## 0.1 Main algorithm

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Algorithm 1: main algorithm

Input: Demand matrix DM

Output: Routing \phi_{best} (Fraction of demand)

1 Initialization : D = \{DM\}, DM_{worst} = 0;

// D is the set of demand matrix

2 while not meet the cutoff condition do

3 D = D \cup DM_{worst};

4 Find the best routing \phi of D; // The First model

5 \phi_{best} = \phi;

6 Find the worst demand matrix DM_{worst} of \phi_{best} // The second model

7 end

8 return \phi_{best}
```

## 0.2 Original model in the paper

$$\begin{aligned} \min \alpha \\ \phi \text{ is a PD routing} \\ \forall \text{ edges } e = (u, v) \\ \forall \text{ DMs D} \in D \text{ with } OPTU(D) = r : \\ \sum_{(s,t)} \frac{d_{st} f_{st}(u) \phi_t(e)}{c(e)} \leq \alpha r \end{aligned} \tag{1}$$

$$OPTU(D) = \min_{\phi \mid \phi \text{ is a PD routing}} MxLU(\phi, D)$$
 (2)

## 0.3 Model 1: The first model

Find the best routing  $\phi$  of the given set of demand matrix to minimize maximum link ultilization

 $\min \alpha$  (Minimize maximum link ultilization)

$$\sum_{e=(v,u)\in E} \phi_{st}(e) - f^i_{st}(u) = \begin{cases} 1 & \text{if } u = s \\ -1 & \text{if } u = t \\ 0 & \text{others} \end{cases}$$

$$, \forall (s,t) \in V \times V, 1 \leq i \leq |D|, \forall u \in V \qquad \qquad \text{(Flow conservation)}$$

$$\sum_{(s,t)} \frac{d^i_{st} f^i_{st}(u) \phi_{st}(e)}{c(e)} \leq \alpha \ , \ \forall e \in E, 1 \leq i \leq |D| \qquad \text{(Link capacity constraint)}$$

$$0 \leq \phi_{st}(e) \leq 1 \ , \forall e \in E \qquad \qquad \text{(Decision variable constraint)}$$

$$0 \leq \alpha \leq 1 \qquad \qquad \text{(Decision variable constraint)}$$

$$(3)$$

## 0.4 Model 2: The second model

Find the worst demand matrix with given routing  $\phi$  to maximize maximum link ultilization

 $\max \rho$  (Maximize minimum link ultilization)

$$\sum_{e=(v,u)\in E} \phi_{st}(e) - f_{st}(u) = \begin{cases} 1 & \text{if } u = s \\ -1 & \text{if } u = t \text{ ,} \forall s,t,u \in V \\ 0 & \text{others} \end{cases}$$
 (Flow conservation) 
$$\rho \leq \sum_{(s,t)} \frac{d_{st} f_{st}(u) \phi_{st}(e)}{c(e)} \leq 1$$
 (Link capacity constraint) 
$$0 \leq \rho \leq 1$$
 (Decision variable constraint) 
$$(4)$$