

-07-'19

UNIT-I: EMBEDDED COMPUTING

- Introduction

- Complex Systems and Microprocessors

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- Complex Algorithms
- User Interfaces (GPS), Complex Interface
- Real Time (deadline)
- Multivariate
- Manufacturing Cost
- Power and Energy

Characteristics

we need to tackle care
of deadlines, improve
performance

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- How much h/w do we need?
- How do we meet deadlines?
- How do we minimize power consumption?
- How do we design for upgradability?
- Complex Testing?
- Limited observability and controllability?
- Restricted development environment?

Challenges

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Performance

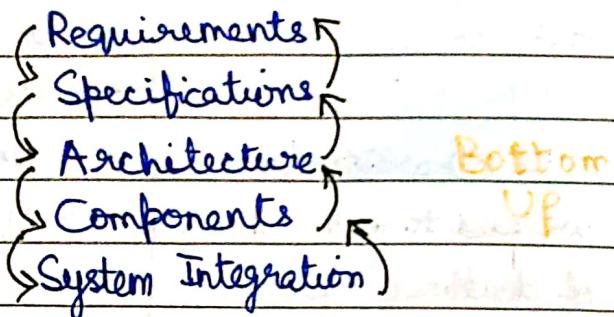
- CPU
- Platform
- Program
- Task
- Multiprocessor

- The Embedded System Design Process

Normal Purpose - General Purpose

Embedded → which includes a microprocessor
but not used for general purpose, i.e.,
specific purpose

ES Design Process



- We will do all the programs using microcontrollers.
- A microcontroller is application specific: General Purpose, Only for one purpose, Only one operation
- A microprocessor is general purpose whereas a microcontroller is in built, one system contains all components.
- First Microprocessor - Intel 4004
- Calculator Operation - Calculator Application
- In a microprocessor, all connections are external.
- When we are upgrading something, we need to think about cost vs performance.
- We need to generate embedded system in such a way that programming should take less power and energy.
- Why do we use microprocessor for Embedded System?
This is because any app can be run or implemented on microprocessor.
- Here: Components = Component Design

Challenges:

i) ES comes with hardware

- Optimize the hardware, in case also, we use ES (embedded system), many vehicles, application specific.

ii) Real Time

- Cost vs Performance

iii) Eg: Battery in Mobile Phones

Before it was very large, now it is small and large utilization

★ iv) ES installed software - we are using it

with little changes. We need to make

upgradability. It should execute a new program when it is upgraded.

v) } Write program on common desktop,

vi) } Same Concept

vii) } compile and execute

How to improve performance of ES?

- It depends on processor (CPU)
- Depends on wires, buses, (platform), connection media
- Divide program into smaller one (program)
- Multitasking (Single processor doing many tasks)
- Multiprocessor (it contains many processors)

In ES, Hexcode is not written manually.

- We will be writing C and C++ programs
- Hexfile will go to ES. Before dumping, we need to check all the possibilities when doing in ES (it is also one of the challenges).
- Testing on desktop.
- We cannot change it in ES, so before dumping, we only need to check. It is one of the disadvantages.
- In ES, we can read hardly 30 characters, but in other processors, it can read all characters, i.e., 1550.



