CSC240 Winter 2024 Homework Assignment 3

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1. Let \mathcal{F} be the set of all functions from D to D, where D is a nonempty set. Consider the following two predicates with domain $\mathcal{F} \times \mathcal{F}$:

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\begin{array}{lcl} P(f,g) & = & \exists y \in D. \forall x \in D. [f(g(x)) \neq y] \text{ and} \\ Q(f,g) & = & \exists v \in D. [\forall u \in D. (f(u) \neq v) \text{ OR } \forall u \in D. (g(u) \neq v)]. \end{array}
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Formally prove that $\forall f \in \mathcal{F}. \forall g \in \mathcal{F}. (P(f,g) \text{ IMPLIES } Q(f,g)).$

Remember to number all lines, indent properly, and justify all your steps, including references to the appropriate line numbers, as described in Proof Outlines. Only do one step of the proof per line.

Proof

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Let f \in \mathcal{F} be arbitrary
1
          Let g \in \mathcal{F} be arbitrary
2
              Assume P(f,g)
3
              \exists y \in D. \forall x \in D. [f(g(x)) \neq y]; definition of P(f,g), L3
              Let i \in D be such that \forall x \in D.[f(g(x)) \neq i]; instantiation, L4
5
              R 	ext{ OR NOT}(R); tautology
6
              Let S = \exists a \in D.(f(a) = i)
7
              \exists a \in D.(f(a) = i) \text{ OR NOT}(\exists a \in D.(f(a) = i)); \text{ substitution of all } R \text{ by } S, \text{ L6}
8
                  Assume \exists a \in D.(f(a) = i)
9
                  Let r \in D be such that f(r) = i; instantiation, L9
10
                  Let v = r
11
                  v \in D
12
                      To obtain a contradiction, assume NOT(\forall u \in D.(g(u) \neq v))
13
                      \exists u \in D.(q(u) = v); negation of quantifiers, L13
14
                      Let k \in D be such that g(k) = v; instantiation, L14
15
                      f(g(k)) = f(v); property of function, L15
16
                      f(g(k)) = f(r); substitute v = r from L11 to L16
17
                      f(q(k)) = i; substitute f(r) = i from L10 to L17
18
                      f(g(k)) \neq i; specialization, L5, L15
19
                      This is a contradiction: L18, L19
20
                  \forall u \in D.(g(u) \neq v); proof by contradiction, L13, L20
21
                  \forall u \in D.(f(u) \neq v) \text{ OR } \forall u \in D.(g(u) \neq v); \text{ proof of disjunction, L21}
22
                  \exists v \in D. [\forall u \in D. (f(u) \neq v) \text{ OR } \forall u \in D. (g(u) \neq v)]; \text{ construction, L11, L12, L22}
23
              \exists a \in D.(f(a) = i) \text{ IMPLIES } \exists v \in D. [\forall u \in D.(f(u) \neq v) \text{ OR } \forall u \in D.(g(u) \neq v)];
24
              direct proof, L9, L23
                  Assume NOT(\exists a \in D.(f(a) = i))
25
                  \forall a \in D.(f(a) \neq i); negation of quantifiers, L25
26
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\forall u \in D.(f(u) \neq i); substitute u for quantified a L26
27
                    Let v = i
28
                    v \in D
29
                    \forall u \in D.(f(u) \neq v); substitute v = i from L28 to L27
30
                    \forall u \in D.(f(u) \neq v) \text{ OR } \forall u \in D.(g(u) \neq v); \text{ proof of disjunction, L30}
31
                    \exists v \in D. [\forall u \in D. (f(u) \neq v) \text{ OR } \forall u \in D. (g(u) \neq v)]; \text{ construction, L28, L29, L31}
32
                NOT(\exists a \in D.(f(a) = i)) \text{ IMPLIES } \exists v \in D.[\forall u \in D.(f(u) \neq v) \text{ OR } \forall u \in D.(g(u) \neq v)];
33
                direct proof, L25, L32
                \exists v \in D. [\forall u \in D. (f(u) \neq v) \text{ OR } \forall u \in D. (g(u) \neq v)]; \text{ proof by cases, L8, L24, L33}
34
                Q(f,g); definition of Q, L34
35
            P(f,g) IMPLIES Q(f,g); direct proof L3, L35
36
       \forall g \in \mathcal{F}.(P(f,g) \text{ IMPLIES } Q(f,g)); \text{ generalization, L2, L36}
38 \forall f \in \mathcal{F}. \forall g \in \mathcal{F}. (P(f,g) \text{ IMPLIES } Q(f,g)); \text{ generalization, L1, L37}
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2. Recall that, if p is a polynomial of degree $m \ge 1$, then there exist coefficients a_i for $0 \le i \le m$ such that $a_m \neq 0$ and, for all numbers n,

$$p(n) = \sum_{i=0}^{m} a_i n^i.$$

Give a well-structured informal proof that, for any polynomial p of degree at least 1 whose coefficients are natural numbers, there is a natural number n such that p(n) is not prime.

Proof:

We will prove by contradiction.

• Assume there is a polynomial p of degree $m \geq 1$ whose coefficients are natural numbers, where for all natural number n, p(n) is a prime.

Let
$$p(n) = a_0 + \sum_{i=1}^{m} a_i n^i$$
, where $a_m \neq 0$, $m \geq 1$, and $a_i \geq 0$ for $0 \leq i \leq m$.

By our assumption,
$$p(0) = a_0$$
 is a prime (every term containing n goes to 0).
Since $a_0 \in \mathbb{N}$, $p(a_0) = a_0 + \sum_{i=1}^m a_i \cdot (a_0)^i = a_0 \cdot \left(1 + \sum_{i=1}^m a_i \cdot (a_0)^{i-1}\right)$ is also a prime.

By the definition of a prime, $p(a_i)$ only has 1 and itself as factors, therefore:

- case 1: $a_0 = 1$

Since we have shown a_0 is a prime under the assumption, we have $a_0 \neq 0$ and $a_0 \neq 1$ since 0 and 1 are not prime. There is a contradiction.

- case 2: $(1 + \sum_{i=1}^{m} a_i \cdot (a_0)^{i-1}) = 1$. Immediately, we have: $\sum_{i=1}^{m} a_i \cdot (a_0)^{i-1} = 0$ Since for $0 \le i \le m$, $a_i \ge 0$, we have each term $a_i \cdot (a_0)^{i-1} \ge 0$. Since $a_0 \ge 2$ and $i-1 \geq 0$, the latter term $(a_0)^{i-1} \geq 1$. To keep the summation 0, we must have $a_i = 0$ for $1 \le i \le m$ $(a_1 = ... = a_m = 0)$, which contradicts with $a_m \ne 0$

Therefore, $p(a_i)$ cannot be prime. This contradicts to $p(a_i)$ is a prime.

Because of the contradiction, we conclude there is NO polynomial p of degree m > 1 whose coefficients are natural numbers, where for all natural number n, p(n) is a prime. This is equivalent to say for any polynomial p of degree at least 1 whose coefficients are natural numbers, there is a natural number n such that p(n) is not prime.