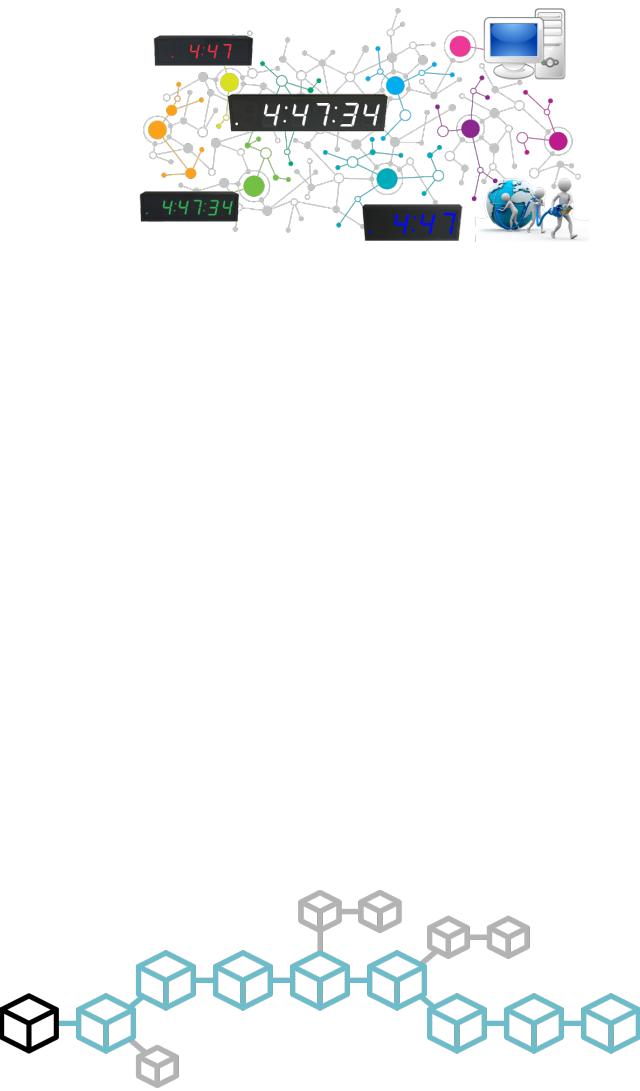
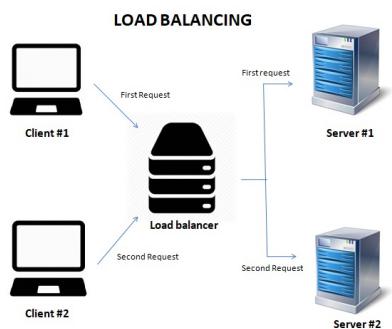
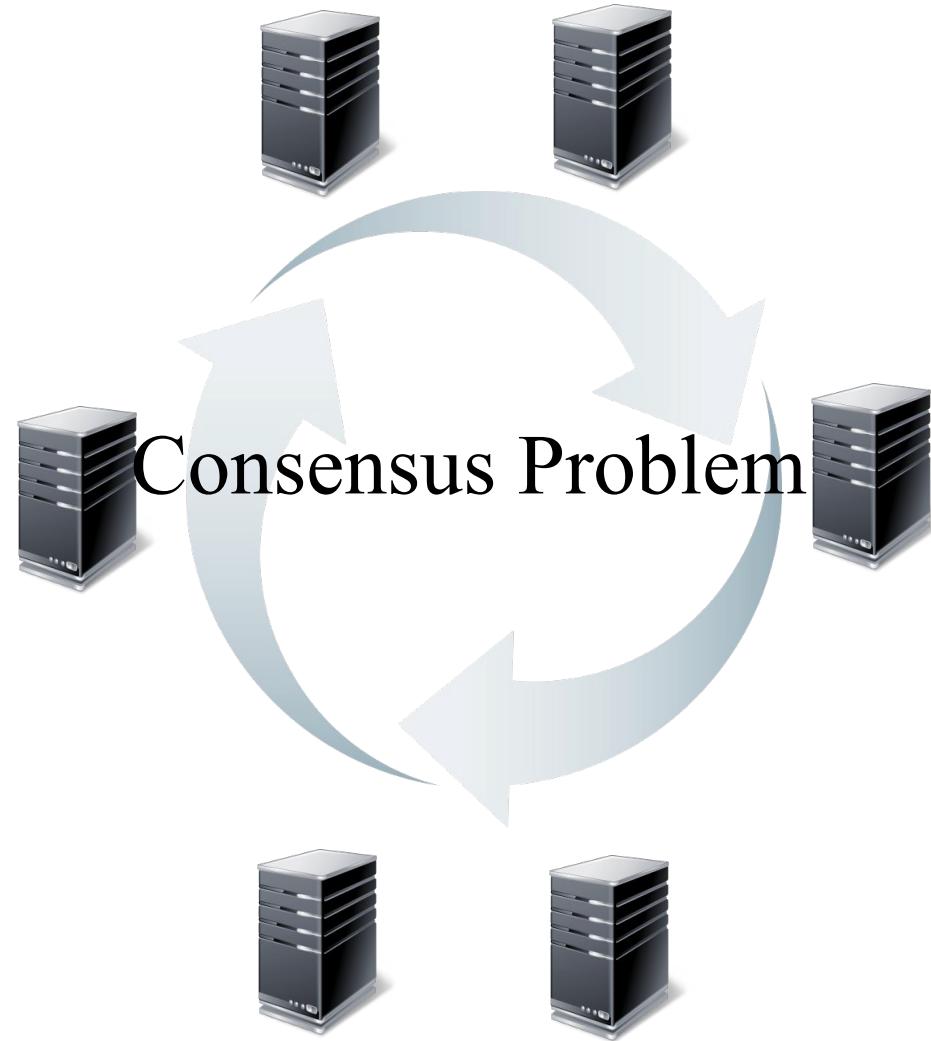
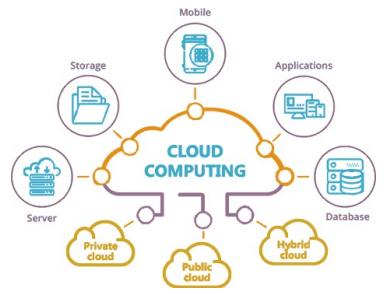


