

## Tutorial - 3

Ans-1 Pseudocode for linear search is

```
int linear ( int * arr, int n, int key ) {  
    for ( int i = 0; i < n; i++ ) {  
        if ( arr[i] == key ) {  
            return i;  
        }  
    }  
    return -1;  
}
```

Ans-2 Pseudocode of insertion sort

```
void insertion ( int arr[], int n ) {  
    for ( int i = 1; i < n; i++ ) {  
        int key = arr[i];  
        int j = i - 1;  
        while ( j >= 0 && arr[j] > key ) {  
            arr[j+1] = arr[j];  
            j--;  
        }  
        arr[j+1] = key;  
    }  
}
```

Ans-3 Average case complexity of sorting algos

- Bubble sort  $\rightarrow O(n^2)$
- insertion sort  $\rightarrow O(n^2)$
- Selection sort  $\rightarrow O(n^2)$
- Merge sort  $\rightarrow O(n \log n)$
- Quick sort  $\rightarrow O(n \log n)$
- heap sort  $\rightarrow O(n \log n)$

Ans-4

Stable  $\rightarrow$  (appears in same order)

Bubble, insertion, merge

Inplace  $\rightarrow$  (using constant space)

Bubble, selection, insertion, heap

Ans-5

Pseudocode for Binary Search

```
int BS(int arr[], int n, int key) {
```

```
    int low = 0, high = n-1;
```

```
    while (low <= high) {
```

```
        int mid =  $\left\lfloor \frac{low + high}{2} \right\rfloor$ 
```

```
        if (key == arr[mid]) {
```

```
            return mid;
```

```
        }
```

```
        else if (key > arr[mid]) {
```

```
            start = mid + 1;
```

```
        }
```

```
        else {
```

```
            end = mid - 1;
```

```
        }
```

```
    }
```

```
    return -1;
```

```
}
```

TC  $\rightarrow O(\log n)$

SC  $\rightarrow O(1)$

Ans-6

Recurrence relation of Binary Search is

$$T(n) = T(n/2) + 1$$

Ans-7 Assuming given array is sorted then we can find sum in linear complexity.

```
int sum (int arr[], int target) {  
    int i = 0; j = n-1;  
    while (i < j) {  
        if (arr[i] + arr[j] == target)  
            return 1;  
        else if (arr[i] + arr[j] > target)  
            j--;  
        else  
            i++;  
    }  
    return -1;  
}
```

Ans-8 Quick sort is the best sorting algorithm in practical use as it follows the locality of reference and also its best case time complexity is  $O(n \log n)$ .

Ans-9 No of inversion tells us how far the array is from being sorted.

if ( $arr[i] > arr[j]$  &&  $i < j$ )

→ 4, 7, 21, 31, 8, 10, 1, 20, 6, 45

$$\begin{aligned} \text{inversion} &= 4 + 7 + 7 + 4 + 4 + 3 + 2 \\ &= \underline{31} \end{aligned}$$

Ans-10 Quick Sort

Best Case → When array is totally unsorted.

Worst Case → When array is sorted or reverse sorted.

Ans-11

Merge Sort

Best

$$2T(n/2) + O(n)$$

Worst

$$2T(n/2) + O(n)$$

Quick Sort

$$T(n) = T(k) + T(n-k-1) + O(n)$$

$$T(n) = T(n-1) + O(n)$$

Similarity → Both are based on divide and conquer technique.

Differences → Worst case TC of merge sort is  $O(n \log n)$  while of quick sort is  $O(n^2)$

Ans-12

Void Stable Selection Sort(int arr[], int n) {

for(int i=0; i<n-1; i++) {

int min = i;

for(int j=i+1; j<n; j++) {

if(arr[min] > arr[j]) {

min = j;

}

}

int key = arr[min];

while(min > i) {

arr[min] = arr[min-1];

min--;

}

arr[i] = key;

}

}

Ans-13

### Optimised Bubble Sort

```
for (int i=0; i<n; i++) {  
    bool swap = false;  
    for (j=0; j<n-i-1; j++) {  
        if (arr[j] < arr[j+1]) {  
            swap(arr[j], arr[j+1]);  
            swap = true;  
        }  
    }  
    if (!swap)  
        return;  
}
```

Ans-14 In such case, merge sort would be efficient as it is an External Sorting algorithm. i.e. data is divided into chunks and then sorted using merge sort.

→ Sorted data is dumped into files.

- Internal Sorting → It is a type of sorting technique in which whole sorting takes place in main memory of computer.

