CSE 222: Data Structures and Algorithms

()

Homework #4 Time Complexity Analysis

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Problem 1: checkIfValidUsername

(? points)

```
static boolean checkIfValidUsername(String username) {
       // checks string validation and lenght
         if(username = null \mid | username.length() = 0)  {
a .
           System.out.println("the username is invalid, it should be at least 1 character
           return false;
       // checks first char is letter or not
         if(isLetter(username.charAt(0))) {
b.
           // if last letter valid string is valid
           if(username.length() == 1)
           return true;
           // checks other letters
           return checkIfValidUsername(username.substring(1));
       System.out.println("the username is invalid, it shouldcontains only letters");
       return false;
   }
```

- (a) {This part have basic compression and println, there is no loop all parts have O(1) time complexity }
- (b) {There is no loop in this part is Letter(O(1), charat(O(1)) .lenght()(O(1)) and a recursive call with substring function that O(N).)}

(Conc.) { This method recursively checks the rest of the characters with creating substrings and have O(N) per step. The recursion continues until either the end of the string is reached that's why this make the method $O(N^2)$ timecomplexity)}

(Solution)

 $check If Valid Username\ have\ O(N^2) time complexity (length to f the username)$

Problem 2: isBalancedPassword

(? points)

```
static boolean isBalancedPassword(String password1) {
                                Stack < Character > stack = new Stack < >();
  a .
                          boolean isAnyBracket = false, isAnyLetter = false;
                               for (char c : password1.toCharArray()) {
  b.
                                       if (c == '(' | c == '{' | c == 
                                                    stack.push(c);
                                                    isAnyBracket = true;
                                       \{ else\ if\ (c == ')' \mid |c == '\}' \mid |c == ']' \}
                                                    if (stack.isEmpty()) {
                                                                 System.out.println("the password1 is inv
                                                                 return false;
                                                    char top = stack.pop();
                                                    if ((c == ')' && top != '(') || (c == '}' &&
                                                                 System.out.println("the password1 is inv
                                                                 return false;
                                       } else if(isLetter(c)) {
                                                    isAnyLetter = true;
                                           else {
                                                    System.out.println("the password1 is invalid
                                                    return false;
                                       }
                                 if (!isAnyBracket)
C .
                                       System.out.println("the passwordl is invalid, it
                          if (!isAnyLetter)
                                       System.out.println("the password1 is invalid, it
                          if (! stack.isEmpty())
                                       System.out.println("the passwordl is invalid, it
                          return (stack.isEmpty() && isAnyBracket && isAnyLett
             }
(a) {Initialization of Stack and boolean variables has a time complexity of O(1).
```

(b) { for loop iterates over each character of the password1.toCharArray(), so

the time complexity of the loop is O(N), where N is the length of the password. and to CharArray method have O(N) time complexity that calls once $\}$

(b.inloop) {In the loop there are push, pop, isLetter, isEmpty and println methods all of these methods have O(1) time complexity in worst case }

(c) {There is only compressions and println function that have O(1) time complexity }

(Conc.) { Therefore, the time complexity of the isBalancedPassword function is O(N), where N is the length of the password.)}

(Solution) is BalancedPassword function is O(N) (length of the password.)

Problem 3: containsUserNameSpirit

(? points)

```
static\ boolean\ contains User Name Spirit (String\ username,\ String\ username)
           Stack < Character > stack = new Stack < Character > ();
a .
         if(password1.length() < 8) {
             System.out.println("the passwordl is invalid, it
        }
         // Push each character of the password onto the stac
          for (char c : password1.toCharArray()) {
b .
           stack.push(lowerCase(c));
        }
          for (char c: username.toCharArray()) {
 c .
             if (stack.contains(lowerCase(c)))
                  return true;
        }
           System.out.println("the passwordl is invalid, it s
 d.
         return false:
```

- (a) {Initialization of stack, compression in if and println have a time complexity of O(1). }
- (b) { for loop iterates over each character of the password1.toCharArray() and pushes to stack, so the time complexity of the loop is O(N). toCharArray method have O(N) time complexity that calls once, and push function have O(1). At the end for loop has O(N) time complexity, where N is the length of password1. }
- (c) { for loop iterates over each character of the password1.toCharArray() and checks stacks for it's lowercase version, so the time complexity of the loop is O(N*M). toCharArray method have O(M) time complexity that calls once, lowercase function have O(1), and contains method has O(N). At the end for loop has O(N*M) time complexity, where N is the length of password1 and M is length of username. }
 - (d) {There is only println function that have O(1) time complexity }
- (Conc.) { In the worst case, if both the length of password1 and the length of the username are equal to the maximum possible length, then the time complexity

