Welcome To My Presentation

My Presentation Topic is

Radix Sort

Presented By

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Content

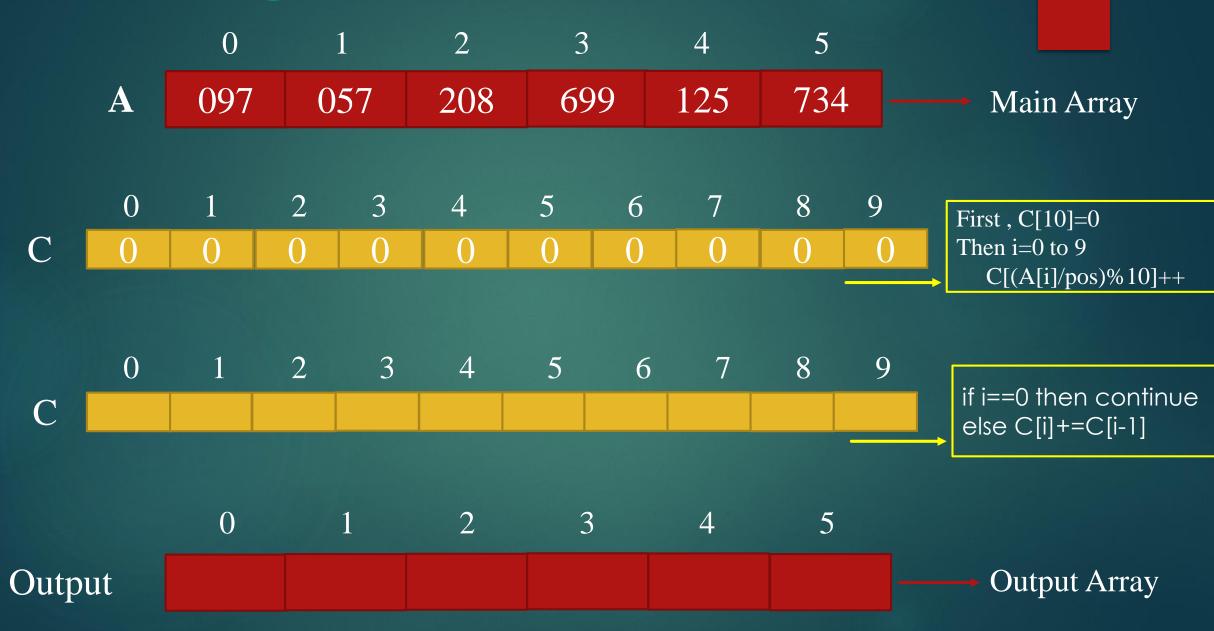
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- ► Radix Sort Algorithm
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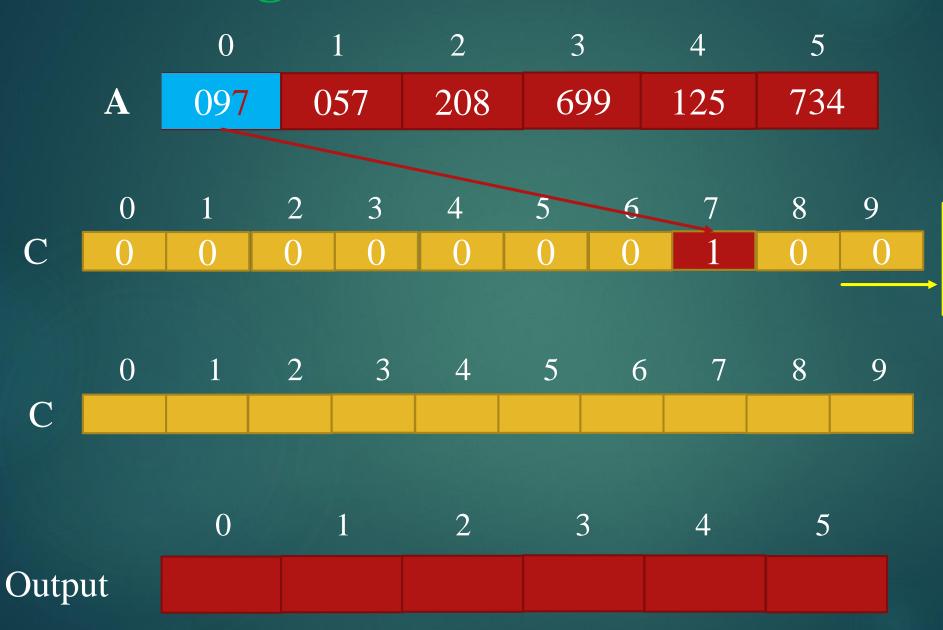
Introduction

- ▶ Radix sort is *a sorting algorithms* that sorts the elements by first grouping the individual digits of the same **place value**.
- ▶ Then, sort the elements according to their increasing/decreasing order.
- ▶ We have seen many sorting algorithms but this one is different method.
- ▶ It is interesting because it requires the absolute minimum amount of apace and the minimum amount of data movement and most amazing off all, it does no comparisons.

Radix Sort Algorithm

- ▶ Find the largest element in the array, i.e. max. Let X be the number of digits in max . X is calculated because we have to go through all the significant places of all elements.
- Now, go through each significant place one by one. Use any stable sorting technique to sort the digits at each significant place. We have used counting sort for this
- ▶ First, we will sort elements based on the value of the unit place. Then, we will sort elements based on the value of the tenth place. This process goes on until the last significant place.



