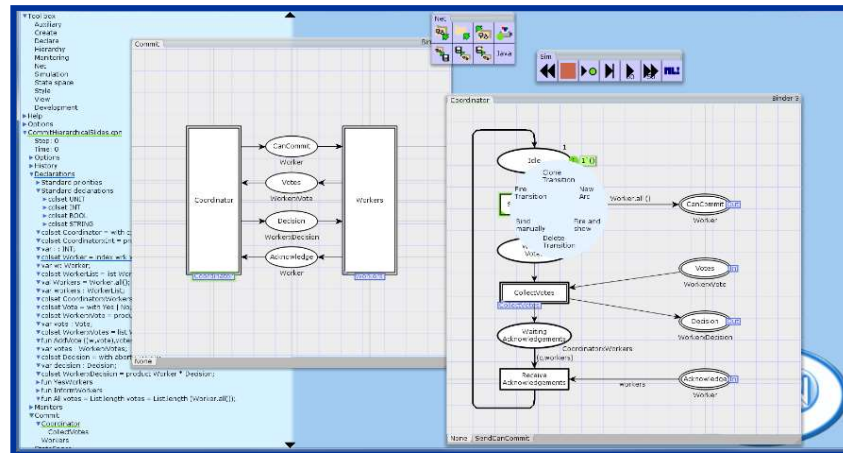


## Lecture 4

# Hierarchical Coloured Petri Nets with Modules



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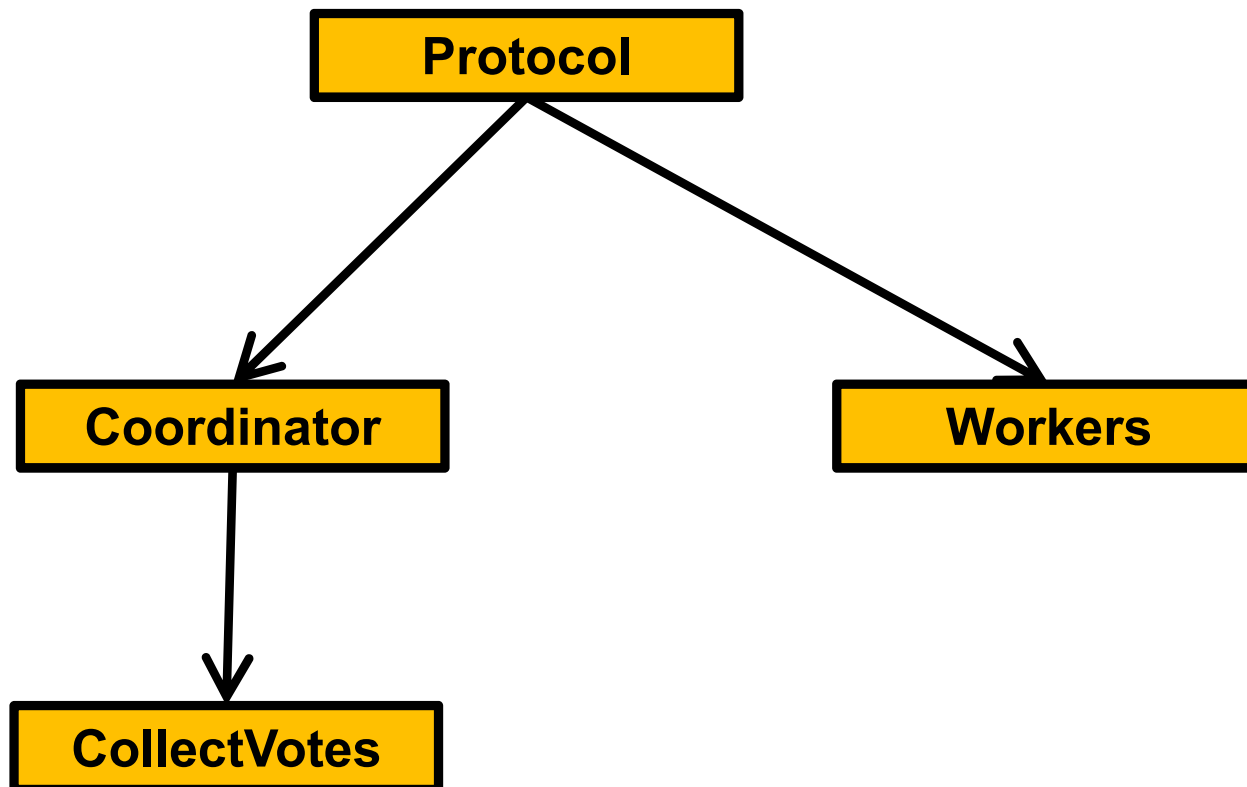
Email: [lmkr@hvl.no](mailto:lmkr@hvl.no) / WWW: [home.hib.no/ansatte/lmkr](http://home.hib.no/ansatte/lmkr)