Blockchain Security and Scalability

Tuyet Duong



Consensus Problem

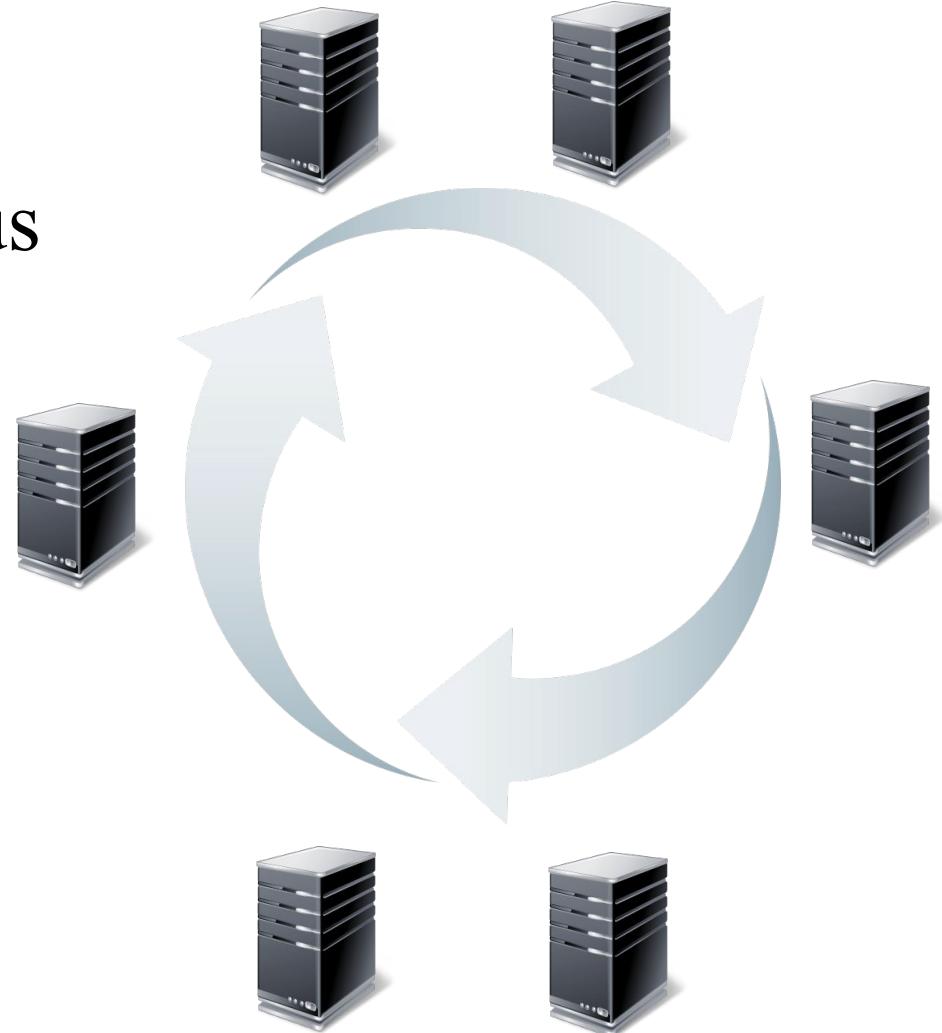
Problem

N servers/nodes, $f < N$ are malicious

Propose and agree on one value

Termination : everyone decides

Agreement : all non-malicious
agree



Classical Consensus Protocol



Classical Consensus Protocol

✗ Assumption of known identity set

✗ Bandwidth limited

- $O(n^2)$ messages (e.g. PBFT)
- Work for a small network (e.g $n < 100$)

How to Build a Secure and Scalable Consensus Protocol?

Blockchain Consensus Protocol

Bitcoin

- Nakamoto 2008
- Maintained by network participants (miners)



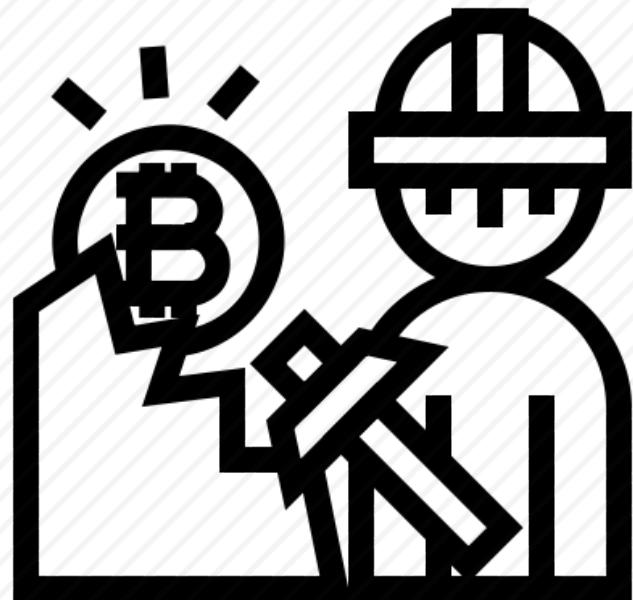
Bitcoin Network

Bitcoin Blockchain Consensus Protocol

How to pick a leader?



Bitcoin Blockchain Consensus Protocol



How to pick a leader?

Proof of Work Puzzle

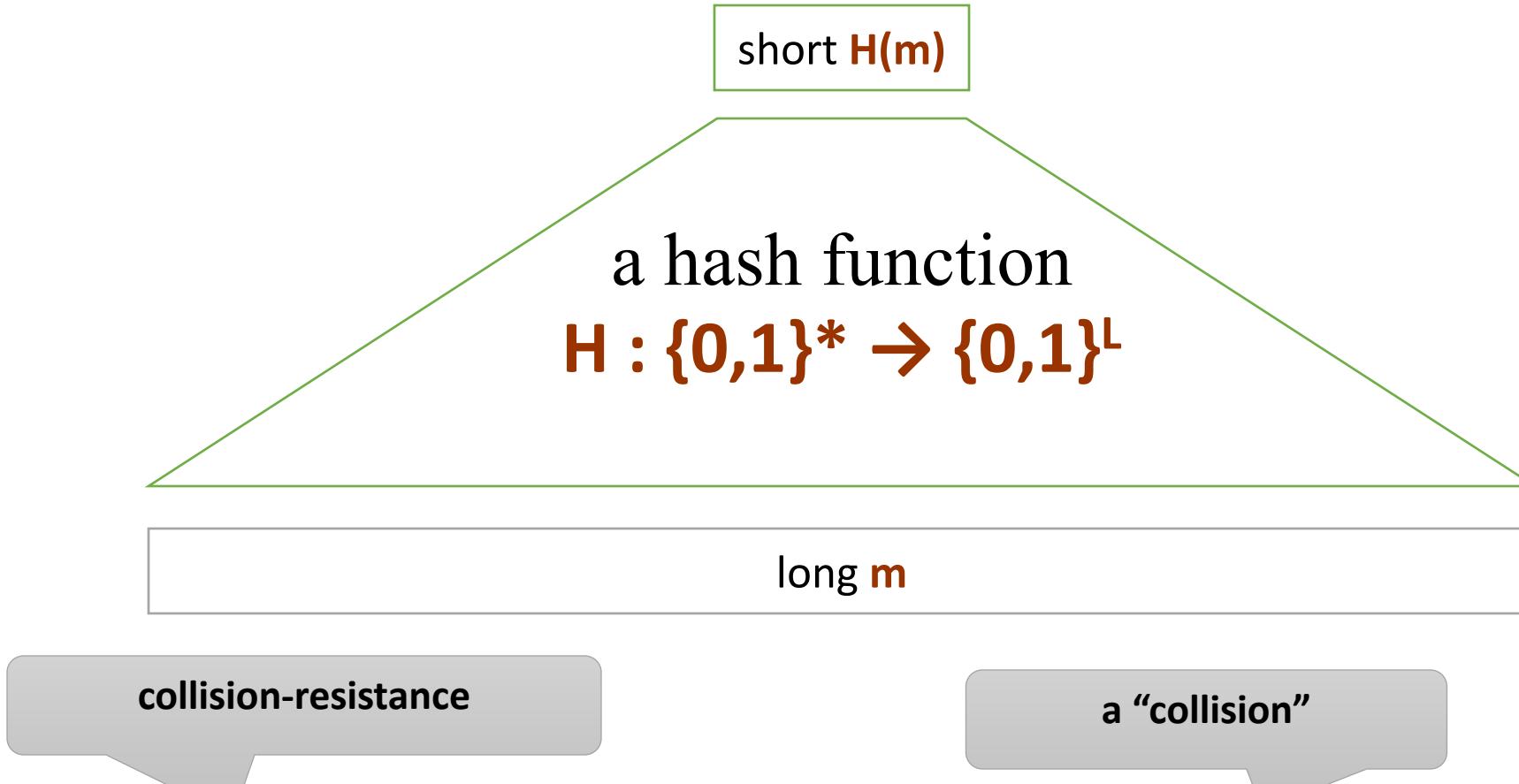
by Cynthia Dwork and Moni Naor in 1993

Recap

- Two different directions to solve the consensus problem
 - Committee-based
 - **Leader-based**