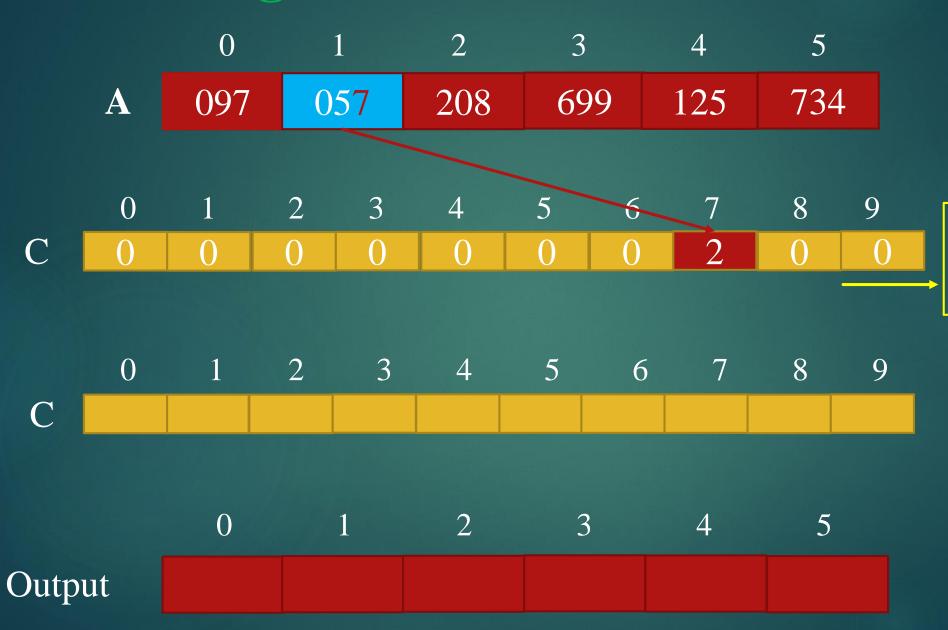


If i=0 & pos=1

C[(097/1)%10]++

Then, C[7]=1

=C[7]++

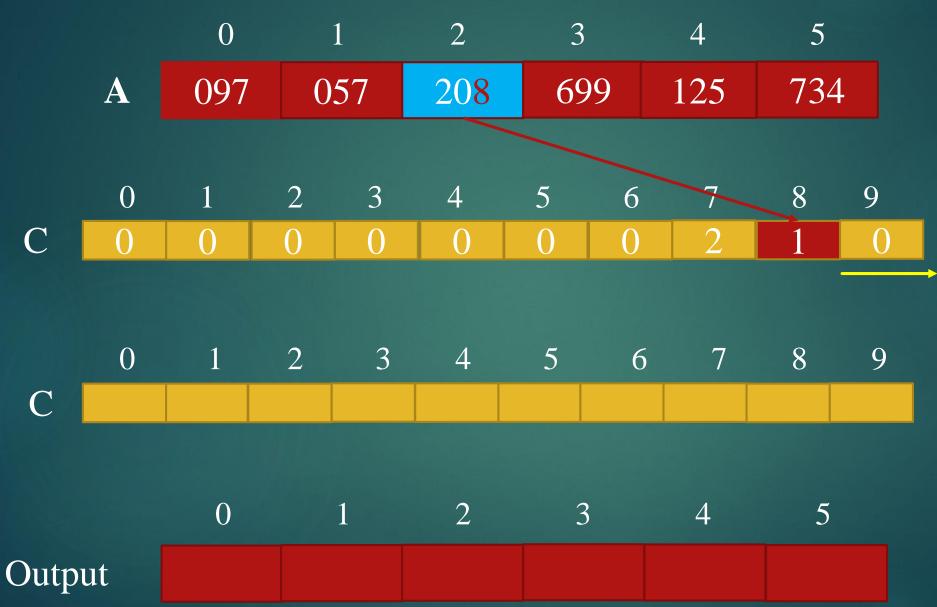


If i=1 & pos=1

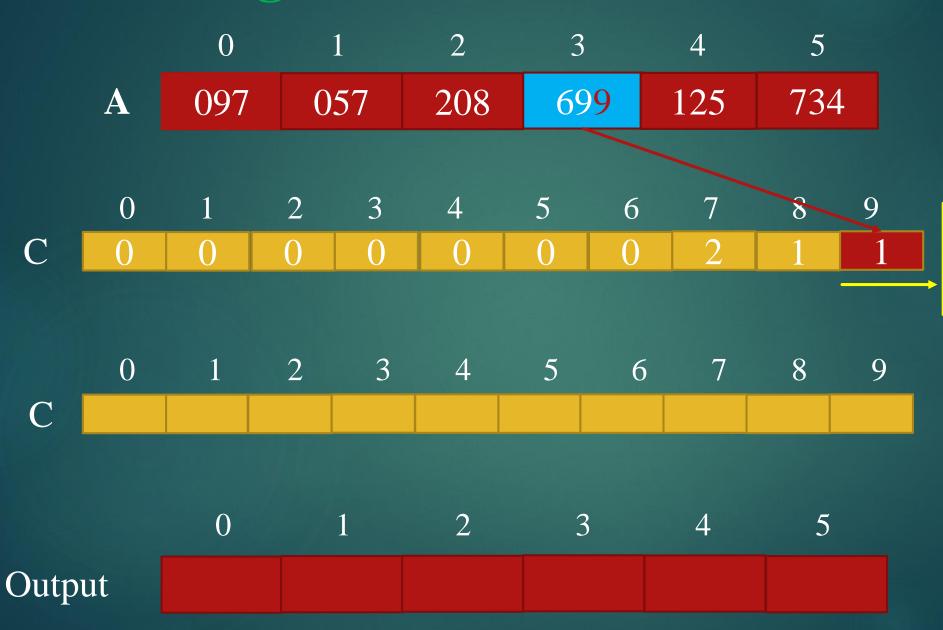
C[(057/1)%10]++

Then, C[7]=2

=C[7]++



If i=2 & pos=1 C[(208/1)%10]++ =C[8]++ Then, C[8]=1

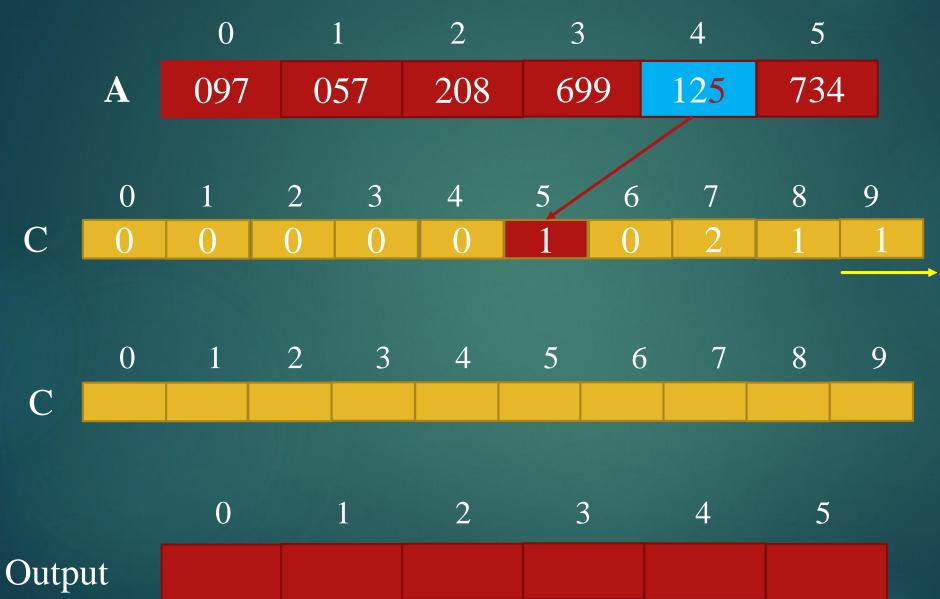


If i=3 & pos=1

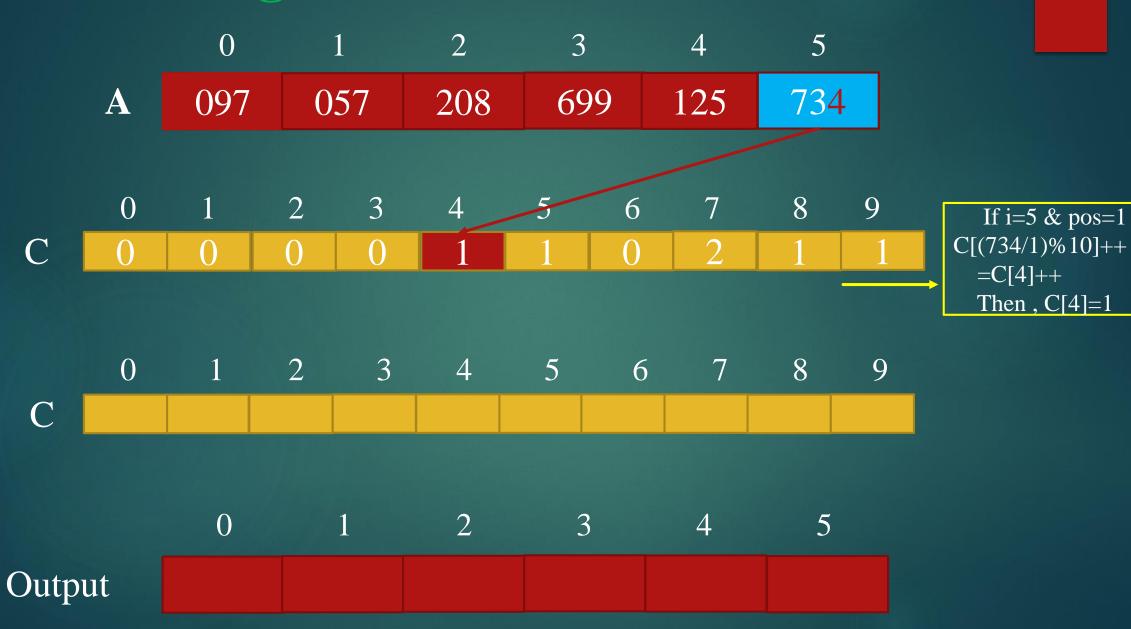
C[(699/1)%10]++

Then, C[9]=1

=C[9]++



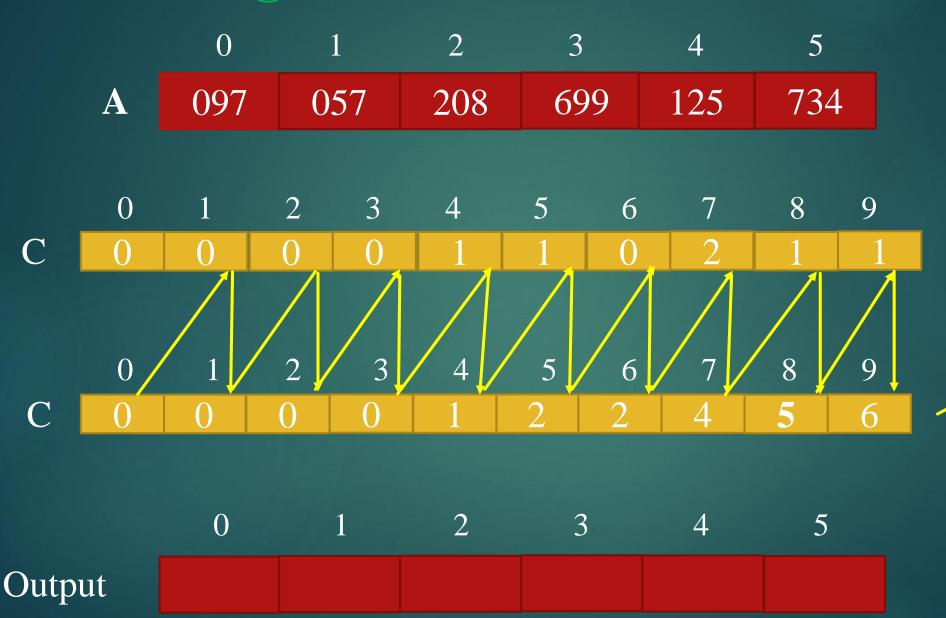
If i=4 & pos=1 C[(125/1)%10]++ =C[5]++ Then, C[5]=1



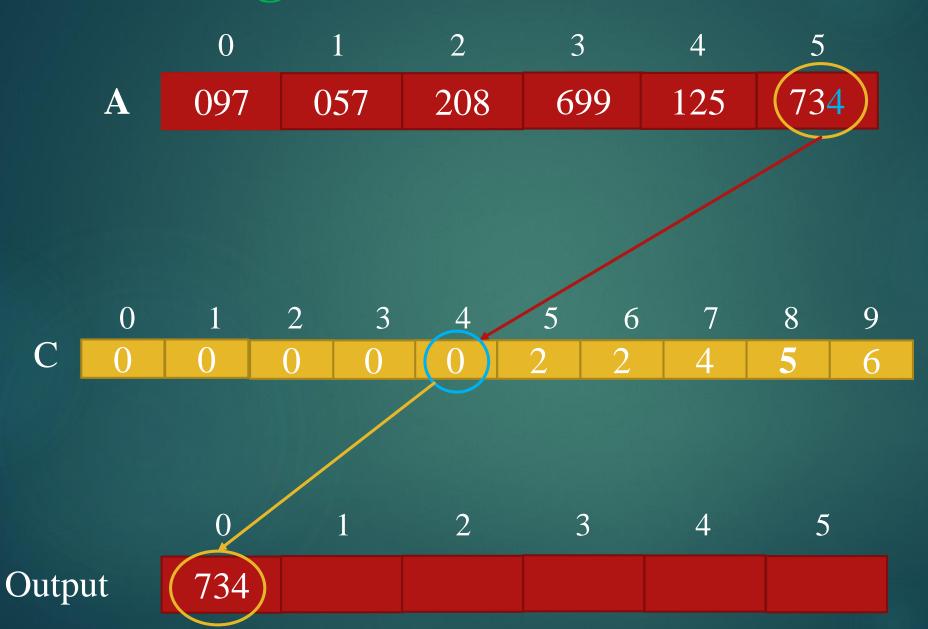
If i=5 & pos=1

Then, C[4]=1

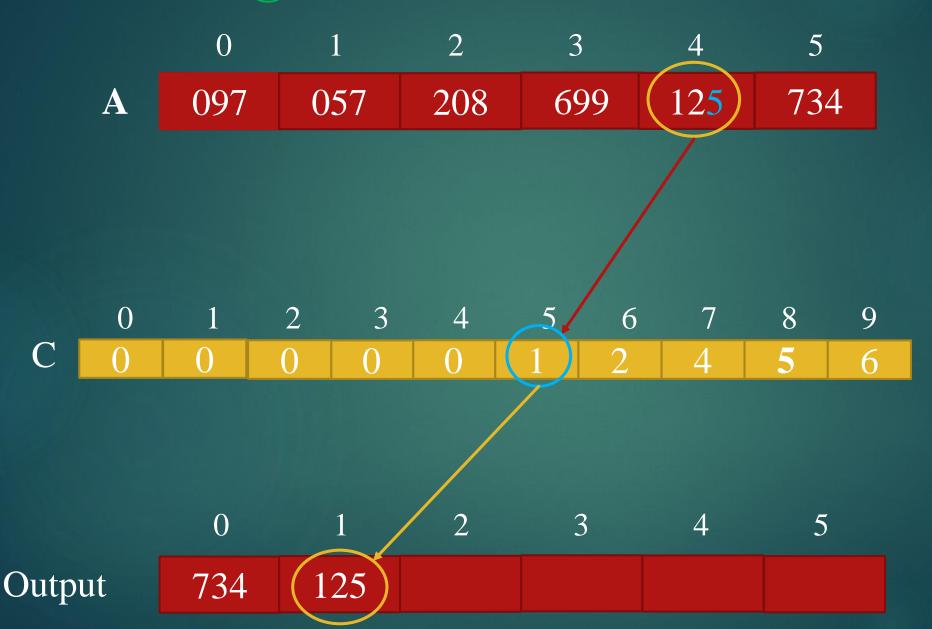
=C[4]++



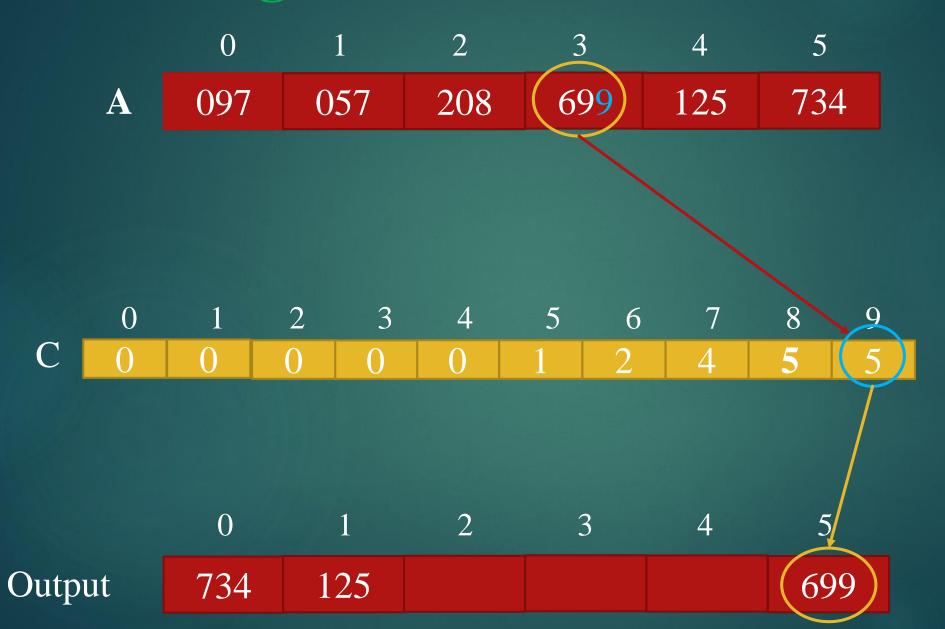
If I == 0 then continue Else c[i]+=C[i-1]