# Introduction

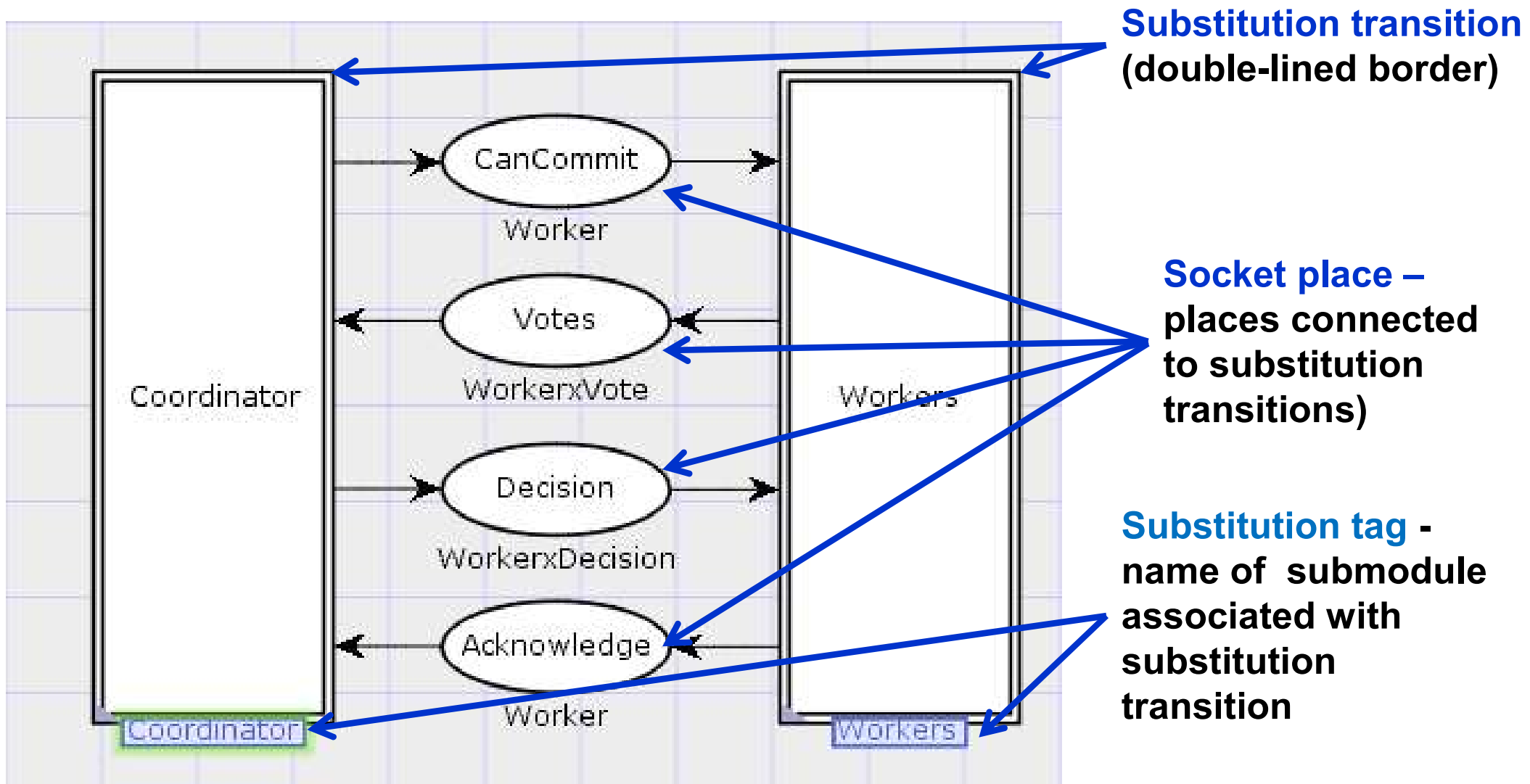
- **Important to be able to split a large CPN models into a set of modules with interfaces**
  - To support construction of large CPN models
  - To support reuse of modules and maintainability
  - To support abstraction and management of details
- **Key concepts**
  - A **module** exchange tokens with its environment using input/output **port places**
  - **Substitution transitions** have associated **submodules**
  - **Port-socket relation** associates socket places of substitution transitions with the port places in the submodule