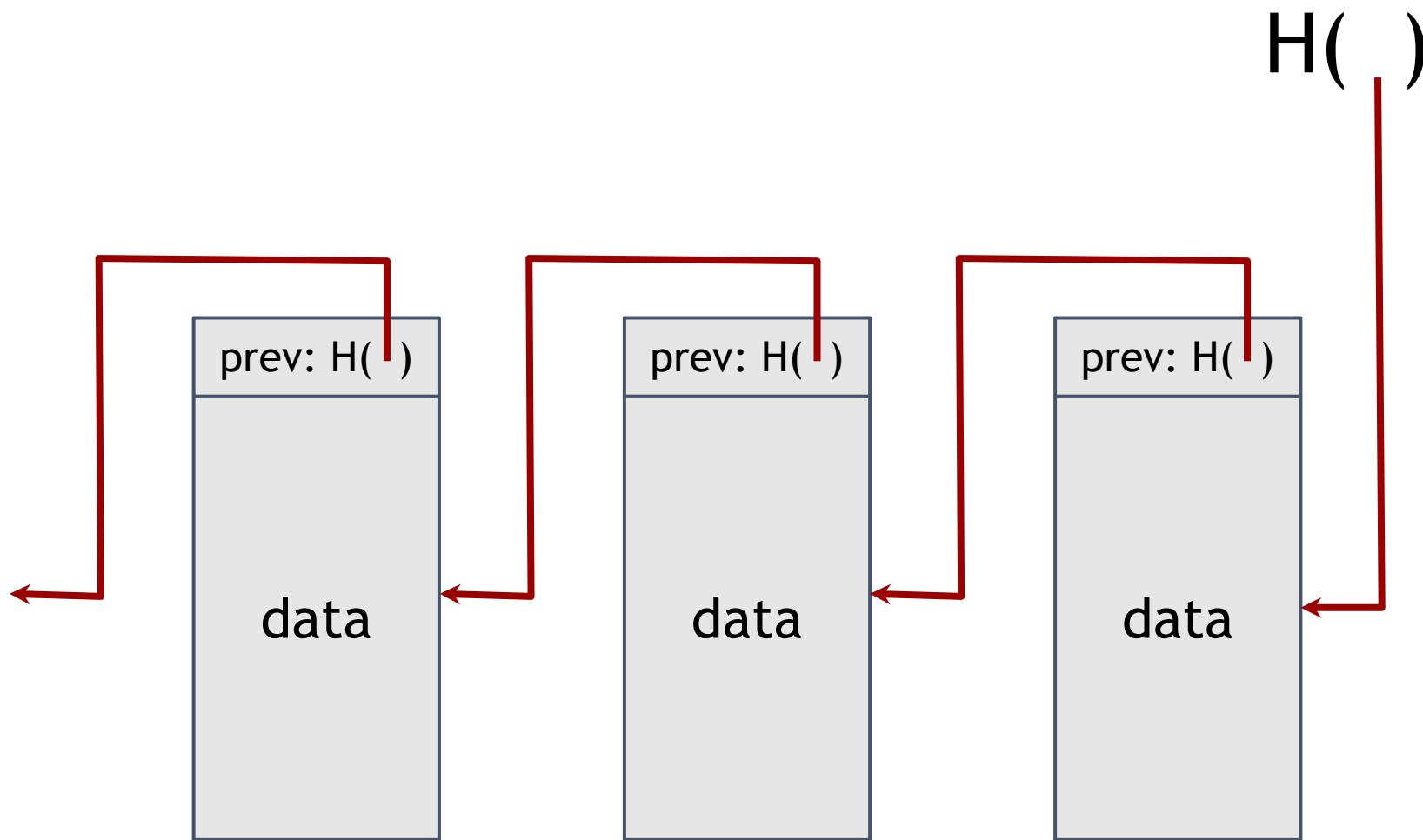
Hash Function and Blockchain Structure

Hash Function



Requirement: it should be hard to find a pair (m, m') such that
 $H(m) = H(m')$

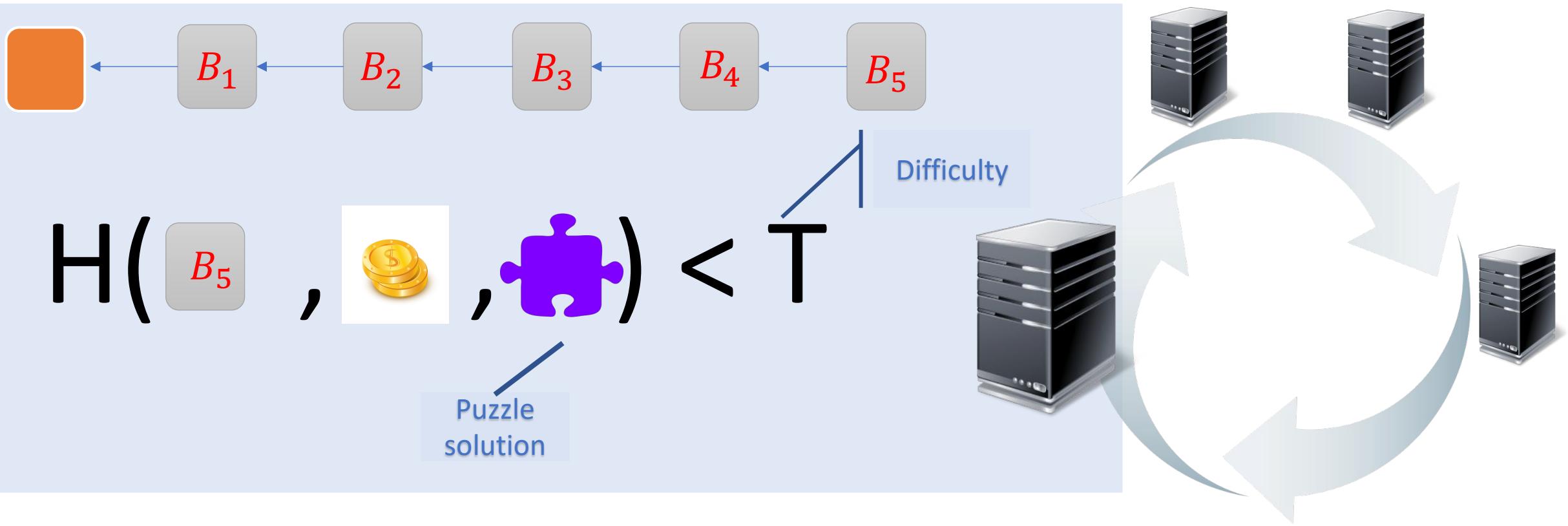
linked list with hash pointers = “block chain”