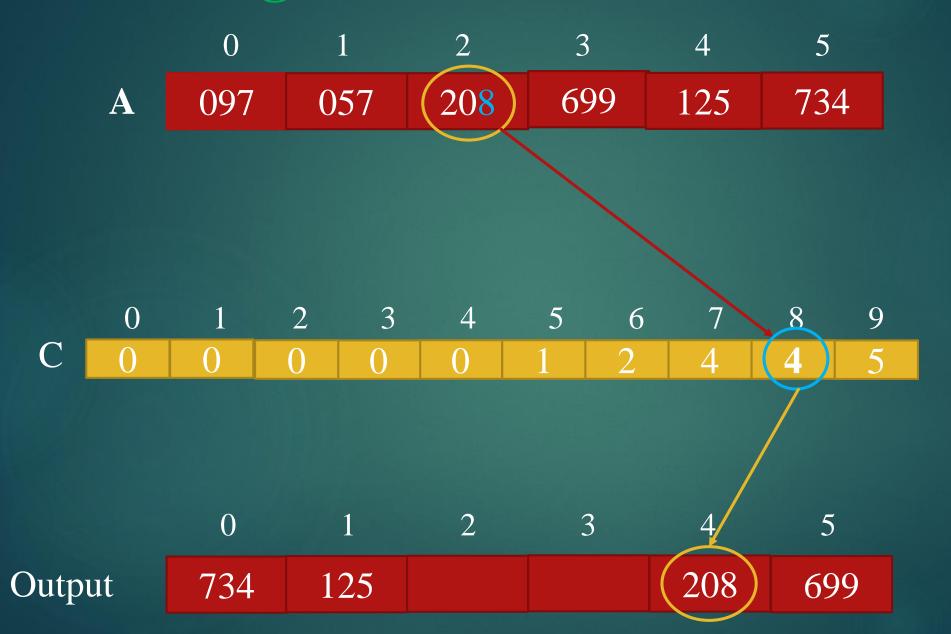
If i=5 & pos=1 C[(734/1)% 10]--=C[4]-1,C[4]=0 Output[0]=734



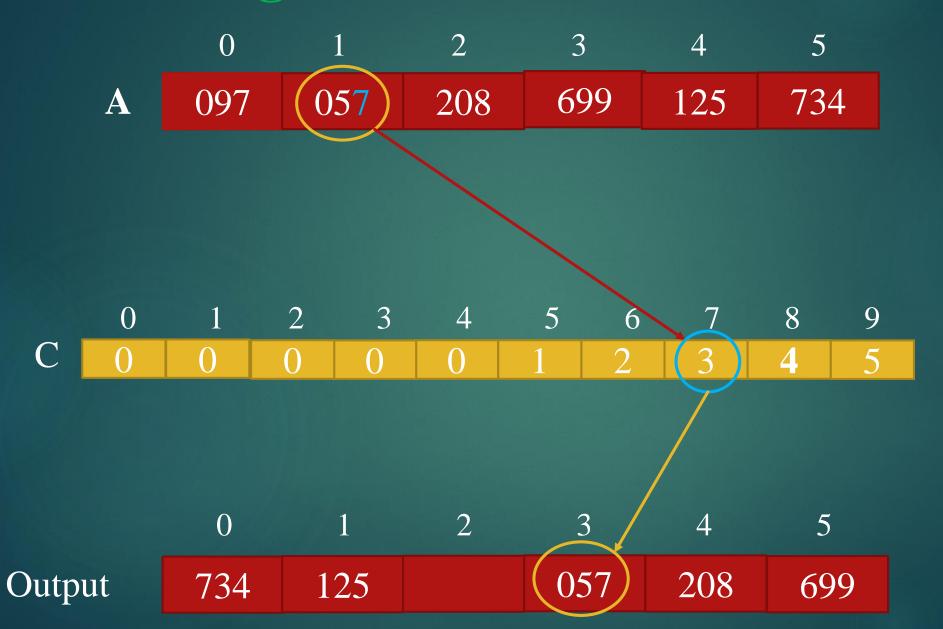
If i=4 & pos=1 C[(125/1)% 10]--=C[5]-1,C[5]=1 Output[1]=125



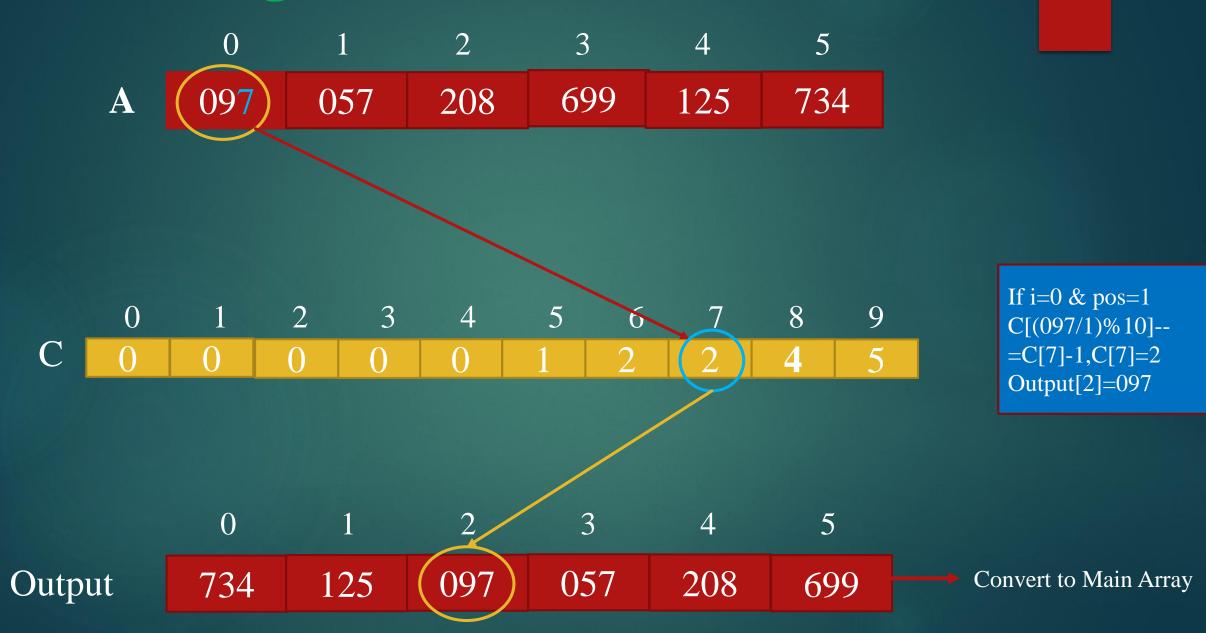
If i=3 & pos=1 C[(699/1)% 10]--=C[9]-1,C[9]=5 Output[5]=699

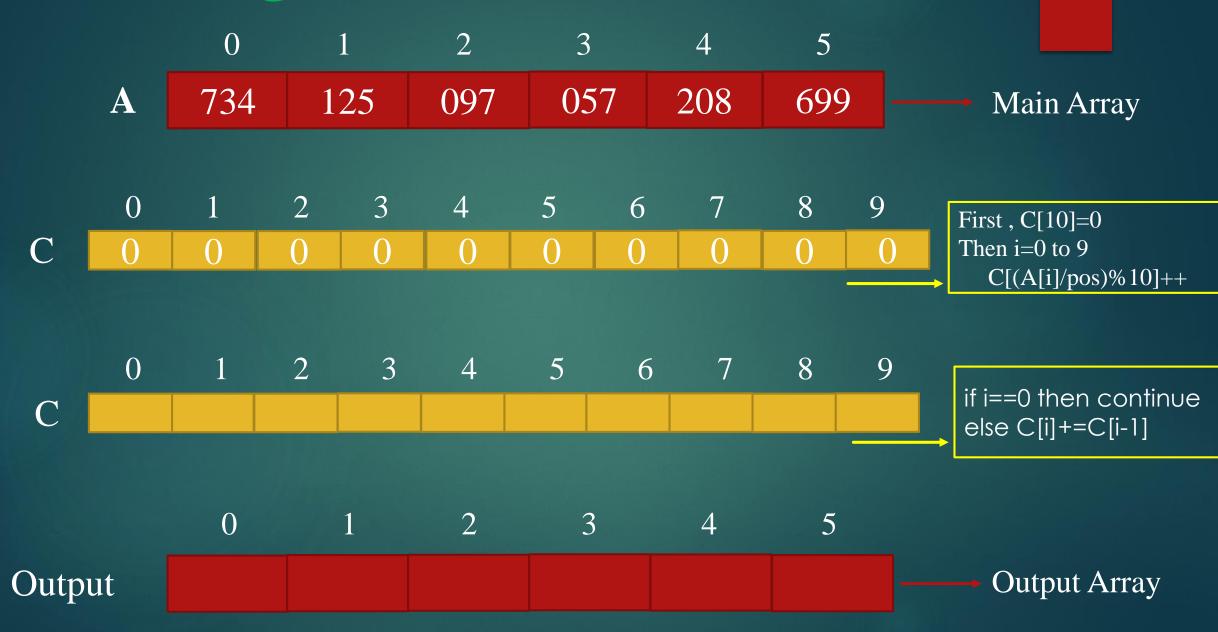


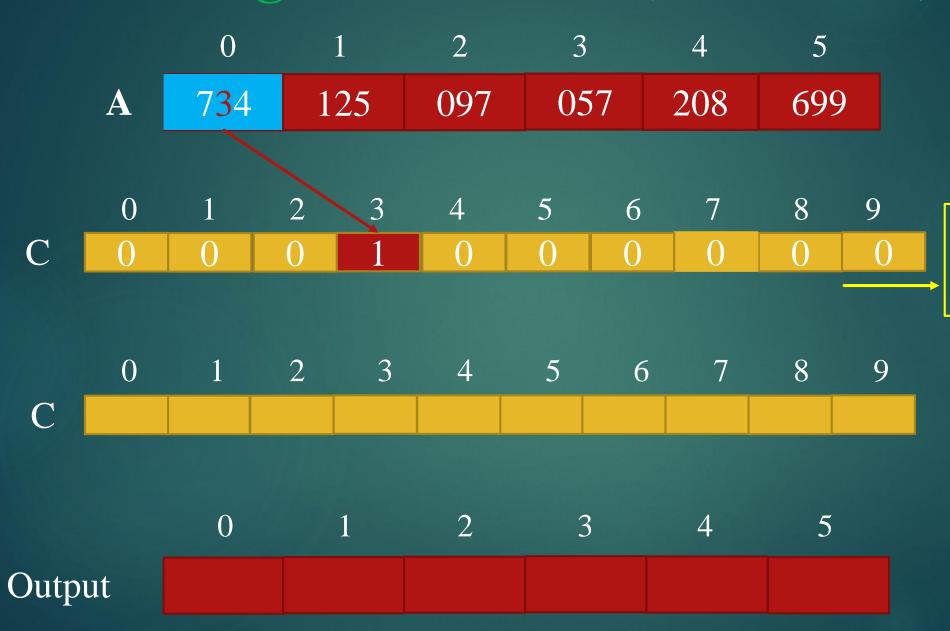
If i=2 & pos=1 C[(208/1)%10]--=C[8]-1,C[8]=4 Output[4]=208



