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## 1 Preface

# 2 Introduction to quivers and category theory

# 3 Datatype convention of catreps

Since the goal of this thesis is a translation of the package catreps by Peter Webb et al. into CAP, this section is a short overview of the package catreps.

In this package a category is stored as a concrete category (i.e. a category where the objects are sets and morphisms are maps of sets). A category is stored as a record (cat, say) with fields cat.objects, cat.generators, cat.domain, cat.codomain. Each object in the list cat.object is a set, and each morphism in the list of generator morphisms cat.generators is stored as a mapping of sets, which we notate as the list of its values. ([5])

```
Example

gap> c3c3 := ConcreteCategory( [ [2,3,1], [4,5,6], [,,,5,6,4] ] );

rec( codomain := [ 1, 2, 2 ], domain := [ 1, 1, 2 ],

generators := [ [ 2, 3, 1 ], [ 4, 5, 6 ], [ ,,, 5, 6, 4 ] ],

objects := [ [ 1, 2, 3 ], [ 4, 5, 6 ] ], operations := rec( ) )
```

The list of values as seen in the example above may be easy to type in, but does have its disadvantages: If for example you want to store the morphism that maps the set  $\{9\}$  to itself, i.e. the identity morphism  $1_{\{9\}}$ , you first have to write the eight commas that are not part of that morphism definition  $[\ ,,,,,,,,9]$  and you might make a mistake by forgetting one comma. Another issue is that the source object of a morphism gen is only implicitly given by those list entries i for which IsBound(gen[i]) = true.

Using instead MapOfFinSets in CAP solves both of these issues, and it lets us use a different model for concrete categories in CAP, i.e. that of a subcategory of FinSets, for which we already have an implementation in CAP. Another advantages of this method is that a MapOfFinSets can cache known properties about itself:

```
gap> S := FinSet( [1,2,3] );
<An object in FinSets>
gap> T := FinSet( [4,5,6] );
<An object in FinSets>
gap> map1 := MapOfFinSets( S, [ [1,1], [2,2], [3,3] ], S );
<A morphism in FinSets>
gap> IsAutomorphism( map1 );
true
gap> map1;
<An automorphism in FinSets>
```

Going further in the cited example,

"The following constructs a representation:" ([5])

```
Example
  gap> one:=One(GF(3));;
 gap> d:=[[1,1,0,0,0],[0,1,1,0,0],[0,0,1,0,0],[0,0,0,1,1],[0,0,0,0,1]]*one;;
 gap> e:=[[0,1,0,0],[0,0,1,0],[0,0,0,0],[0,1,0,1],[0,0,1,0]]*one;;
 gap> f:=[[1,1,0,0],[0,1,1,0],[0,0,1,0],[0,0,0,1]]*one;;
 gap> nine:=CatRep(c3c3,[d,e,f],GF(3));
category := rec( generators := [ [ 2, 3, 1 ], [ 4, 5, 6 ], [ ,,, 5, 6, 4 ] ]
, operations := rec(), objects := [[1, 2, 3], [4, 5, 6]],
domain := [ 1, 1, 2 ], codomain := [ 1, 2, 2 ] ),
genimages := [ [ [Z(3)^0, Z(3)^0, 0*Z(3), 0*Z(3), 0*Z(3)],
[ 0*Z(3), Z(3)^0, Z(3)^0, 0*Z(3), 0*Z(3) ],
[ 0*Z(3), 0*Z(3), Z(3)^0, 0*Z(3), 0*Z(3) ],
[ 0*Z(3), 0*Z(3), 0*Z(3), Z(3)^0, Z(3)^0 ],
[ 0*Z(3), 0*Z(3), 0*Z(3), 0*Z(3), Z(3)^0 ] ],
[ [ 0*Z(3), Z(3)^0, 0*Z(3), 0*Z(3) ], [ 0*Z(3), 0*Z(3), Z(3)^0, 0*Z(3) ]
, [ 0*Z(3), 0*Z(3), 0*Z(3), 0*Z(3) ],
[ 0*Z(3), Z(3)^0, 0*Z(3), Z(3)^0 ],
[ 0*Z(3), 0*Z(3), Z(3)^0, 0*Z(3) ] ],
[ [ Z(3)^0, Z(3)^0, 0*Z(3), 0*Z(3) ], [ 0*Z(3), Z(3)^0, Z(3)^0, 0*Z(3) ]
, [ 0*Z(3), 0*Z(3), Z(3)^0, 0*Z(3) ],
[0*Z(3), 0*Z(3), 0*Z(3), Z(3)^0]], field := GF(3),
dimension := [5, 4]
```

we see that catreps works with GAP matrices directly whereas with CAP we use HomalgMatrix and RingsForHomalg which lets us delegate computation to faster computer algebra systems like Singular or Magma. What is also noticable is the big chunk of output we get as a result of CatRep(c3c3,[d,e,f],GF(3)). In CAP we hide the output and give a short description of the resulting object or morphism, and use the Display function to display the whole result.

