

Question 1

1. $P(X \geq a) \leq \mathbb{E}(X)/a = 1/a$.

2. 不等式在 $a \in (0, 1)$ 时显然成立. 设 $a > 1$.

$$P(X \geq a) = P(X - \mathbb{E}(X) \geq (a - 1)\sigma(X)) \leq 1/(a - 1)^2.$$

3. 考虑计算 $\mathbb{E}(e^{tX})$. 下面的计算要求 $t < 1$.

$$\begin{aligned} \mathbb{E}(e^{tX}) &= \int_0^{+\infty} e^{-x} \cdot e^{tx} dx \\ &= \int_0^{+\infty} e^{-(1-t)x} dx \\ &= \frac{1}{t-1} e^{-(1-t)x} \Big|_0^{+\infty} = \frac{1}{1-t} \\ &= \sum_{k=0}^{\infty} t^k. \end{aligned}$$

于是 $\mathbb{E}(X^k) = k!$, $P(X \geq a) = P(X^k \geq a^k) \leq \frac{\mathbb{E}(X^k)}{a^k} = k!/a^k$.

4. 要求 $\min_t \frac{1}{1-t} e^{-ta}$. 记该函数为 $f(t)$.

$$\begin{aligned} f'(t) &= \frac{1}{(1-t)^2} e^{-ta} + \frac{1}{1-t} \cdot (-a)e^{-ta} \\ &= \frac{1}{1-t} e^{-ta} \cdot \left(\frac{1}{1-t} - a \right). \end{aligned}$$

于是取 $t = 1 - a^{-1}$ 即得 $P(X \geq a) \leq ae^{1-a}$.

5. 所求为

$$P(X \geq a) = \int_a^{+\infty} e^{-x} dx = e^{-a}.$$

6. 令 $K \sim \pi(ta)$, $f(k) = \mathbb{E}(X^k)/a^k$.

$$\begin{aligned} \mathbb{E}(f(K)) &= \sum_{k=0}^{+\infty} \frac{(ta)^k}{k!} e^{-ta} \cdot \frac{k!}{a^k} \\ &= e^{-ta} \sum_{k=0}^{+\infty} t^k \\ &= \frac{e^{-ta}}{1-t}. \end{aligned}$$

取 $t = 1 - a^{-1}$ 即有 $\mathbb{E}(f(K)) = \lambda_{\text{Chernoff}}$. 于是矩函数给出的界较 Chernoff 更紧.

Question 2

1. 对于 $\varepsilon < 1/2$, 有

$$\lim_{n \rightarrow \infty} P(|X_n - X| < \varepsilon) = \lim_{n \rightarrow \infty} \frac{1}{n} = 0,$$

于是 $\{X_n\}$ 依概率收敛到 0. 相对的

$$P\left(\bigcup_{m=n}^{\infty} P(|X_m - X| < \varepsilon)\right) = \sum_{m=n}^{\infty} \frac{1}{m}$$

求和发散, 于是 $\{X_n\}$ 不几乎必然收敛到 0.

2. 利用 Chernoff-Hoeffding 不等式:

$$\begin{aligned} P(|Y_m - p| > \varepsilon) &= P\left(\left|\sum_{i=1}^m X_i - mp\right| > m\varepsilon\right) \\ &\leq 2 \cdot \exp\left(-\frac{2m^2\varepsilon^2}{m}\right) \\ &= 2 \exp(-2m\varepsilon^2). \end{aligned}$$

于是

$$\begin{aligned} P\left(\bigcup_{m=n}^{\infty} |Y_m - p| > \varepsilon\right) &\leq 2 \sum_{m=n}^{+\infty} \exp(-2m\varepsilon^2) \\ &= 2 \exp(-2n\varepsilon^2) \cdot \frac{1}{1 - \exp(-2\varepsilon^2)} \rightarrow 0. \end{aligned}$$

Question 3

观察到若 A 输出“合数”, 则输入为合数; 若 B 输出“质数”, 则输入为质数.

考虑如下算法: 循环执行 AB , 直至某一程序产生如上输出. 下面证明这一算法运行时间期望为 $O(T)$.

由对称性不妨设输入为质数, 则 C 在 B 输出“质数”时停止. 于是

$$\mathbb{E}(T(C)) \leq 2\mathbb{E}\left(G\left(\frac{1}{2}\right)\right) \cdot O(T) = O(T).$$

Question 4

1. 有

$$\begin{aligned} M_X(t) &= (M_{B(1,p)}(t))^n = (pe^t + 1 - p)^n \\ &= \exp(n \cdot \ln(pe^t + 1 - p)) \\ &\leq \exp(n \cdot (pe^t - p)) \\ &= \exp(\mathbb{E}(X) \cdot (e^t - 1)). \end{aligned}$$