If i=1 & pos=1 C[(057/1)% 10]--=C[7]-1,C[7]=3 Output[3]=057





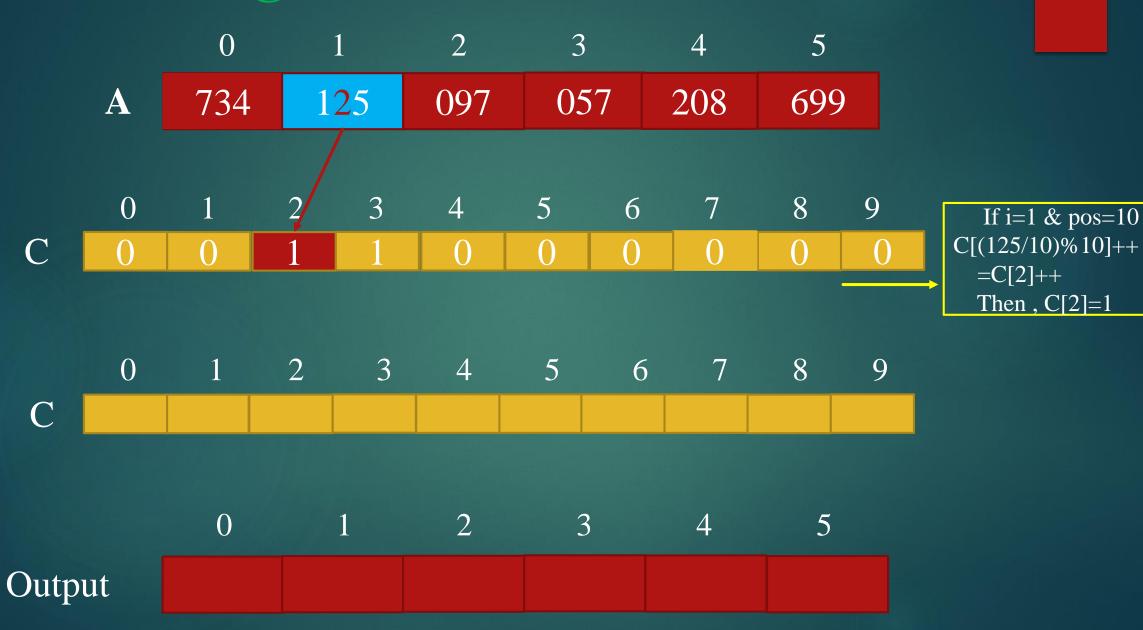


If i=0 & pos=10

C[(734/10)%10]++

Then, C[3]=1

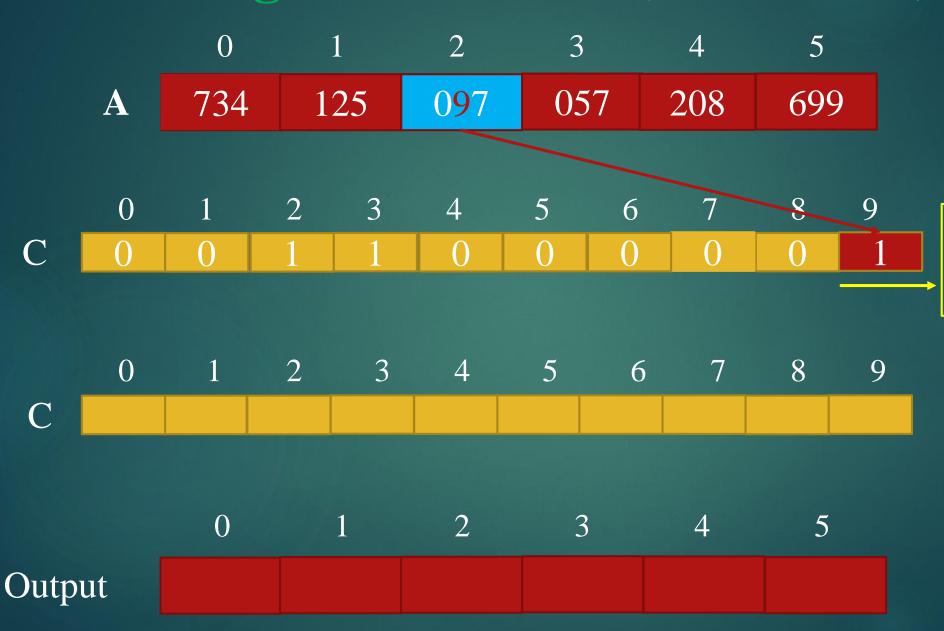
=C[3]++



If i=1 & pos=10

=C[2]++

Then, C[2]=1

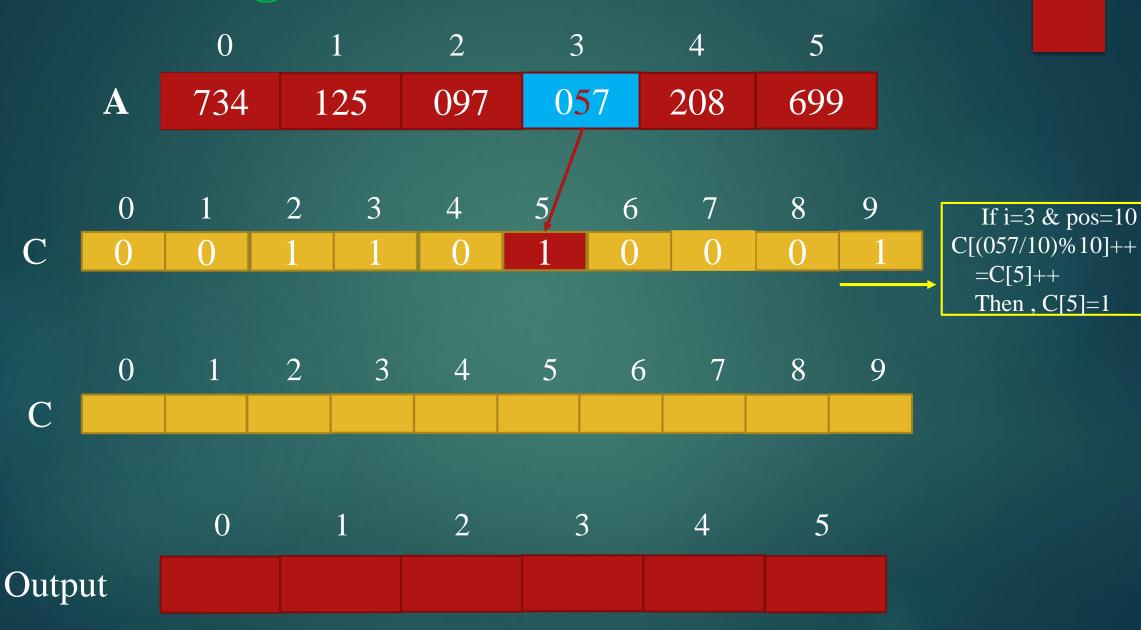


If i=2 & pos=10

C[(097/10)%10]++

Then, C[9]=1

=C[9]++



If i=3 & pos=10

=C[5]++

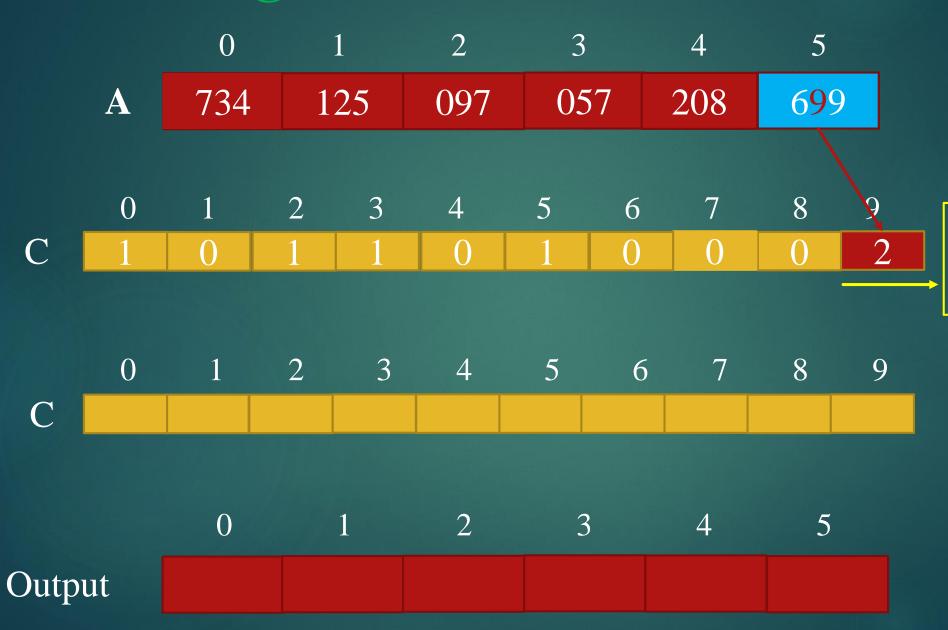
Then, C[5]=1



If i=4 & pos=10

=C[0]++

Then, C[0]=1

