
UNIT 1 SOLUTION OF LINEAR ALGEBRAIC EQUATIONS

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1.0 INTRODUCTION

In Block 1, we have discussed various numerical methods for finding the approximate roots of an equation $f(x) = 0$. Another important problem of applied mathematics is to find the (approximate) solution of systems of linear equations. Such systems of linear equations arise in a large number of areas, both directly in the modelling physical situations and indirectly in the numerical solution of other mathematical models. Linear algebraic systems also appear in the optimization theory, least square fitting of data, numerical solution of boundary value problems of ODE's and PDE's etc.

In this unit we will consider two techniques for solving systems of linear algebraic equations – Direct method and Iterative method.

These methods are specially suited for computers. Direct methods are those that, in the absence of round-off or other errors, yield the exact solution in a finite number of elementary arithmetic operations. In practice, because a computer works with a finite word length, direct methods do not yield exact solutions.

Indeed, errors arising from round-off, instability, and loss of significance may lead to extremely poor or even useless results. The fundamental method used for direct solution is Gauss elimination.

Iterative methods are those which start with an initial approximations and which, by applying a suitably chosen algorithm, lead to successively better approximations. By this method, even if the process converges, we can only hope to obtain an approximate solution. The important advantages of iterative methods are the simplicity and uniformity of the operations to be performed and well suited for computers and their relative insensitivity to the growth of round-off errors.

So far, you know about the well-known Cramer's rule for solving such a system of equations. The Cramer's rule, although the simplest and the most direct method, remains a theoretical rule since it is a thoroughly inefficient numerical method where even for a system of ten equations, the total number of arithmetical operations required in the process is astronomically high and will take a huge chunk of computer time.

1.1 OBJECTIVES

After going through this unit, you should be able to:

- obtain the solution of system of linear algebraic equations by direct methods such as Cramer's rule, and Gauss elimination method;
- use the pivoting technique while transforming the coefficient matrix to upper triangular matrix;
- obtain the solution of system of linear equations, $A\mathbf{x} = \mathbf{b}$ when the matrix A is large or sparse, by using one of the iterative methods – Jacobi or the Gauss-Seidel method;
- predict whether the iterative methods converge or not; and
- state the difference between the direct and iterative methods.

1.2 PRELIMINARIES

Let us consider a system of n linear algebraic equations in n unknowns

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n &= b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n &= b_2 \\ a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n &= b_n \end{aligned} \quad (1.2.1)$$

Where the coefficients a_{ij} and the constants b_i are real and known. This system of equations in matrix form may be written as

$$\begin{aligned} A\mathbf{x} &= \mathbf{b} \quad \text{where } A = (a_{ij})_{n \times n} \\ \mathbf{x} &= (x_1, x_2, \dots, x_n)^T \text{ and } \mathbf{b} = (b_1, b_2, \dots, b_n)^T. \end{aligned} \quad (1.2.2)$$

A is called the coefficient matrix.

We are interested in finding the values x_i , $i = 1, 2, \dots, n$ if they exist, satisfying Equation (3.3.2).

We now give the following

Definition 1: A matrix in which all the off-diagonal elements are zero, i.e. $a_{ij} = 0$ for $i \neq j$ is called a diagonal matrix; e.g., $A = \begin{bmatrix} a_{11} & 0 & 0 \\ 0 & a_{22} & 0 \\ 0 & 0 & a_{33} \end{bmatrix}$ is a 3×3 diagonal matrix.

A square matrix is said to be upper – triangular if $a_{ij} = 0$ for $i > j$, e.g.,

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ 0 & a_{22} & a_{23} \\ 0 & 0 & a_{33} \end{bmatrix}$$

Definition 2: A system of linear equations (3.3.2) is said to be consistent thus exists a solution. The system is said to be inconsistent if no solution exists. The system of equations (3.3.2) is said to be homogeneous if vector $\underline{b} = \underline{0}$, that is, all $b_i = 0$, otherwise the system is called non-homogeneous.

We state the following useful result on the solvability of linear systems.

Theorem 1: A non-homogeneous system of n linear equations in n unknown has a unique solution if and only if the coefficient matrix A is non singular ($\det A \neq 0$) and the solution can be expressed as $\mathbf{x} = A^{-1}\mathbf{b}$.

1.3 DIRECT METHODS

In schools, generally Cramer's rule/method is taught to solve system of simultaneous equations, based on the evaluation of determinants. This is a direct method. When n is small (say, 3 or 4), this rule is satisfactory. However, the number of multiplication operations needed increases very rapidly as the number of equations increases as shown below:

Number of equations	Number of multiplication operations
2	8
3	51
4	364
5	2885
.	.
.	.
.	.
10	359251210

Hence a different approach is needed to solve such a system of equations on a computer. Thus, Cramer's rule, although the simplest and the most direct method, remains a theoretical rule and we have to look for other efficient direct methods. We are going to discuss one such direct method – Gauss' elimination method next after stating Cramer's Rule for the sake of completeness.

1.3.1 Cramer's Rule

In the system of equation (3.3.2), let $\Delta = \det(A)$ and $\mathbf{b} \neq 0$. Then the solutions of the system is obtained as $x_i = \frac{\Delta_i}{\Delta}$, $i = 1, 2, \dots, n$

where Δ_i is the determinant of the matrix obtained from A by replacing the i^{th} column of Δ by vector \mathbf{b} .

1.3.2 Gauss Elimination Method

In Gauss's elimination method, one usually finds successively a finite number of linear systems equivalent to the given one such that the final system is so simple that its solution may be readily computed. In this method, the matrix A is reduced to the form U (upper triangular matrix) by using the elementary row operations like

- (i) interchanging any two rows
- (ii) multiplying (or dividing) any row by a non-zero constant
- (iii) adding (or subtracting) a constant multiple of one row to another row.

If any matrix A is transformed to another matrix B by a series of row operations, we say that A and B are equivalent matrices. More specifically we have.

Definition 3: A matrix B is said to be row-equivalent to a matrix A , if B can be obtained from A by a using a finite number of row operations.

Two linear systems $A\mathbf{x} = \mathbf{b}$ and $A'\mathbf{x} = \mathbf{b}'$ are said to be equivalent if they have the same solution. Hence, if a sequence of elementary operations on $A\mathbf{x} = \mathbf{b}$ produces the new system $A'\mathbf{x} = \mathbf{b}'$, then the systems $A\mathbf{x} = \mathbf{b}$ and $A'\mathbf{x} = \mathbf{b}'$ are equivalent.