2. 对任意 $\varepsilon > 0$,

$$\begin{aligned} P(X \geq (1 + \varepsilon)\mathbb{E}(X)) &\leq M_X(t) \cdot \exp(-t \cdot (1 + \varepsilon)\mathbb{E}(X)) \\ &\leq \exp(\mathbb{E}(X) \cdot (e^t - 1 - t(1 + \varepsilon))). \end{aligned}$$

对 $f(t) := e^t - 1 - t(1 + \varepsilon)$ 求导:

$$f'(t) = e^t - (1 + \varepsilon),$$

取 $t = \ln(1 + \varepsilon)$ 即有

$$\begin{aligned}
P(X \geq (1 + \varepsilon)\mathbb{E}(X)) &\leq \exp(\mathbb{E}(X) \cdot (\varepsilon - \ln(1 + \varepsilon)(1 + \varepsilon))) \\
&= \left(\frac{e^\varepsilon}{(1 + \varepsilon)^{1+\varepsilon}} \right)^{\mathbb{E}(X)}.
\end{aligned}$$

对任意 $0 < \varepsilon < 1$,

$$\begin{aligned}
P(X \leq (1 - \varepsilon)\mathbb{E}(X)) &\leq P(e^{-tX} \geq e^{-t(1-\varepsilon)\mathbb{E}(X)}) \\
&\leq M_X(-t) \cdot \exp(t(1 - \varepsilon)\mathbb{E}(X)) \\
&= \exp(\mathbb{E}(X) \cdot (e^{-t} - 1 + t(1 - \varepsilon))).
\end{aligned}$$

对 $g(t) := e^{-t} - 1 + t(1 - \varepsilon)$ 求导:

$$g'(t) = -e^{-t} + 1 - \varepsilon,$$

取 $t = -\ln(1 - \varepsilon)$ 即有

$$\begin{aligned}
P(X \leq (1 - \varepsilon)\mathbb{E}(X)) &\leq \exp(\mathbb{E}(X) \cdot (1 - \varepsilon - 1 - \ln(1 - \varepsilon)(1 - \varepsilon))) \\
&= \exp(\mathbb{E}(X) \cdot (-\varepsilon - \ln(1 - \varepsilon)(1 - \varepsilon))) \\
&= \left(\frac{e^{-\varepsilon}}{(1 - \varepsilon)^{1-\varepsilon}} \right)^{\mathbb{E}(X)}.
\end{aligned}$$

3. 由 Union Bound 只需证明 $X \sim B(n, 1/n)$ 时 $P(X \geq c_1 \ln n / \ln \ln n) \leq 1/n^2$.

$$\begin{aligned}
\ln P(X \geq c_1 \ln n / \ln \ln n) &\leq \ln \left(\frac{e^{\frac{c_1 \ln n}{\ln \ln n} - 1}}{\left(\frac{c_1 \ln n}{\ln \ln n} \right)^{\frac{c_1 \ln n}{\ln \ln n}}} \right) \\
&= \frac{c_1 \ln n}{\ln \ln n} - 1 - \left(\frac{c_1 \ln n}{\ln \ln n} \right) (\ln c_1 + \ln \ln n - \ln \ln \ln n) \\
&\leq \left[\frac{c_1 \ln n}{\ln \ln n} - \frac{c_1 \ln c_1 \ln n}{\ln \ln n} \right] - c_1 \ln n + \frac{c_1 \ln n \ln \ln \ln n}{\ln \ln n} \\
&= \ln n \cdot \left(\underbrace{\frac{c_1(1 - \ln c_1)}{\ln \ln n}}_{A} - c_1 + \underbrace{\frac{c_1 \ln \ln \ln n}{\ln \ln n}}_{B} \right)
\end{aligned}$$

其中 $A, B \xrightarrow{n \rightarrow \infty} 0$. 于是当 n 充分大时 $\ln P(X \geq c_1 \ln n / \ln \ln n) \leq -c_1 \ln n$, 取 $c_1 = 3$ 即证.

4. 有

$$\begin{aligned}
\mathbb{E}(Y) &\leq c_1 \ln n / \ln \ln n + n \cdot P(Y \geq c_1 \ln n / \ln \ln n) \\
&\leq c_1 \ln n / \ln \ln n + 1 \\
&\leq (c_1 + 1) \ln n / \ln \ln n.
\end{aligned}$$

Question 5

1. 考虑由 $\{y_i\}, \{-y_i\}, \{\mathbf{0}\}$ 组成的 $\{x_i\}$, 则不等式中的向量是 $\{x_i - x_j\}$ 的子集, 结论由题设成立.
2. 假定 (1) 中全部命题成立, 则有

$$\begin{aligned}
\left\langle \frac{1}{\sqrt{k}}Ay_i, \frac{1}{\sqrt{k}}Ay_j \right\rangle &= \frac{1}{2} \left[\left\| \frac{1}{\sqrt{k}}A(y_i + y_j) \right\|^2 - \left\| \frac{1}{\sqrt{k}}Ay_i \right\|^2 - \left\| \frac{1}{\sqrt{k}}Ay_j \right\|^2 \right] \\
&\leq \frac{1}{2} \left[\left(1 + \frac{\varepsilon}{4}\right) \|y_i + y_j\|^2 - \left(1 - \frac{\varepsilon}{4}\right) \|y_i\|^2 - \left(1 - \frac{\varepsilon}{4}\right) \|y_j\|^2 \right] \\
&= \left(1 + \frac{\varepsilon}{4}\right) \langle y_i, y_j \rangle + \frac{\varepsilon}{4} (\|y_i\|^2 + \|y_j\|^2),
\end{aligned}$$

