

Petri Nets

Lesson 4

Special thanks to J.P. Bahsoun

Petri nets: Model parallel activities

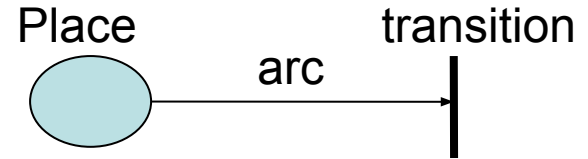
Definition:

- Oriented Bipartite Graph, consists of two sets of vertices (place, transitions) and one set of arcs.

Representation:

- An arc (arrow) links a place (resp. a transition) to a transition (resp. a place).
- A place expresses a state, an activity, a condition, a precondition or a postcondition
- A transition expresses an event

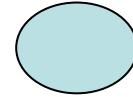
Places and transitions are numbered or labeled



Place

A place has a name and expresses its semantics according to the presence of tokens

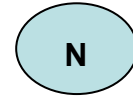
Empty canal



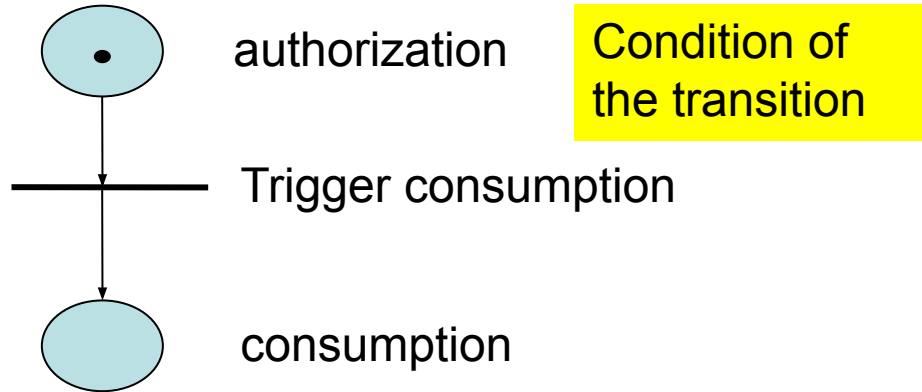
Canal of capacity 3



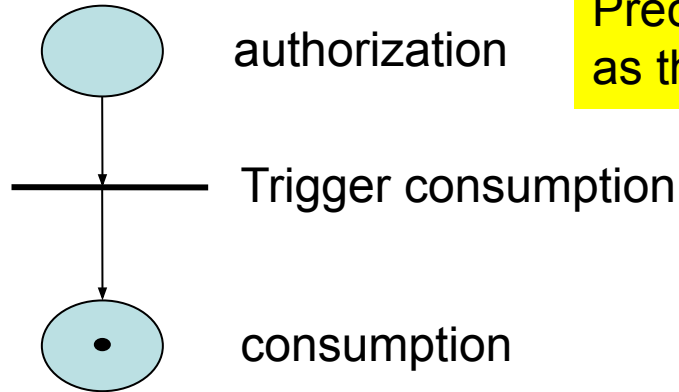
Canal of capacity N



Example of Transition



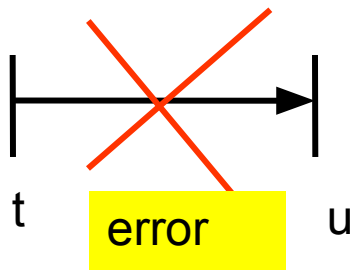
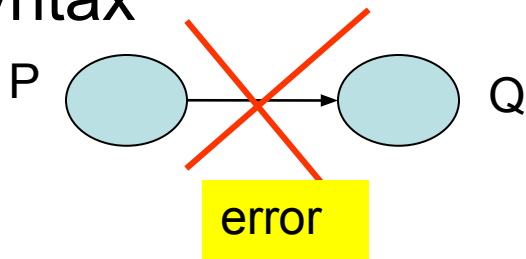
Example of Transition