Let us illustrate (Naive) Gauss elimination method by considering a system of three equations:

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 &= b_1 \\ a_{21}x_1 + a_{22}x_2 + a_{23}x_3 &= b_2 \\ a_{31}x_1 + a_{32}x_2 + a_{33}x_3 &= b_3. \end{aligned} \quad (1.3.1)$$

Let $a_{11} \neq 0$. We multiply first equation of the system by $-\frac{a_{22}}{a_{11}}$ and add

to the second equation. Then we multiply the first equation by $-\frac{a_{31}}{a_{11}}$ and add to the third equation. The new equivalent system (first derived system) then becomes

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 &= b_1 \\ a_{22}^{(1)}x_2 + a_{23}^{(1)}x_3 &= b_2^{(1)} \\ a_{32}^{(1)}x_2 + a_{33}^{(1)}x_3 &= b_3^{(1)} \end{aligned} \quad (1.3.2)$$

where

$$a_{22}^{(1)} = a_{22} - \frac{a_{21}}{a_{11}} \cdot a_{12}, \quad a_{23}^{(1)} = a_{23} - \frac{a_{21}}{a_{11}} \cdot a_{13},$$

$$b_2^{(1)} = b_2 - \frac{a_{21}}{a_{11}} \cdot b_1, \text{ etc.}$$

Next, we multiply the second equation of the derived system provided $a_{22}^{(1)} \neq 0$, by $-\frac{a_{32}^{(1)}}{a_{22}^{(1)}}$ and add to the third equation of (3.4.2). The system becomes

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 &= b_1 \\ a_{22}^{(1)}x_2 + a_{23}^{(1)}x_3 &= b_2^{(1)} \\ a_{33}^{(2)}x_3 &= b_3^{(2)} \end{aligned} \quad (1.3.3)$$

where

$$a_{33}^{(2)} = a_{33}^{(1)} - \frac{a_{32}^{(1)}}{a_{22}^{(1)}} \cdot a_{23}^{(1)}$$

and

$$b_3^{(2)} = b_3^{(1)} - \frac{a_{32}^{(1)}}{a_{22}^{(1)}} b_2^{(1)}.$$

This system is an upper-triangular system and can be solved using back substitutions method provided $a_{33}^{(2)} \neq 0$. That is, the last equation gives $x_3 = \frac{b_3^{(2)}}{a_{33}^{(2)}}$; then substituting

this value of x_3 in the last but one equation (second) we get the value of x_2 and then substituting the obtained values of x_3 and x_2 in the first equation we compute x_1 . This process of solving an upper-triangular system of linear equations is often called **back substitution**. We illustrate this by the following example:

Example 1: Solve the following system of equations consisting of four equations.

$$\begin{aligned} \text{(Equation 1)} \quad E_1: \quad x_1 + x_2 + 0x_3 + 3x_4 &= 4 \\ E_2: \quad 2x_1 + x_2 - x_3 + x_4 &= 1 \\ E_3: \quad 3x_1 - x_2 - x_3 + 2x_4 &= -3 \\ E_4: \quad -x_1 + 2x_2 + 3x_3 - x_4 &= 4. \end{aligned}$$

Solution: The first step is to use first equation to eliminate the unknown x_1 from second, third and fourth equation. This is accomplished by performing $E_2 - 2E_1$, $E_3 - 3E_1$ and $E_4 + E_1$. This gives the derived system as

$$\begin{aligned} E'_1: & x_1 + x_2 + 0x_3 + 3x_4 = 4 \\ E'_2: & -x_2 - x_3 + 5x_4 = -7 \\ E'_3: & -4x_2 - x_3 - 7x_4 = -15 \\ E'_4: & 3x_2 + 3x_3 + 2x_4 = 8. \end{aligned}$$

In this new system, E'_2 is used to eliminate x_2 from E'_3 and E'_4 by performing the operations $E'_3 - 4E'_2$ and $E'_4 + 3E'_2$. The resulting system is

$$\begin{aligned} E''_1: & x_1 + x_2 + 0x_3 + 3x_4 = 4 \\ E''_2: & -x_2 - x_3 + 5x_4 = -7 \\ E''_3: & 3x_3 + 13x_4 = 13 \\ E''_4: & -13x_4 = -13. \end{aligned}$$

This system of equation is now in triangular form and can be solved by back substitution. E''_4 gives $x_4 = 1$, E''_3 gives

$$x_3 = \frac{1}{3}(13 - 13x_4) = \frac{1}{3}(13 - 13 \times 1) = 0.$$

E''_2 gives $x_2 = -(-7 + 5x_4 + x_3) = -(-7 + 5 \times 1 + 0) = 2$ and E''_1 gives $x_1 = 4 - 3x_4 - x_2 = 4 - 3 \times 1 - 2 = -1$.

The above procedure can be carried out conveniently in matrix form as shown below:

We consider the Augmented matrix $[A|b]$ and perform the elementary row operations on the augmented matrix.

$$\begin{aligned} [A|b] &= \left[\begin{array}{cccc|c} 1 & 2 & 0 & 3 & 4 \\ 2 & 1 & -1 & 1 & 1 \\ 3 & -1 & -1 & 2 & -3 \\ -1 & 2 & 3 & -1 & 4 \end{array} \right] \quad \begin{array}{l} R_2 - 2R_1, R_3 - 3R_1 \text{ and} \\ R_4 + R_1 \text{ gives} \end{array} \\ &= \left[\begin{array}{cccc|c} 1 & 2 & 0 & 3 & 4 \\ 0 & -1 & -1 & -5 & -7 \\ 0 & -4 & -1 & -7 & -15 \\ 0 & 3 & 3 & 2 & 8 \end{array} \right] \quad R_3 - 4R_2, R_4 + 3R_2 \text{ gives} \\ &= \left[\begin{array}{cccc|c} 1 & 2 & 0 & 3 & 4 \\ 0 & -1 & -1 & -5 & -7 \\ 0 & 0 & 3 & 13 & 13 \\ 0 & 0 & 0 & -13 & -13 \end{array} \right] \end{aligned}$$

This is the final equivalent system:

$$\begin{aligned} x_1 + x_2 + 0x_3 + 3x_4 &= 4 \\ -x_2 - x_3 - 5x_4 &= -7 \\ 3x_3 + 13x_4 &= 13 \\ -13x_4 &= -13. \end{aligned}$$

The method works with the assumption that none of the elements $a_{11}, a_{22}^{(1)}, \dots,$

$a_{n-1,n-1}^{(n-2)}, a_{n,n}^{(n-1)}$ is zero. This does not necessarily mean that the linear system is not solvable, but the following technique may yield the solution:

Suppose $a_{kk}^{(k-1)} = 0$ for some $k = 2, \dots, n-2$. The k th column of $(k-1)$ th equivalent system from the k th row is searched for the first non zero entry. If $a_{pk}^{(k)} \neq 0$ for some p ,

$k + 1 \leq p \leq n$, then interchange R_k by R_p to obtain an equivalent system and continue the procedure. If $a_{pk}^{(k)} = 0$ for $p = k, k + 1, \dots, n$, it can be shown that the linear system does not have a unique solution and hence the procedure is terminated.