All in all, there are plenty of reasons to change to CAP. In the meantime, in order to still support inputs in the convention of catreps, I wrote a converter function ConvertToMapOfFinSets.

# 4 The category FinSets

There are algorithms whose sole purpose is to convert data structures, so they are not of much interest to the mathematical theory, and then there are algorithms that implement our category theoretical calculations, so they are important to our theory.

The algorithms ConvertToMapOfFinSets,
ConcreteCategoryForCAP and RightQuiverFromConcreteCategory
are more of the data structure conversion type, while RelationsOfEndomorphisms,
Algebroid,
EmbeddingOfSubRepresentation
and WeakDirectSumDecomposition are also important to our theory.

## 4.1 MapOfFinSets

### Algorithm 1: ConvertToMapOfFinSets

**Input :** a list *objects* of objects in FinSets and a morphism *gen* given as a list of images in the convention of catreps

**Output:** the corresponding map of finite sets from source S to target T

```
1 let T be the first object O \in objects such that gen \cap O \neq \emptyset;
```

```
2 if gen \cap O = \emptyset \ \forall O \in objects then
```

**3** Error "unable to find target set"

4 end

**5** let fl be the flattening of *objects* as a list;

6 let S be the sublist of fl according to positions i such that gen[i] is bound;

**7** set S to be the first object  $O \in objects$  such that O = S;

s if  $S \neq O \forall O \in objects$  then

9 | Error "unable to find source set"

10 end

11 let G be the list of pairs  $[i, gen[i]], i \in S$ ;

12 return MapOfFinSets(S, G, T);

# 5 The categories Functor-Categories and Cat-Reps

### 6 Conclusion

### References

- [1] https://web.northeastern.edu/martsinkovsky/p/Parnu2019/slides-facchini.pdf
- [2] https://www.math.uni-bielefeld.de/~sek/kau/leit4.pdf
- [3] Jan Geuenich. https://hss.ulb.uni-bonn.de/2017/4681/4681.pdf
- [4] Mohamed Barakat, Julia Mickisch and Fabian Zickgraf, FinSetsForCAP, https://github.com/mohamed-barakat/FinSetsForCAP/
- [5] Peter Webb, Dan Christensen, Fan Zhang, and Moriah Elkin, catrepstutorialMarch11, http://www-users.math.umn.edu/~webb/GAPfiles/catrepstutorial.html

## A Implementation in Cap

# Listings

1	ConvertToMapOfFinSets
2	ConcreteCategoryForCAP
3	RightQuiverFromConcreteCategory
4	RelationsOfEndomorphisms 6
5	Algebroid
6	EmbeddingOfSubRepresentation
7	WeakDirectSumDecomposition

### Procedure 1: ConvertToMapOfFinSets

```
function( objects, gen )
   local 0, T, f1, S, G, i;
   T := First( objects, 0 -> Length(
   Intersection( gen, AsList( 0 ) ) > 0 );
   if T = fail then
       Error( "unable to find target set\n" );
   fl := Flat( List( objects, 0 -> AsList( 0 ) ));
   S := fl{PositionsProperty( fl , i -> IsBound( gen[i] ))};
   S := First(objects, O \rightarrow AsList(O) = S);
   if S = fail then
       Error( "unable to find source set\n" );
   fi;
   G := [];
   G := List(S, i \rightarrow [i, gen[i]]); # gen[i] is sure to be bound
   return MapOfFinSets( S, G, T );
end );
```

Back to Algorithm 1. Back to Index

Procedure 2: ConcreteCategoryForCAP

```
function( L )
  local C, c, objects;

DeactivateCachingOfCategory( FinSets );
CapCategorySwitchLogicOff( FinSets );
DisableSanityChecks( FinSets );
```

```
C := Subcategory( FinSets, "A finite concrete category" :
   overhead := false, FinalizeCategory := false );
   DeactivateCachingOfCategory( C );
   CapCategorySwitchLogicOff( C );
   DisableSanityChecks( C );
   SetFilterObj( C, IsFiniteConcreteCategory );
   AddIsAutomorphism(C,
     function( alpha )
       return IsAutomorphism( UnderlyingCell( alpha ) );
   end );
   AddInverse( C,
     function( alpha )
       return Inverse( UnderlyingCell( alpha ) ) / CapCategory( alpha );
   end );
   c := ConcreteCategory( L );
   C!.ConcreteCategoryRecord := c;
   objects := List( c.objects, FinSet );
   SetSetOfObjects( C, List( objects, o -> o / C ) );
   SetSetOfGeneratingMorphisms( C, List(
   c.generators, g -> ConvertToMapOfFinSets( objects, g ) / C ) );
   Finalize( C );
   return C;
end );
```