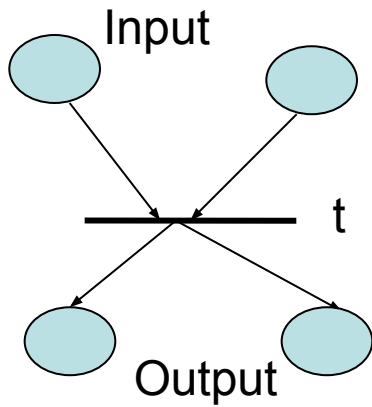
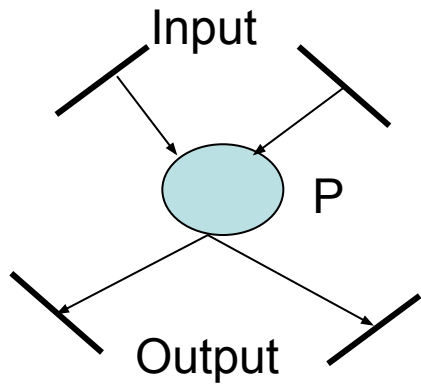
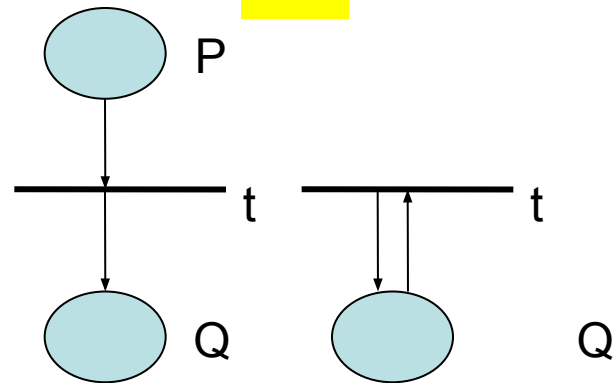
Precondition not fulfilled
as there is no token

Disarmed transition, cannot
be activated

Syntax



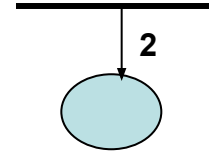
OK



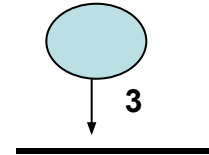
Syntax

An arc has a default weight of 1.

Generates 2
tokens

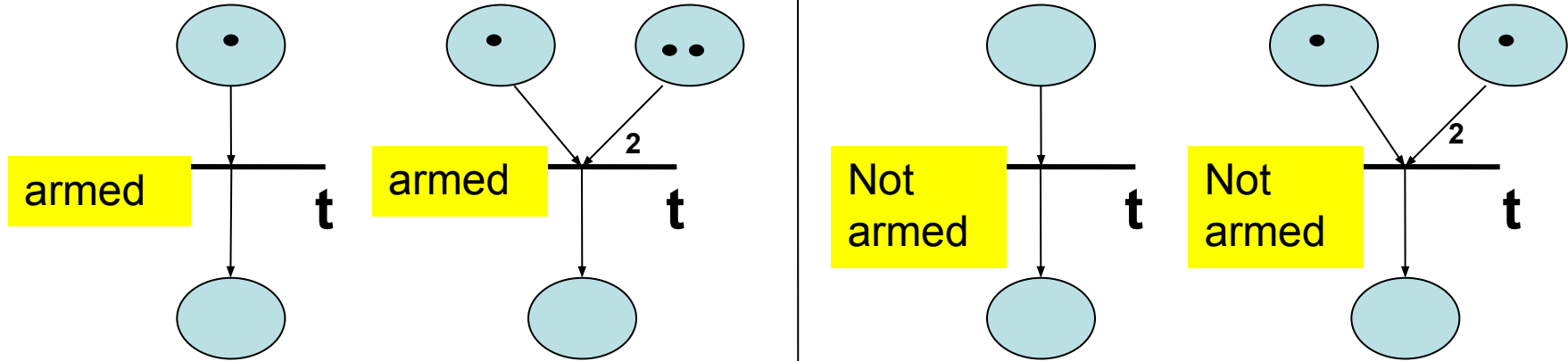


At least 3 tokens needed to arm
the transition



Semantic

A transition **t** is armed if each input place **P_i** contains tokens of number \geq the weight of the output arc of **P_i** that link it to **t**