ES Design Process

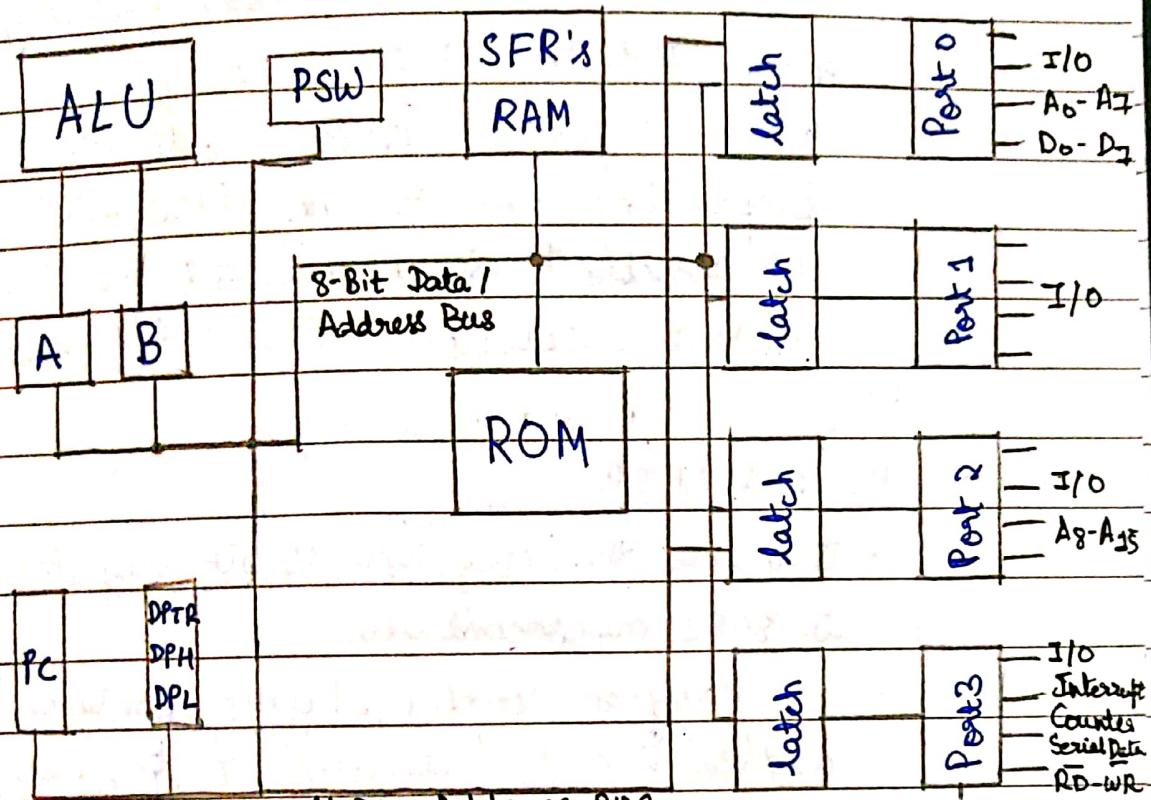
(5marks, explain with applicat)

- i) It is a 5 step process
It basically uses top-down approach.
- I/O, O/P
Cost
- Performance of System
- ii) Specifications
- iii) Architecture - what it does
Organization - How it does
(I/O, O/P, how many different hardware, how many sensors)
- iv) Each and every component
(which component for what purpose)
- v) Combining via System Integration

EA
ALE
PSEN
XTAL1
XTAL2
RESET
Vec
GND

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8051 Architecture

16 Bit Address BUS
Internal RAM Structure

EA	System	Bit / Byte Addresses	SFR's	
ALE	Timing		IE	
PSEN	System		IP	
XTAL1	Interrupt		PCON	
XTAL2	Timers		SBUF	
RESET	Data Buffer	Register	SCON	
Vcc	Memory	Block 3 Bank	ICON	
GND	Control	Registers	IMOP	
		Block 2 Bank	TLO	
			THO	
			TL1	Port 3
			TH1	Memory
		Registers		Control
		Block 1 Bank		
		Registers		
		Block 0 Bank		

→ 8051 Oscillator, Clock

- 8051 consists of 2 internal oscillators and 3 external oscillators, whose frequency ranges from 1MHz to 16MHz.

- We fix crystal frequency to 12MHz, to use proper time; i.e., to make time 1μs. $\left[\frac{16\text{MHz}}{12\text{MHz}} < 1\mu\text{s} \right]$
- Not possible to synchronize all the tasks without oscillator. } Time $\propto 1/\text{frequency}$

→ PC and DPTR

- They are the only two 16 bit registers available in 8051 microcontroller.
- PC: Program counter always contains 16-bit address of next instruction to be executed.
- DPTR: Data Pointer is used to access external addresses (16bit address = 2×8 bit (DPH + DPL)) of external devices attached to our device.

→ A and B CPU registers

- All operations performed such as arithmetic and logical operations are performed by default we use 'A' and 'B' registers and also store data in them. They are CPU Registers
- Data transfer between external and internal devices are done using 'A' register
- A = 8 bit register, B = 8-bit register

