

Introduction

- AI: It is the study of how to make computers do things which at the moment people do better.
- The AI problems are classified into 3 tasks:
- Mundane tasks: Daily routine tasks.
 - Perception: Vision, speech
 - Natural language: Understanding, Generation, Translation
 - Common sense Reasoning
 - Formal tasks:
 - Games: chess (computer copies and predicts our moves after a few turns)
 - Mathematics: Geometry, logic, Integral Calculus, proving properties of program
 - Expert tasks:
 - Engineering: Design, fault finding, manufacturing & planning, Medical diagnosis, financial analysis

Intelligence: Thinking, decision making, problem solving, learning from experience, understanding, adaptation, reasoning, reaction to situations.

There is no parameter to measure intelligence and hence it cannot be defined.

- Fields of AI:
- Speech recognition
 - Image processing
 - Knowledge discovery
 - Natural language understanding
 - Quantitative reasoning
 - Knowledge representation
 - Machine learning
 - * Search models
 - Memory
 - Adverse Serial Learning
 - Handling uncertainty
- Computer vision
- Computerisation
- Neural Networks
- Data mining
- Pattern recognition
- NW generator
- Fuzzy sets
- Logic
- Ontology
- Semantics
- Robot control
- Planning

Technique: AI problems span a very broad spectrum. They appear to have very little in common except that they are hard (Time complexity is exponential) in terms of knowledge:

It is voluminous - Hard to categorize accurately
Constantly changing - It differs from data by being organized in a way that corresponds to the way it will be used.

Technique in terms of knowledge: We are forced to include that an AI technique is a method that exploits knowledge that should be represented in such a way that-

The knowledge captures generalization. In other words it is not necessary to represent separately each individual situation. Instead situations that share important properties are grouped together.

- 1. It can be understood by people who must provide it.
- 2. It can easily be modified to correct errors and to reflect changes in the world and in our world view.
- 3. It can be used in great many situations even if it is not totally accurate or complete.
- 4. It can be used to help us share bulk by helping to narrow the range of possibilities that must usually be considered.

Tic Tac Toe: 3 scenarios will be compared on basis of complexity, generalization, clarity of knowledge and extensibility of approach.

2 Ds: (Program 1)

Board: A 9 element vector representing the board, where the elements of the vector corresponds to the board position as follows:

0 → blank 1 - x 2 - 0

Move table: It is a large vector of $19683 (3^9)$ elements,

Each square \rightarrow 3 possibilities

of squares

Total possibilities $\rightarrow 3^9$

each element of which is a 9 element vector:

Algorithm: To make a move, following steps:

1. View the board as ternary (base 3) no. Convert it to a decimal no.
2. Use the no. computed in step 1 as an index into move table and access the vector stored there.
3. The vector selected in step 2 represents the way the board will look after the move that should be made.

Less time consuming, specialized approach, not much knowledge used, not extensible to 3D as no. of possibilities increase to 3^{27} (space complexity \uparrow)

Program 2: PS wed:

Board: 9 element vector but instead of using 0, 1, 2 we use 2 \rightarrow blank, 3 \rightarrow X, 5 \rightarrow O

Turn: It is an integer indicating which move of the game is about to be played. 1 \rightarrow 1st move 9 \rightarrow last move

Algorithm: 3 sub procedures:

1. Move 2: Returns 5 if the centre square of the board is blank. Otherwise it will return any blank non corner square.

2. Posswin (p): It returns 0 if player p cannot win on his next move; otherwise it returns a number of squares that constitutes a winning move.

This function will enable the program both to win and to block the opponent's win.

Posswin operates by checking 1 at a time each of the rows, columns and diagonals. Because $3 \times 3 \times 2 = 5 \times 5 \times 2$ possibility of wins.

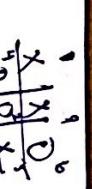
3. Go(n): It makes a move in-square n. It sets Board(n) = 3 if turn is odd

" " 5 " " " even

It also increments turn "by 1".