# Hierarchical Modules

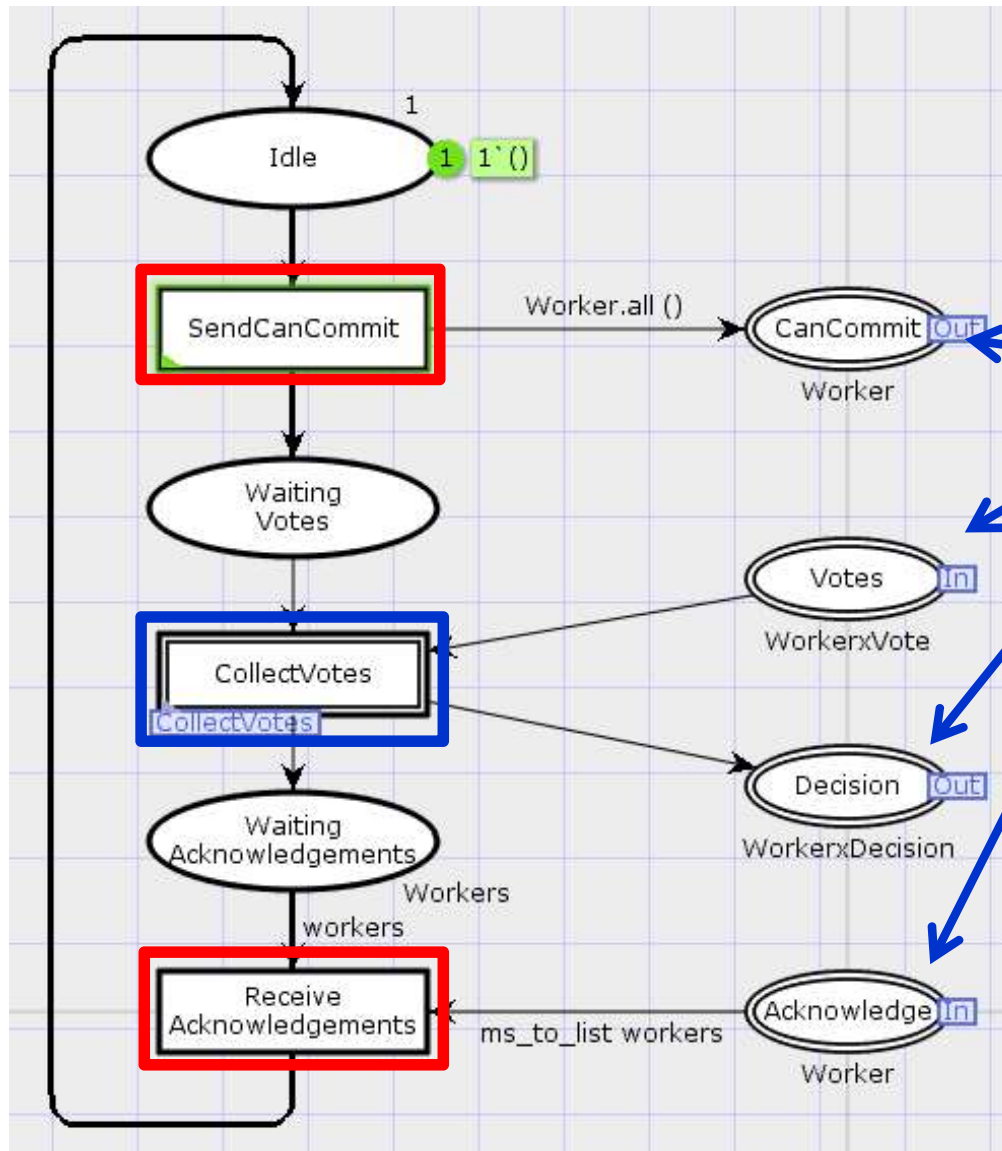
- Model is comprised of collection of **modules** that are hierarchically organised into levels
- **Example:** two-phase commit protocol



# Top-level: Protocol Module



# Coordinator Module



**Port place** - used for exchanging tokens with the upper-level module (IN,OUT,IN/OUT).

**SendCanCommit** and **ReceiveAcknowledgement** are ordinary transitions.

**CollectVotes** is a substitution transition

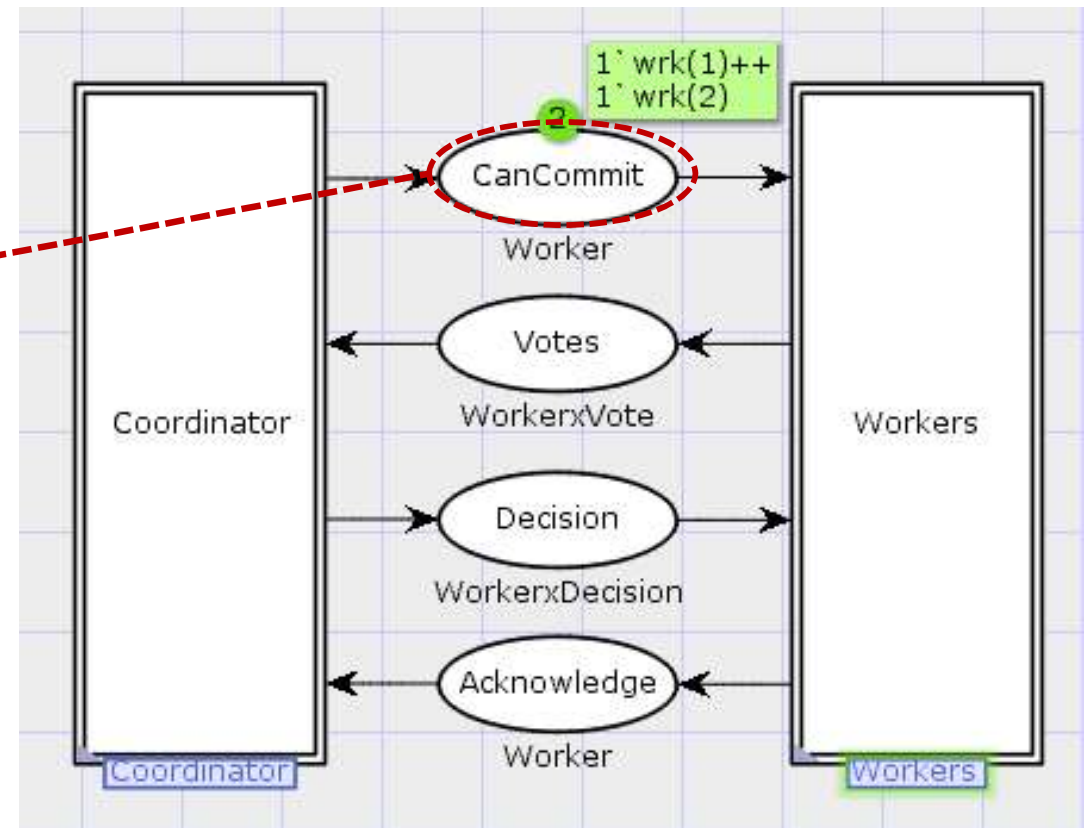
# Port and Socket Places

- Tokens added (removed) on a port place are added (removed) on the **associated socket place**