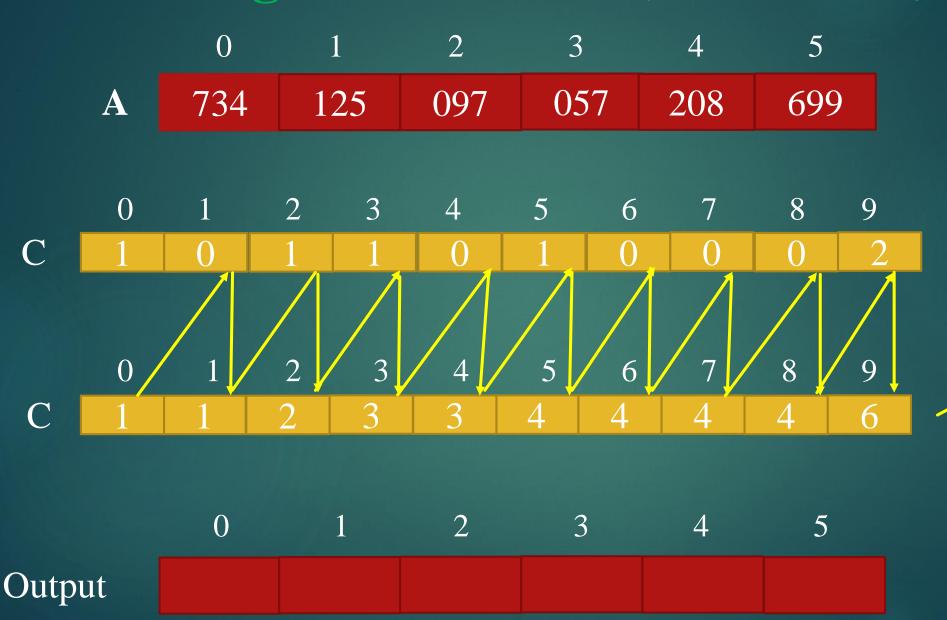


If i=5 & pos=10

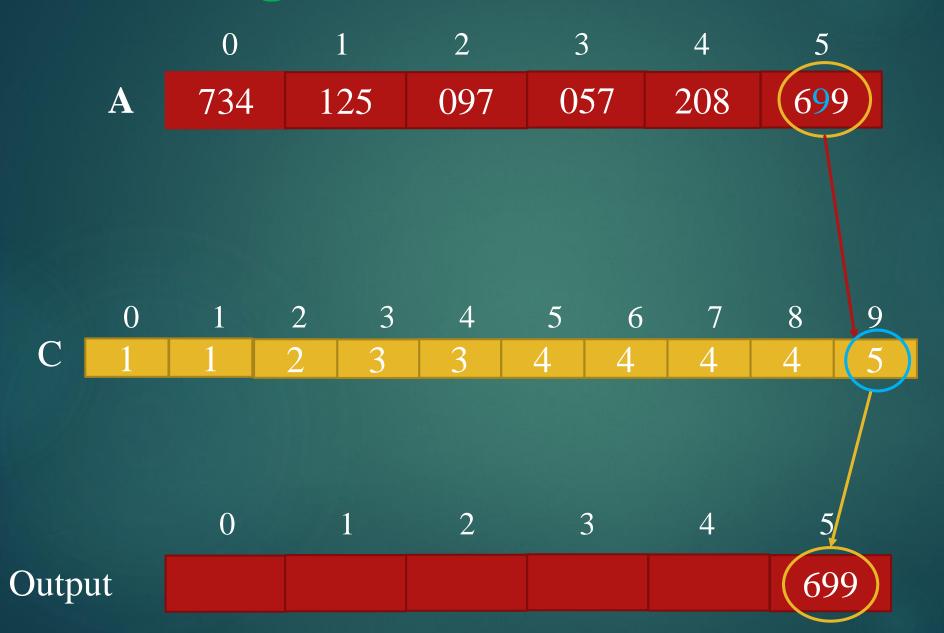
C[(699/10)%10]++

Then, C[9]=2

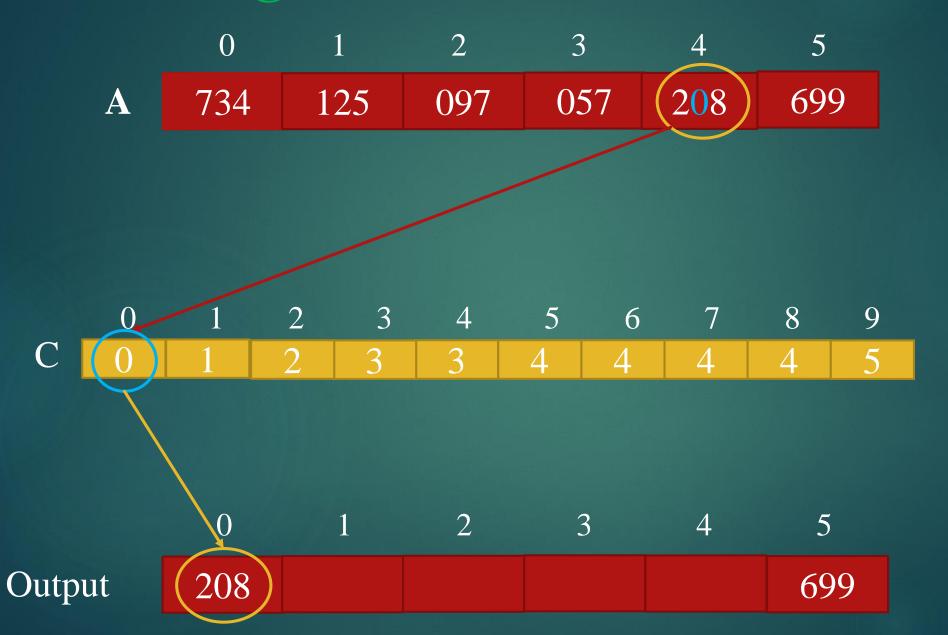
=C[9]++



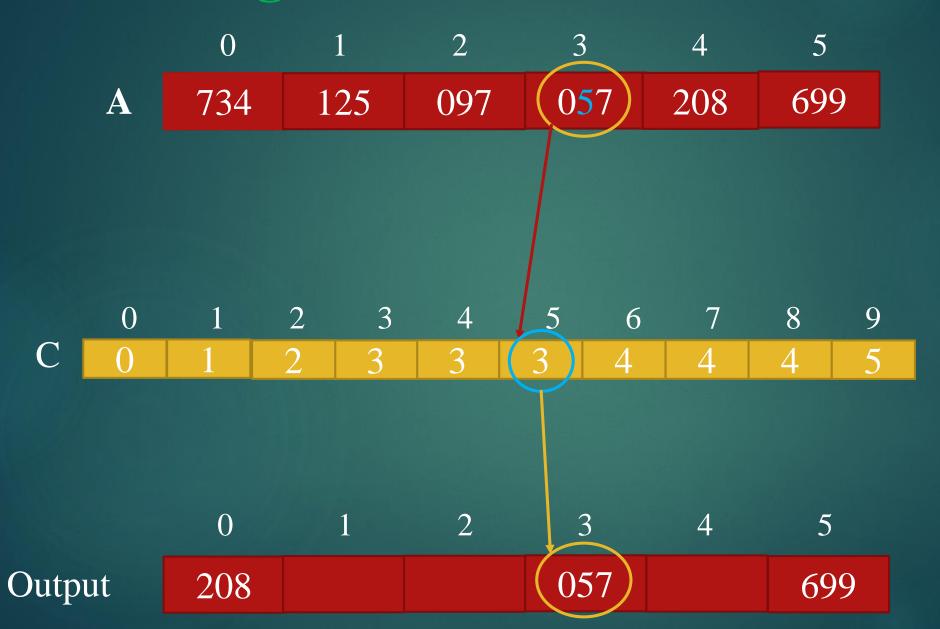
If I == 0 then continue Else c[i]+=C[i-1]



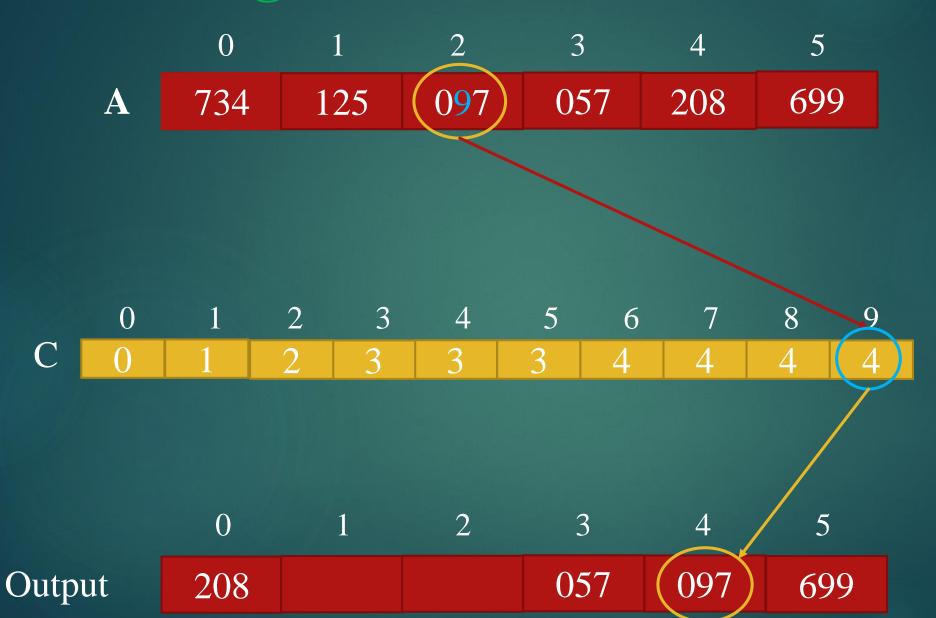
If i=5 & pos=10 C[(699/10)%10]--=C[9]-1,C[9]=5 Output[5]=699



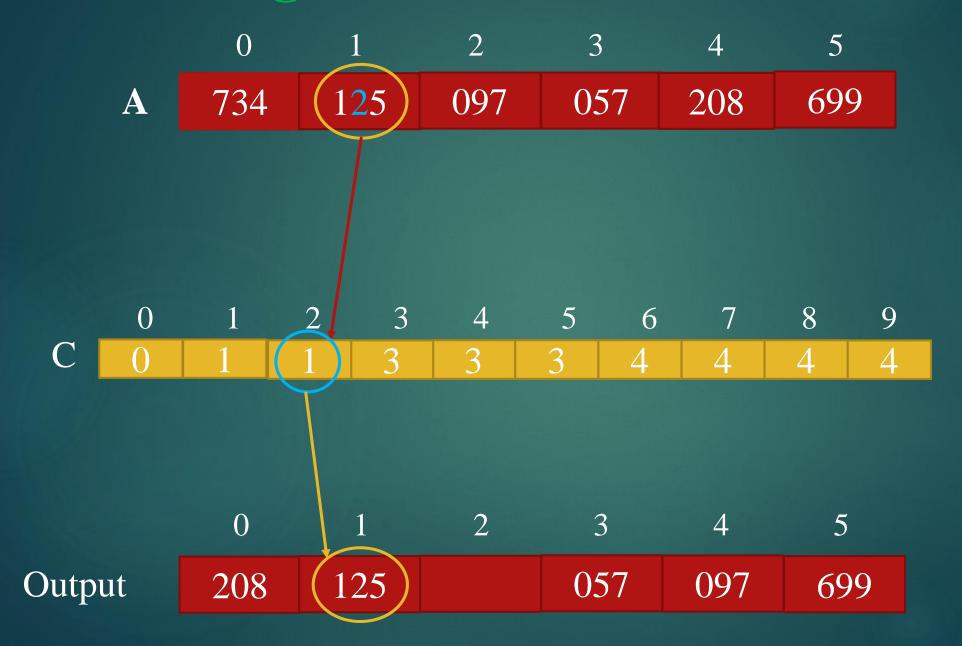
If i=4 & pos=10 C[(208/10)%10]--=C[0]-1,C[0]=0 Output[0]=208