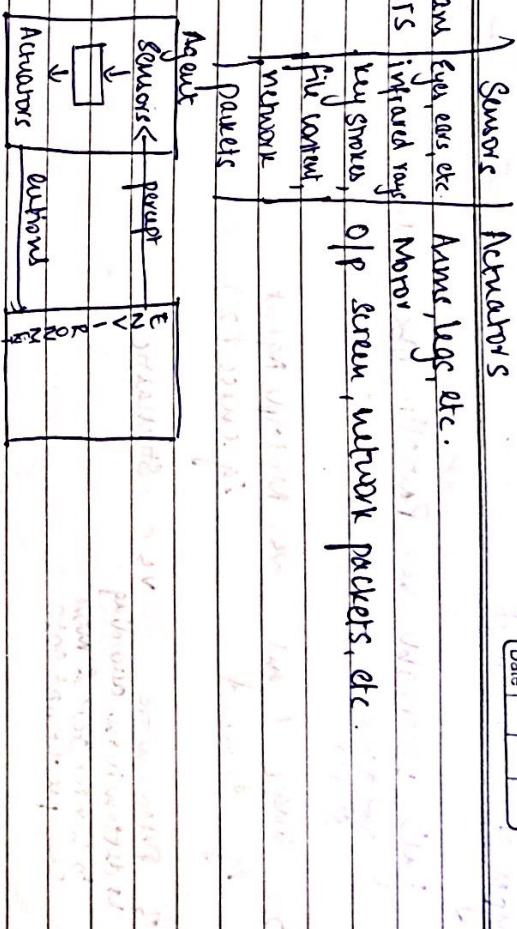
Page No.	_____
Date	_____

In your turn at
random any non-
empty blank square



Page No.	_____
Date	_____

- Turn(1) Use(1). (Assuming X is O)
 Turn(2) Use(5) if blank otherwise Use(11)
 Turn(3) Board(9) if blank Use(9) else Use(3)
 Turn(4) If Posswin(x) != 0 then Use(Posswin(x))
 Use Use(0) max(2)
 Turn(5) If Posswin(x) is not 0, Use(Posswin(x))
 otherwise If Posswin(0) insert 0, Use(Posswin(0))
 else if
 Board(7) is blank Use(7) else Use(13)
 Turn(6) If Posswin(0) is not 0 Use(Posswin(0))
 otherwise if Posswin(x) is not 0
 Use Use(Posswin(x))
 else
 Use(0)



Page No.	_____
Date	_____

- Turn(7) If Posswin(x) is not 0, then Use(Posswin(x))
 Use if Posswin(0) is not 0 Use(Posswin(0))
 we find any position which is blank
 else if Posswin(x) is not 0 Use " (x)"
 else if Posswin(0) is not 0 Use " (0)"
 else if Posswin(x) is not 0 Use " (x)"
 else if Posswin(0) is not 0 Use " (0)"

Rational Agent: It is the one that does the right thing.
 things are important:
 • True performance measures that determine the intention of success
 • The agent's prior knowledge of the environment
 • The actions that agent can perform
 • The agent's percept sequence to date

- Turn(8) Game is 7:
 Turn decides who will win.
 Now more complex, less space, specialised, creating of
 knowledge easier to extend to 3D compared to
 program 2, program 3, many conclusions → from the book

- AGENTS: An agent is anything that can be viewed as perceiving
 its environment through sensors and acting upon that environment
 through actuators.

• Omnidirectional Agent: It knows the actual outcome of
 its actions and can act accordingly.

Type of Environment:

1. Fully Observable vs. Partially Observable
↳ everything can be sensed
2. Single Agent vs. Multiple Agent
↳ cross world
eg. chess (2)
3. Deterministic vs. Stochastic
↳ deterministically according to current state & move can predict next state
4. Episodic vs. Sequential
↳ information about current state depends upon previous parts (episodes), individual parts
5. Static vs. Dynamic
↳ finite no. of
absence states
↳ infinite no. of
eg. taxi driving
6. Discrete vs. Continuous
↳ discrete states
↳ continuous states

agent should maintain some sort of internal state that depends on the percept history and thereby reflects about some of the unobserved aspects of the current state.