You may now solve the following exercises:

- E1) Solve the system of equations
 $3x_1 + 2x_2 + x_3 = 3$
 $2x_1 + x_2 + x_3 = 0$
 $6x_1 + 2x_2 + 4x_3 = 6$
 using Gauss elimination method. Does the solution exist?
- E2) Solve the system of equations
 $16x_1 + 22x_2 + 4x_3 = -2$
 $4x_1 - 3x_2 + 2x_3 = 9$
 $12x_1 + 25x_2 + 2x_3 = -11$
 using Gauss elimination method and comment on the nature of the solution.
- E3) Solve the system of equations by Gauss elimination.
 $x_1 - x_2 + 2x_3 - x_4 = -8$
 $2x_1 - 2x_2 + 3x_3 - 3x_4 = -20$
 $x_1 + x_2 + x_3 + 0.x_4 = -2$
 $x_1 - x_2 + 4x_3 + 3x_4 = 4$
- E4) Solve the system of equations by Gauss elimination.
 $x_1 + x_2 + x_3 + x_4 = 7$
 $x_1 + x_2 + 0.x_3 + 2x_4 = 8$
 $2x_1 + 2x_2 + 3x_3 + 0.x_4 = 10$
 $-x_1 - x_2 - 2x_3 + 2x_4 = 0$
- E5) Solve the system of equation by Gauss elimination.
 $x_1 + x_2 + x_3 + x_4 = 7$
 $x_1 + x_2 + 2x_4 = 5$
 $2x_1 + 2x_2 + 3x_3 = 10$
 $-x_1 - x_2 - 2x_3 + 2x_4 = 0$

It can be shown that in Gauss elimination procedure and back substitution

$$(2n^3 + 3n^2 - 5n) \Big/ 6 + \frac{n^2 + n}{2} \text{ multiplications/divisions and } \frac{n^3 - n}{3} + \frac{n^2 - n}{2}$$

additions/subtractions are performed respectively. The total arithmetic operation

$$\text{involved in this method of solving a } n \times n \text{ linear system is } \frac{n^3 + 3n^2 - n}{3}$$

$$\text{multiplication/divisions and } \frac{2n^3 + 3n^2 - 5n}{6} \text{ additions/subtractions.}$$

Definition 4: In Gauss elimination procedure, the diagonal elements $a_{11}, a_{22}^{(1)}, a_{33}^{(2)}$, which have been used as divisors are called pivots and the corresponding equations, are called pivotal equations.

1.3.3 Pivoting Strategies

If at any stage of the Gauss elimination, one of these pivots say $a_{ii}^{i-1} (a_{11}^{(0)} = a_{11})$, vanishes then we have indicated a modified procedure. But it may also happen that the pivot $a_{ii}^{(i-1)}$, though not zero, may be very small in magnitude compared to the

remaining elements ($\geq i$) in the i th column. Using a small number as divisor may lead to growth of the round-off error. The use of large multipliers like

$$\frac{-a_{i+1}^{(i-1)}, i}{a_{ii}^{(i-1)}}, \frac{a_{i+2,i}^{(i-1)}}{a_{ii}^{(i-1)}}$$

etc. will lend to magnification of errors both during the elimination phase and during the back substitution phase of the solution procedure. This can be avoided by rearranging the remaining rows (from i th row up to n th row) so as to obtain a non-vanishing pivot or to choose one that is largest in magnitude in that column. This is called pivoting strategy.

There are two types of pivoting strategies: partial pivoting (maximal column pivoting) and complete pivoting. We shall confine to simple partial pivoting and complete pivoting. That is, the method of scaled partial pivoting will not be discussed. Also there is a convenient way of carrying out the pivoting procedure where instead of interchanging the equations all the time, the n original equations and the various changes made in them can be recorded in a systematic way using the augmented matrix $[A|b]$ and storing the multipliers and maintaining pivotal vector. We shall just illustrate this with the help of an example. However, leaving aside the complexities of notations, the procedure is useful in computation of the solution of a linear system of equations.

If exact arithmetic is used throughout the computation, pivoting is not necessary unless the pivot vanishes. But, if computation is carried up to a fixed number of digits (precision fixed), we get accurate results if pivoting is used.

The following example illustrates the effect of round-off error while performing Gauss elimination:

Example 2: Solve by the Gauss elimination the following system using four-digit arithmetic with rounding.

$$\begin{aligned} 0.003000x_1 + 59.14x_2 &= 59.17 \\ 5.291x_1 - 6.130x_2 &= 46.78. \end{aligned}$$

Solution: The first pivot element $a_{11}^0 = a_{11} = 0.0030$ and its associated multiplier is

$$\frac{5.291}{0.0030} = 1763.66 \approx 1763$$

Performing the operation of elimination of x_1 from the second equation with appropriate rounding we got

$$\begin{aligned} 0.003000x_1 + 59.14x_2 &= 59.17 \\ -104300x_2 &= -104400 \end{aligned}$$

By backward substitution we have

$$x_2 = 1.001 \text{ and } x_1 = \frac{59.17 - (59.14)(1.001)}{0.00300} = -10.0$$

The linear system has the exact solution $x_1 = 10.00$ and $x_2 = 1,000$.

However, if we use the second equation as the first pivotal equation and solve the system, the four digit arithmetic with rounding yields solution as $x_1 = 10.00$ and $x_2 = 1.000$. This brings out the importance of partial or maximal column pivoting.

Partial pivoting (Column Pivoting)

In the first stage of elimination, instead of using $a_{11} \neq 0$ as the pivot element, the first column of the matrix A ($[A|b]$) is searched for the largest element in magnitude and this largest element is then brought at the position of the first pivot by interchanging first row with the row having the largest element in magnitude in the first column.

