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Sparse Block-Encodings for Linear Combinations of Ladder Operators

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In this work, we detail the construction of quantum circuit oracles that create block-encodings for observables described as a linear combination of products of ladder operators acting on fermionic, antifermionic, and bosonic modes. We refer to this construction as LOBE (Ladder Operator Block-Encoding) and show how it can be used to simulate Hamiltonians involving interactions between these different types of particles. Our work builds off of similar sparse-oracle constructions in the literature, but generalizes prior works to establish a clear connection with block-encoding methods that are commonly refered to as LCU (Linear Combination of Unitaries). In addition to extnding LCU to more general observables that are given as a Linear combination of products of ladder operators, we also demonstrate how these oracles can be extended to include bosonic ladder operators. To our knowledge, this is the first block-encoding construction that allows for interactions between fermions, antifermions, and bosons, paving the way for simulation of more complicated quantum systems such as those that arise in high-energy physics.

I. INTRODUCTION

Wop, wop, wop, wop, Dot, fuck 'em up Wop, wop, wop, wop, I'ma do my stuff [1]

II. THEORY

Give background of ladder operators and constructions of realistic Hamiltonians/Observables from ladder operators

A. Encoding

Here we'll discuss how we encode the physical states we are interested in terms of qubits/registers.

Question: do we want to also cost-out the unary encoding or just the "compact" encoding that we've been primarily working with?

B. Ladder Operators

@Gus, you probably have much better language to define all of this stuff. I just needed to write something down so I could reference it in the circuit construction. Don't hesitate to scrap anything in here.

1. Feromons and Antiferomons

Define action of fermionic ladder operators.

Fermions (and antifermions) obey the Pauli-exclusion principle citation and therefore the occupation of a (anti)fermionic mode can only be occupied ($|1\rangle$) or unoccupied ($|0\rangle$). Fermionic (and antifermionic) ladder operators only act non-trivially on the qubits encoding the mode that the ladder operator acts on and we define their action as follows.

The fermionic creation operator is given by:

$$b_i^{\dagger} | n_{b_i} \rangle = \begin{cases} (-1)^{\sum_{j < i} b_j} | 1 \rangle & when | n_{b_i} \rangle is | 0 \rangle \\ 0 & when | n_{b_i} \rangle is | 1 \rangle \end{cases}$$
 (1)

where b_i denotes a fermionic ladder operator on the i^{th} mode, the † indicates a creation operator, and $|n_{b_i}\rangle$ is the occupation of the i^{th} fermionic mode. An antifermionic creation operator is defined as above with the symbol d to denote that the operator acts on antifermions.

For a fermionic creation operator, if the mode being acted upon is unoccupied, then the creation operator "creates" a fermion in that mode and applies a phase determined by the parity of the occupation of the previous modes. Therefore the ordering of the

modes in the encoding has an implication on the action of the operator that must be accounted for. Since fermionic modes can only be either occupied or unoccupied, then if the mode is already occupied the operator zeroes the amplitude of the quantum state, thereby "destoying" that portion of the quantum state.

The fermionic annihilation operator is given by:

$$b_i |n_{b_i}\rangle = \begin{cases} (-1)^{\sum_{j < i} b_j} |0\rangle & when |n_{b_i}\rangle = |1\rangle \\ 0 & when |n_{b_i}\rangle = |0\rangle \end{cases}$$
 (2)

and the antifermionic annihilation operator is likewise defined for d instead of b.

The action of the annihilation operators is similar (and opposite) to the creation operators. If the mode is already occupied, then the annihilation operator "annihilates" the fermion at that mode by setting the occupation to zero and applies a phase based on the parity of the occupation of the preceding modes. If the mode is unoccupied before the operator is applied, then the annihilation operator zeroes the amplitude.