Coordinator module



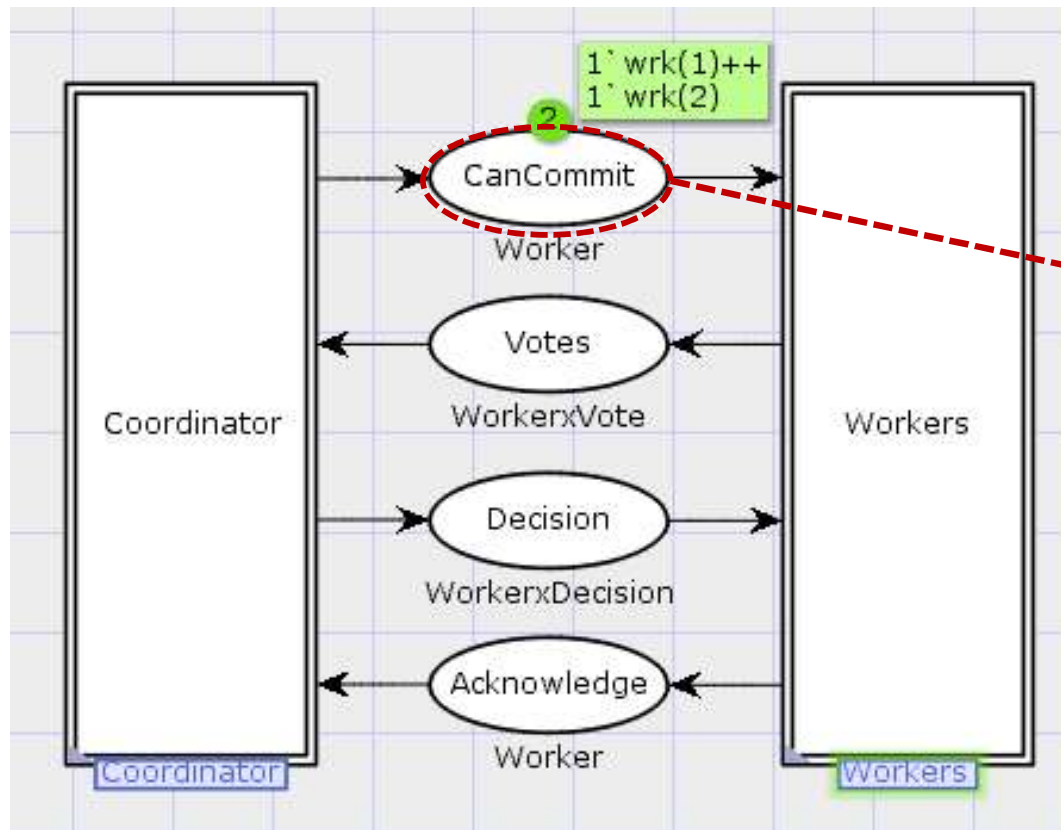
Protocol module



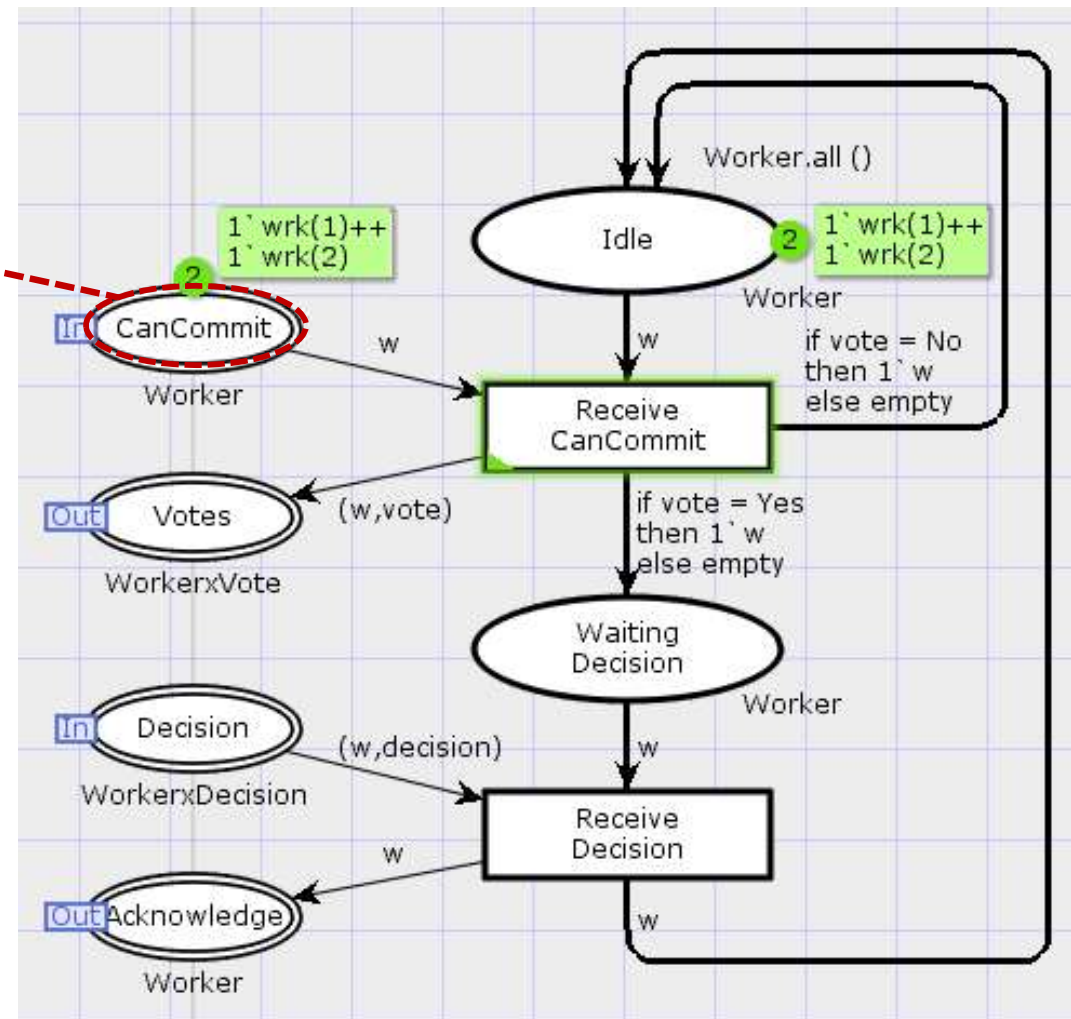
- Associated port and socket places constitute a **compound place**.