1. Goal based agent: As a current state description, the agents need some sort of goal information that describes situations that are desirable.
2. Learning Agent: All the agents can improve their performance by learning. \Rightarrow constant feedback

Utility based Agent: Goals alone are not enough to generate high quality behaviour in most environments.
↳ tools just provide a crude binary distinction between happy & unhappy states. A more general performance measure allows a comparison of different world states according to exactly how happy they would make the agent.

Learning Agent: All the agents can improve their performance by learning. \Rightarrow constant feedback

Chapter -2
Problem Solving Criteria

To build a system to solve a particular problem, we need to do 4 things:

1. Define the problem specifically / precisely : Definition must include precise specification of what the initial

situations will be as well as what final situations constitute acceptable solutions to the problem.

2. Analyse the problem: A very few important features can have an immense effect on the appropriateness

of various possible techniques for solving the problem.

3. Isolate & represent the relevant knowledge that is necessary to solve the problem.

4. Choose the best problem solving techniques and apply it to the particular problems.

1. Simple reflex agents: These agents select actions on the basis of current percept ignoring rest of percept history.
2. Model based reflex agents: The most effective way to handle partial observability is for the agent to keep track of the part of the world it can't see now, i.e., the

enumerating all of the states it contains.

1. (x,y) If $x < y \rightarrow (y,x)$
2. $(x,y) \rightarrow$ If $y < 3 \rightarrow (x,3)$
3. $(x,y) \rightarrow (x,y)$

4. (x,y) If $y > 0 \rightarrow (x,y-1)$

5. $(x,y) \rightarrow (0,y)$

6. $(x,y) \rightarrow (x,0)$

7. $(x,y) \rightarrow (4-y-1, y-x)$

8. (x,y)

If $(x+y) \geq 3$ and $x > 0 \rightarrow (x-(3-y), 3)$

9. $(x,y) \rightarrow (x+y, 0)$

If $x+y \leq 4$ and $y > 0$

10. $(x,y) \rightarrow (0, x+y)$

If $(x+y) \leq 3$ and $x > 0$

11. $(0,2) \rightarrow (2,0)$

12. $(2,y) \rightarrow (0,y)$

(0,0)

Rule 2. $(0,0) \rightarrow (0,3)$

9. $(0,3) \rightarrow (3,0)$

2. $(3,0) \rightarrow (3,3)$

7. $(3,3) \rightarrow (4,2)$

5. $(4,2) \rightarrow (0,2)$

9. $(0,2) \rightarrow (2,0)$

Learning Outcomes : (Mentioned for implementation of state space search)

1. Define a state space that contains all the possible configurations of the relevant object. It is of course possible to define this state space without explicitly

1. Specify 1 or more states within that space that describes possible situations from which the problem solving process may start. These states are called initial states.
2. Specify 1 or more states that would be acceptable as solutions to the problem. These states are called goal states.
3. Specify a set of rules that describe the actions available.
4. Production System: It consists of set of rules each consisting of left side that determines the applicability of the rule and a right side that describes the operation to be performed if the rule is applied.
5. One or more knowledge DB that contains whatever info is of priority for the particular task. Some parts of the DB may be permanent while other parts may pertain only to the solution of the current problem.
6. The info in these DBs may be structured in any appropriate way.

3. A control strategy that specifies the order in which the rules will be compared to the DB and a way of resolving the conflicts that arise when several rules match at once.

Right. Of a good control strategy:

1. It first comes motion (should move forward)
2. Initial
3. Goal

Move operation

Generate & Test()

1. While more candidates exist

do generate a candidate

Test whether it is the goal state

- return Failure

River Crossing Puzzle:

A man needs to transport a lion, a goat & a cabbage across a river. He has a boat in which he can take only 1 of them at a time. It is only his presence that prevents the lion from eating goat & goat from eating cabbage. He can neither leave the goat alone with lion nor cabbage with goat. Now how will he take them all across the river.