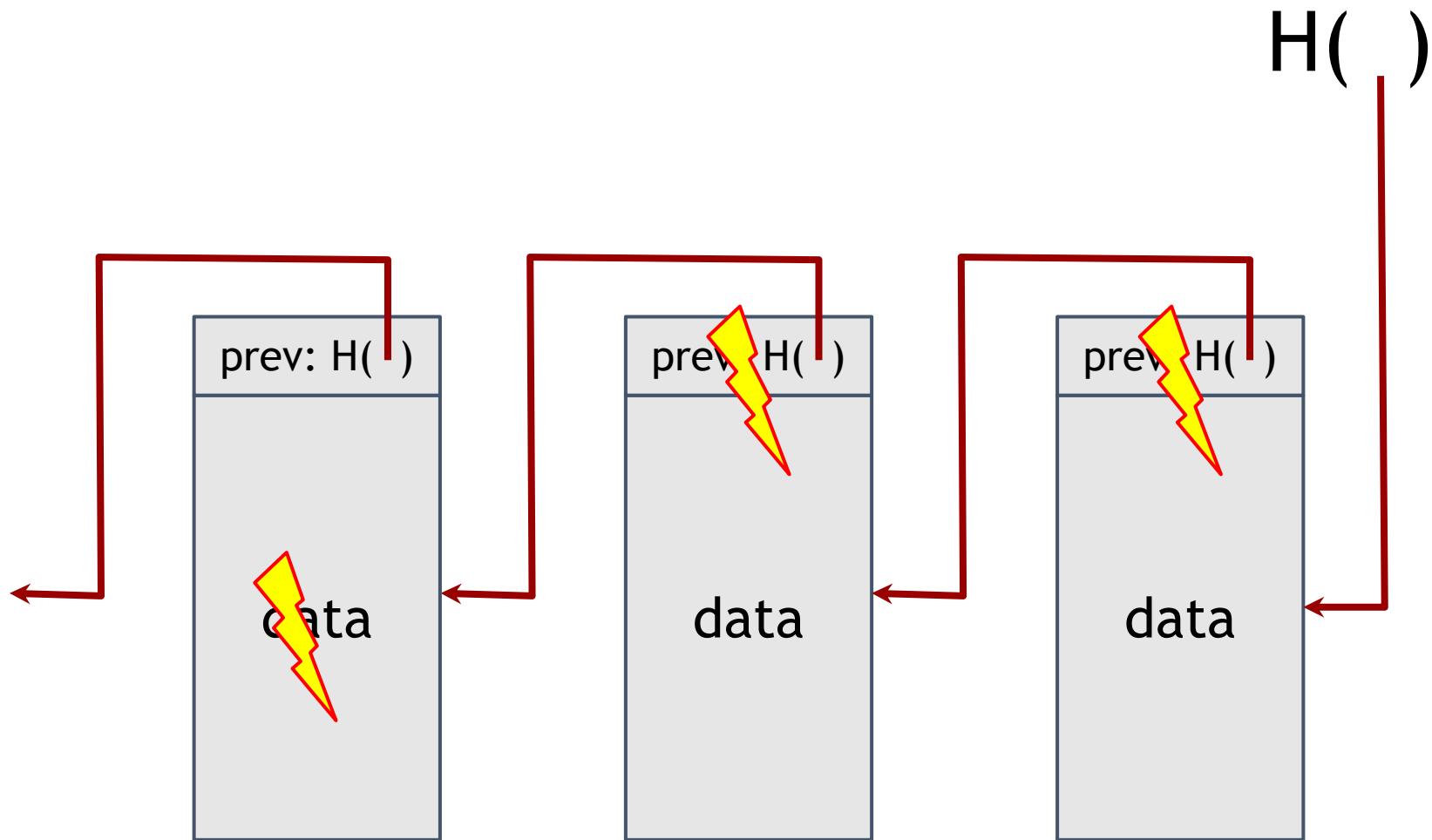
Proof-of-Work Puzzle

Bitcoin Blockchain Consensus Protocol

[Nak08]



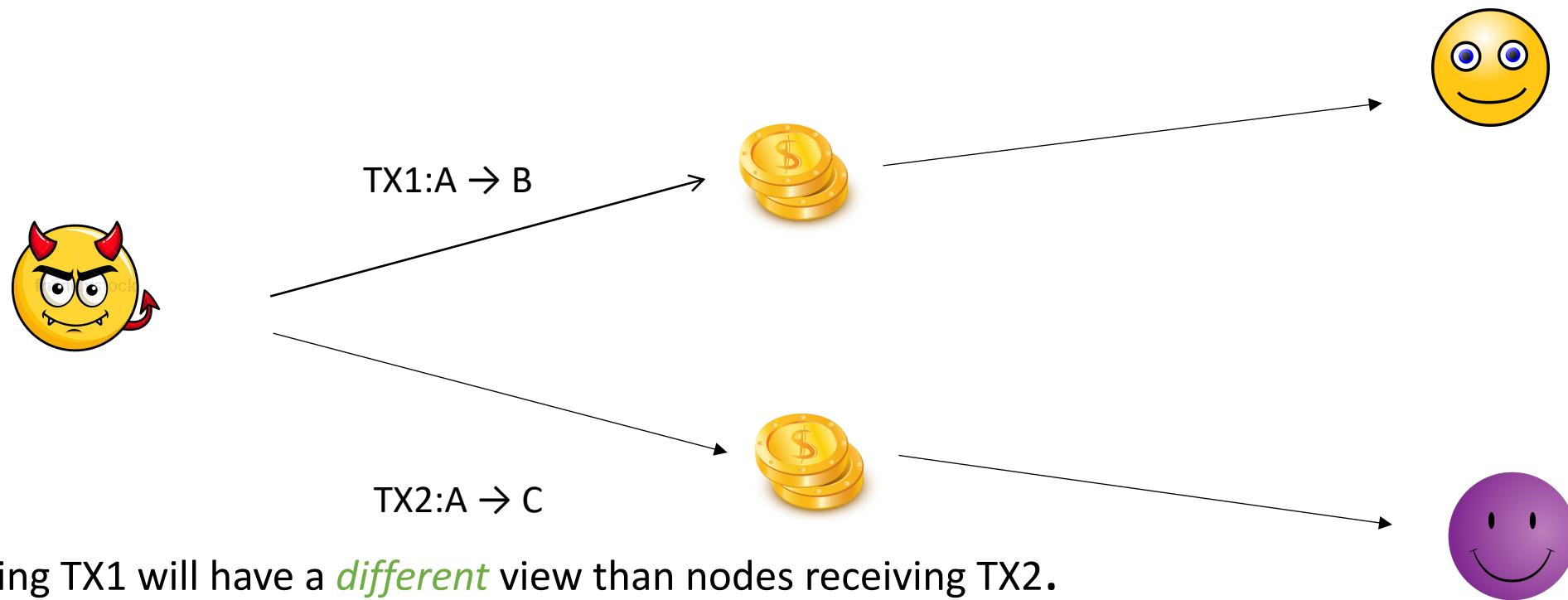
detecting tampering



use case: tamper-evident log

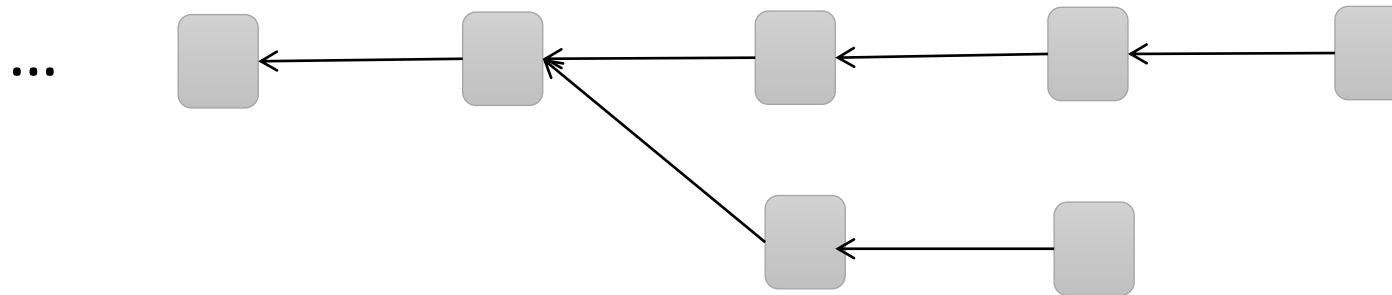
Bitcoin Blockchain Consensus Protocol: Fork

[Nak08]



Bitcoin Blockchain Consensus Protocol

[Nak08]



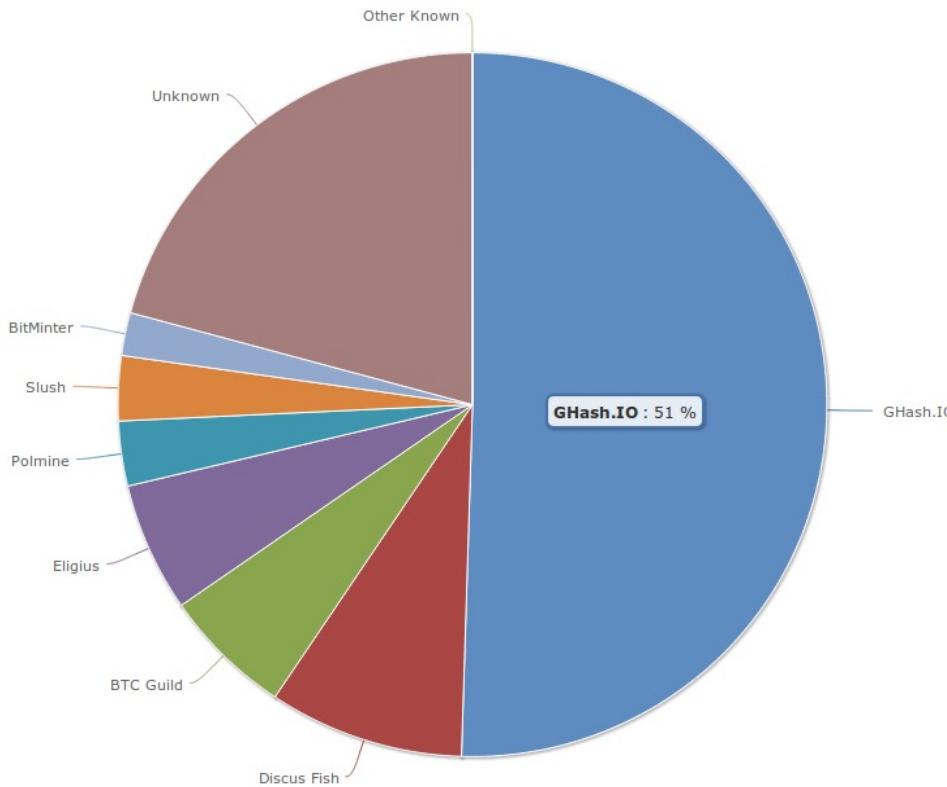
Honest nodes “believe” the longest chain

