

CSED601

Dependable Computing

Lecture 13

Jong Kim
Dept. of CSE
POSTECH

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Review of Previous Lecture

- Reliability Evaluation Techniques
 - Markov modeling
 - Discrete Parameter Markov Chain
 - Continuous Parameter Markov Chain

Fault Trees

- Concept
 - A dual of reliability block diagram
 - Logic failure diagram
 - Think in terms of logic where
 - 0 = operating, 1 = failed
- Diagram
 - AND gate
 - All inputs must fail for the gate to fail
 - OR gate
 - Any input failure causes the gate to fail
 - K-of-n gate
 - K or more input failures cause gate to fail

Fault Trees

- Example
 - Triplex Bus Guardian
 - Three serially connected pass transistor
 - Two failure states
 - Failed Active : signal pass but does not function as a guardian
 - Failed Passive : signal cut
 - Modeling
 - Active mode
 - If three switches are failed active, then the system fail → AND gate of Failed Active
 - Passive mode
 - Cut off with any single unit failure → OR gate

Fault Trees

- Example
 - Total failure of Bus guardian
 - Caused by either active or passive mode
 - Draw your diagram here
 - K-of-N

Petri Nets

- Refer the slide made by Dr. Axel Krings
 - <http://www2.cs.uidaho.edu/~krings/CS449/Notes.S13/449-13-11.pdf>

GSPN

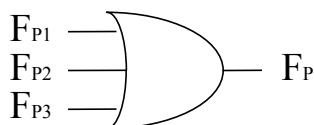
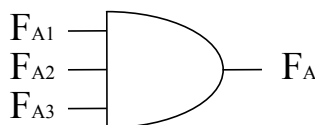
- General Stochastic Petri Net
- Refer the slide made by Dr. Axel Krings
 - <http://www2.cs.uidaho.edu/~krings/CS449/Notes.S13/449-13-12.pdf>

Fault Trees

- ◆ Fault Trees
 - dual of Reliability Block Diagram
 - logic failure diagram
 - think in terms of logic where
 - » 0 = operating, 1 = failed
- ◆ AND Gate
 - all inputs must fail for the gate to fail
- ◆ OR Gate
 - any input failure causes the gate to fail
- ◆ k-of-n Gate
 - k or more input failures cause gate to fail

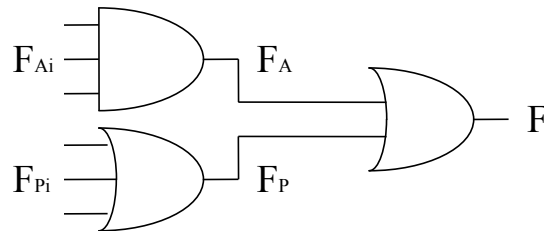
e.g. Triplex Bus Guardian

- ◆ Active mode
 - M₁ and M₂ and M₃ fail =>
 - AND Gate
- ◆ Passive Mode
 - “cutoff” with any single unit failure =>
 - OR Gate



e.g. Triplex Bus Guardian

- ◆ Total Failure
 - caused by either active or passive mode

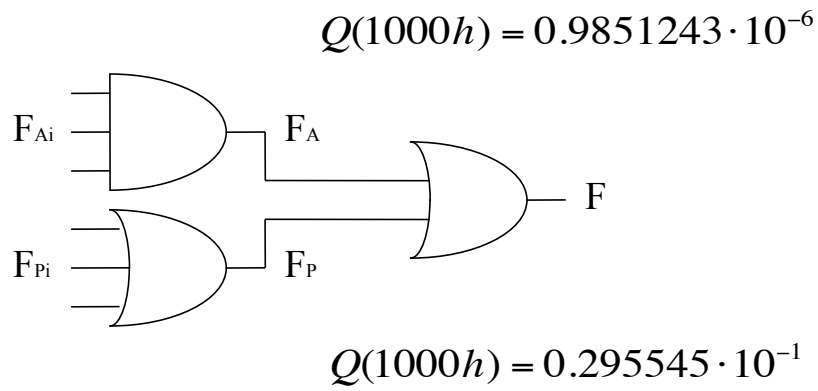


e.g. Triplex Bus Guardian

- ◆ How can one use the fault tree effectively to isolate those parts of the system that need reliability considerations?