* 8 bit general purpose registers $\times 4$

* $\alpha \times 16$ bit registers : PC, DPTR

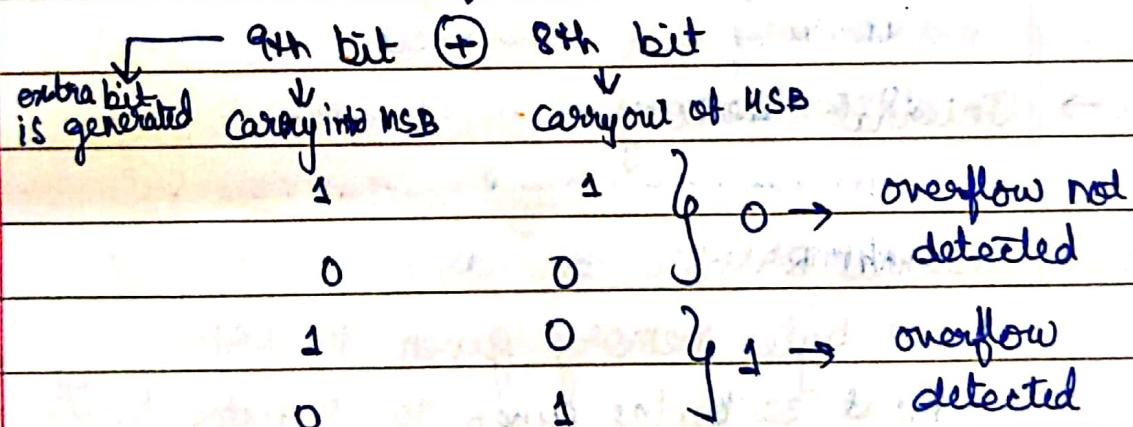
→ Flags and PSW

- Seven flags are present.
 - Four special purpose flags, modified when arithmetic and logical operations are performed
- i) Carry (cy) - Set when extra bit during addition, else reset
 - ii) Auxiliary Carry (Ac) - Used in BCD addition, set when carry is transferred from lower nibble to higher level, otherwise reset
 - iii) Overflow (ov) - Cannot store 256 in 8 bit register then overflow ov is set, else reset.

Overflow is detected with: 9bit carry (carry into MSB)

XOR 8th bit carry (carry out of MSB)

$$= 1 \xrightarrow{\text{set}} \text{overflow detected} \quad = 0 \xrightarrow{\text{not set}} \text{no overflow detected}$$



iv) Parity (P) - Always odd parity

$$\text{No. of 1s in 8 bits} + \text{Parity bit} = \text{Odd no}$$

↳ Parity Set

$$\text{No. of 1s in 8 bits} + \text{Parity bit} = \text{Even no}$$

↳ Parity Reset

- Three special purpose flags, can be access only by programmer, not by operations

- i) GFO (General Flag 0)
- ii) GF 1 (General Flag 1)
- iii) FO (Flag 0)

- PSW is a special purpose register.

Applications:

- To check the status of flags.
- To modify the status of flags.
- To check which registers are under use.

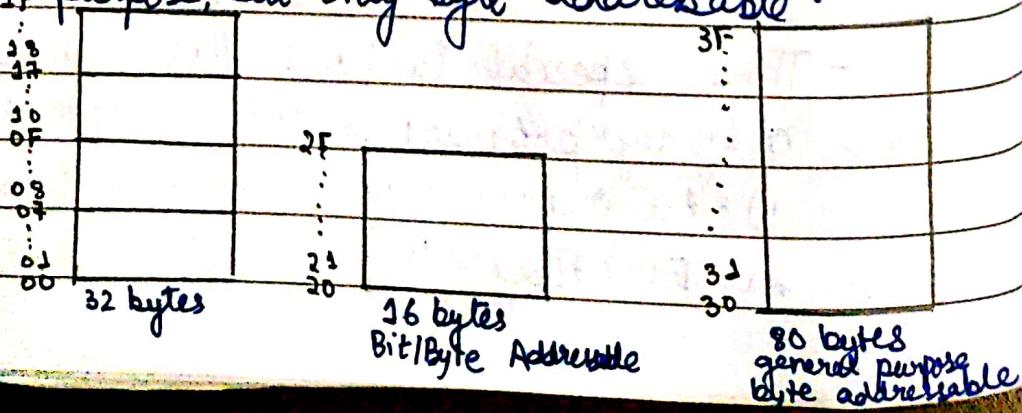


RS0	RS1	Registers Bank
0	0	R.B. 0
0	1	R.B. 1
1	0	R.B. 2
1	1	R.B. 3

→ Internal Memory

Internal RAM:

- 128 bytes memory given to RAM.
- First 32 bytes given to register bank's, each consists of 8 bits (RBO + RB1 + RB2 + RB3 = 8 + 8 + 8 + 8) ($4 \times 8 = 32$)
- After that, 16 bytes memory are accessed by bit addresses / byte addresses.
- Remaining 80 bytes are used for general purpose, but only byte addressable.