Next, after elimination of x_1 , the second column of the derived system is searched for the largest element in magnitude among the $(n - 1)$ element leaving the first element. Then this largest element in magnitude is brought at the position of the second pivot by interchanging the second row with the row having the largest element in the second column of the derived system. The process of searching and interchanging is repeated in all the $(n - 1)$ stages of elimination. For selecting the pivot we have the following algorithm:

For $i = 1, 2, \dots, n$ find j such that

$$\left| a_{ji}^{(i-1)} \right| = \max_{i \leq k \leq n} \left| a_{ki}^{(i-1)} \right| \quad \left(a_{ji}^0 = a_{ji} \right)$$

Interchange i th and j th rows and eliminate x_i .

Complete Pivoting

In the first stage of elimination, we look for the largest element in magnitude in the entire matrix A first. If the element is a_{pq} , then we interchange first row with p th row and interchange first column with q th column, so that a_{pq} can be used as a first pivot. After eliminating x_q , the process is repeated in the derived system, more specifically in the square matrix of order $n - 1$, leaving the first row and first column. Obviously, complete pivoting is quite cumbersome.

Scaled partial pivoting (Scaled column pivoting)

First a scale factor d_i for each row i is defined by $d_i = \max_{i \leq j \leq n} |a_{ij}|$

If $d_i = 0$ for any i , there is no unique solution and procedure is terminated. In the first stage choose the first integer k such that

$$\left| a_{k1} \right| / d_k = \max_{i \leq j \leq n} \left| a_{j1} / d_j \right|$$

interchange first row and k th row and eliminate x_1 . The process is repeated in the derived system leaving aside first row and first column.

We now illustrate these pivoting strategies in the following examples.

Example 3: Solve the following system of linear equations with partial pivoting

$$\begin{aligned} x_1 - x_2 + 3x_3 &= 3 \\ 2x_1 + x_2 + 4x_3 &= 7 \\ 3x_1 + 5x_2 - 2x_3 &= 6 \end{aligned}$$

$$[A|b] = \left(\begin{array}{ccc|c} 1 & -1 & 3 & 3 \\ 2 & 1 & 4 & 7 \\ 3 & 5 & -2 & 6 \end{array} \right) \quad R_1 - \frac{1}{3}R_3, R_2 - \frac{2}{3}R_3$$

$$= \left(\begin{array}{ccc|c} 0 & -\frac{8}{3} & \frac{11}{3} & 1 \\ 0 & -\frac{7}{3} & \frac{16}{3} & 3 \\ 3 & 5 & -2 & 6 \\ \hline 0 & -\frac{8}{3} & \frac{11}{3} & 1 \\ 0 & 0 & \frac{51}{24} & \frac{17}{8} \\ 3 & 5 & -2 & 6 \end{array} \right) \quad R_2 - \frac{7}{3} \cdot \frac{3}{8} R_1$$

Re-arranging the equations (3rd equation becomes the first equation and first equation becomes the second equation in the derived system), we have

$$\begin{aligned} 3x_1 + 5x_2 - 2x_3 &= 6 \\ -\frac{8}{3}x_2 + \frac{11}{3}x_3 &= 1 \\ \frac{51}{24}x_3 &= \frac{17}{8} \end{aligned}$$

Using back substitution we have $x_1 = 1$, $x_2 = 1$ and $x_3 = 1$.

You may now solve the following exercises:

- E6) Solve the system of linear equation given in the Example 3 by complete pivoting.
- E7) Solve the system of linear equation given in Example 3 by scaled partial pivoting.
- E8) Solve the system of equations with partial (maximal column) pivoting.

$$\begin{aligned} x_1 + x_2 + x_3 &= 6 \\ 3x_1 + 3x_2 + 4x_3 &= 20 \\ 2x_1 + x_2 + 3x_3 &= 13 \end{aligned}$$

1.4 ITERATIVE METHODS

Consider the system of equations

$$A\mathbf{x} = \mathbf{b} \quad \dots (1.4.1)$$

Where A is an $n \times n$ non-singular matrix. An iterative technique to solve the $n \times n$ linear system (1.4.1) starts with an initial approximation $\underline{x}^{(0)}$ to the solution \mathbf{x} , and generates a sequence of vectors $\{\underline{x}^k$ that **converges** to \mathbf{x} , the actual solution vector (When $\max_{1 \leq i \leq n} |x_i^{(k)} - x_i| < \varepsilon$ for some k when ε is a given small positive numbers.).

Most of these iterative techniques entails a process that converts the system $A\mathbf{x} = \mathbf{b}$ into an equivalent system of the form $\mathbf{x} = T\mathbf{x} + \mathbf{c}$ for some $n \times n$ matrix T and vector \mathbf{c} . In general we can write the iteration method for solving the linear system (3.5.1) in the form

$$\mathbf{x}^{(k+1)} = T\mathbf{x}^{(k)} + \mathbf{c} \quad k = 0, 1, 2, \dots,$$

T is called the iteration matrix and depends on A, \mathbf{c} is a column vector which depends on A and \mathbf{b} . We illustrate this by the following example.

Iterative methods are generally used when the system is large (when $n > 50$) and the matrix is sparse (matrices with very few non-zero entries).

Example 4: Convert the following linear system of equations into equivalent form $\mathbf{x} = \mathbf{T}\mathbf{x} + \mathbf{c}$.

$$\begin{aligned} 10x_1 - x_2 + 2x_3 &= 6 \\ -x_1 + 11x_2 - x_3 + 3x_4 &= 25 \\ 2x_1 - x_2 + 10x_3 - x_4 &= -11 \\ 3x_2 - x_3 + 8x_4 &= 15 \end{aligned}$$

Solution: We solve the i th equation for x_i (assuming that $a_{ii} \neq 0 \forall i$. If not, we can interchange equations so that is possible)

$$x_1 = +\frac{1}{10}x_2 - \frac{1}{5}x_3 + \frac{3}{5}$$

$$x_2 = \frac{1}{11}x_1 + \frac{1}{11}x_3 - \frac{3}{11}x_4 + \frac{25}{11}$$

$$x_3 = -\frac{1}{5}x_1 + \frac{1}{10}x_2 + \frac{1}{10}x_4 - \frac{11}{10}$$

$$x_4 = -\frac{3}{8}x_2 + \frac{1}{8}x_3 + \frac{15}{8}$$

$$\text{Here } \mathbf{T} = \begin{bmatrix} 0 & \frac{1}{10} & -\frac{1}{5} & 0 \\ \frac{1}{11} & 0 & \frac{1}{11} & -\frac{3}{11} \\ -\frac{1}{5} & \frac{1}{10} & 0 & \frac{1}{10} \\ 0 & -\frac{3}{8} & \frac{1}{8} & 0 \end{bmatrix} \quad \text{and } \mathbf{c} = \begin{bmatrix} \frac{3}{5} \\ \frac{25}{11} \\ -\frac{11}{10} \\ \frac{15}{8} \end{bmatrix}$$