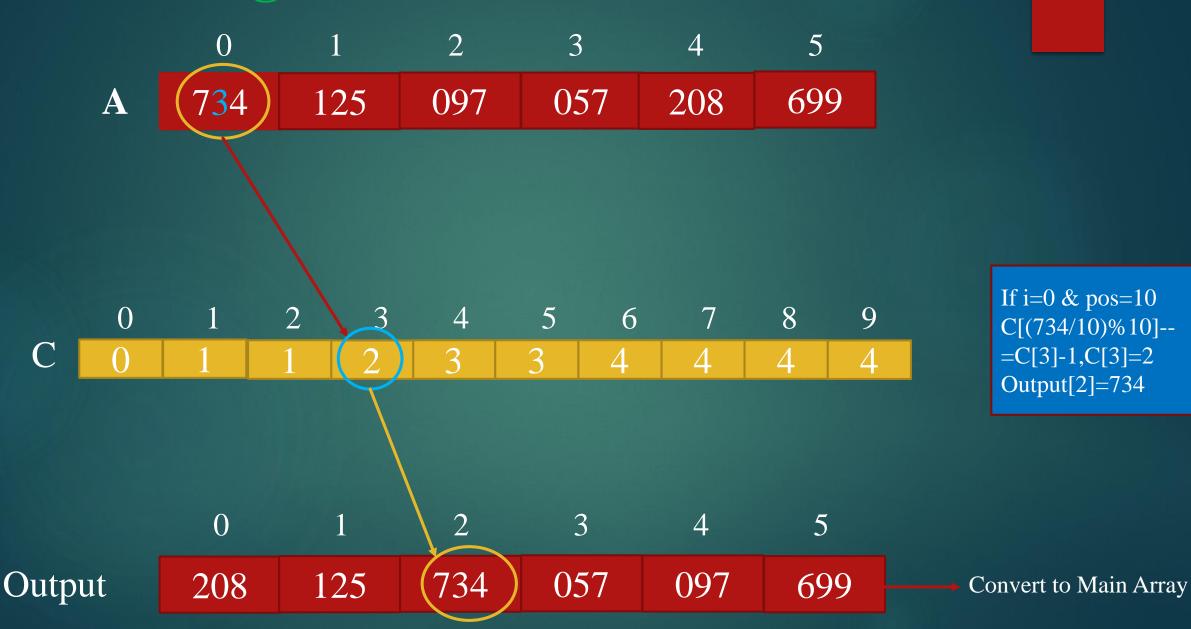
If i=3 & pos=10 C[(057/10)%10]--=C[5]-1,C[5]=3 Output[3]=057



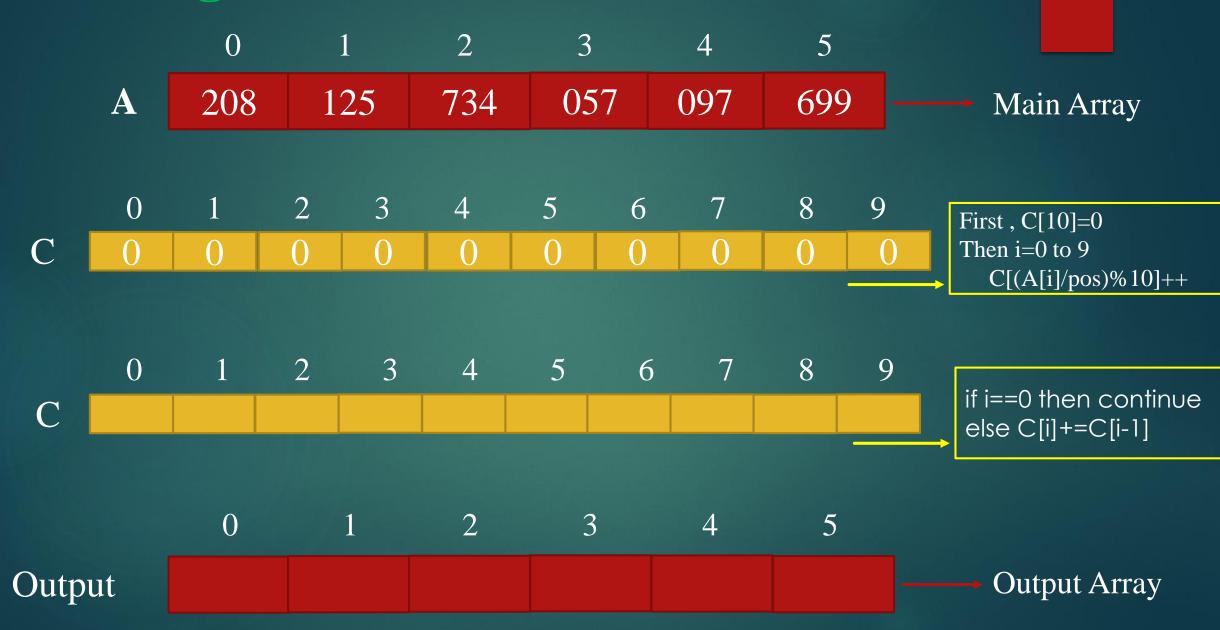
If i=2 & pos=10 C[(097/10)%10]--=C[9]-1,C[9]=4 Output[4]=097



If i=1 & pos=10 C[(125/10)%10]--=C[2]-1,C[2]=1 Output[1]=125



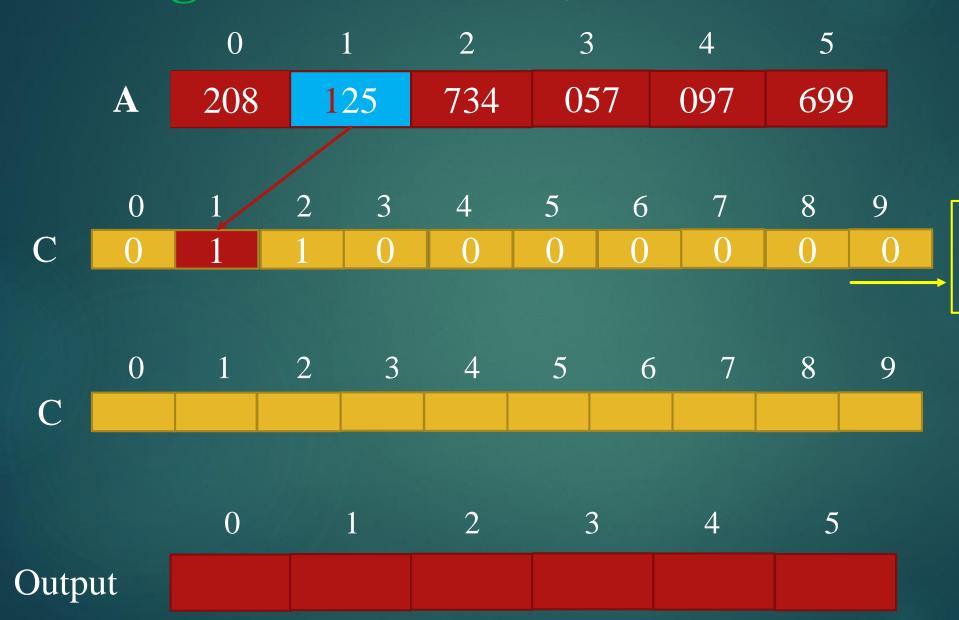
Working of Radix Sort(Hundreds Place)



Working of Radix Sort(Hundreds Place)



Working of Radix Sort(Hundreds Place)