Assumption: *honest nodes control the majority of computing power*

Security and Scalability Problems and my Contributions

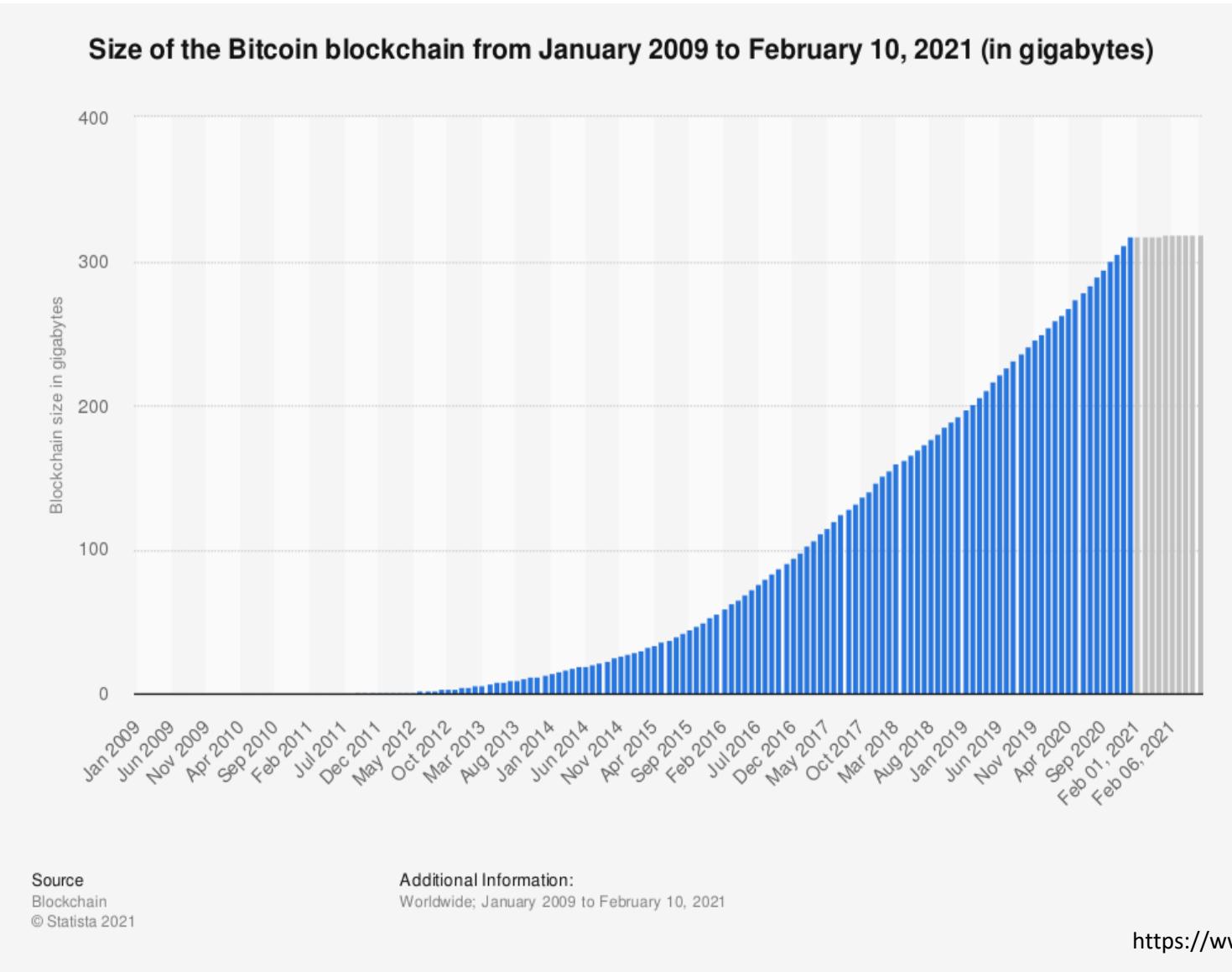
Problem 1 : Security

**51% Honest Mining Power Assumption
could be challenged**



June 12, 2014
GHash.IO large mining pool crisis
<https://blockchain.info/pools>

Problem 2: Storage Scalability



Problem 2: Throughput Scalability



3-7 TXs per second

Demand from Practice: 1,200 - 50,000 TXs/s

PayPal™



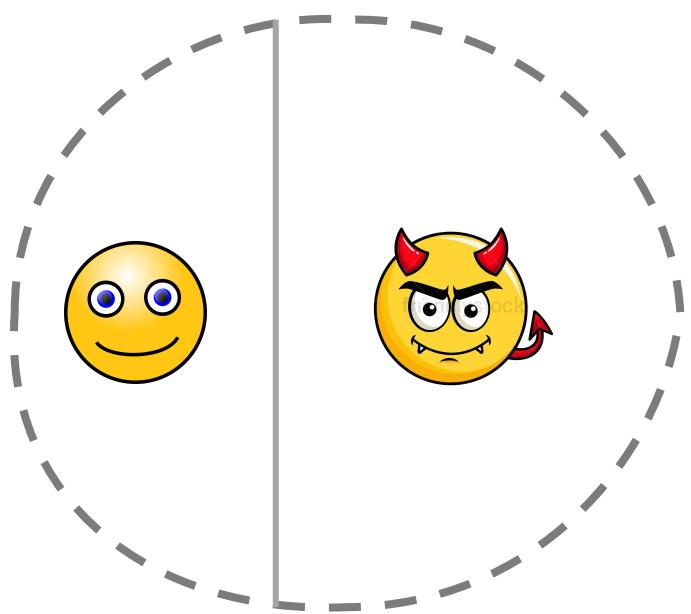