2. Bosos

$$a_{i}^{\dagger} | n_{a_{i}} \rangle = \begin{cases} \sqrt{n_{a_{i}} + 1} | n_{a_{i}} + 1 \rangle & when | n_{a_{i}} \rangle \neq | \Omega \rangle \\ 0 & when | n_{a_{i}} \rangle = | \Omega \rangle \end{cases}$$

$$(3)$$

where a_i denotes a bosonic ladder operator on the i^{th} mode, the \dagger indicates a creation operator, $|n_{a_i}\rangle$ is the occupation of the i^{th} bosonic mode, and Ω is the maximum allowable bosonic occupation.

$$a_{i} | n_{a_{i}} \rangle = \begin{cases} \sqrt{n_{a_{i}}} | n_{a_{i}} - 1 \rangle & when | n_{a_{i}} \rangle \neq | 0 \rangle \\ 0 & when | n_{a_{i}} \rangle = | 0 \rangle \end{cases}$$

$$(4)$$

3. Commutation Rules

C. Observables

1. Products of Ladder Operators (Terms)

We define a term(T) as a product of ladder operators that can act on fermionic, antifermionic, and bosonic modes:

$$T = \prod_{m=0}^{M-1} c_m {5}$$

where M is the number of ladder operators in the term and $c_m \in \{b_i, b_i^\dagger, d_i, d_i^\dagger, a_i, a_i^\dagger\}$.

The ladder operators (c_m) can be reordered arbitrarily with the introduction of additional terms due to the commutation rules described in II B 3. In this work, we will have a preference for *normal ordering* of the operators to obey the following structure:

$$T = \left(\prod_{i} (\delta_{b_i^{\dagger}} b_i^{\dagger}) (\delta_{b_i} b_i)\right) \left(\prod_{i} (\delta_{d_i^{\dagger}} d_i^{\dagger}) (\delta_{d_i} d_i)\right) \left(\prod_{i} (\delta_{a_i^{\dagger}} (a_i^{\dagger})^r) (\delta_{a_i} (a_i)^s)\right)$$
(6)

where δ takes the value 0 or 1 to denote if the operator is active in the term and the values r and s are positive integers $\in [0, \Omega]$ and denote the exponent of the bosonic ladder operators acting on that particular bosonic mode.

Describe number/occupation operators acting on fermions/antifermions and bosons and rewrite previous equation for T including these operators.

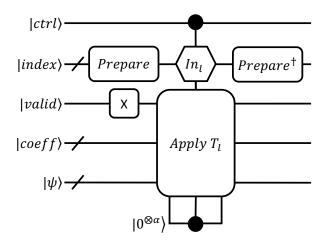


FIG. 1. Ladder Operator Block-Encoding.

2. Linear Combinations of Terms

We can write Hamiltonians (or observables) in the form of linear combinations of terms:

$$H = \sum_{l=0}^{L-1} \alpha_l T_l \tag{7}$$

where L is the total number of terms and α_l is a real-valued coefficient associated with the term T_l .

III. LADDER OPERATOR BLOCK-ENCODING (LOBE)

In this section of the text, we'll do the following:

A. Defining "Block-Encoding"

B. Prior Works

- 1. Linear Combination of Unitaries
- 2. Sparse Block-Encoding of Pairing Hamiltonians

stuff from Liu et al Might ask you to fill this out @Gus

C. Rescaling of Coefficients Due to Bosonic Terms

D. Circuit Construction

In Figure 1, we define the LOBE circuit in terms of generic oracles. Disregarding the (optional) control qubit ($|ctrl\rangle$), the LOBE circuit makes use of 5 qubit registers: $|index\rangle$, $|valid\rangle$, $|coeff\rangle$, $|\psi\rangle$, and $|0^{\otimes \alpha}\rangle$.

The register denoted $|index\rangle$ is referred to as the index register and is used to index the terms in the Hamiltonian as is done in LCU constructions. The integer representations of the computational basis states of the index register corresponds to the indices l in Eq. 7.

The register denoted $|valid\rangle$ consists of a single qubit and is referred to as the validation qubit. It serves the same purpose as in [2] which is to denote whether or not the term at the current index (T_l) will annihilate the quantum state. If the term will annihilate the state, then the validation qubit remains in the $|1\rangle$ state such that the branch of the wavefunction stays outside the

desired subspace of the block-encoding. If the term will not annihilate the state, then the validation qubit gets flipped to the $|0\rangle$ state for the term T_l .