### Procedure 3: RightQuiverFromConcreteCategory

```
];
i := i+1;
od;
q := RightQuiver( "q(1)[a]", Length( objects ), arrows );
return q;
end );
```

### Procedure 4: RelationsOfEndomorphisms

```
function( k, C )
 local objects, gmorphisms, q, kq, relation_of_endomorphism,
       arrows, endos, vertices, i, mor, mpowers, m, npowers, n, foundEqual,
       relsEndo:
 objects := SetOfObjects( C );
 gmorphisms := SetOfGeneratingMorphisms( C );
 q := RightQuiverFromConcreteCategory( C );
 kq := PathAlgebra( k, q );
 relation_of_endomorphism := function(kq, a, m, n)
     local rel, one;
     rel := [];
     if m = 0 then
         one := Source( a );
         rel := PathAsAlgebraElement( kq, a )^n
               - PathAsAlgebraElement( kq, one );
     else
         rel := PathAsAlgebraElement( kq, a )^(m+n)
               - PathAsAlgebraElement( kq, a )^m;
     fi;
     return rel;
 end:
 arrows := Arrows( q );
 endos := Filtered( arrows, a -> Source( a ) = Target( a ) );
 vertices := Collected( List( endos, Source ) );
 if ForAny( vertices, 1 -> 1[2] > 1 ) then
     Error( "we assume at most 1 generating endomorphism per vertex\n" );
 fi:
 relsEndo := [];
 for i in [ 1 .. Length( gmorphisms ) ] do
     mor := gmorphisms[i];
     if not IsEndomorphism( mor ) then
         continue;
     fi;
     mpowers := [];
```

```
m := 0;
       # sigma lemma
       foundEqual := false;
       while not mor^m in mpowers do
          n := 1;
          npowers := [];
          while not mor^(m+n) in npowers and
            not foundEqual do
              if IsCongruentForMorphisms( mor^(m+n), mor^m ) then
                  Add( relsEndo,
                      relation_of_endomorphism( kq, arrows[i], m, n ) );
                  foundEqual := true;
              fi;
              Add( npowers, mor^(m+n) );
              n := n+1;
          Add( mpowers, mor^m );
          m := m+1;
       od;
   od;
   return relsEndo;
end );
```

### Procedure 5: Algebroid

```
function( k, C )
 local objects, gmorphisms, q, kq, relEndo, A, F, vertices, rel,
       func, st, s, t, homST, list, p, pos;
 objects := SetOfObjects( C );
 gmorphisms := SetOfGeneratingMorphisms( C );
 q := RightQuiverFromConcreteCategory( C );
 kq := PathAlgebra( k, q );
 relEndo := RelationsOfEndomorphisms( k, C );
 A := Algebroid( kq, relEndo );
 kq := UnderlyingQuiverAlgebra( A );
 F := CapFunctor( A, objects, gmorphisms, C );
 vertices := List( SetOfObjects(A), UnderlyingVertex );
 rel := [];
 func :=
   function( p, 1 )
     return ForAny( 1, p1->
                   IsCongruentForMorphisms(
                          ApplyToQuiverAlgebraElement(F, p),
                          ApplyToQuiverAlgebraElement( F, p1 ) )
                  );
 end;
 for st in Cartesian(vertices, vertices) do
```

```
s := st[1];
       t := st[2];
       if s = t then
          continue;
       fi:
       homST := BasisPathsBetweenVertices( kq, s, t );
       homST := List( homST, p -> PathAsAlgebraElement( kq, p ) );
       list := [];
       for p in homST do
          pos := PositionProperty( list, 1->func(p,1) );
          if IsInt(pos) then
              Add( list[pos], p );
          else
              Add(list, [p]);
          fi;
       list := List( list, l-> List( l, p -> p!.representative ) );
       Append( rel, list );
   od;
   rel := Filtered( rel, 1 -> Length(1)>1 );
   rel := List( rel, 1 -> List( 1{[ 2 .. Length(1) ]}, p -> 1[1]-p ) );
   rel := Flat( rel );
   rel := Concatenation( relEndo, rel );
   kq := PathAlgebra( kq ) / rel;
   kq := PathAlgebra( kq ) / GroebnerBasis( IdealOfQuotient( kq ) );
   kq := Algebroid( kq );
   SetUnderlyingCategory( kq, C );
   SetIsLinearClosureOfACategory( kq, true );
   return kq;
end );
```

#### Procedure 6: EmbeddingOfSubRepresentation

```
function( eta, F )
  local kq, objects, morphisms, subrep, embedding;

kq := Source( CapCategory( F ) );

eta := List( eta, function( eta_o ) if IsMonomorphism( eta_o ) then
    return eta_o; fi; return ImageEmbedding( eta_o ); end );

objects := List( eta, Source );
morphisms := List(
```

### Procedure 7: WeakDirectSumDecomposition