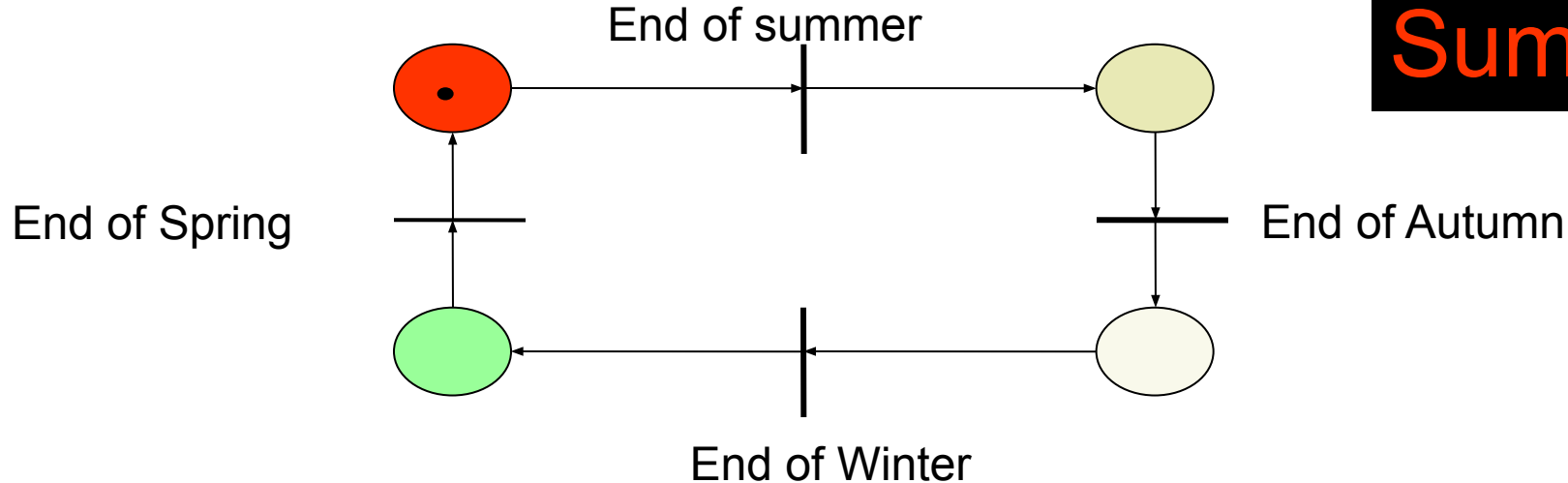
Dynamic modeling



One transition is triggered at a time

The triggering of a transition:

- eventually leads to the circulation of tokens
- eventually leads the RoP from one state to another