# Workers Module

## Protocol module



## Workers module





# CPN Tools Demo

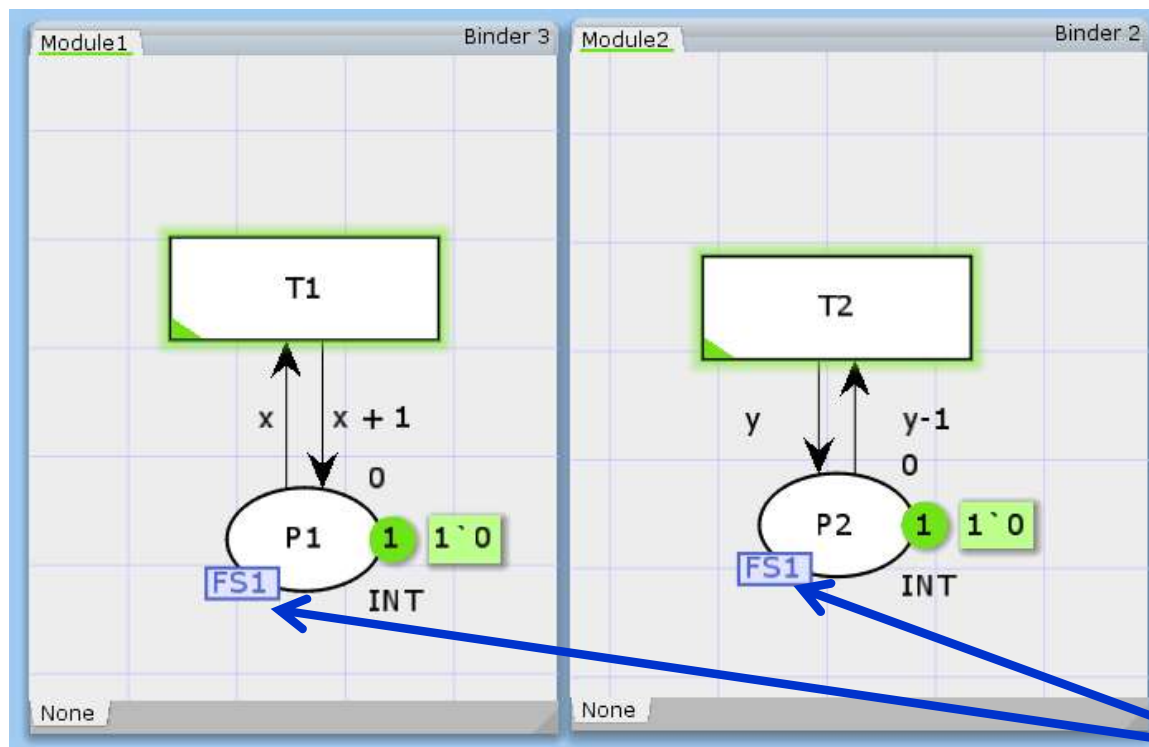
- **Hierarchical CPN models**
  - Navigating hierarchical models
  - Simulation of hierarchical models
  - Editing of modules: top-down and bottom-up development





# Place Fusion Sets

- Group of places to be treated as one conceptual (global) place



Any change in the marking of P1 will be reflected on P2 (and vice versa).

