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Tutorial - 1

Ans - 1 Asymptotic notation discribe the algorithim efficiency and Performance in a meaningful way. it describe the behaviour of time or space Complexity for large instance characteristics. They are mathematical tool to represent the time complexity of algorithim for asypotic analysis.

There are mainly three asymptotic notations

1 Big - O Notation

The Big O notation defines an upper bound of an algorithim, it bounds a function only from above. for ex - insertion sort. It takes linear time in best case & quadratic time in worst case is $O(n^2)$ So we can say that TC of insertion sort is $O(n^2)$.

$$f(n) = O g(n)$$

$g(n)$ is tight upper bound of $f(n)$

$$f(n) = O(g(n))$$

if $f(n) < C.g(n)$

$\forall \; n \geq n_0$, some constant $C > 0$

2 Omega notation (Ω - notation)

Omega notation Represent the lower bound of the running time of an algorithim. It can be usefull when we have lower bound on time complexity of an algorithim.

Ex → The time complexity of insertion sort can be written as $\Omega(n)$, but is not a very useful information about insertion.

3 Theta notation (Θ notation)

The Theta notation bounds a function from above and below so it define exact asyptotic behaviours

Ex $3n^3 + 6n^2 + 6000 = \Theta(n^3)$

Ans-2 $O(\log n)$

Ans-3 $T(n) = 3T(n-1)$ if $n>0$ otherwise 1

$= 3(3T(n-2))$

$= 3^2 T(n-2)$

$= 3^3 T(n-3)$

$= 3^n T(n-n)$

$T(n) = 3^n$

Ans-4 $T(n) = 2T(n-1)-1$ if $n>0$ otherwise 1

$= 2(2T(n-2)-1)-1$

$= 2^2(T(n-2))-2-1$

$= 2^3(T(n-3))-2^2-2^1-2^0$

$= 2^n(T(n-n))-2^{n-1}-2^{n-2}---2^0$

$T(0) = 1$

$= 2^n-(2^n-1)$

$TC = O(2^n)$

Ans-5 $TC = O(\sqrt{n})$

Ans-6 $TC = O(\sqrt{n})$

Ans-7 $TC = n/2 \times \log n \times \log n$

$= O(n \log^2 n)$

Ans-8 $TC = O(n^3)$

Ans-9

i	j
1	n times
2	n/2 times
3	n/3 times
n	n/n times

TC $\in O(n \log n)$

Ans-10 Since Polynomial grow Smaller than exponential n^k has an asmyptotic upper bound of $O(a^n)$.

for $a = 2$, $n_0 = 2$