1.4.1 The Jacobi Iterative Method

This method consists of solving the i th equation of $A\mathbf{x} = \mathbf{b}$ for x_i , to obtain

$$x_i = \sum_{\substack{j=1 \\ j \neq i}}^n \left(\frac{-a_{ij}x_j}{a_{ii}} \right) + \frac{b_i}{a_{ii}} \quad \text{for } i = 1, 2, \dots, n$$

provided $a_{ii} \neq 0$.

We generate $\mathbf{x}^{(k+1)}$ from $\mathbf{x}^{(k)}$ for $k \geq 0$ by

$$x_i^{(k+1)} = \sum_{\substack{j=1 \\ j \neq i}}^n \left(\frac{-a_{ij}x_j^{(k)} + b_i}{a_{ii}} \right) \quad i = 1, 2, \dots, n \quad (1.4.2)$$

We state below a sufficient condition for convergence of the Jacobi Method.

Theorem

If the matrix A is strictly diagonally dominant, that is, if

$$|a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}}^n |a_{ij}|, \quad i = 1, 2, \dots, n.$$

then the Jacobi iteration method (3.5.2) converges for any initial approximation $\mathbf{x}^{(0)}$.

Generally $\mathbf{x}^{(0)} = \mathbf{0}$ is taken in the absence of any better initial approximation.

Example 5: Solve the linear system $A\mathbf{x} = \mathbf{b}$ given in previous example (Example 4) by Jacobi method rounded to four decimal places.

Solution: Letting $\mathbf{x}^{(0)} = (0, 0, 0, 0)^T$, we get

$$\mathbf{x}^{(1)} = (0.6000, 2.2727 - 1.1000, 1.8750)^T$$

$$\mathbf{x}^{(2)} = (1.0473, 1.7159, -0.8052, 0.8852)^T \text{ and}$$

$$\mathbf{x}^{(3)} = (0.9326, 2.0533, -1.0493, 1.1309)^T$$

Proceeding similarly one can obtain

$$\mathbf{x}^{(5)} = (0.9890, 2.0114, -1.0103, 1.0214)^T \text{ and}$$

$$\mathbf{x}^{(10)} = (1.0001, 1.9998, -0.9998, 0.9998)^T.$$

The solution is $\mathbf{x} = (1, 2, -1, 1)^T$. You may note that $\mathbf{x}^{(10)}$ is a good approximation to the exact solution compared to $\mathbf{x}^{(5)}$.

You also observe that A is strictly diagonally dominant (since $10 > 1 + 2$, $11 > 1 + 1 + 3$, $10 > 2 + 1 + 1$ and $8 > 3 + 1$).

Now we see how $A\mathbf{x} = \mathbf{b}$ is transformed to an equivalent system $\mathbf{x} = T\mathbf{x} + \mathbf{c}$.

The matrix can be written as

$$A = D + L + U \quad \text{where}$$

$$D = \begin{bmatrix} a_{11} & 0 & \dots & 0 \\ 0 & a_{22} & \dots & 0 \\ 0 & \dots & 0 & a_{nn} \end{bmatrix}$$

$$U = \begin{bmatrix} 0 & a_{12} & \dots & a_{1n} \\ 0 & 0 & a_{23} & \dots & a_{2n} \\ 0 & 0 & 0 & \dots & a_{n-1, n} \end{bmatrix}$$

$$L = \begin{bmatrix} 0 & 0 & \dots & \dots & 0 \\ a_2 & 0 & \dots & \dots & 0 \\ a_3 & a_{32} & 0 & \dots & 0 \\ a_n & a_{n2} & \dots & a_{n, n-1} \end{bmatrix}$$

Since $(D + L + U)\mathbf{x} = \mathbf{b}$

$$D\mathbf{x} = - (L + U)\mathbf{x} + \mathbf{b}$$

$$\mathbf{x} = -D^{-1}(L + U)\mathbf{x} + D^{-1}\mathbf{b}$$

$$\text{i.e. } T = -D^{-1}(L + U) \text{ and } \mathbf{c} = D^{-1}\mathbf{b}.$$

In Jacobi method, each of the equations is simultaneously changed by using the most recent set of \mathbf{x} -values. Hence the Jacobi method is called method of simultaneous displacements.

You may now solve the following exercises:

- E9) Perform five iterations of the Jacobi method for solving the system of equations.

$$\begin{bmatrix} 5 & -1 & -1 & -1 \\ -1 & 10 & -1 & -1 \\ -1 & -1 & 5 & -1 \\ -1 & -1 & -1 & 10 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} -4 \\ 12 \\ 8 \\ 34 \end{bmatrix}$$

Starting with $\mathbf{x}^{(0)} = (0,0,0,0)$. The exact solution is $\mathbf{x} = (1,2,3,4)^T$. How good $\mathbf{x}^{(5)}$ as an approximation to \mathbf{x} ?

- E10) Perform four iterations of the Jacobi method for solving the following system of equations.

$$\begin{bmatrix} 2 & -1 & -0 & -0 \\ -1 & 2 & -1 & 0 \\ 0 & -1 & 2 & -1 \\ 0 & 0 & -1 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} -1 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

With $\mathbf{x}^{(0)} = (0.5, 0.5, 0.5, 0.5)^T$. Here $\mathbf{x} = (1, 1, 1, 1)^T$. How good $\mathbf{x}^{(5)}$ as an approximation to \mathbf{x} ?