Similar to global variables  
- and should be used with care

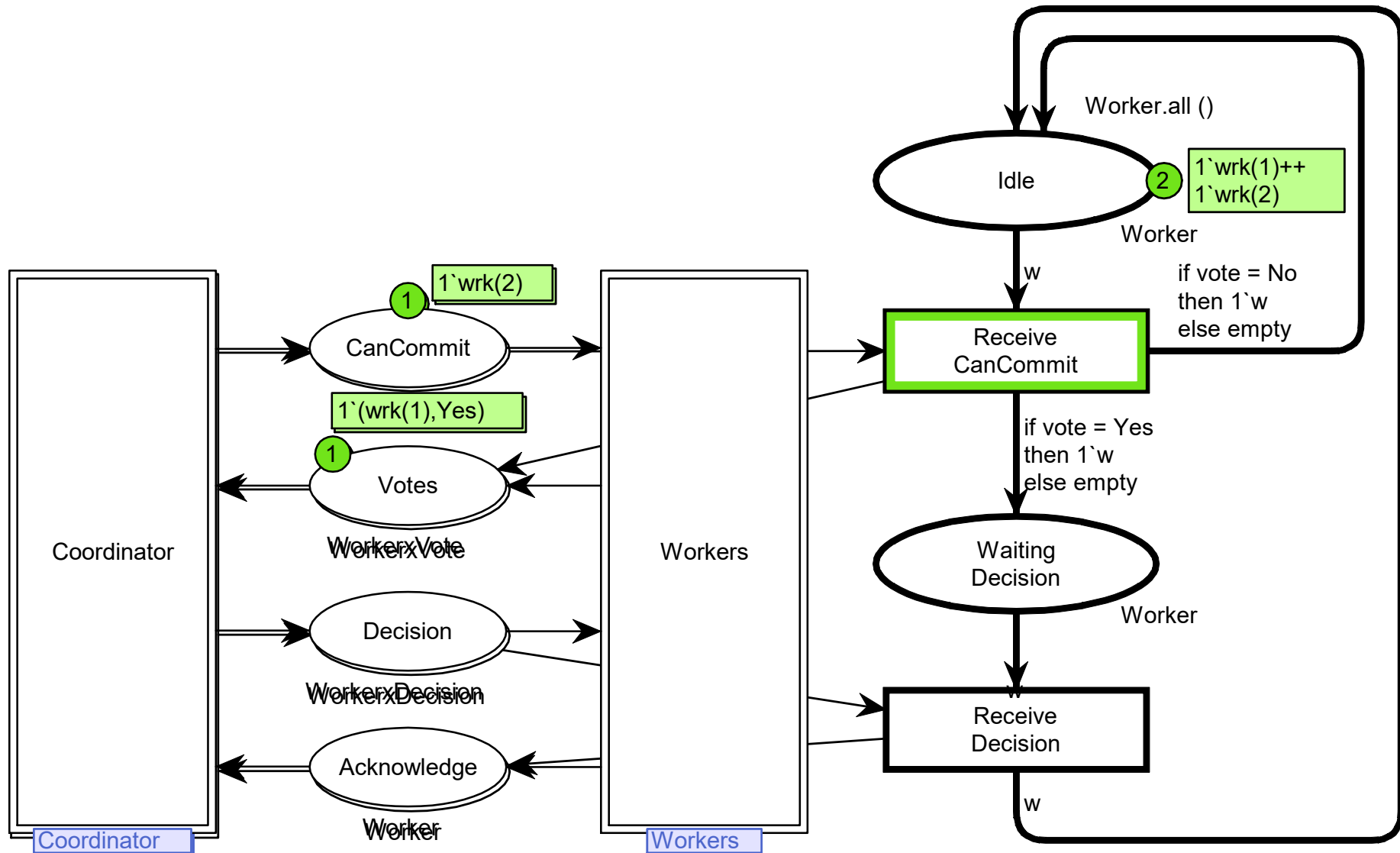
P1 and P2 are fusion places belonging to fusion set FS1.



