

- 1) The Imitation Game
- 2) Utilizăm alg. lui Euclid cu împărțiri repetate

$$\begin{array}{r} 101000110101 : 10000111011 = 1 \text{ rest } 110111010 \\ \underline{10000111011} \\ 110111010 \end{array}$$

$$\begin{array}{r} 10000111011 : 110111010 = 100 \text{ rest } 110010011 \\ \underline{110111010} \\ 110010011 \\ \underline{0} \\ 110010011 \\ \underline{0} \\ 110010011 \end{array}$$

$$\begin{array}{r} 110111010 : 110010011 = 1 \text{ rest } 100111 \\ \underline{110010011} \\ 100111 \end{array}$$

$$\begin{array}{r} 110010011 : 100111 = 1010 \text{ rest } 1101 \\ \underline{100111} \\ 10110 \\ \underline{0} \\ 101101 \\ \underline{100111} \\ 1101 \\ \underline{0} \\ 1101 \end{array}$$

$$\begin{array}{l} 100111 \cdot 1101 = ? \\ \left. \begin{array}{l} \frac{100111}{1101} \Rightarrow 100111 \equiv 11010 \pmod{1101} \\ \text{Dar } 11010 = 10 \cdot 1101 \Rightarrow 11010 \equiv 0 \pmod{1101} \end{array} \right\} \Rightarrow 100111 \equiv 0 \pmod{1101} \Rightarrow \text{cmmdc}((101000110101)_2, (100001111011)_2) = (1101)_2 \end{array}$$

Verificăm rez. în baza 10.

$$\begin{aligned} (101000110101)_2 &= 2^0 + 2^2 + 2^4 + 2^5 + 2^9 + 2^{11} = 1 + 4 + 16 + 32 + 512 + 2048 = 53 + 2560 = (2613)_{10} \\ (100001111011)_2 &= 2^0 + 2^1 + 2^3 + 2^4 + 2^5 + 2^6 + 2^{11} = 1 + 2 + 8 + 16 + 32 + 64 + 2048 = 123 + 2048 = (2171)_{10} \\ (1101)_2 &= 2^0 + 2^2 + 2^3 = 1 + 4 + 8 = 13 \end{aligned}$$

$$\begin{array}{r} 2613 : 2171 = 1 \text{ rest } 442 \\ \underline{2171} \\ 442 \end{array}$$

$$\begin{array}{r} 2171 : 442 = 4 \text{ rest } 403 \\ \underline{1768} \\ 403 \end{array}$$

$$\begin{array}{r} 442 : 403 = 1 \text{ rest } 39 \\ \underline{403} \\ 39 \end{array}$$

$$\begin{array}{r} 403 : 39 = 10 \text{ rest } 13 \\ \underline{390} \\ 13 \end{array}$$

$$39 : 13 = 3 \text{ rest } 0 \Rightarrow \text{cmmdc}((2163)_{10}, (2171)_{10}) = (13)_{10} = (1101)_2$$

3) $0 \neq N$ are k biți $\Leftrightarrow (N)_2 = (\overline{a_{k-1} \dots a_0})_2 \Leftrightarrow (N)_6 = \sum_{i=0}^{k-1} (a_i \cdot 2^i) \Leftrightarrow 2^{k-1} \leq N \leq 2^k$

Dacă știm $(N)_2$, at. conversia $(N)_2 \mapsto (N)_6$ este de complexitate $O(k)$, adică liniară, deoarece la fiecare pas al sumei are loc câte o înmulțire și o adunare.

Dacă știm $(N)_6$, at. conversia $(N)_6 \mapsto (N)_2$ este de complexitate $O(\log_2 N)$, dar cum $N \leq 2^k$ at. $O(\log_2 N) = O(\log_2 2^k) = O(k)$ complexitate liniară.

5 39)

a) $(101011)_2 = 1 \cdot 2^0 + 1 \cdot 2^1 + 1 \cdot 2^3 + 1 \cdot 2^5 = 1 + 2 + 8 + 32 = (43)_{10}$

b) $(3C)_{16} = 12 \cdot 16^0 + 3 \cdot 16^1 = 12 + 48 = (60)_{10}$

c) $(12Z)_6 = 2 \cdot 6^0 + 2 \cdot 6^1 + 1 \cdot 6^2 = 2 + 12 + 36 = (50)_{10} = (302)_4$

\wedge
 $50 : 4 = 12 \text{ rest } 2$
 $12 : 4 = 3 \text{ rest } 0$
 $3 : 4 = 0 \text{ rest } 3$

d) $(44)_7 : (11)_7 = 4 \text{ rest } 0$

6 39) $7^{58} \pmod{59}$

$$7^{58} = (7^2)^{29} = 49^{29} = 49 \cdot (49)^{28} = 49 \cdot (2 \cdot 49)^{14} = 49 \cdot (4 \cdot 49)^7 = 49 \cdot (16 \cdot 49)^4 = 49 \cdot (29)^4 = \frac{49 \cdot 29 \cdot (84)^3}{=1421} \equiv 5 \cdot (15)^3 = \frac{5 \cdot 15 \cdot 225}{=75} = \frac{16 \cdot 48}{=768} \equiv 1 \pmod{59}$$