1.4.2 The Gauss-Seidel Iteration Method

In this method, we can write the iterative scheme of the system of equations $\mathbf{Ax} = \mathbf{b}$ as follows:

$$\begin{aligned} a_{11}x_1^{(k+1)} &= -a_{12}x_2^{(k)} - a_{13}x_3^{(k)} - \dots - a_{1n}x_n^{(k)} + b_1 \\ a_{21}x_1^{(k+1)} + a_{22}x_2^{(k+1)} &= -a_{23}x_3^{(k)} - \dots - a_{2n}x_n^{(k)} + b_2 \\ &\vdots \\ a_{n1}x_1^{(k+1)} + a_{n2}x_2^{(k+1)} \dots + a_{nn}x_n^{(k+1)} &= b_n \end{aligned}$$

In matrix form, this system can be written as $(\mathbf{D} + \mathbf{L}) \mathbf{x}^{(k+1)} = -\mathbf{U} \mathbf{x}^{(k)} + \mathbf{b}$ with the same notation as adopted in Jacobi method.

From the above, we get

$$\begin{aligned} \mathbf{x}^{(k+1)} &= -(\mathbf{D} + \mathbf{L})^{-1} \mathbf{U} \mathbf{x}^{(k)} + (\mathbf{D} + \mathbf{L})^{-1} \mathbf{b} \\ &= \mathbf{T} \mathbf{x}^{(k)} + \mathbf{c}_n \end{aligned}$$

i.e. $\mathbf{T} = -(\mathbf{D} + \mathbf{L})^{-1} \mathbf{U}$ and $\mathbf{c} = (\mathbf{D} + \mathbf{L})^{-1} \mathbf{b}$

This iteration method is also known as the method of successive displacement.

For computation point of view, we rewrite $(\mathbf{A} \mathbf{x})$ as

$$x_i^{(k+1)} = -\frac{1}{a_{ii}} \left[\sum_{j=1}^{i-1} a_{ij} x_j^{(k+1)} + \sum_{j=i+1}^n a_{ij} x_j^{(k)} - b_i \right]$$

$i = 1, 2, \dots, n$

Also in this case, if A is diagonally dominant, then iteration method always converges. In general Gauss-Seidel method will converge if the Jacobi method converges and will converge at a faster rate. You can observe this in the following example. We have not considered the problem: How many iterations are needed to have a reasonably good approximation to \mathbf{x} ? This needs the concept of matrix norm.

Example 6: Solve the linear system $A\mathbf{x} = \mathbf{b}$ given in Example 4 by Gauss-Seidel method rounded to four decimal places. The equations can be written as follows:

$$\begin{aligned}x_1^{(k+1)} &= \frac{1}{10}x_2^{(k)} - \frac{1}{3}x_3^{(k)} + \frac{3}{5} \\x_2^{(k+1)} &= \frac{1}{11}x_1^{(k+1)} + \frac{1}{11}x_3^k - \frac{3}{11}x_4^{(k)} + \frac{25}{11} \\x_3^{(k+1)} &= -\frac{1}{3}x_1^{(k+1)} + \frac{1}{10}x_2^{(k+1)} + \frac{1}{10}x_4^{(k)} - \frac{11}{10} \\x_4^{(k+1)} &= -\frac{3}{8}x_2^{(k+1)} + \frac{1}{8}x_3^{(k+1)} + \frac{15}{8}.\end{aligned}$$

Letting $\mathbf{x}^{(0)} = (0, 0, 0, 0)^T$ we have from first equation

$$\begin{aligned}x_1^{(1)} &= 0.6000 \\x_2^{(1)} &= \frac{0.6000}{3} + \frac{25}{11} = 2.3273 \\x_3^{(1)} &= -\frac{0.6000}{3} + \frac{1}{10}(2.3273) - \frac{11}{10} = -0.1200 + 0.2327 - 1.1000 = -0.9873 \\x_4^{(1)} &= -\frac{3}{8}(2.3273) + \frac{1}{8}(-0.9873) + \frac{15}{8}\end{aligned}$$

$$= -0.8727 - 0.1234 + 1.8750$$

$$= 0.8789$$

Using $\mathbf{x}^{(1)}$ we get

$$\mathbf{x}^{(2)} = (1.0300, 2.037, -1.014, 0.9844)^T$$

and we can check that

$$\mathbf{x}^{(5)} = (1.0001, 2.0000, -1.0000, 1.0000)^T$$

Note that $\mathbf{x}^{(5)}$ is a good approximation to the exact solution. Here are a few exercises for you to solve.

You may now solve the following exercises:

E11) Perform four iterations (rounded to four decimal places) using Jacobi Method and Gauss-Seidel method for the following system of equations.

$$\begin{bmatrix} -8 & 1 & 1 \\ 1 & -5 & -1 \\ 1 & 1 & -4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 1 \\ 16 \\ 7 \end{bmatrix}$$

With $\mathbf{x}^{(0)} = (0, 0, 0)^T$. The exact solution is $(-1, -4, -3)^T$. Which method gives better approximation to the exact solution?

E12) For linear system given in E10), use the Gauss Seidel method for solving the system starting with $\mathbf{x}^{(0)} = (0.5, 0.5, 0.5)^T$ obtain $\mathbf{x}^{(4)}$ by Gauss-Seidel method and compare this with $\mathbf{x}^{(4)}$ obtained by Jacobi method in E10).

1.4.3 Comparison of Direct and Iterative Methods

Both the methods have their strengths and weaknesses and a choice is based on the particular linear system to be solved. We mention a few of these below:

Direct Method

1. The direct methods are generally used when the matrix A is dense or filled, that is, there are few zero elements, and the order of the matrix is not very large, say $n < 50$.
2. The rounding errors may become quite large for ill conditioned equations (If at any stage during the application of pivoting strategy, it is found that all values of $\left\{ \left| a_{mk} \right| \right\}$ for $m = k + 1$, to n are less than a pre-assigned small quantity ε , then the equations are ill-conditioned and no useful solution is obtained). Ill-conditioned matrices are not discussed in this unit.

Iterative Method

1. These methods are generally used when the matrix A is sparse and the order of the matrix A is very large say $n > 50$. Sparse matrices have very few non-zero elements.
2. An important advantage of the iterative methods is the small rounding error. Thus, these methods are good choice for ill-conditioned systems.
3. However, convergence may be guaranteed only under special conditions. But when convergence is assured, this method is better than direct.

With this we conclude this unit. Let us now recollect the main points discussed in this unit.

1.5 SUMMARY

In this unit we have dealt with the following:

1. We have discussed the direct methods and the iterative techniques for solving linear system of equations $A\mathbf{x} = \mathbf{b}$ where A is an $n \times n$, non-singular matrix.
2. The direct methods produce the exact solution in a finite number of steps provided there are no round off errors. Direct method is used for linear system $A\mathbf{x} = \mathbf{b}$ where the matrix A is dense and order of the matrix is less than 50.
3. In direct methods, we have discussed Gauss elimination, and Gauss elimination with partial (maximal column) pivoting and complete or total pivoting.