-Stack Memory: From internal memory, some part of RAM is used for stack pointer, increment the stack pointer by 1 if data is inserted, and if you decrement stack pointer by 1, data will be deleted.

- More priority will be for stack, it has capacity to override data, so care must be taken to give initial address to stack.
- Stack is 468 bit.

Special Function Registers (SFR's)

- They will have special names as well as special addresses also.
- Among the special function registers, a few are bit addressable also (can access point by point), others are byte addressable only (can access entire byte).
- Bit Addressable are also byte addressable, but not vice versa.

ROM

- ROM is of two types:
 - i) Internal ROM
 - ii) External ROM

* Internal Memory

- * Internal RAM
- * Stack and SP
- * Special Function Registers (SFR's)
- * Internal ROM

→ 8051 Microcontroller Architecture - 40 Pin Diagram

Port 1 Bit 0	0	RES	40	+5V
Port 1 Bit 1	1	AD0	39	Port 0 Bit 0
Port 1 Bit 2	2	AD1	38	Port 0 Bit 1
Port 1 Bit 3	3	AD2	37	Port 0 Bit 2
Port 1 Bit 4	4	AD3	36	Port 0 Bit 3
Port 1 Bit 5	5	AD4	35	Port 0 Bit 4
Port 1 Bit 6	6	AD5	34	Port 0 Bit 5
Port 1 Bit 7	7	AD6	33	Port 0 Bit 6
Reset Input	8	AD7	32	Port 0 Bit 7
Port 3 Bit 0	10	RxD	(Vpp)	EA
Port 3 Bit 1	11	TxD	(PROG)	ALE
Port 3 Bit 2	12	INT0		PSEN
Port 3 Bit 3	13	INT1		Port 2 Bit 1
Port 3 Bit 4	14	To		Port 2 Bit 6
Port 3 Bit 5	15	T1		Port 2 Bit 5
Port 3 Bit 6	16	WR		Port 2 Bit 4
Port 3 Bit 7	17	RD		Port 2 Bit 3
Crystal Input 1	18	XTAL2		Port 2 Bit 2
Crystal Input 2	19	XTAL1		Port 2 Bit 1
	20	GND		Port 2 Bit 0

- Port 0, Port 1, Port 2, Port 3 - All pins can be used as I/O pins or O/P pins.
- Port 0, Port 2 can be used for dual purpose, i.e., accessing address and accessing data.
- Port 0 is used to access lower byte address ($A_0 - A_7$)
- Port 2 is used to access higher byte address ($A_8 - A_{15}$)

- Port 3 has control pins Read, Write, Define Data, T transmit Data.

- RD = 0 → Performing Read Operation
- WR = 0 → Performing Write Operation
- $8 \times 4 = 32$ Port Pins
- Other Pins :
 - V_{cc} - Ground Control
 - V_{ss} - Frequency / Voltage
 - Crystal Input 1 - } To provide frequency
 - Crystal Input 2 - } to device
 - Reset Input - To reset the microcontroller
 - EA } Access external RAM, ROM
 - ALE }
 - PSEN }

- Addresses of Special Function Registers

A - 0E0	P0 - 80	PCON - 87	TCON - 88
B - 0F0	P1 - 90	PSW - 0D0	TLO - 8A
DPH - 83	P2 - A0	SCON - 98	TH0 - 8C
DPL - 82	P3 - 0B0	SBUF - 99	TL1 - 8B
IE - 0A8		SP - 81	TH1 - 8D
IP - 0B8		TMOD - 89	

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→ Counters and Timers

→ TCON

TF1	TR1	TF0	TR0	IE1	IT1	IE0	IT0
7	6	5	4	3	2	1	0
001B	000B			0013 P3.3		0003 P3.3	
				INT1		INT0	

- Timers are used to calculate / count internal clock pulses.

- Clocks are used to calculate / count external clock pulses.

- Processor can also do all these, but if a processor has to do its own tasks and also counting, burden on processor is more.

- So, to reduce burden on processors, two hardware chips, externally added 16 bit registers, timers and counters are used.

- They are used as 2x8 bit registers

$$TL0 + TH0 = \text{Timer 0}$$

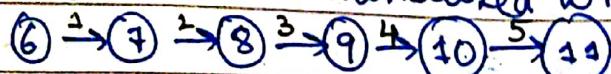
$$TL1 + TH1 = \text{Timer 1}$$

L=Lower, H=Higher

- To find 'n' number of clock pulses, timer must be initialized with max-n+1 value.

Eg: Maximum = 10, we want to count 5 events,

timer must be initialized with $10-5+1=6$



When overflow occurs, flag is set.