即

$$\left| \left\langle \frac{1}{\sqrt{k}}Ay_i, \frac{1}{\sqrt{k}}Ay_j \right\rangle - \langle y_i, y_j \rangle \right| \leq \frac{\varepsilon}{2} + \frac{\varepsilon}{4} < \varepsilon.$$

3. 将 x_i 正则化为 $\hat{x}_i = x_i / \|x_i\|$ 即可. $x_i = \mathbf{0}$ 时显然.
4. 考虑 \mathbb{R}^n 的单位正交基 e_1, \dots, e_n , 构造如上随机矩阵 $B \in M_{k \times n}$, $k = O(\log n / \varepsilon^2)$, ε 取 0.1. 利用上一问结论有

$$\left| \left\langle \frac{1}{\sqrt{k}}Be_i, \frac{1}{\sqrt{k}}Be_j \right\rangle - \langle e_i, e_j \rangle \right| \leq \varepsilon.$$

构造矩阵 $A := \frac{1}{k}B^\top B$, $A_{ij} - I_{ij}$ 即满足如上不等式, 于是 A 符合要求.

Question 6

1. 显然 Y_i 的密度函数为 $g(x) = \frac{2}{\pi(x^2+1)}$, $x > 0$.
2. 考虑对 Y_i 进行截断. 令 $Y_i^{(t)} := \min(Y_i, t)$, $Z_i := Y_i - Y_i^{(t)} = \max(Y_i - t, 0)$, 则

$$P\left(\sum Y_i > c_1 n^2\right) \leq P\left(\sum Z_i > 0\right) + P\left(\sum Y_i^{(t)} > c_1 n^2\right).$$

- 先证 $P(\sum Z_i > 0) < 1/6$. 由 Union Bound 只需要保证 $P(Z_i > 0) < 1/6n$. 已知

$$P(Z_i > 0) = P(Y_i > t) = \frac{2}{\pi} \left(\frac{\pi}{2} - \arctan t \right) = 1 - \frac{2}{\pi} \arctan t,$$

洛必达法则告诉我们 $t \rightarrow +\infty$ 时 $\pi/2 - \arctan t$ 和 $1/t$ 同阶, 于是取 $t = O(n)$ 即可满足要求.

- 接下来证 $P\left(\sum Y_i^{(t)} > c_1 n^2\right) < 1/6$. 有

$$\begin{aligned}
\mathbb{E}(Y_i) &= \int_0^t x \cdot \frac{2}{\pi(x^2+1)} dx + \left(1 - \frac{2}{\pi} \arctan t\right) \cdot t \\
&= \frac{2}{\pi} \int_0^t \frac{x}{x^2+1} dx + \left(1 - \frac{2}{\pi} \arctan t\right) \cdot t \\
&= \frac{1}{\pi} \ln(t^2+1) + \left(1 - \frac{2}{\pi} \arctan t\right) \cdot t.
\end{aligned}$$

取 $t = O(n)$ 时由 Markov 不等式

$$\begin{aligned}
P\left(\sum Y_i^{(t)} > c_1 n^2\right) &\leq \frac{\frac{1}{\pi} \ln(t^2+1) - \left(1 - \frac{2}{\pi} \arctan t\right) \cdot t}{c_1 n} \\
&\leq \frac{O(\ln n) - O(\frac{1}{n}) \cdot O(n)}{c_1 n} = O\left(\frac{\ln n}{n}\right),
\end{aligned}$$

选取足够大的 c_1 即可保证概率不超过 $1/6$. 明所欲证.

3. 利用和上文类似的截断方式可知

$$\mathbb{E}\left(\sum Y_i\right) = n \cdot \mathbb{E}(Y_i) = O(n \ln t),$$

利用 Chernoff-Hoeffding 不等式即有

$$\begin{aligned} P\left(\sum Y_i^{(t)} \leq c_2 n\right) &= P\left(\sum Y_i^{(t)} \leq \mathbb{E}(Y_i) - O(n \ln t)\right) \\ &\leq \exp\left(-O\left(\frac{n \ln^2 t}{t^2}\right)\right). \end{aligned}$$

令 $t = O(\sqrt{n})$ 并利用 $P(\sum Y_i < c_2 n) < P\left(\sum Y_i^{(t)} < c_2 n\right)$ 即证.