4. We have discussed two iterative methods, Jacobi method and Gauss-Seidel method and stated the convergence criterion for the iteration scheme. The iterative methods are suitable for solving linear systems when the matrix is sparse and the order of the matrix is greater than 50.

1.6 SOLUTION/ANSWERS

E1) [A1b]
$$\left[\begin{array}{ccc|c} 3 & 2 & 1 & 3 \\ 2 & 1 & -1 & 0 \\ 6 & 2 & 4 & 6 \end{array} \right]$$

$a_{11} \neq 0$
 \longrightarrow
$$\left[\begin{array}{ccc|c} 3 & 2 & 1 & 3 \\ 0 & \frac{1}{3} & -\frac{1}{3} & -2 \\ 0 & -2 & 2 & 0 \end{array} \right]$$

$a_{22}^{(1)} \neq 0$
 \longrightarrow
$$\left[\begin{array}{ccc|c} 3 & 2 & 1 & 3 \\ 0 & -\frac{1}{3} & -\frac{1}{3} & -2 \\ 0 & 0 & 0 & 12 \end{array} \right]$$

This system has no solution since x_3 cannot be determined from the last equation. This system is said to be inconsistent. Also note that $\det(A) = 0$.

E2) [A1b]
$$\left[\begin{array}{ccc|c} 16 & 22 & 4 & -2 \\ 4 & -3 & 2 & 9 \\ 12 & 25 & 2 & -11 \end{array} \right]$$

$a_{11} \neq 0$
 \longrightarrow
$$\left[\begin{array}{ccc|c} 16 & 22 & 4 & 2 \\ 0 & -\frac{17}{2} & 1 & \frac{19}{2} \\ 0 & \frac{17}{2} & -1 & -\frac{19}{2} \end{array} \right]$$

$a_{22}^{(1)} \neq 0$
 \longrightarrow
$$\left[\begin{array}{ccc|c} 16 & 22 & 4 & 2 \\ 0 & -\frac{17}{2} & 1 & \frac{19}{2} \\ 0 & 0 & 0 & 0 \end{array} \right] \Rightarrow x_3 = \text{arbitrary value, } x_2 =$$

$$\frac{-2}{17} \left(\frac{19}{2} - x_3 \right) \text{ and } x_3 = \frac{1}{6} (-2 - 22x_3 - 22x_3)$$

This system has infinitely many solutions. Also you may check that $\det(A) = 0$.

E3) Final derived system:

$$\left[\begin{array}{cccc|c} 1 & -1 & 2 & -1 & -8 \\ 0 & 2 & -1 & 1 & 6 \\ 0 & 0 & -1 & -1 & -4 \\ 0 & 0 & 0 & 2 & 4 \end{array} \right] \text{ and the solution is } x_4 = 2, x_3 = 2$$

$x_2 = 3, x_1 = -7.$

E4) Final derived system:

$$\left[\begin{array}{cccc|c} 1 & -1 & 1 & 1 & 7 \\ 0 & 0 & -1 & 1 & 1 \\ 0 & 0 & 1 & -2 & -4 \\ 0 & 0 & 0 & 1 & 3 \end{array} \right] \quad \text{and the solution is}$$

$$x_4 = 3, \quad x_3 = 2, \quad x_2 \text{ arbitrary and } x_1 = 2 - x_2.$$

Thus this linear system has infinite number of solutions.

E5) Final derived system:

$$\left[\begin{array}{cccc|c} 1 & 1 & 1 & 1 & 7 \\ 0 & 0 & -1 & 1 & -2 \\ 0 & 0 & 1 & -2 & -4 \\ 0 & 0 & 0 & 1 & 3 \end{array} \right] \quad \text{and the solution does not}$$

exist since we have $x_4 = 3, x_3 = 2$ and third equation $\phi -x_3 + x_4 = -2$ implies $1 = -2$, leading to a contradiction.

E6) Since $|a_{32}|$ is maximum we rewrite the system as

$$\left[\begin{array}{ccc|c} 5 & 3 & -2 & 6 \\ 1 & 2 & 4 & 7 \\ -1 & 1 & 3 & 3 \end{array} \right] \quad \begin{array}{l} \text{by interchanging } R_1 \text{ and } R_3 \text{ and } C_1 \text{ and } C_2 \\ R_2 - \frac{1}{5}R_1, R_3 + \frac{1}{5}R_1 \text{ gives} \end{array}$$

$$\left[\begin{array}{ccc|c} 5 & 3 & -2 & 6 \\ 0 & \frac{7}{5} & \frac{22}{5} & \frac{29}{5} \\ 0 & \frac{8}{5} & \frac{13}{5} & \frac{21}{5} \end{array} \right] \rightarrow \left[\begin{array}{ccc|c} 5 & -2 & 3 & 6 \\ 0 & \frac{22}{5} & \frac{7}{5} & \frac{29}{5} \\ 0 & \frac{13}{5} & \frac{8}{5} & \frac{21}{5} \end{array} \right] \quad \begin{array}{l} \text{by inter-} \\ \text{changing} \\ C_2 \text{ and } C_3 \end{array}$$

Since $|a_{23}|$ is maximum –

By $R_3 - \frac{5}{12}x \frac{13}{15}R_2$ we have

$$\left[\begin{array}{ccc|c} 5 & -2 & 3 & 6 \\ 0 & \frac{22}{5} & \frac{7}{5} & \frac{29}{5} \\ 0 & 0 & \frac{17}{22} & \frac{43}{22} \end{array} \right]$$

$$5x_2 + 3x_1 - 2x_3 = 6$$

$$\frac{22}{5}x_3 + \frac{7}{5}x_2 = \frac{29}{5}$$

$$\frac{17}{22}x_2 = \frac{17}{22}$$

$$\text{We have } x_2 = 1, \frac{22}{5}x_3 = \frac{29}{5} - \frac{7}{5} = \frac{22}{5} \Rightarrow x_3 = 1$$

$$3x_1 = 6 - 5 + 2 \Rightarrow x_1 = 1$$

E7) For solving the linear system by sealed partial pivoting we note that $d_1 = 3$, $d_2 = 4$ and $d_3 = 5$ in

$$W = [A|b] = \begin{bmatrix} 1 & -1 & 3 & 3 \\ 2 & 1 & 4 & 7 \\ 3 & 5 & -2 & 6 \end{bmatrix} \quad p = [a, 2, 3]^T$$

Since $\max \left\{ \frac{1}{3}, \frac{2}{4}, \frac{3}{5} \right\} = \frac{3}{5}$, third equation is chosen as the first pivotal equation.