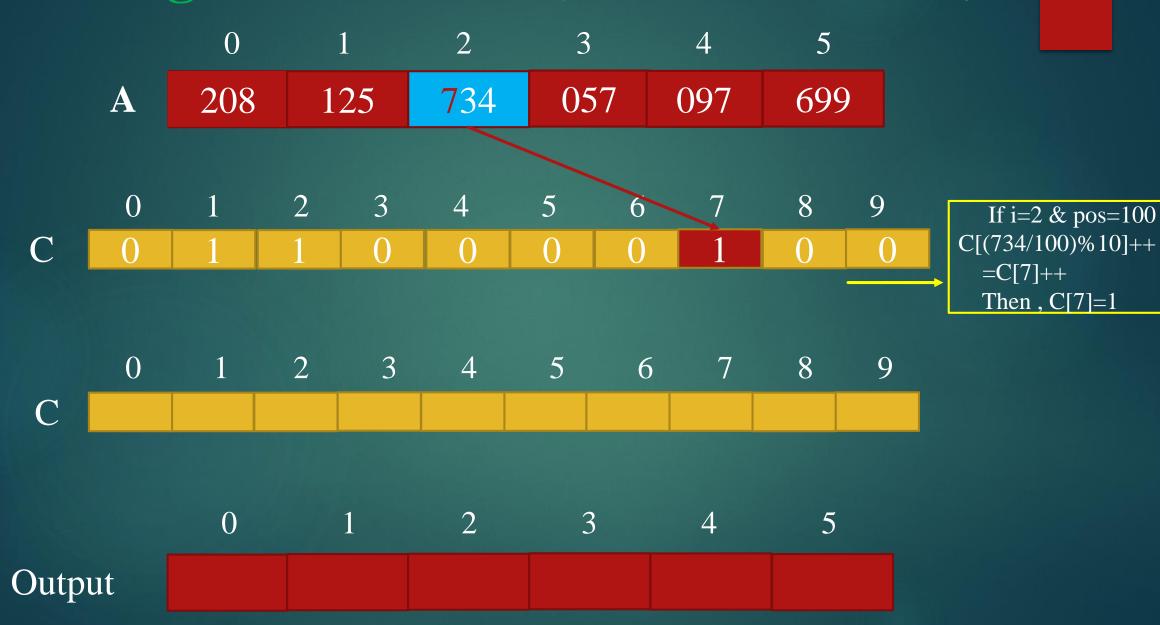


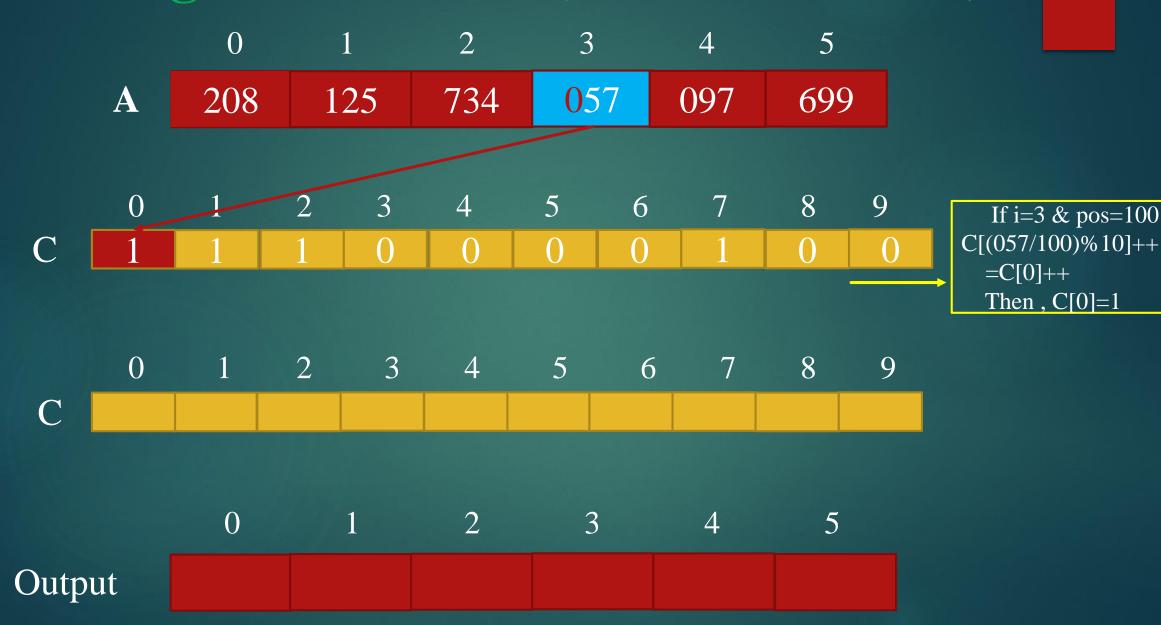
If i=1 & pos=100

C[(125/100)%10]++

Then, C[1]=1

=C[1]++

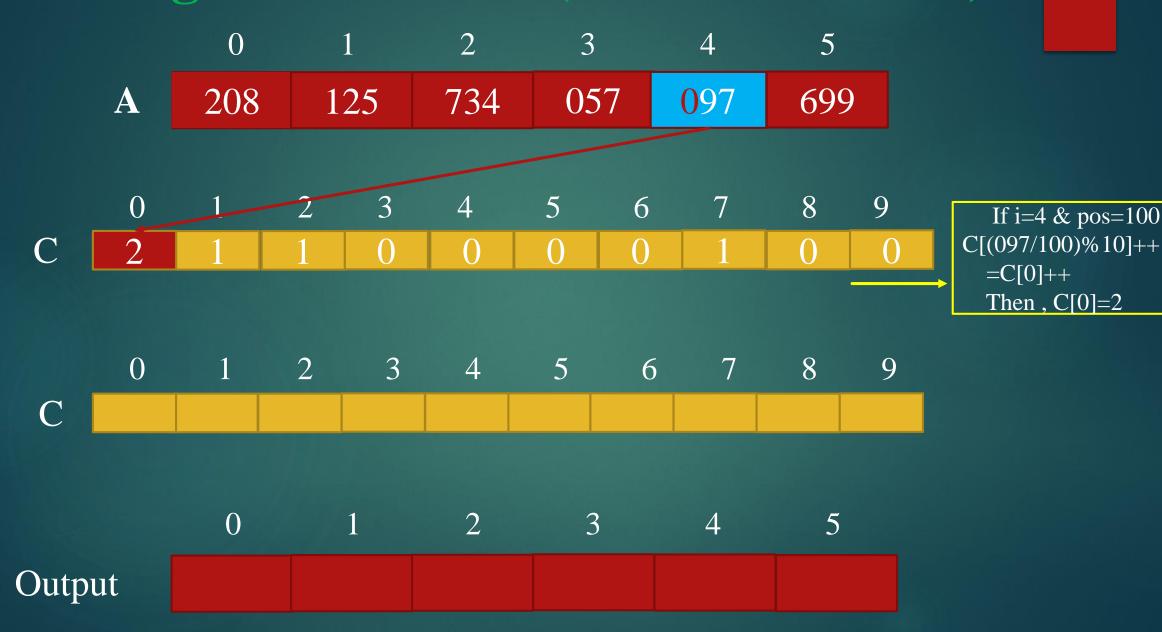




If i=3 & pos=100

=C[0]++

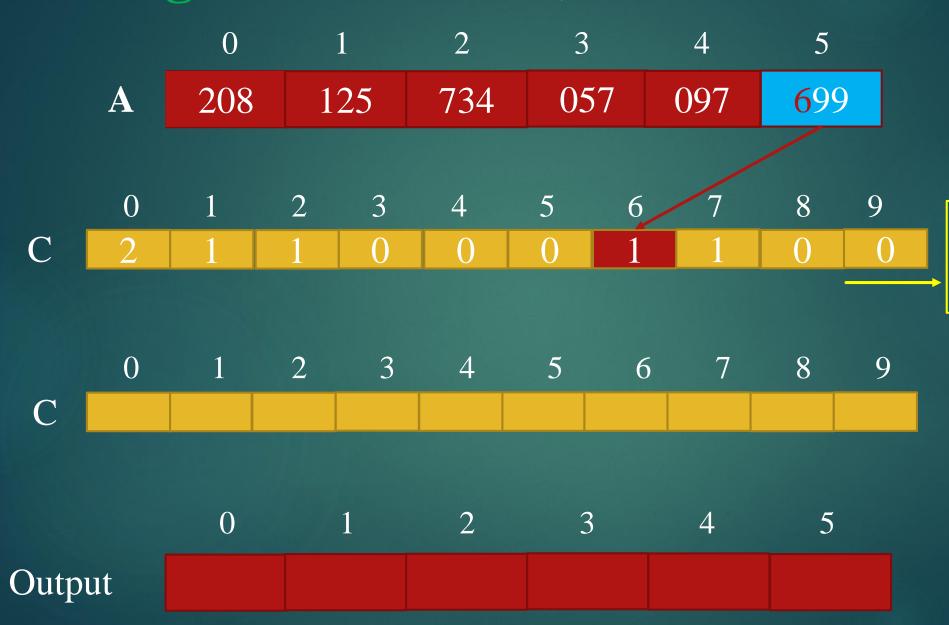
Then, C[0]=1



If i=4 & pos=100

=C[0]++

Then, C[0]=2

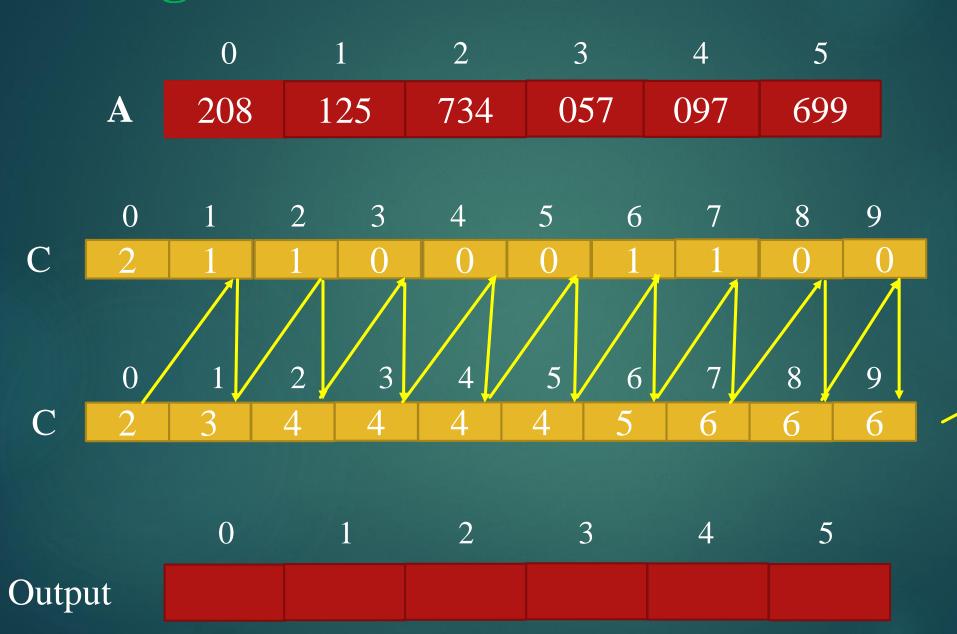


If i=5 & pos=100

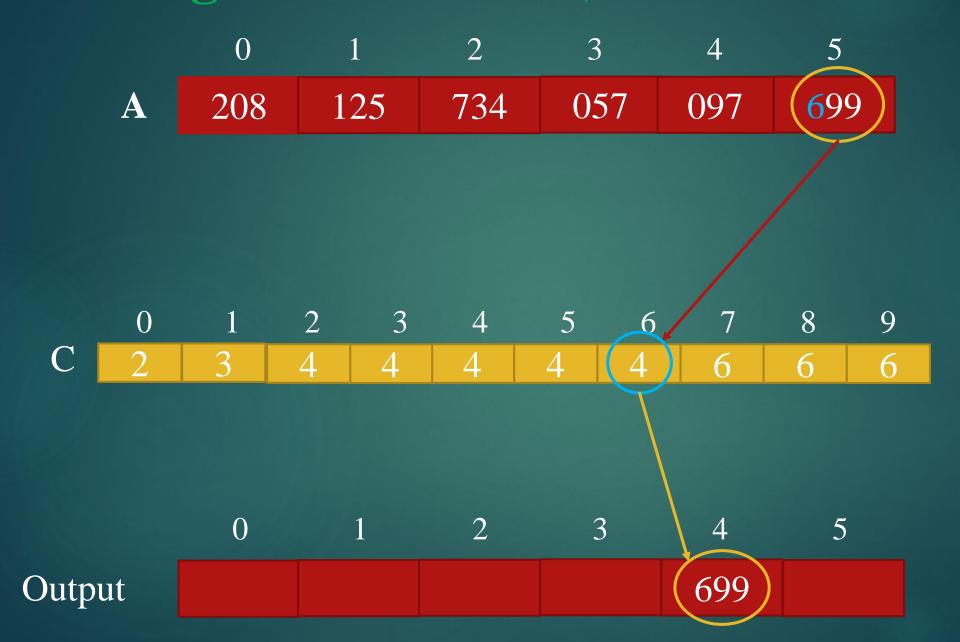
C[(699/100)%10]++

Then, C[6]=1

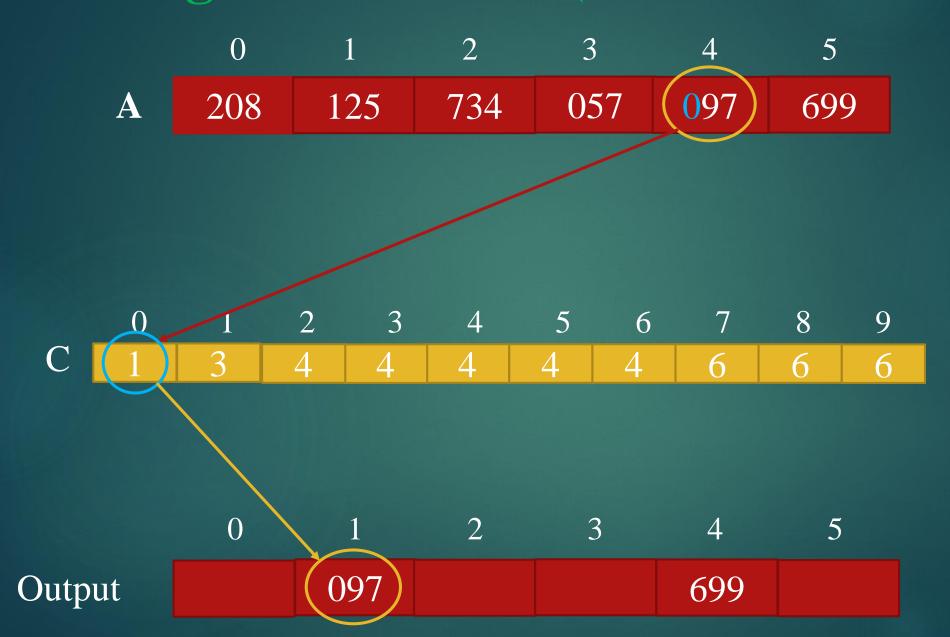
=C[6]++



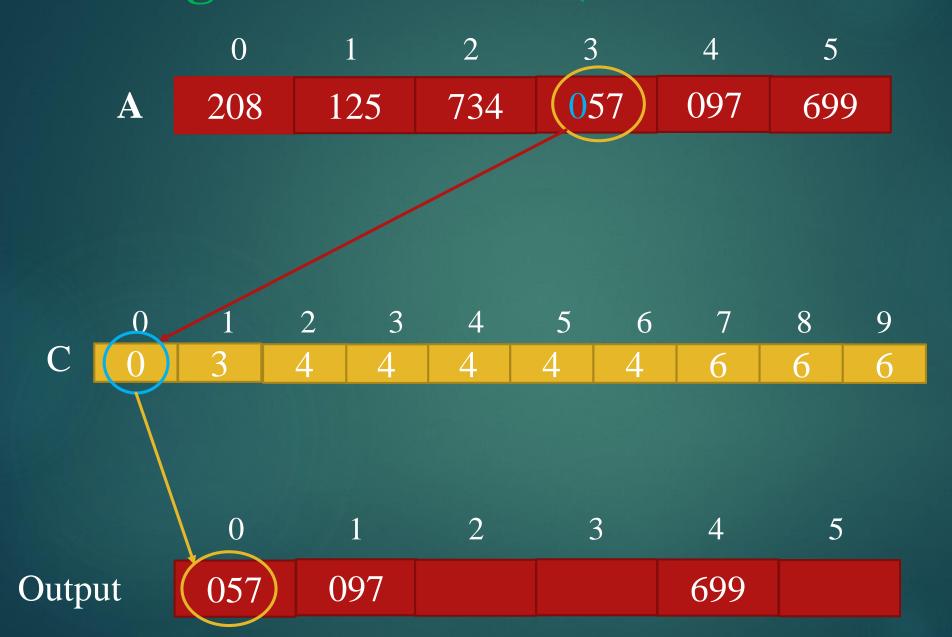
If I == 0 then continue Else c[i]+=C[i-1]



