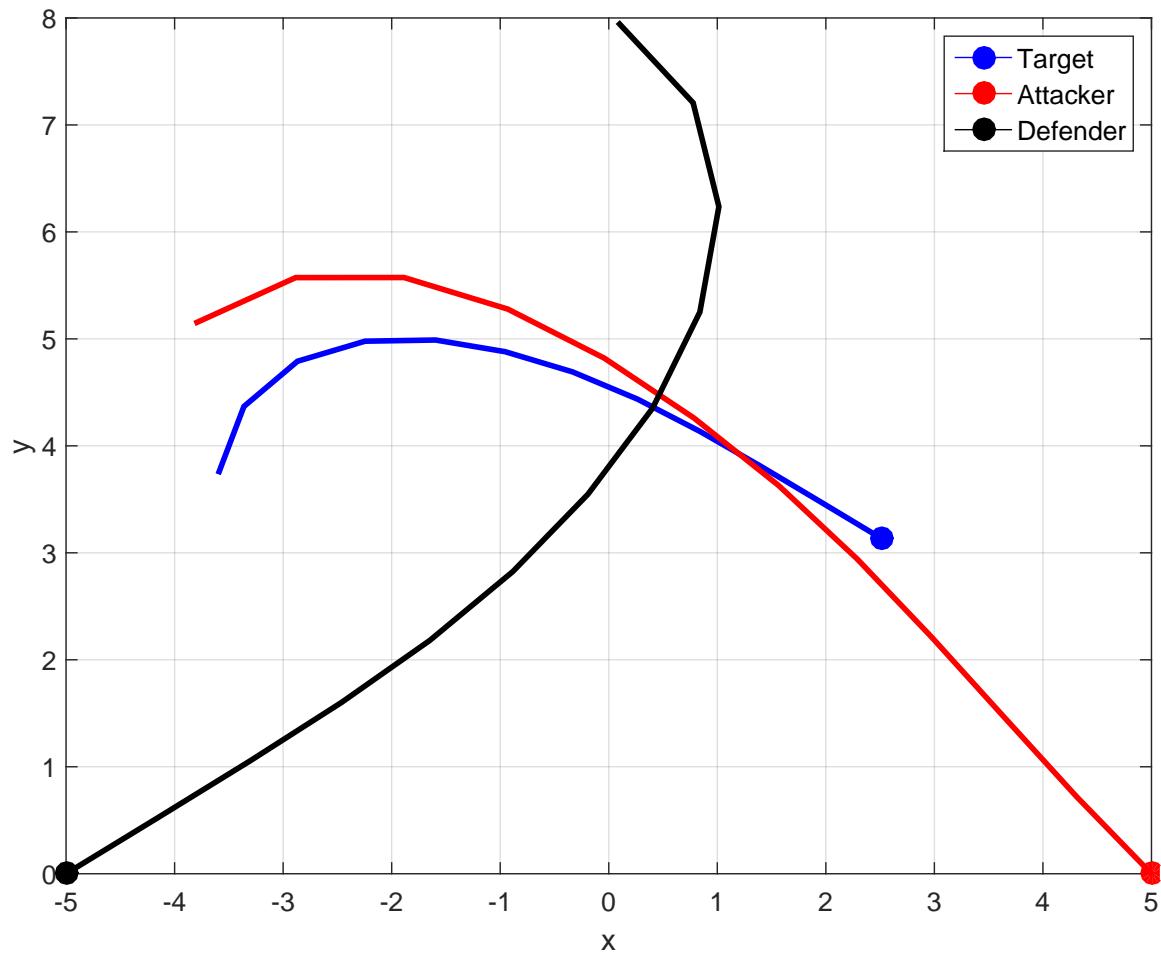


Figure 10.8: Optimal trajectories

Part IV

CONCLUSION AND FUTURE WORK

Chapter 11: CONCLUSION AND FUTURE WORK

This thesis offered a unified analytic solution of the *TAD* problem in which an Attacker is pursuing a Target while attempting to evade a Defender. The thesis reviews and extends the work that has recently appeared in [1, 2, 3], beside making the following new contributions:

1. The thesis covers all cases for the ratio γ of the Attacker's speed w.r.t the Defender's speed. It treats the case of a slow Defender $\gamma > 1$ for the first time, and presents this case along with the case of a fast Defender ($\gamma < 1$) discussed in [2] and the case of a similar Defender ($\gamma = 1$), discussed earlier in [1, 3].
2. The thesis demonstrates a striking increase in complexity when $\gamma \neq 1$ compared with the case $\gamma = 1$. It also demonstrates some sort of *duality* between the two cases of ($\gamma < 1$) and ($\gamma > 1$).
3. The thesis develops novel analytic expressions for the Voronoi diagrams for bordering the escape regions when ($\gamma < 1$) or ($\gamma > 1$). These expressions are more complex than the ones obtained in [3] for ($\gamma = 1$), and reduce to it as a limiting case.
4. The thesis offers a tutorial exposition of the *TAD* problem, uses simple arguments of plane geometry to develop the necessary Apollonius circle, utilizes equalities rather than inequalities in developing Voronoi diagram, and pays careful attention to the inadvertent inclusion of extraneous solutions so as to justify their subsequent rejection.
5. The thesis supplements its analysis with extensive computations for the critical speed ratio, Voronoi diagrams, and the optimal interception points. The results obtained encompass all possible values of γ , and they reduce to the already available results for $\gamma = 1$. Results for the trajectories and optimal interception points obtained agree with those obtained by the numerical solution of a two-point boundary value problem (TPBVP) utilizing Pontryagin's Maximum Principle [2].

Some possible extensions of the current work that warrant further exploration include:

1. Further analysis of the quartic equation obtained for the Voronoi diagram when $\gamma \neq 1$, with an aim to *split* it into two factors representing the rejected and accepted branches of the diagram.
2. Investigation of the sixth-degree complex polynomial equation for the optimal interception angle to get some *insight* about its six roots, and to find a better way for *selecting* desirable root.
3. Relaxation of some of the assumption used in this study. In particular, it is very interesting to consider the possibility of *variable* rather than constant speeds the three agents.
4. Addition of an element of *uncertainty* to the computation. For example, we might assume that the initial positions of the three agents are not deterministic but *stochastic* or *fuzzy*.

5. Extension of the current work to a more general situation involving several Targets, several Attackers and/or several Defenders.

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