Eliminating x_1 we have

d

$$\begin{bmatrix} 3 \\ 4 \\ 5 \end{bmatrix} \quad W^1 = \begin{bmatrix} \frac{1}{3} & \frac{8}{3} & \frac{11}{3} & 1 \\ -\frac{2}{3} & -\frac{7}{3} & \frac{16}{3} & 3 \\ 3 & 5 & -2 & 6 \end{bmatrix}$$

where we have used a square to enclose the pivot element and

in place of zero entries, after elimination of x_1 from 1st and 2nd equation, we have

stored the multipliers. Here $m_{11} = \frac{a_{11}}{a_{31}} = \frac{1}{3}$ and $m_{2,1} = \frac{a_{21}}{a_{31}} = \frac{2}{3}$

Instead of interchanging rows (here R_1 and R_3) we keep track of the pivotal equations being used by the vector $p = [3, 2, 1]^T$

In the next step we consider $\max \left\{ \frac{7}{3}, \frac{1}{4}, \frac{8}{3}, \frac{1}{3} \right\} = \frac{8}{9}$

So the second pivotal equation is the first equation.

i.e. $p = [3, 1, 3]^T$ and multiplier is $-\frac{7}{3} - \frac{8}{3} = -\frac{15}{3} = -5 = m_{2,2}$

$$\text{and } W^{(2)} = \begin{bmatrix} \frac{1}{3} & \frac{8}{3} & \frac{11}{3} & 1 \\ -\frac{2}{3} & -\frac{7}{3} & \frac{16}{3} & 3 \\ 3 & 5 & -2 & 6 \end{bmatrix} \quad p = [3, 1, 2]^T$$

The triangular system is as follows:

$$3x_1 + 5x_2 - 2x_3 = 6$$

$$-\frac{8}{3}x_2 + \frac{11}{3}x_3 = 1$$

$$\frac{17}{8}x_3 = \frac{17}{8}$$

By back substitution, this yields $x_1 = 1$, $x_2 = 1$ and $x_3 = 1$.

Remark: The p vector and storing of multipliers help solving the system $Ax = b'$ where b is changed b' .

$$\begin{aligned}
 \text{E8)} \quad [A, \mathbf{b}] &= \begin{bmatrix} 1 & 1 & 1 & 6 \\ 3 & 3 & 4 & 20 \\ 2 & 1 & 3 & 13 \end{bmatrix} \rightarrow \begin{bmatrix} 3 & 3 & 4 & 20 \\ 1 & 1 & 1 & 6 \\ 2 & 1 & 3 & 13 \end{bmatrix} \\
 &\xrightarrow{R_2 - \frac{1}{3}R_1, R_3 - \frac{2}{3}R_1} \begin{bmatrix} 3 & 3 & 4 & 20 \\ 0 & 0 & -\frac{1}{3} & -\frac{2}{3} \\ 0 & -1 & \frac{1}{3} & -\frac{1}{3} \end{bmatrix} \rightarrow \begin{bmatrix} 13 & 3 & 4 & 20 \\ 0 & -1 & \frac{1}{3} & -\frac{1}{3} \\ 0 & 0 & -\frac{1}{3} & -\frac{2}{3} \end{bmatrix}
 \end{aligned}$$

Since the resultant matrix is in triangular form, using back substitution we get $x_3 = 2$, $x_2 = 1$ and $x_1 = 3$.

E9) Using $\mathbf{x}^{(0)} = [0, 0, 0, 0]^T$ we have

$$\mathbf{x}^{(1)} = [-0.8, 1.2, 1.6, 3.4]^T$$

$$\mathbf{x}^{(2)} = [0.44, 1.62, 2.36, 3.6]^T$$

$$\mathbf{x}^{(3)} = [0.716, 1.84, 2.732, 3.842]^T$$

$$\mathbf{x}^{(4)} = [0.8823, 1.9290, 2.8796, 3.9288]^T$$

E10) Using $\mathbf{x}^{(0)} = [0.5, 0.5, 0.5, 0.5]^T$, we have

$$\mathbf{x}^{(1)} = [0.75, 0.5, 0.5, 0.75]^T$$

$$\mathbf{x}^{(2)} = [0.75, 0.625, 0.625, 0.75]^T$$

$$\mathbf{x}^{(3)} = [0.8125, 0.6875, 0.6875, 0.8125]^T$$

$$\mathbf{x}^{(4)} = [0.8438, 0.75, 0.75, 0.8438]^T$$

E 11) By Jacobi method we have

$$\mathbf{x}^{(1)} = [-0.125, -3.2, -1.75]^T$$

$$\mathbf{x}^{(2)} = [-0.7438, -3.5750, -2.5813]^T$$

$$\mathbf{x}^{(3)} = [-0.8945, -3.8650, -2.8297]^T$$

$$\mathbf{x}^{(4)} = [-0.9618, -3.9448, -2.9399]^T$$

Where as by Gauss-Seidel method, we have

$$\mathbf{x}^{(1)} = [-0.125, -3.225, -2.5875]^T$$

$$\mathbf{x}^{(2)} = [-0.8516, -3.8878, -2.9349]^T$$

$$\mathbf{x}^{(3)} = [-0.9778, -3.9825, -2.9901]^T$$

$$\mathbf{x}^{(4)} = [-0.9966, -3.9973, -2.9985]^T$$

E12) Starting with the initial approximation

$$\mathbf{x}^{(0)} = [0.5, 0.5, 0.5, 0.5]^T, \text{ we have the following iterates:}$$

$$\mathbf{x}^{(1)} = [0.75, 0.625, 0.5625, 0.7813]^T$$

$$\mathbf{x}^{(2)} = [0.8125, 0.6875, 0.7344, 0.8672]^T$$

$$\mathbf{x}^{(3)} = [0.8438, 0.7891, 0.8282, 0.9141]^T$$

$$\mathbf{x}^{(4)} = [0.8946, 0.8614, 0.8878, 0.9439]^T$$

Since the exact solution is $\mathbf{x} = [1, 1, 1, 1]^T$, the Gauss, Seidel method gives better approximation than the Jacobi method at fourth iteration.