e.g. Triplex Bus Guardian

◆ Combined fault model



Examples

- ◆ Simple Passive TMR (no diagnosis)
 - RBD = (2 of 3): 2 operable \Rightarrow System operable
 - F-Tree = (2 of 3): 2 failed \Rightarrow System failed
- ◆ Simple TMR with *Benign* failures
 - RBD = (1 of 3): 1 operable \Rightarrow System operable
 - F-Tree = (3 of 3): 3 failed \Rightarrow System failed
- ◆ Summary
 - Parallel \Rightarrow AND
 - Series \Rightarrow OR
 - K-of-N \Rightarrow (n-k+1 of n)

Petri Nets

- ◆ Part of this discussion is based on the paper
 - *Petri Nets: Properties, Analysis and Applications*
 - by Tadao Murata, Proc. IEEE, Vol. 77, No. 4, April 1989.
- ◆ Petri Nets
 - graphical and mathematical modeling tool
 - tool for describing systems characterized as being:
 - » concurrent, asynchronous, distributed, parallel, nondeterministic and/or stochastic

Petri Nets

- ◆ History
 - **1962:** Carl Adam Petri's submitted his dissertation at the Uni. Darmstadt, Germany
 - **1970:** early development was published by A.W. Host and in the records of the 1970 Project MAC Conference on Concurrent Systems and Parallel Computation
 - **1970-75:** Computation Structure Group and MIT was most active
 - **1975:** conference on Petri Nets and Related Methods at MIT
 - **1979:** 135 researchers assembled in Hamburg, Germany, for 2-week advanced course on General Net Theory of Processes and Systems
 - **1980:** first European Workshop on Applications and Theory of Petri Nets, Strasbourg, France.
 - check out Murata's paper for the extensive literature discussion

Petri Nets


◆ General:

- directed, weighted, bipartite graph
- two kinds of nodes (Places P, Transitions T)
- arcs from P to T or from T to P
- arcs have integer weights
- non-negative Place weights are called tokens


Petri Nets

◆ A Petri Net is a 5-tuple $PN = \{P, T, A, W, M_0\}$

◆ Place Set $P = \{p_1, p_2, \dots, p_m\}$

- finite set of places
- condition = place
- one condition or set of atomic conditions
- symbol 

◆ Transition Set $T = \{t_1, t_2, \dots, t_n\}$

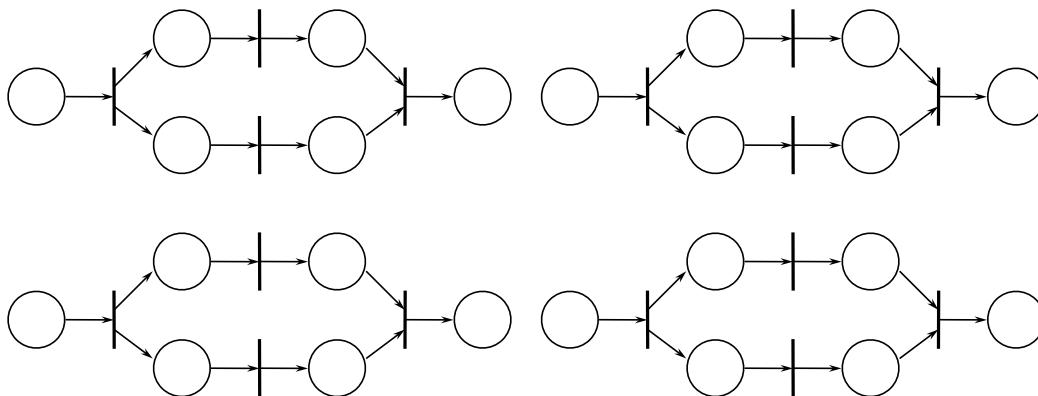
- finite set of transitions
- action = transition
- one action or set of atomic transitions
- symbol 