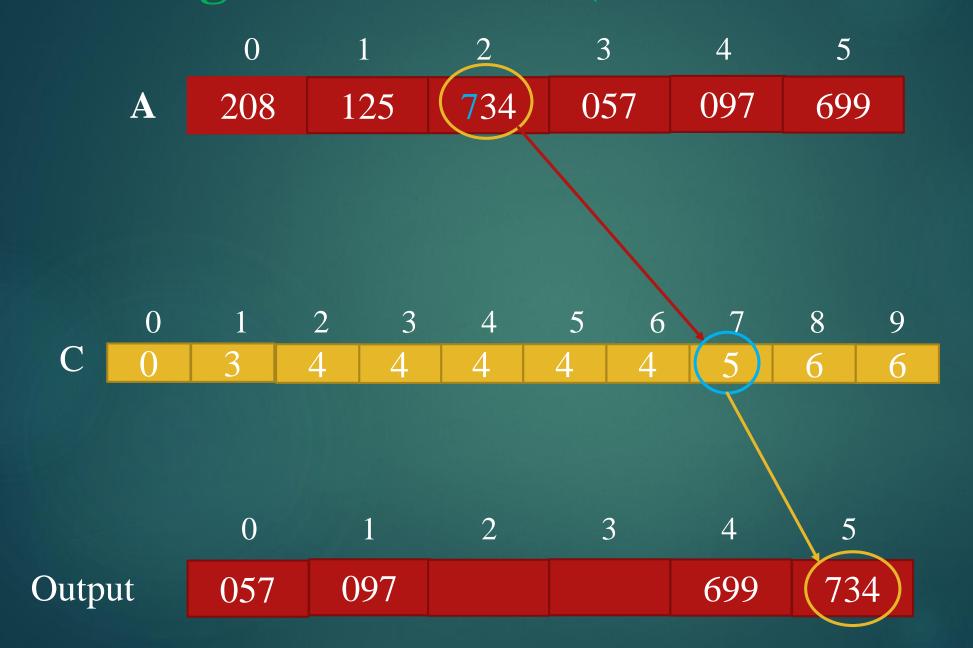
If i=5 & pos=100 C[(699/100)%10]--=C[6]-1,C[6]=4 Output[4]=699



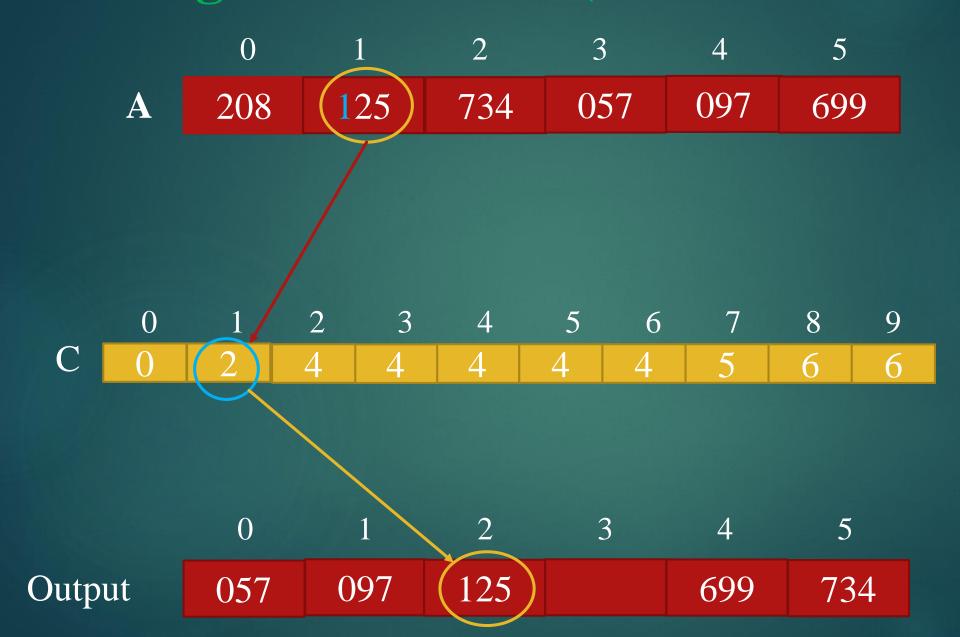
If i=4 & pos=100 C[(097/100)%10]--=C[0]-1,C[0]=1 Output[1]=097



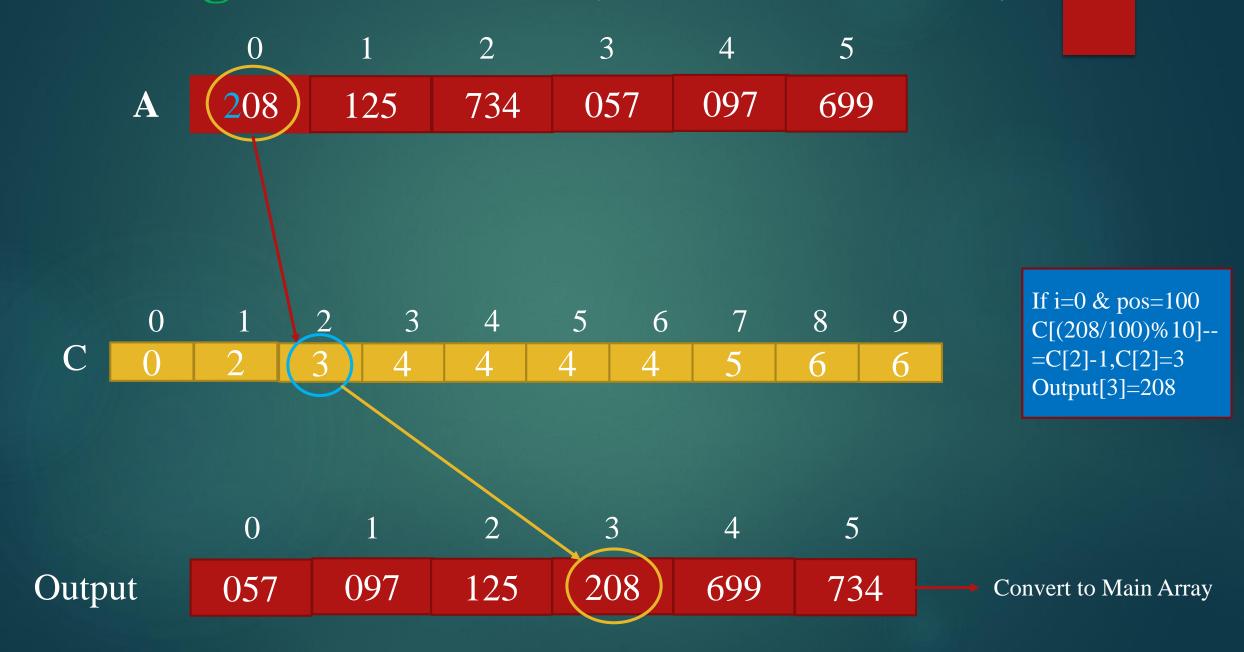
If i=3 & pos=100 C[(057/100)%10]--=C[0]-1,C[0]=0 Output[0]=057



If i=2 & pos=100 C[(734/100)%10]--=C[7]-1,C[7]=5 Output[5]=734

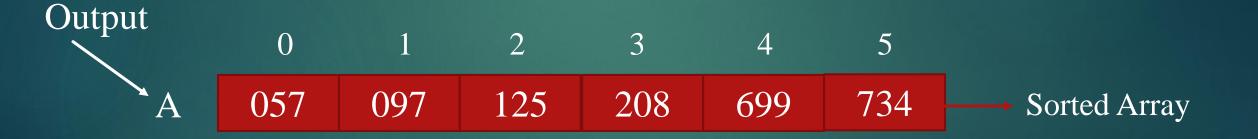


If i=1 & pos=100 C[(125/100)%10]--=C[1]-1,C[1]=2 Output[2]=125



Working of Radix Sort





Pseudocode

```
Quick Selection Sort.cpp × Queue.cpp × Margesort.cpp × Counting Sort.cpp
                                                                         Radix_Sort.cpp 

                                                                                        B_Sale.cpp >
                                                                                                    A Cabbages.cpp X B Bouzu Mekuri.cpp X B Boo
      void counting_sort(if A[], if n, if pos)
          ll ctlen = 10, i;
           11 count[ctlen], Output[n];
          for (i = 0; i < ctlen; i++)count[i] = 0;</pre>
          for (i = 0; i < n; i++)count[(A[i] / pos) % 10]++;</pre>
 94
          for (i = 1; i < ctlen; i++)count[i] += count[i - 1];</pre>
           for (i = n - 1; i \ge 0; i--)Output[--count[(A[i] / pos) % 10]] = A[i];
           for (i = 0; i < n; i++)A[i] = Output[i];</pre>
      void radix_sort(ll A[], ll n)
100
          11 i, mx = 0, pos;
          Forn(i, n)mx = max(mx, A[i]);
102
103
           for (pos = 1; mx / pos > 0; pos *= 10)counting sort(A, n, pos);
104
                                                              C:\WINDOWS\system32\cmd.exe - pause
                                                                                                                                  X
                                                                                                                            int main()
                                                              Enter the value of n : 6
                                                              Enter the element of A: 97 57 208 699 125 734
          ll n, i;
                                                              Output: 57 97 125 208 699 734
           cout << "Enter the value of n : ";</pre>
109
                                                             Press any key to continue . . .
110
           cin >> n;
111
           11 A[n];
112
           cout << "Enter the element of A : ";</pre>
113
           Forn(i, n)cin >> A[i];
114
           radix sort(A, n);
115
           cout << "Output : ";</pre>
           Forn(i, n)cout << A[i] << " ";</pre>
116
117
           cout << endl;</pre>
118
119
           biday;
120
121
```

Time Complexity

* The time complexity of radix sort is given by the formula,

$$T(n) = O(d*(n+b))$$

Where,

- d is the number of digits in the given listn is the number of elements in the list
- > n is the number of elements in the list
- b is the base or bucket size used, which is normally base 10 for decimal representation.

Thank You