Recap

- Problems:
 - Security Assumption
 - Storage and Throughput Scalability

Solving Problem 1 : Security in the presence
of the 51% Attack

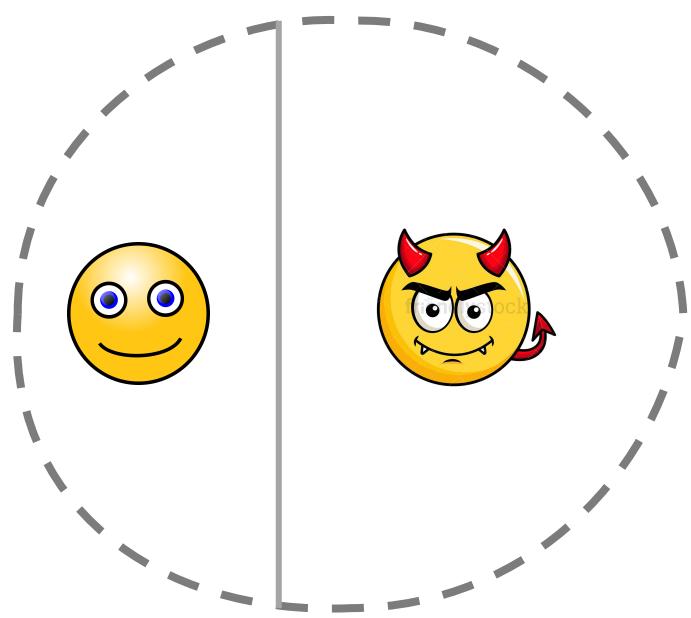
2-hop Protocol: Idea

Computing power

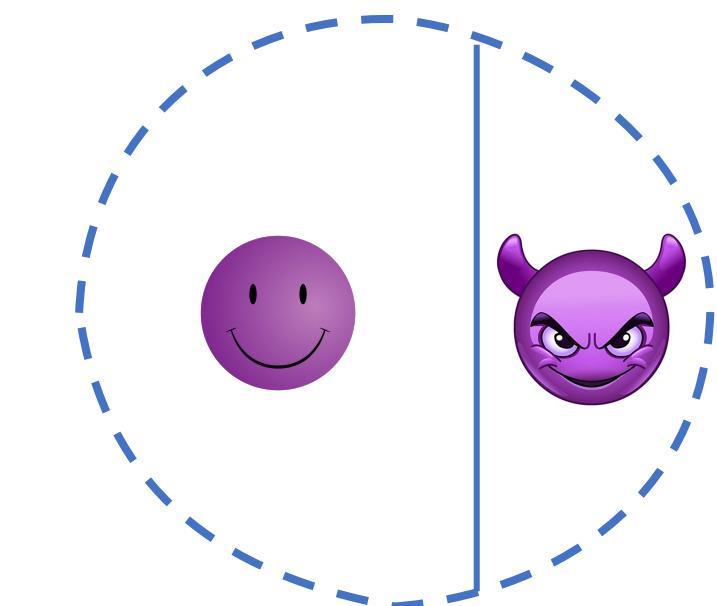


2-hop Protocol: Idea

Computing power

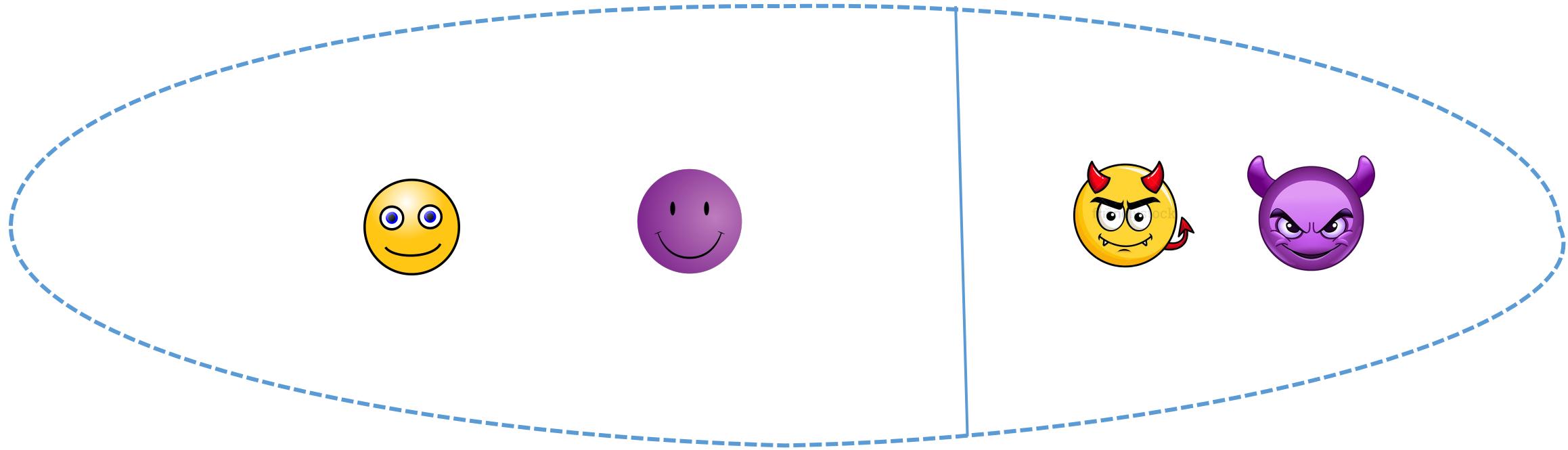


Stake



2-hop Protocol: Idea

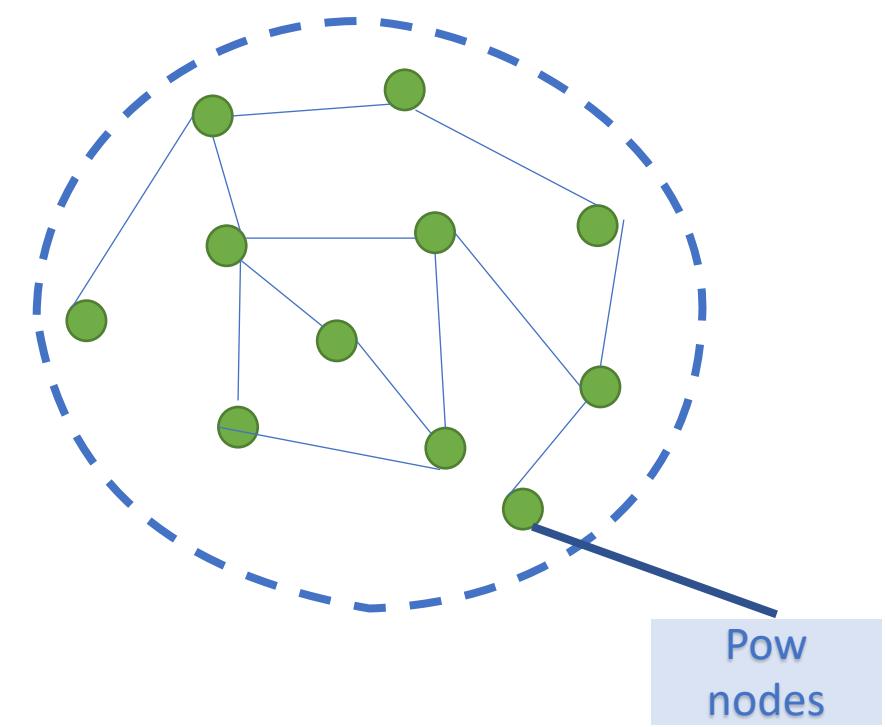
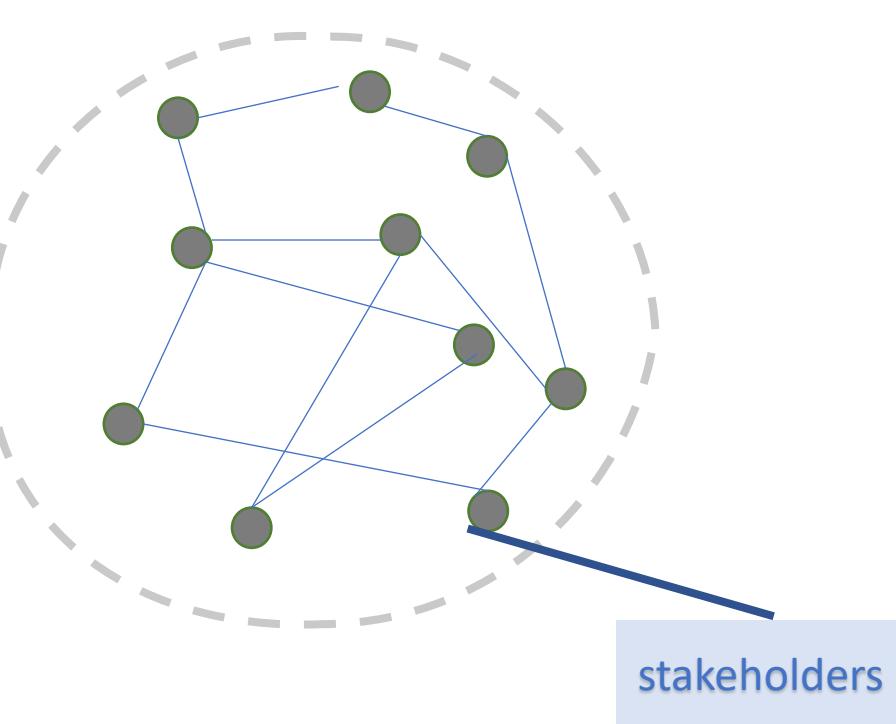
Computing power and Stake