The register denoted $|\mathit{coeff}\rangle$ is referred to as the coefficient register and is used to apply the coefficients associate with the term T_l . These coefficients include both the coefficients of the terms in the linear combination (α_l) as well as the coefficients associated with the bosonic ladder operators. One qubit is used to apply the α_l coefficient while a separate qubit will be required for each bosonic operator (defined as a creation, annihilation, or occupation operator acting with a positive integer-valued exponent) in the term.

The register denoted $|\psi\rangle$ is referred to as the system register and is used to encode the state of the system. This register can be broken up into three subsequent registers: the fermionic system $|\psi_b\rangle$, the antifermionic system $(|\psi_d\rangle)$, and the bosonic system $(|\psi_a\rangle)$. The encoding studied in this work is outlined in more detail in Subsection II A.

Finally, the register denoted $|0^{\otimes \alpha}\rangle$ is referred to as the clean ancillae register. This register includes ancillae qubits that are promised to begin in the $|0\rangle$ state and are returned to this state at the end of the block-encoding circuit.

The LOBE circuit is structured similarly to an LCU circuit in that it involves sequential application of a *Prepare* oracle, a *Select* oracle, and then uncomputing the *Prepare* oracle using the daggered circuit. In addition to these oracles, a single X gate is applied to the $|valid\rangle$ qubit in order to prepare this qubit in the $|1\rangle$ state. This single X gate is *not* repeated on subsequent applications of the block-encoding.

1. Prepare Oracle

The *Prepare* oracle is applied to the index register to prepare a superposition over the computational basis states of this register. **Uniform State Preparation**

The simplest implementation of the *Prepare* oracle is to prepare a uniform superposition over the computational basis states using a series of Hadamard gates on each qubit in the register. The rescaling condition for this implementation is that the coefficients must be rescaled such that they all have magnitude ≤ 1 . Let the largest coefficient magnitude be given by $\alpha^* = \max |\alpha_l|$. Then, this rescaling imposes a constant rescaling factor on the Hamiltonian of $1/\alpha^*$ We refer to these rescaled coeffifients as $\tilde{\alpha}_l$. An additional rescaling factor of $2^{-\lceil \log_2 L \rceil}$ from the uniform superposition results in an overall rescaling of the Hamiltonian by a constant factor of:

$$\lambda_{usp} = \frac{1}{2^{\lceil \log_2 L \rceil} \alpha^*} \tag{8}$$

The quantum resource requirements of this implementation of *Prepare* are negligible as no ancillae are required and only Clifford operations (Hadamards) are performed. However, this implementation of *Prepare* requires that the rescaled coefficients of the terms ($\tilde{\alpha}_l$) be incorporated within the *Select* oracle which will be described in Subsubsection III D 2.

Arbitrary State Preparation

An alternative implementation of *Prepare* is to use the same implementation traditionally used in LCU circuits. The coefficients are first rescaled by their L1 - norm:

$$\lambda_{asp} = \sum_{l=0}^{L-1} |\alpha_l| \tag{9}$$

Then, a state preparation routine is applied that performs the following operation:

$$\left|0^{\otimes \lceil \log_2 L \rceil}\right\rangle \to_{Prepare} \sum_{l=0}^{L-1} \sqrt{\left|\alpha_l\right|/\lambda} \left|l\right\rangle$$
 (10)

This results in a weighted superposition over the computational basis states where the squared amplitudes of the basis states are equal to the magnitude of the associated coefficients.

For Hamiltonians that have structure among the coefficients of the terms, implementations of the *Prepare* oracle can be constructed that leverage this structure. In certain cases, this can drastically reduce the cost of *Prepare* such as is done for the Fermi-Hubbard model in [3]. However, when a certain structure cannot be assumed, the Grover-Rudolph algorithm from [4] gives a formulaic routine to generate these *Prepare* circuits.