Start state : ((M,G)) () (Two lists)

Goal state : (()) (M,G)

movement : → specific problem

Initialise set of successors C to an empty set.

Add m to the complement of the given state to get the new states.

Given state has left (L) turn add right to S.R else add left

If legal(s) then add s to the set of successors C.

For each other validity check make a copy of s.

Add E to S'.

if legal(s') then add s' to C.

else return to C

SimpleSearch()

1. Open $\leftarrow \{ \text{start} \}$

2. while open is not empty

do pick a node from open

if Open $\leftarrow \text{open} \setminus \{ \text{node} \}$ & Goal(node) = TRUE then

return node

else

Open $\leftarrow \text{open} \cup \text{Movement}(\text{node})$

return Failure

SimpleSearch2()

1. Open $\leftarrow \{\text{start}\}$

2. Closed $\leftarrow \{\}$ (keeps track of visited nodes)

3. while Open is not empty

4. do pick some node n from open

Open $\leftarrow \text{Open} \setminus \{ \text{n} \}$

Closed $\leftarrow \text{closed} \cup \{ \text{n} \}$

If GoalTest(n) = TRUE

return n

else

Open $\leftarrow \text{Open} \cup \text{Movement}(\text{n})$

Open $\leftarrow \text{Open} \setminus \{ \text{n} \}$

Closed $\leftarrow \text{closed} \cup \{ \text{n} \}$

Open $\leftarrow \text{Open} \cup \text{Movement}(\text{n})$

Open $\leftarrow \text{Open} \setminus \{ \text{n} \}$

Closed $\leftarrow \text{closed} \cup \{ \text{n} \}$

Open $\leftarrow \text{Open} \cup \text{Movement}(\text{n})$

Open $\leftarrow \text{Open} \setminus \{ \text{n} \}$

Closed $\leftarrow \text{closed} \cup \{ \text{n} \}$

Open $\leftarrow \text{Open} \cup \text{Movement}(\text{n})$

Open $\leftarrow \text{Open} \setminus \{ \text{n} \}$

Closed $\leftarrow \text{closed} \cup \{ \text{n} \}$

keeping track of path we refine?

Simple Search 3()

open $\leftarrow \{ \text{start} \}$

closed $\leftarrow \{ \}$

while open is not empty

do pick some node n from open
 $n \leftarrow \text{Head}(n)$

if GoalTest(n) == True

Return Reverse(n),

else closed $\leftarrow \text{closed} \cup \{ n \}$

successor $\leftarrow \text{MoveGen}(n) \setminus \text{closed}$

open $\leftarrow \text{open} \cup \text{successor}$

for each s in successor

do AddCons(s, n), to open

return Failure

open closed

(S,) (S, S,) (S)

(AS, BS, CS), (S)

(BS, LS) (A, S)

(AS, BS, CS) (A, S)

(A, S)

Modifying Simple Search 3

DFS

1. open $\leftarrow (\text{start NIL})$

2. closed $\leftarrow ()$

3. while not null(open)

do nodepair $\leftarrow \text{Head}(\text{open})$

node $\leftarrow \text{Head}(\text{nodepair})$

if GoalTest(node) == True

then return

Reconstruct(nodepair, closed)

else

closed $\leftarrow \text{cons}(\text{nodepair}, \text{closed})$

children $\leftarrow \text{MoveGen}(\text{node})$

noloops $\leftarrow \text{Reconstruct}(\text{nodepair}, \text{closed})$

new $\leftarrow \text{makePair}(\text{noloops}, \text{node})$

EP AS BS CS

C Simple Search 3()