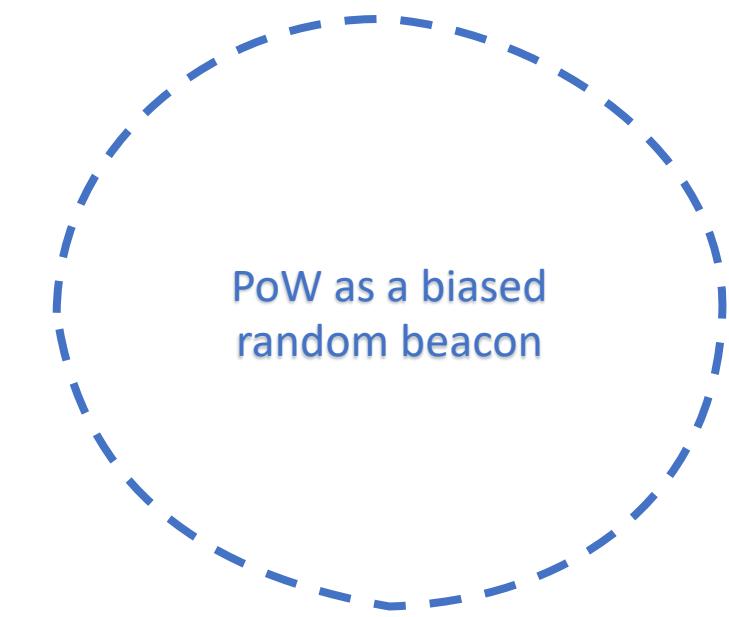
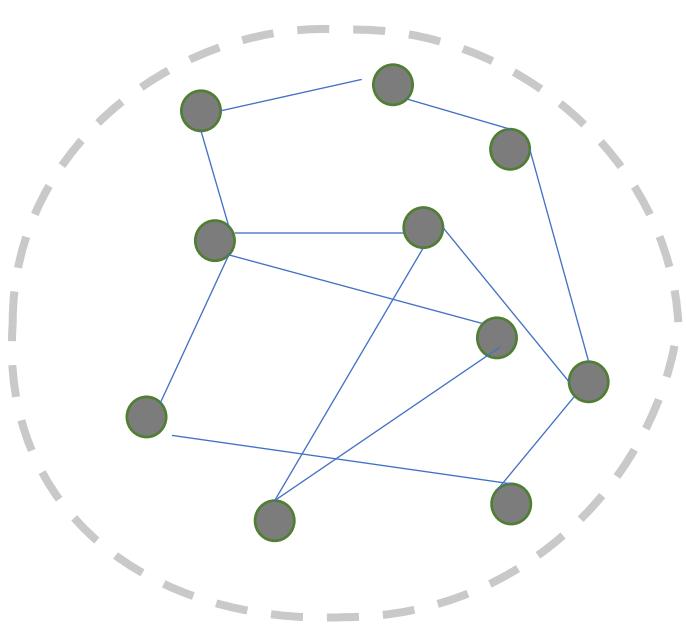
2-hop Protocol: Challenge

How to design a Proof-of-Stake Puzzle?

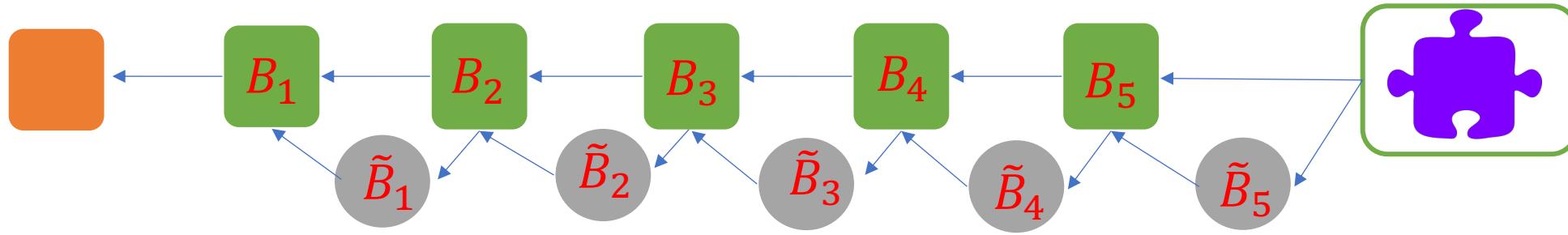
2-hop Protocol: Challenge



2-hop Protocol: Challenge



2-hop Protocol: Blockchain extension

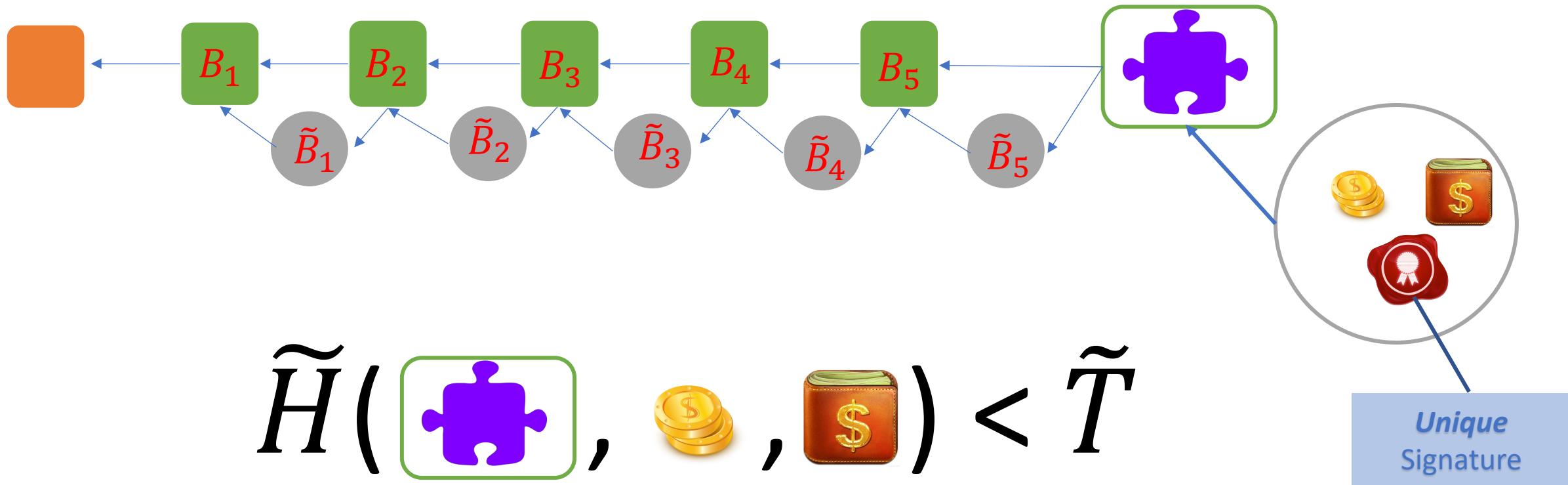


$$\tilde{H}(\boxed{\text{puzzle piece}}, \text{coins}, \text{wallet}) < \tilde{T}$$