Gover-Rudolph requires L-1 rotations, with 1 being uncontrolled and the others controlled. The implementation also requires several series of multiplexing. I think we can make the implementation use fewer Toffolis by doing some smart multiplexing. The total number of Toffolis (if my assumption is correct) is given by:

$$N_{toff} = \lfloor \frac{L}{2} \rfloor \tag{11}$$

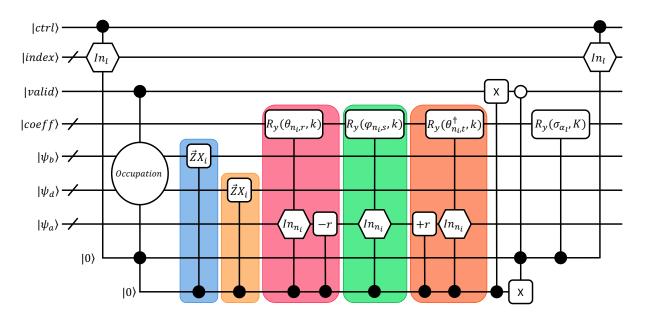


FIG. 2. Ladder Operator Term Oracle.

and the total number of ancillae is given by:

$$Q_{asp} = \lceil \log_2 L \rceil - 2 \tag{12}$$

check this. It should be noted that these ancillae, which will be returned to $|0\rangle$ after the oracle call completes, will likely be reused in the *Select* oracle and therefore will not contribute to the total qubit requirement of the block-encoding.

2. Select Oracle

E. Analytical Cost Analysis

The minimum number of qubits required for the index register is given by:

$$Q_{index} = \lceil \log_2 L \rceil \tag{13}$$

If we let K denote the maximum number of bosonic operators within a single term, then the minimum number of qubits in the coefficient register is given by:

$$Q_{coeff} = K + 1 \tag{14}$$

Using the qubit-efficient encoding described in Subsection II A, the minimum number of qubits required for the system registers is given by:

$$\begin{aligned} Q_{\psi_b} &= \Lambda \\ Q_{\psi_d} &= \Lambda \\ Q_{\psi_a} &= \Lambda \lceil \log_2 \left(\Omega + 1 \right) \rceil \end{aligned} \tag{15}$$

There are numerous space-time tradeoffs that one can make to either reduce the number of ancillae qubits at the cost of more gates or vice versa. In this work, we opt for compilations that minimize the number of non-Clifford operations at the expense of more ancillae qubits. With this choice, the number of clean ancillae required is given by:

$$\alpha = \lceil \log_2 L \rceil + (B+1) + \lceil \log_2 (\Omega+1) \rceil \tag{16}$$

where $\lceil \log_2 L \rceil$ qubits are used for multiplexing over the index register, $\lceil \log_2 (\Omega + 1) \rceil$ qubits are used for multiplexing over the bosonic occupation registers, and B denotes the maximum number of fermionic and antifermionic ladder operators within

a single term. And we assume that any left-elbow with N controls and one ancilla qubit storing the quantum boolean can be decomposed into N-1 left-elbows, each using 2 controls at the expense of N-1 total ancillae.

Therefore, the total qubit requirement under these compilation choices - disregarding the control qubit - is given by:

$$N = Q_{index} + 1 + Q_{coeff} + Q_{\psi_b} + Q_{\psi_d} + Q_{\psi_a} + \alpha$$

= $2\lceil \log_2 L \rceil + 2\Lambda + (\Lambda + 1)\lceil \log_2 (\Omega + 1) \rceil + B + K + 3$ (17)

F. Example

Step-by-step example (intention is to move this to an appendix)

IV. RESULTS / NUMERICAL BENCHMARKING

here we'll "benchmark" (aka numerically compute the cost - number of qubits and gates - of creating a lobe block-encoding) for some systems

ideas of systems to benchmark on:

- Fermi-Hubbard
- Something with just bosons
- Something with fermions, antifermions, and bosons

V. CONCLUSIONS

[1] K. Lamar, Not Like Us, 1st ed. (Interscope Records, 2024).

^[2] D. Liu, W. Du, L. Lin, J. P. Vary, and C. Yang, An efficient quantum circuit for block encoding a pairing hamiltonian, arXiv preprint arXiv:2402.11205 (2024).

^[3] R. Babbush, C. Gidney, D. W. Berry, N. Wiebe, J. McClean, A. Paler, A. Fowler, and H. Neven, Encoding electronic spectra in quantum circuits with linear t complexity, Physical Review X 8, 041015 (2018).

^[4] L. Grover and T. Rudolph, Creating superpositions that correspond to efficiently integrable probability distributions, arXiv preprint quantph/0208112 (2002).