1. open $\leftarrow \{ \text{start} \}$

2. closed $\leftarrow \{ \}$

3. while open is not empty

do pick some node n from open

$h \leftarrow \text{Head}(n)$

if $\text{GoalTest}(h) == \text{True}$

Return Reverse(n)

else

closed $\leftarrow \text{closed} \cup \{ h \}$

successor $\leftarrow \text{MoreGen}(h)$ / closed

successor $\leftarrow \text{successor} / \text{mapcar}(\text{open})$; open $\leftarrow \text{open} / n$

for each s in successor

do add con(n, s) to open

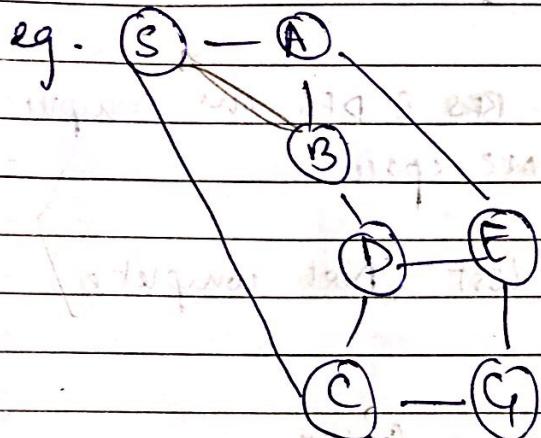
return failure

l/mapcar(open)
gives head
of nodes
in open.

Open	Closed	Successor
$\{ A, S \}$	$\{ \}$	A, B, C
$\{ A, S, B, S, L, S \}$	$\{ A, S \}$	A, B, C

new node will be added in front of tail

Open \leftarrow Append (new, Tail(open)) new node inserted in the front of list
return "No solution"



open	closed
(SNil)	
(AS, BS, CS)	(Tail)
(EA, BS, CS)	(AS, SNil)
(UE, DE, BS, CS)	(EA, AS, SNil)

- For BFS, modify as
- open \leftarrow Append (Tail(open), new) \rightarrow new node inserted at the end.
- Branching factor = 3 \Rightarrow every node will generate 3 nodes

4 factors on which BFS & DFS are compared:

- a) completeness
 - b) Time complexity
 - c) Space "
 - d) Quality of solution
- a) Does the algorithm always find a solution where there exists one?
 - b) How much time does the algorithm run before finding a solution? We will measure the time complexity by the no. of nodes the algo picks up before it finds the solution.
 - c) How much space does the algorithm need? we will use the no. of nodes in the open list as a measure of space complexity.

d) A simple criteria for quality might be the length of the path found from the start node to the goal node.

2 scenarios : State space \leftarrow finite.

When the state space is finite both BFS & DFS are complete (there is a goal state in state space).

State space infinite : BFS complete

DFS will get lost (Not complete)

Time complexity :

BFS :
BFS fails
 $d \rightarrow \text{depth}$

DFS :
BFS fails
 $b \rightarrow \text{branching factor}$
 $d \rightarrow \text{depth}$

Branching factor for DES drives down some branch factor

Same

Size of open is linear with depth

DFS : No. of nodes are exponential to depth

Space complexity : In general search tree with a branching factor b DES drives down some branch factor

at each level leaving all nodes in open. When it enters the depth d it has DPS nodes and

$DPS = (b-1)(d-1) + b^d$

Outers the depth d it has DPS nodes and

$DPS = (b-1)(d-1) + b^d$

Size of open is linear with depth

BFS : No. of nodes are exponential to depth

Space complexity of BFS $\geq DPS$

Find the minimum solution among various solutions

Quality of solution : BFS gives better solution even in worst case. Measured in terms of no. of edges bw goal & start states.

BFS & DFS : Blind algorithms, they have no intelligence,

they explore the nodes in pre-defined order.

Average, $N_{DPS} = \frac{(d+1 + b^{d+1})}{b-1} \Big|_2$.

$$= \frac{(d+1)(b-1) + b^{d+1}}{2(b-1)}$$

$$= \frac{db + b - d - 1 + b^{d+1}}{2b} \approx \frac{b^{d+1}}{2b} \quad (\text{dividing by } b \text{ is a good way})$$

$$\approx \frac{b^{d+1}}{2(b-1)} \quad (\text{canceling } b)$$

$$N_{DPS} = \frac{\cancel{b^{d+1}-1} + b^{d+1}-1}{b-1} \Big|_2$$

(+1 ignored)

$$= \frac{b^{d+2} - b^{d+1}}{2(b-1)} \approx \frac{b^d(b+1)}{2(b-1)} \approx \frac{b^d(b+1)}{2b}$$

$$N_{DPS} = \frac{b^d(b+1)}{b^{d+1}} = \frac{b+1}{b}$$

\Rightarrow BFS has slightly more complexity than DFS, otherwise same.