Petri Nets

- ◆ Arc Set $A \subseteq (P \times T) \cup (T \times P)$
 - set of directed arcs
 - edge of graph = arc
 - symbol \longrightarrow
- ◆ Weight Function $W = A \rightarrow \{1, 2, 3, \dots\}$
 - weights are associated with arcs
- ◆ Initial Marking $M_0 = P \rightarrow \{0, 1, 2, \dots\}$
 - the initial assignment of tokens to places

Petri Nets

- ◆ example



Petri Nets

◆ Dynamic Behavior

- during simulation of a petri net the state of the net may change
- change of state:
 - » transitions can be enabled
 - » enabled transitions may fire
 - » firing transition changes the marking of the net
 - » the marking is the “snap-shot” of all the tokens

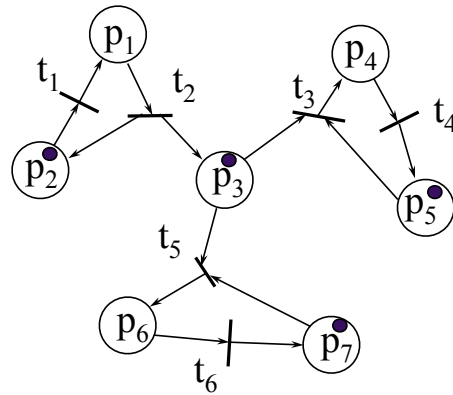
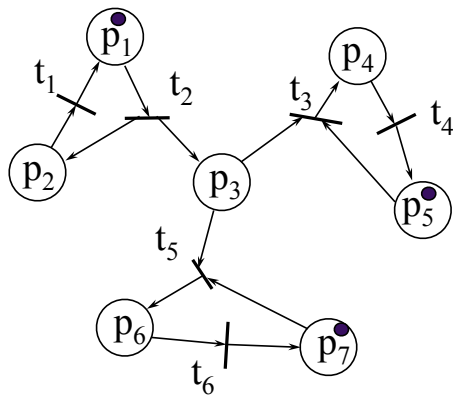
Petri Nets

◆ Firing rules

- A transition T is said to be *enabled* if each input place P is marked with at least $W(P,T)$ tokens
 - » $W(P,T)$ is the weight of the arc from P to T
- An enabled transition may or may not fire (depending on whether or not the event actually takes place).
- A *firing* of an enabled transition T removes $W(P,T)$ tokens from each input place P of T , and adds $W(T,P)$ tokens to each output place P of T
 - » $W(T,P)$ is the weight of the arc from T to P
- Common misconception: When a transition fires, it does **not** move tokens
 - » i.e. the number of tokens in the system is not necessarily constant

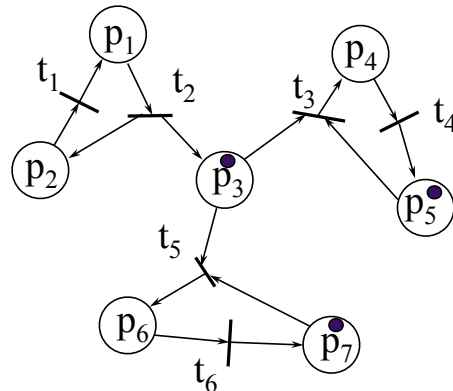
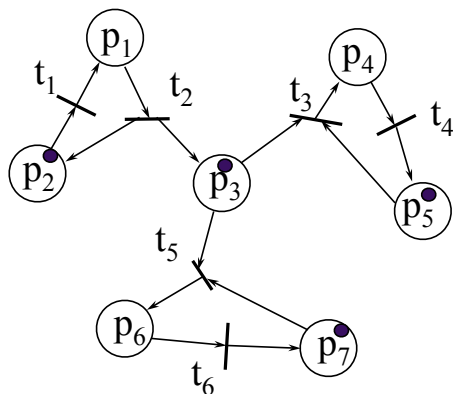
Petri Nets

- ◆ Example: assume the following initial marking
 - Only one transition is enabled, i.e. t_2



