

$$n = p + q.$$

Problem 1

CU's top-secret alien communications program has hired you to decode some radio signals they recently picked up. It turns out two different groups of aliens have been sending us messages for YEARS, and we haven't noticed! You've figured out the problem: the two species' transmissions have gotten mixed up.

Specifically, alien species A has been transmitting the signal $a_1a_2 \dots a_p$, while species B has been transmitting the signal $b_1b_2 \dots b_q$. Each species has been transmitting their message non-stop, but the signals have gotten jumbled: instead of receiving $a_1 \dots a_p a_1 \dots a_p a_1 \dots a_p \dots$ and $b_1 \dots b_q b_1 \dots b_q b_1 \dots b_q \dots$, we've been receiving a mix of the two strings' bits, which might look something like: $a_1a_2b_1b_2a_3a_4a_5b_3a_6b_4 \dots$. The bits are all there, and they're in order - the two strings are just mixed up.

You have an idea what the strings are, but you want to write a program to confirm that you've gotten it right. Write a dynamic program that takes in a received signal $c = c_1c_2 \dots c_n$ and candidate messages $a = a_1a_2 \dots a_p$, $b = b_1b_2 \dots b_q$ and outputs the number of ways those messages could be interleaved to produce the signal.

- Identify a subproblem to solve at the i th bit of the signal c .
- Use your subproblem to define a recurrence. Also state your base cases.
(*hint*: try defining a two-dimensional recurrence, with each dimension corresponding to one of our candidate messages. You may not have to use the whole table...)
- How big is your lookup table? In what order should we fill out our lookup table? How can we use the lookup table to produce a final solution?
- Consider the signal $c = 110110$ and candidate messages $a = 101$, $b = 110$. In how many ways can these be interleaved to create c ?
Use your recurrence to fill out a lookup table.
- Use backtracking on your lookup table to find an interleaving of messages that produces the signal c .

- Identify a subproblem to solve at the i th bit of the signal c .

We need decide whether to attribute the i th bit of c to message a or message b .

$$5 \text{ mod } 5 = 0$$

b. Use your subproblem to define a recurrence. Also state your base cases.

(hint: try defining a two-dimensional recurrence, with each dimension corresponding to one of our candidate messages. You may not have to use the whole table...)

We define $OPT(i, j)$ as # of ways to construct c_1, c_2, \dots, c_{i+j} by interleaving the first i characters of a and the first j characters of b .

$$OPT(i, j) = \begin{cases} OPT(i-1, j) & , c_{i+j} = a_{(i \bmod p)+1} \neq b_{(j \bmod q)+1} \\ OPT(i, j-1) & , c_{i+j} \neq a_{(i \bmod p)+1} = b_{(j \bmod q)+1} \\ OPT(i-1, j) + OPT(i, j-1) & , c_{i+j} = a_{(i \bmod p)+1} = b_{(j \bmod q)+1} \\ 0 & , c_{i+j} \neq a_{(i \bmod p)+1} \neq b_{(j \bmod q)+1} \\ 1 & , c_{1, \dots, j} = b_{(1, \dots, (j \bmod q))}, \bar{i} = 0, \bar{j} \neq 0 \\ 1 & , c_{1, \dots, \bar{i}} = a_{(1, \dots, (\bar{i} \bmod p))}, \bar{i} \neq 0, \bar{j} = 0 \\ 1 & , \bar{i} = 0, \bar{j} = 0 \\ 0 & , \text{otherwise} \end{cases}$$