When count comes to 0000, flag is set.

- We use '1' in max- $n+1$ because when the counter changes from FF to 00, the timer/counter counts that also, and we use that as interrupt.
- Whenever timer becomes 00, we use that as interrupt and can perform a particular task.
Eg. 60 sec green light in traffic signals.
- Two special function registers
Timer Control Register (TCON) and
Timer Mode Register (TMOD)
two 16 bit register control two timers.

TCON:

i) Timer Flag 1 (TF1)	TCON is bit addressable,
ii) Timer Run 1 (TR1)	can change
iii) Timer Flag 0 (TFO)	bit by bit
iv) Timer Run 0 (TR0)	in programming
v) Interrupt Edge 1 (IE1)	and also
vi) Interrupt Signal Timer (IT1)	byte addressable
vii) Interrupt Edge 0 (IE0)	
viii) Interrupt Signal Timer (IT0)	

* Timer 0 to count 'n', whenever Timer 0 moves from FF to 00, i.e., it overflows, TFO becomes 1. Indicates counting completed.

* Timer 1 to count 'n', whenever Timer 1 moves from FF to 00, i.e., it overflows, TF1 becomes 1. Indicates counting completed.

* Both are reset / become 0 whenever microprocessor executes an interrupt service routine (whenever interrupts occur, automatically jump to particular address to execute task) present in address 001B (TF1) and 000B (TF0).

- * TR1 is going to be set by programmer if we want to calculate internal clock pulses, i.e., used for timing purpose.
- * TR1=0 indicates we are not using timing, T1 does not work as timer.

- * TR0 is going to be set by programmer if we want to calculate internal clock pulses, i.e., using T0 as timer, used for timing purpose.
- * TR0=0 indicates we are not using timing, T0 does not work as timer.

* PCON Register

Higher Nibble = Timer Flags
Lower Nibble = Interrupt Flags

- * IE1 - Whenever on P3.3 pin high to low (\overline{L}) pin encountered, i.e., INT1 pin of port 3 encounters change of high to low edge, it indicates external interrupt occurred, and IE1 is set, otherwise reset.

* IED - Whenever on P3.3 pin high to low (\overline{L}) pin edge is encountered; i.e., INT0 pin of Port³ encounters change of high to low edge, it indicates external interrupt has occurred, IED is set.

* IES=0 when process jumps to execute interrupt service routine at address 0013.

* IEO=0 when process jumps to execute interrupt service routine at address 0003.

* IT1 and IT0 are internal interrupts. User can set or reset them based on the task to be performed.

TMOD:

Gate	C/T	M1	M0	Gate	C/T	M1	M0
7	6	5	4	3	2	1	0

- TMOD has two sets of four flags each, one set for Timer0 and another set for Timer1.

- The flags are duplicate ones. We can write only 4 bits.

- The TMOD is byte addressable, not bit addressable.

Flags:

- i) C/T : Counter/Timer
- ii, iii) M0, M1 - Mode Selection Bits
- iv) Gate - Gate Bit

* $C1\bar{T} = \text{Counter 1 Timer}$ if used to know whether $C1\bar{T} = 0$ or used as timer (counter).

$C1\bar{T} = 0 \rightarrow \text{Used as a timer}$

$\hookrightarrow \text{Bit } 6 = 0 \Rightarrow \text{Timer 2 used as timer}$

$\hookrightarrow \text{Bit } 2 = 0 \Rightarrow \text{Timer 0 used as timer}$

$C1\bar{T} = 1 \rightarrow \text{Used as a counter}$

$\hookrightarrow \text{Bit } 6 = 1 \Rightarrow \text{Timer 1 used as clock}$

$\hookrightarrow \text{Bit } 2 = 1 \Rightarrow \text{Timer 1 used as clock}$

* M1 and M0 \rightarrow Mode Selection bits

Used when we use timers

Important when $C1\bar{T} = 0$, does not matter when $C1\bar{T} = 1$.

M1 M0 Timer Mode

0	0	0
---	---	---

0	1	1
---	---	---

1	0	2
---	---	---

1	1	3
---	---	---

• Timer Mode

* Gate Bit

- This device acts as a timer when we encounter conditions.

- When $T0/T1$ works in timer mode, i.e., used to work internal clock pulses:
 - i) Timer 1/0 bit in TCON register must be 1 AND

- ii)
 - a) Either gate bit in TMOD = 0 or

- b) Interrupt 1/0 must be 1.