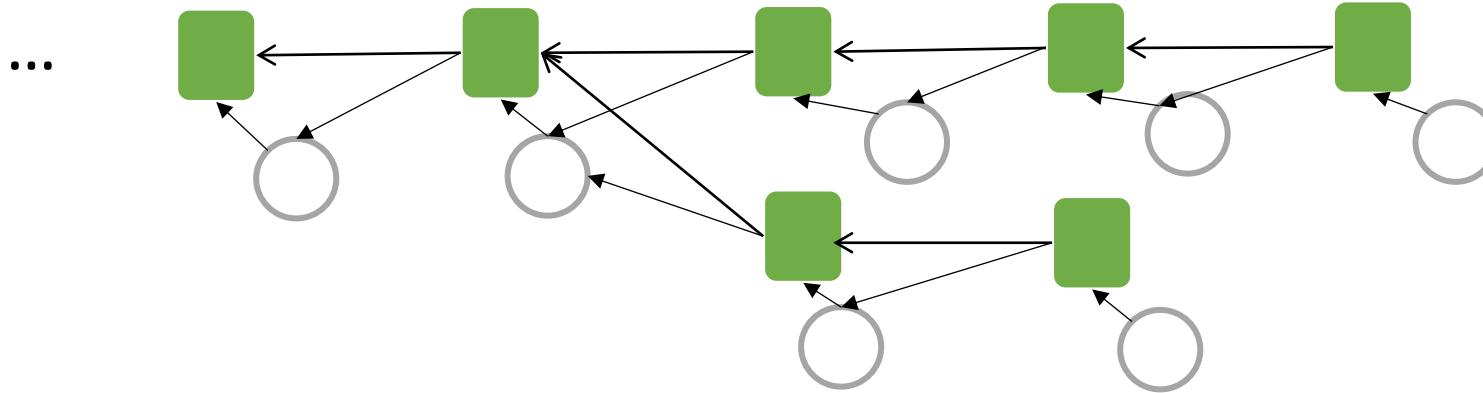
PoW as a biased random beacon

Stakeholder identity

2-hop Protocol: Blockchain Extension

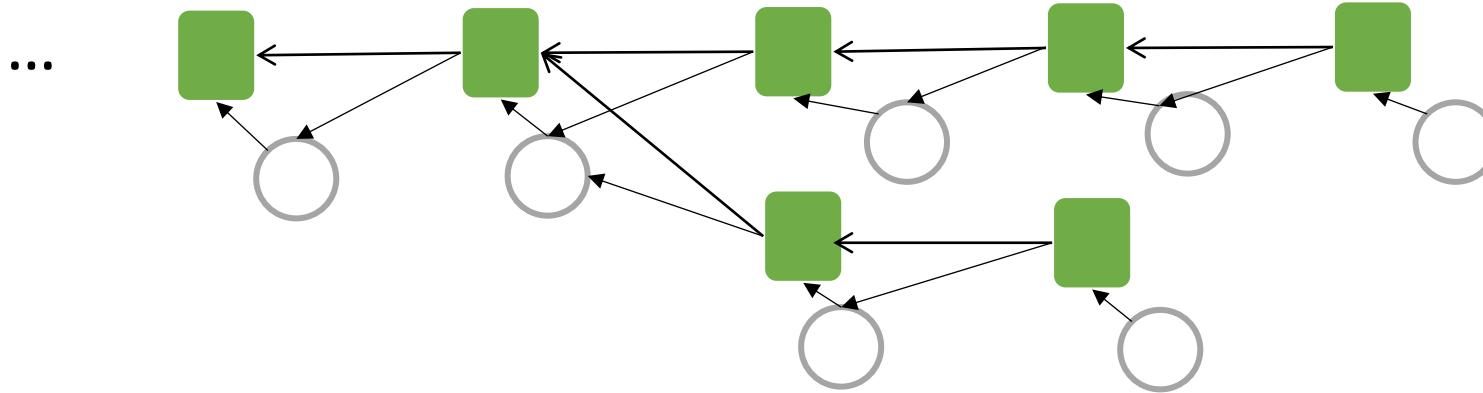


2-hop Protocol: Blockchain Extension



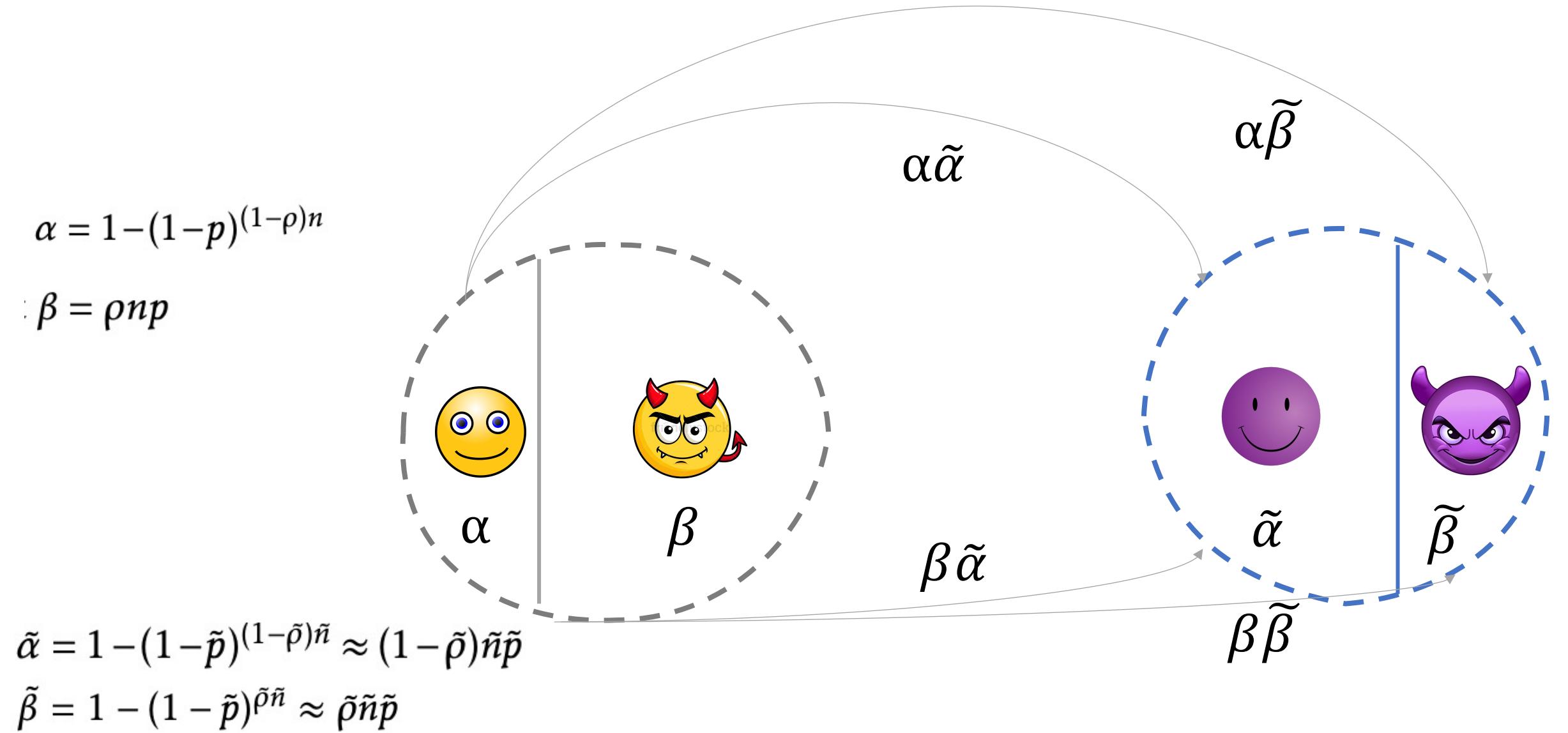
Honest nodes “believe” the chain-pair with the longest PoW chain

2-hop Protocol: Blockchain Extension



Assumption: *honest nodes control the majority of collective resources*

Security Analysis



Security Analysis

Computing power and Stake

$$\alpha \tilde{\alpha}$$

$$\alpha \tilde{\beta} + \beta \tilde{\beta}$$

Security Analysis

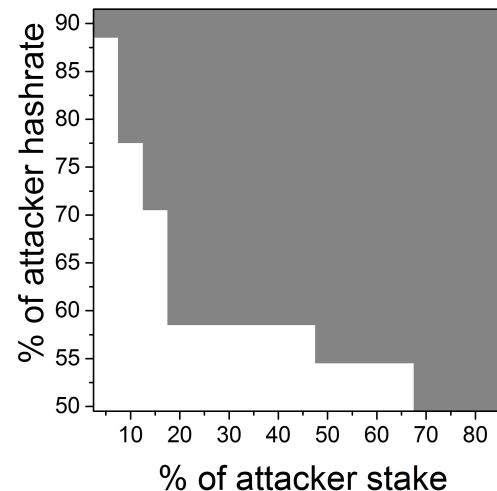
1. Define events of interest, i.e., # honest vs malicious blocks
2. Derive expectations for the events of interest
3. Deploy concentration theory, e.g., Chernoff bound. In t rounds

$$\Pr[Y < (1 - \delta''')\alpha\tilde{\alpha}t] < e^{-\Omega(t)} \quad \Pr[Z > (1 + \delta')(\alpha + \beta)\tilde{\beta}t] < e^{-\Omega(t)}$$

$Y > Z$ with probability at least $1 - e^{-\Omega(t)}$ under the constrain that $\alpha\tilde{\alpha} > \alpha\tilde{\beta} + \beta\tilde{\beta}$

Experiments

Use Scorex framework (<https://github.com/input-output-hk/Scorex>)
<https://bitbucket.org/twincscoin/twincschain>



- 70% of total mining power an adversary also needs for about 20% of total stake to generate a better chain than honest party's.
- 20% of stake is about \$2.3 billion.

Thank you!