$$(i \bmod p) + 1 = (3 \bmod 2) + 1 = 2$$

$$(j \bmod q) + 1 = (6 \bmod 3) + 1 = 1$$

we can run an example to clarify the idea

$a, a_2: x y, p=2$

$b, b_1, b_3: z y k, q=3$

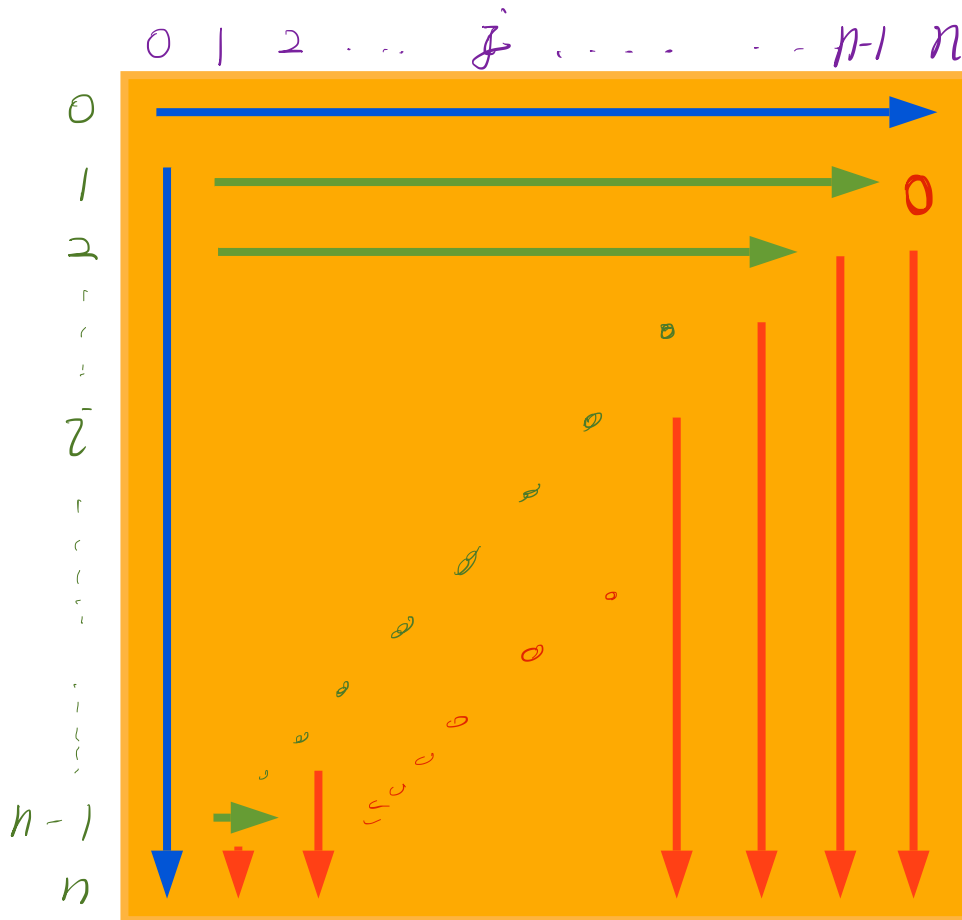
$c_1, c_2, c_3, c_4, c_5, c_6, c_7, c_8, c_9$

$OPT(3, 6) =$

$$\begin{cases} OPT(2, 6) & , c_{10} = a_2 \neq b_1 \\ OPT(3, 5) & , c_{10} = b_1 \neq a_2 \\ OPT(2, 6) + OPT(3, 5) & , c_{10} = a_2 = b_1 \\ 0 & , c_{10} \neq a_2 \neq b_1 \end{cases}$$

c. How big is your lookup table? In what order should we fill out our lookup table? How can we use the lookup table to produce a final solution?

The lookup table is $[1 + \max(i)] \times [1 + \max(j)]$,
 which is $(n+1) \times (n+1)$



- d. Consider the signal $c = 110110$ and candidate messages $a = \overset{\vee}{1}\overset{\vee}{0}1$, $b = 110$. In how many ways can these be interleaved to create c ?

Use your recurrence to fill out a lookup table.

			/	/	0	/	/	0
		0	1	2	3	4	5	6
0		1	1	1	1	1	1	1
/	1	1	2	0	1	2	0	0
0	2	0	2	2	0	2	0	0
/	3	0	2	4	4	0	0	6
/	4	0	2	0	0	0	0	0
0	5	0	0	0	0	0	0	0
/	6	0	0	0	0	0	0	0

They are cases when $i+j = n = 6$,
which means the number of ways to construct (c_1, c_2, \dots, c_6) .

- e. Use backtracking on your lookup table to find an interleaving of messages that produces the signal c .

	0	1	2	3	4	5	6
0	1	1	1	1	1	1	1
1	1	2	0	1	2	0	0
2	0	2	2	0	2	0	0
3	0	2	4	4	0	0	6
4	0	2	0	0	0	0	0
5	0	0	0	0	0	0	0
6	0	0	0	0	0	0	0

$b_1, b_2, b_3, b_4, b_5, b_6$

$b_1, b_2, b_3, b_4, a_1, a_2$

$b_1, b_2, b_3, a_1, b_1, a_2$

$b_1, a_1, a_2, b_2, a_3, b_3$

$b_1, a_1, a_2, a_3, b_2, b_3$

$a_1, b_1, a_2, b_2, a_3, b_3$

$a_1, b_1, a_2, a_3, b_2, a_3$

7 ways