The heuristic function must not be computationally expensive because the idea is to cut down on the computational cost of the search algorithm.

$h(n) \rightarrow$ node

Euclidean distance

$$h(n) = \sqrt{(x_n - x_0)^2 + (y_n - y_0)^2}$$

$$h(n) = \sqrt{(x_n - x_0)^2 + (y_n - y_0)^2}$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Manhattan distance

$$h(n) = |x_n - x_0| + |y_n - y_0|$$

Open \leftarrow sort(h) append(h , tail(h)))

Completeness: Finite: Complete
Infinite: Depends on how good the heuristic function is.

Succinct of solution: Not good.
Sporadic: Heuristic function (Good Perfect) \rightarrow linear
converging " " Non good \rightarrow exponential

Time Complexity: Depends on heuristic function.

Parameters to decide among heuristic functions (which better).

1. Effective branching factor = total nodes expanded \rightarrow more memory required

Nodes in the solution

2. Search space \downarrow perfect. (No. of nodes b/w 0 & 2^n in the path we have)

3. Performance (Inverse of branching factor).

Value tends to $\downarrow \Rightarrow$ perfect.

Or it \downarrow , heuristic function quality \uparrow

Ascent

4. Hill Climbing / Steepest Gradient Descent Algo

Problem of local maxima: Δ

Hill Climbing()

1. node \leftarrow start

2. newnode \leftarrow Head(sort(h (node), node))

do

3. while h (newnode) $<$ h (node) \downarrow wrong

node \leftarrow newnode

newnode \leftarrow Head(sort(h (node), node))

but heuristic function in one with minimum value

2. previous node:

3. break loop:

Best First Search:

sort \rightarrow Used to get nodes in ascending order of heuristic value.

Same as DFS algo but in the end

Propagate : neighbours will be assigned $\text{parent}(m)$ for each $s \in \text{neighbours}$

$\text{newvalue} \leftarrow g(m) + h(m)$

If $\text{newvalue} < g(s)$
then $\text{parent}(s) \leftarrow m$

$g(s) \leftarrow \text{newvalue}$

If s is closed

propagate improvements(s)

1. $\text{open} \leftarrow \text{list}(\text{start})$
2. $(\text{start}) \leftarrow \text{ulstart}$
3. $\text{parent}(\text{start}) \leftarrow \text{NIL}$
4. $\text{closed} \leftarrow \emptyset$
5. while open is not empty
6. do
7. Remove node n from open such that $f(n)$ has the lowest cost
8. Add n to closed
9. If $\text{goaltest}(n) = \text{true}$
10. return ($\text{parent}, \text{path}(n)$)
11. $\text{neighbours} \leftarrow \text{MoveGen}(n)$
12. for each $m \in \text{neighbours}$
13. do $\text{smotar}()$

case $m \notin \text{open}$ & $m \notin \text{closed}$

Add m to open

$\text{parent}(m) \leftarrow n$

$g(m) = g(n) + h(n, m)$

$f(m) = g(m) + h(m)$

case $m \in \text{open}$

If $(g(m) + h(n, m)) < g(m)$

then $\text{parent}(m) \leftarrow n$

$g(m) = g(n) + h(n, m)$

$f(m) = g(m) + h(m)$

case $m \in \text{closed}$

If $(g(m) + h(n, m)) < g(m)$

then $\text{parent}(m) \leftarrow n$

$g(m) = g(n) + h(n, m)$

$f(m) = g(m) + h(m)$

propagate improvements(m)

if failure