 $will\ be\ O(N^2), where N is the maximum length. Otherwise O(N*M) time complexity, where N is the length of password 1 and M is length of username.\}$

(Solution)

contains UserNameSpirit function is O(N*M) (length of the password and username)

Problem 4: isPalindromePossible

(? points)

```
static boolean is Palindrome Possible (String password1) {
           StringBuilder output = new StringBuilder();
a .
         String noBPass;
        for (char c : password1.toCharArray()) {
    if (c != '(' && c != ')' && c != '[' && c != ']'
                 output.append(c);
        }
           Boolean[] table = new Boolean[output.length() + 1]
b.
        for(int \ i = 0; \ i < output.length(); ++i) {
             table[i] = false;
        table[output.length()] = (output.length() \% 2 == 0)
        noBPass = output.toString();
           if (!recursivePolindrome(noBPass, table, 0)) {
c .
             System.out.println("the passwordl is invalid, it
             return false;
         return true;
    }
    private static boolean recursivePolindrome(String passwo
d.
           if(index == password1.length())
         return true;
         if(table[index] == true)
             return recursivePolindrome(password1, table, ind
        char \ c = password1.charAt(index);
          for(int \ i = index + 1; \ i < password1.length(); ++i
e .
             if (c == password1.charAt(i) && !table[i]) {
                 table[i] = true;
                 return recursivePolindrome(password1, table,
        }
```

- (a) $\{string\ builder\ initialization\ have\ O(N)\}$
- (b) { As we analyze in previous questions password1.toCharArray() have O(N) call one,Array initialization has O(N)(default value assignment) or O(1) depend on JVM, for have O(N), inside of for there is only assgiment that is O(1). }
- (c) $\{A \text{ recursive call that we will be analyez in the down, and println that have } O(1)\}$
 - (d) {compression O(1) and continue recursion call }
- (e) {have a loop iterates until find same character and do oen more recursion call that call have
- $O(N^2) inworst case time complexity, and O(N) best case time comlexity, where N is length to find the complexity of th$
- $\textit{(f)} \{It is possible towork one time for each check and continues the recursive vall if work\}$

(Conc.) { In the worst case, have

 $O(N^2)$ time complexity, where N is the length of password 1 and in best case have O(N) time con

(Solution)

 $is Palindrome Possible function \ is \\ O(N^2) worstcase, O(N) bestcase (length of the password)$

Problem 5: isExactDivision

(? points)

```
static boolean is Exact Division (int password2,
[]
    denominations) {
         // Check if password between 10 and 10000
           if(password2 < 10 \mid | password2 > 10000)  {
 a .
             System.out.println("the password2 is invalid, it
             return false;
        }
        // Checks if password2 obtainable by summing denoming
           if (!isExactDivisionRecursive(password2, denominat
  b.
             System.out.println("the password1 is invalid, it
             return false;
         return true;
   private static boolean is Exact Division Recursive (int
            int [] denominations, int currentsum, int inde
password2,
           if(index != -1)
 C .
             currentsum += denominations[index];
         if (currentsum == password2)
             return true;
         if (currentsum > password2)
             return false;
 d.
          for (int i = 0; i < denominations.length; ++i) {
             if (is Exact Division Recursive (password 2, denominat
                 return true;
         return false;
        (a) {Basic compression and print have O(1)}
 (b) { A recursive call that we will be analyze in the down belown, and println
               that have O(1)
```

(**d**) {Each re-

(c) { Basic compressions sum and assignment all of them have O(1)}

cursive call results in another loop that also executes n times, leading to a total of n^n iterations. This is because each recursive call iterates through the entire denominations and the sum of the sum

(Conc.) { In the worst case, have $O(N^N)$ timecomplexity, where N is the length of demonitaions and in best case have O(1) times

(Solution)

 $is Palindrome Possible \ function \ is \\ O(N^N) worst case, O(1) best case (length of the denominations)$