

Problem 1:

Proof. Suppose for contradiction that we have a simple graph $G = (V, E)$, $|V| = n$, where no two vertices have the same degree. Let $V = \{v_1, v_2, \dots, v_n\}$. Observe that $\deg(v_i) \geq 0$ but each vertex can be ~~some~~ adjacent to at most $n-1$ other vertices, so $0 \leq \deg(v_i) \leq n-1$ for $i = 1, \dots, n$. Since $\deg(v_1) \neq \deg(v_2) \neq \dots \neq \deg(v_n)$, ~~we must have there must be a vertex~~ and there are $(n-1) - 0 + 1$ numbers between 0 and $n-1$ inclusive, there must be a vertex v_k with degree 0 and a vertex v_l with degree $n-1$ in G . Contradiction. Since v_l has degree $n-1$, it must be adjacent to v_k . But $\deg(v_k) = 0$ which is impossible. \square

Problem 2.

Proof. Suppose that G is connected ($G=(V,E)$).

(\rightarrow) Suppose G is not complete. We wish to show that we can remove v_1, \dots, v_m such that G becomes disconnected. Observe that G can only be connected but not complete when

$n > 2$. Let $G = G(V \setminus \{v_1, \dots, v_m\}, E')$ where E' is

~~the subgraph of E but with edges incident to v_1, \dots, v_m removed~~

Since G is ~~disconnected~~ ^{not complete}, there exists $v_{m+1}, v_{m+2} \in V$ such that $\{v_{m+1}, v_{m+2}\} \notin E'$. Now remove all vertices and their incident edges except for v_{m+1}, v_{m+2} . We have a disconnected graph.

(\leftarrow) Suppose that we can remove v_1, \dots, v_m to make G disconnected. Let G' be the subgraph of G with v_1, \dots, v_m and their incident edges removed. Since G' is disconnected, we can find $v_{m+1}, v_{m+2} \in G'$ such that no path exists

between them. Thus, $\{v_{m+1}, v_{m+2}\} \notin G'$. But $\{v_{m+1}, v_{m+2}\}$ is not incident with v_1, \dots, v_m , so it cannot be one of the edges we removed to get G' . Therefore, $\{v_{m+1}, v_{m+2}\} \notin G$ and G is not complete as desired. \square