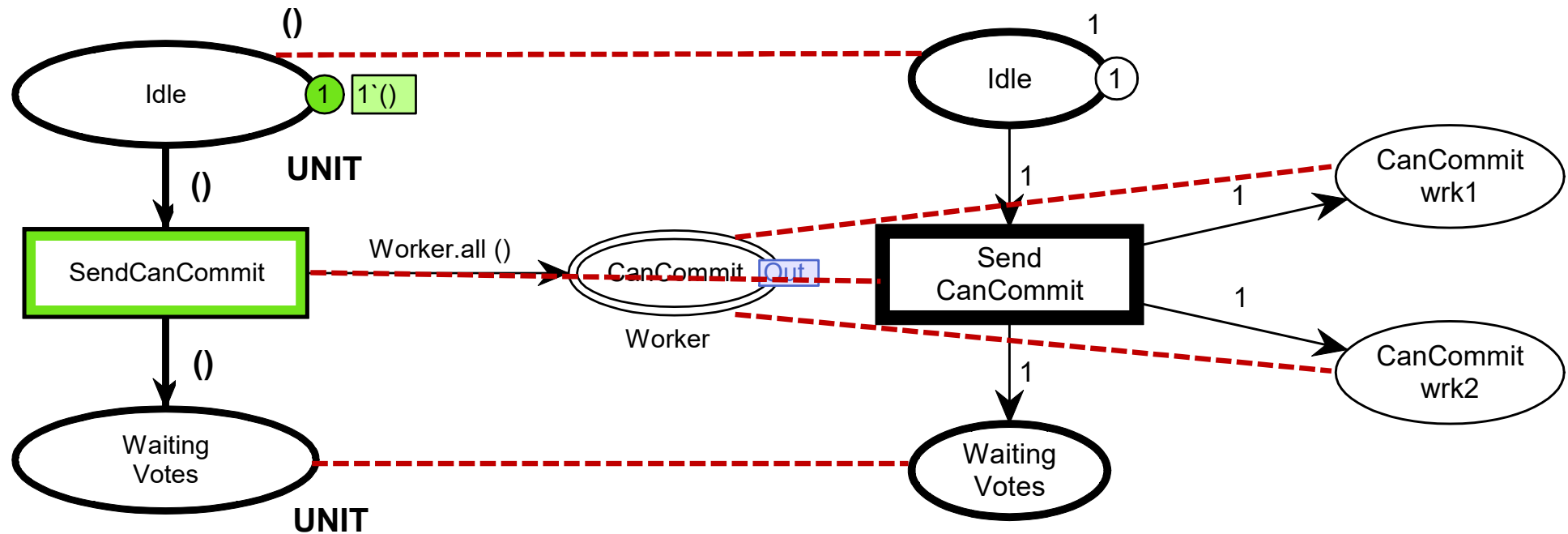
# Unfolding Coloured Petri Nets

- **A hierarchical CPN model can be unfolded to a non-hierarchical Coloured Petri Net**
  - Recursively replace each substitution transition with its associated submodule.
  - Associated port and socket places are merged into a single place.
- **A non-hierarchical Coloured Petri Net can be unfolded into a Place/Transition Net (PTN)**
  - Replace each CPN place with one PTN place for each colour in the colour set of the CPN place.
  - Replace each CPN transition with one PTN transition for each possible binding of the CPN transition.