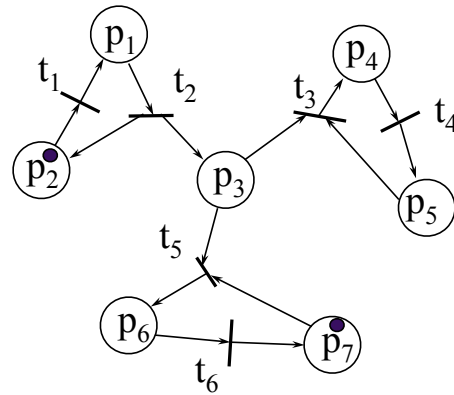
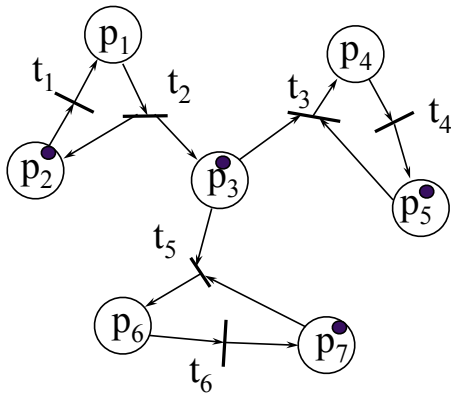
Petri Nets

- Now several transitions are enabled, i.e. t_1 , t_3 and t_5
- if t_1 fires first



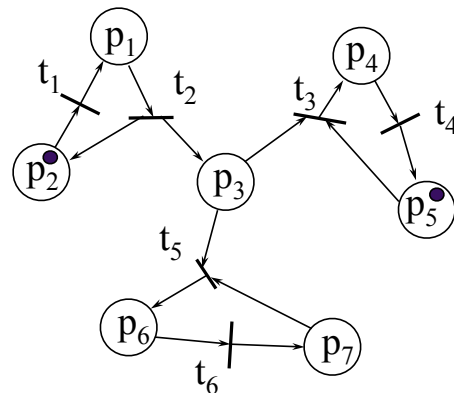
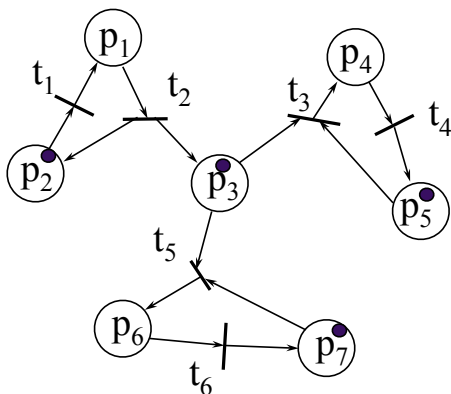
Petri Nets

- if t_3 fires first



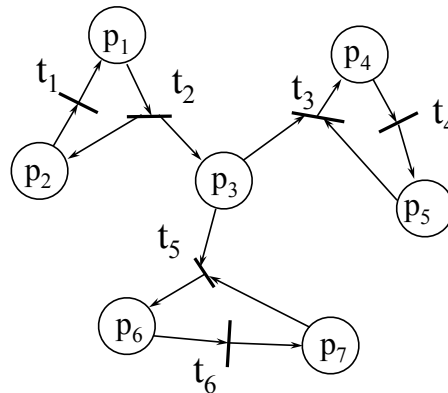
Petri Nets

- if t_5 fires first
- t_3 and t_5 are said to be in conflict



Petri Nets

- what could this Petri net represent?

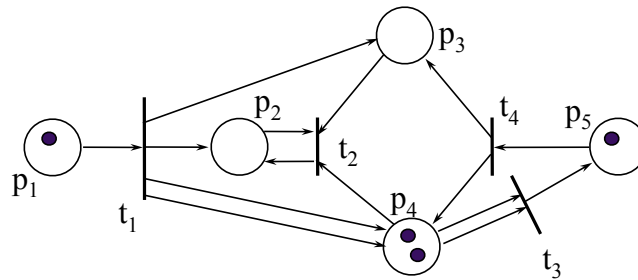


Petri Nets

- ◆ Marking: Number and placement of tokens
 - let $m_i = \#$ of tokens in place p_i
 - then marking
$$M = \{m_1, m_2, \dots, m_n\}$$
 - marking -- system state
 - Advantage: economy of model
 - » e.g. assume net with 6 places
 - we limit each place to maximal 1 token
 - then there are 2^6 possible markings
 - \Rightarrow 64 states
 - thus Petri Nets are a lot smaller than state diagrams, i.e. Markov chains