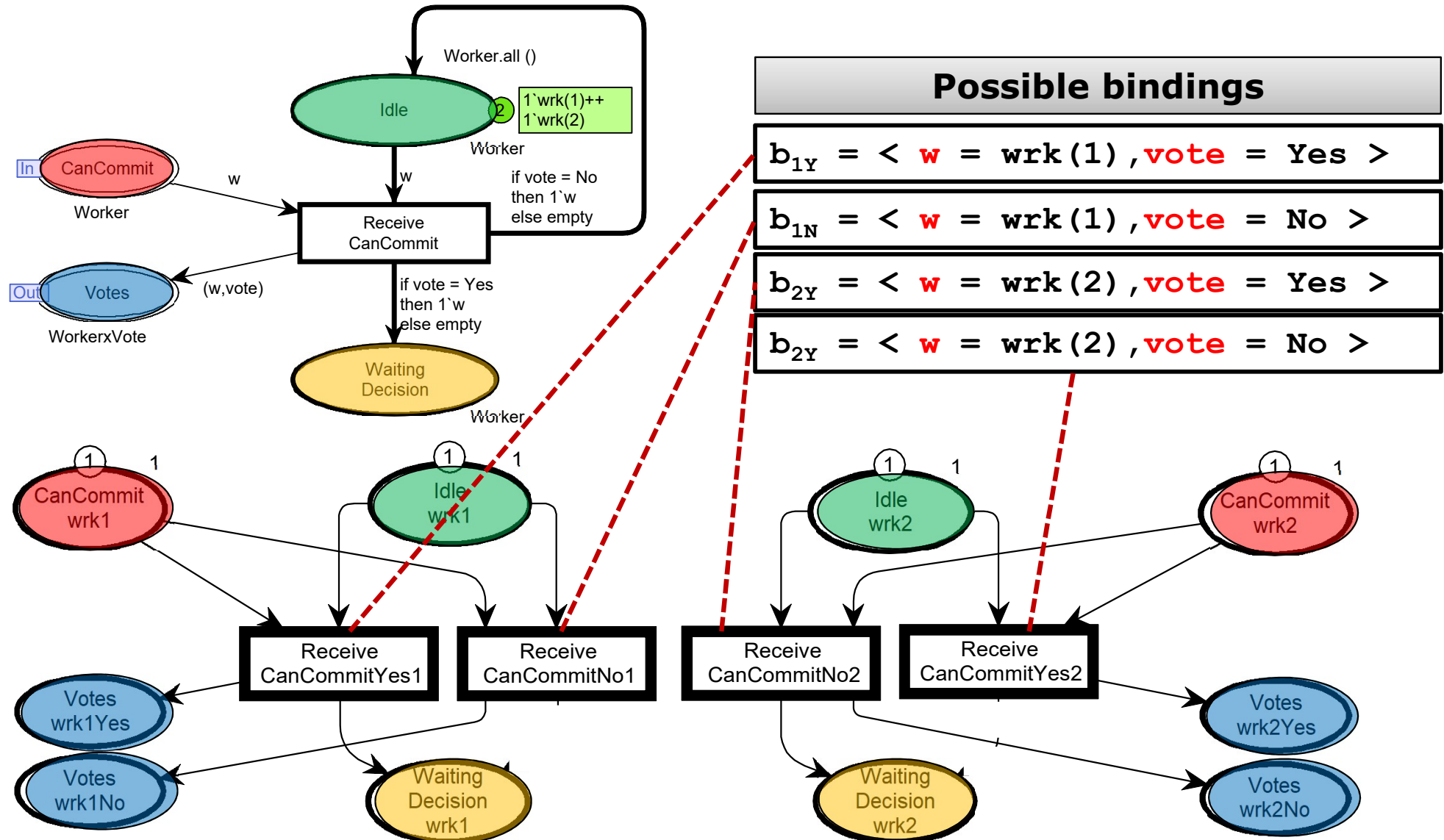
# Unfolding hierarchical CPNs



# Unfolding CPN Places



# Unfolding CPN Transitions

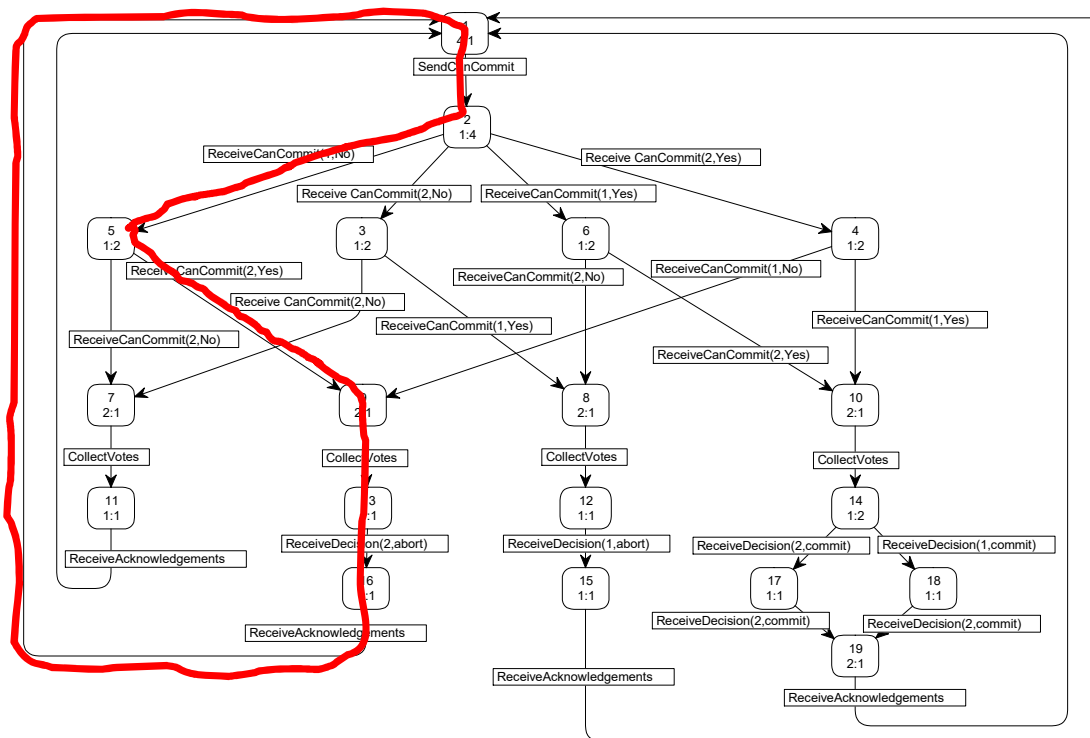






# Verification and Model Checking

- **Formal verification** of CPN models can be conducted using **explicit state space exploration**

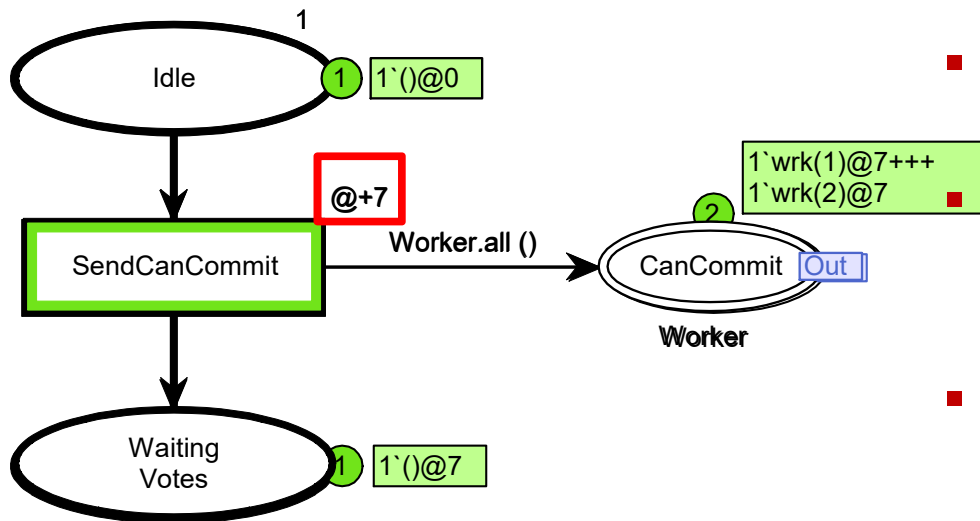


- A state space represents all possible **executions** of the CPN model
- **Standard behavioural properties** can be investigated using the state space report
- **Model-specific properties** can be verified using queries and temporal logic model checking

- Several **advanced techniques** available to alleviate the inherent state explosion problem.

# Performance Analysis

- CPNs include a **concept of time** that can be used to model the timed taken by activities



- A **global clock** representing the **current model time**
  - Tokens carry **time stamps** describing the earliest possible model time at which they can be removed
  - Time inscriptions** on transitions and arcs are used to give time stamps to the tokens produced on output places
- Random distribution functions** can be used in arc expressions (delays, packet loss, ...).
  - Data collection monitors** and batch simulations can be used to compute performance figures.

# Perspectives on CPNs

- **Modelling language combining Petri Nets with a programming language.**
- **The development has been driven by an application-oriented research agenda**
- **Key characteristics**
  - Few but still powerful and expressive modelling constructs
  - **Implicit concurrency** inherited from Petri nets: everything is concurrent unless explicit synchronised
  - **Verification** and **performance analysis** supported by the same modelling language