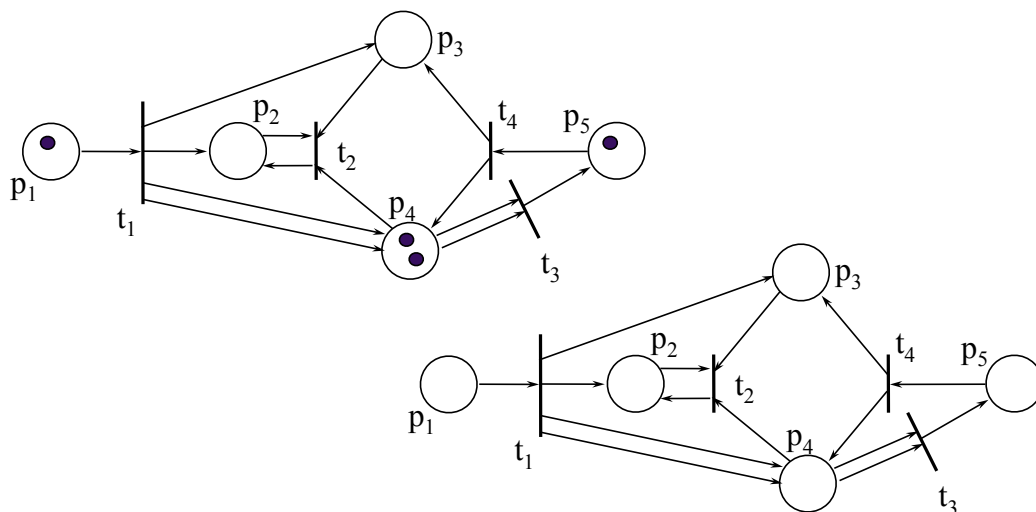
Petri Nets

- ◆ Firing rules
 - transition 1,3 and 4 are enabled



Petri Nets

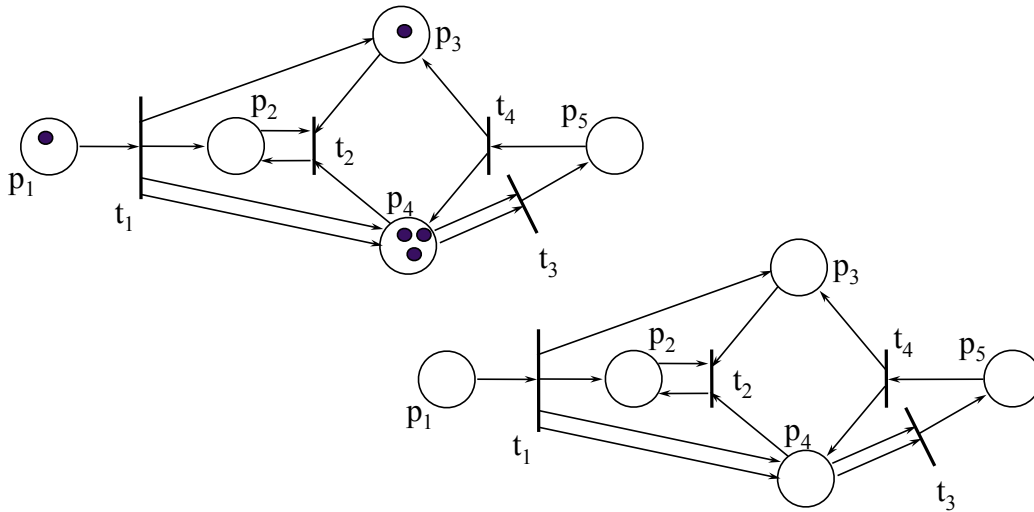
- ◆ Firing rules
 - transition 4 fires



Petri Nets

◆ Firing rules

- transition 1 fires



Petri Nets

◆ Firing rules

- transition 3 fires