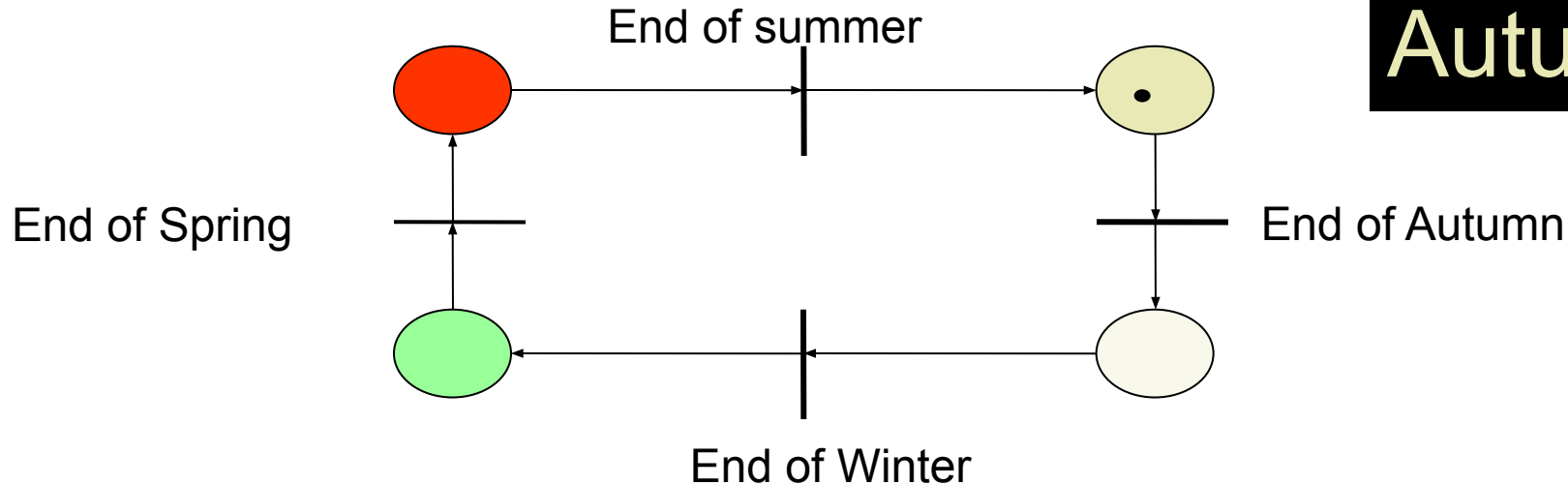
Dynamic modeling



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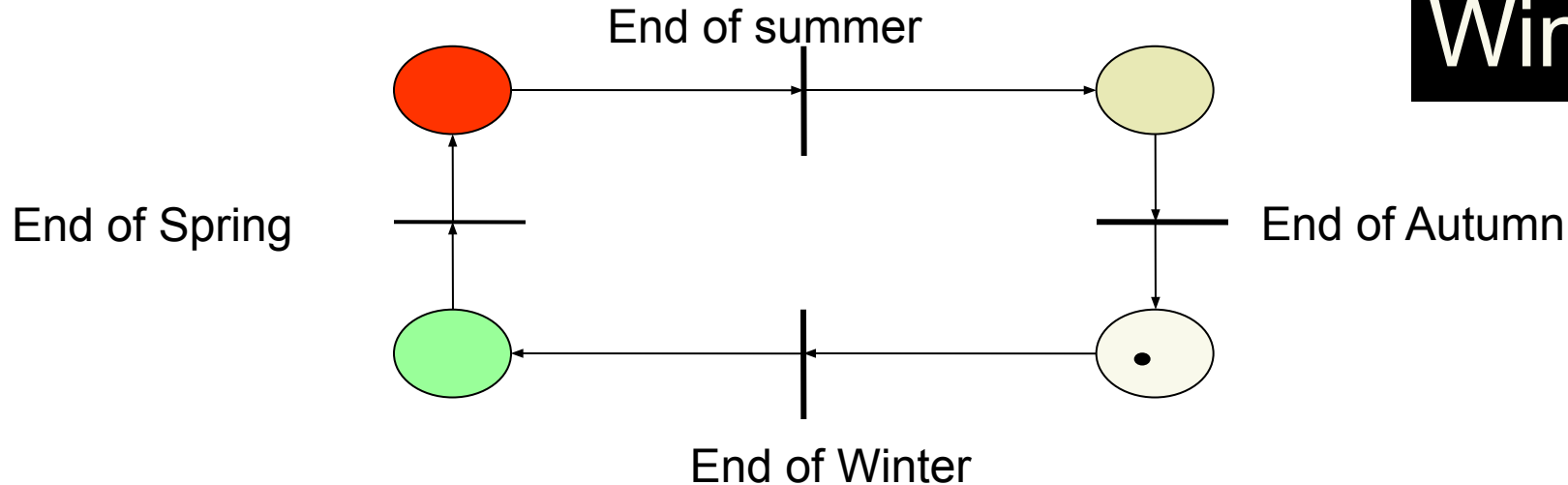
Dynamic modeling



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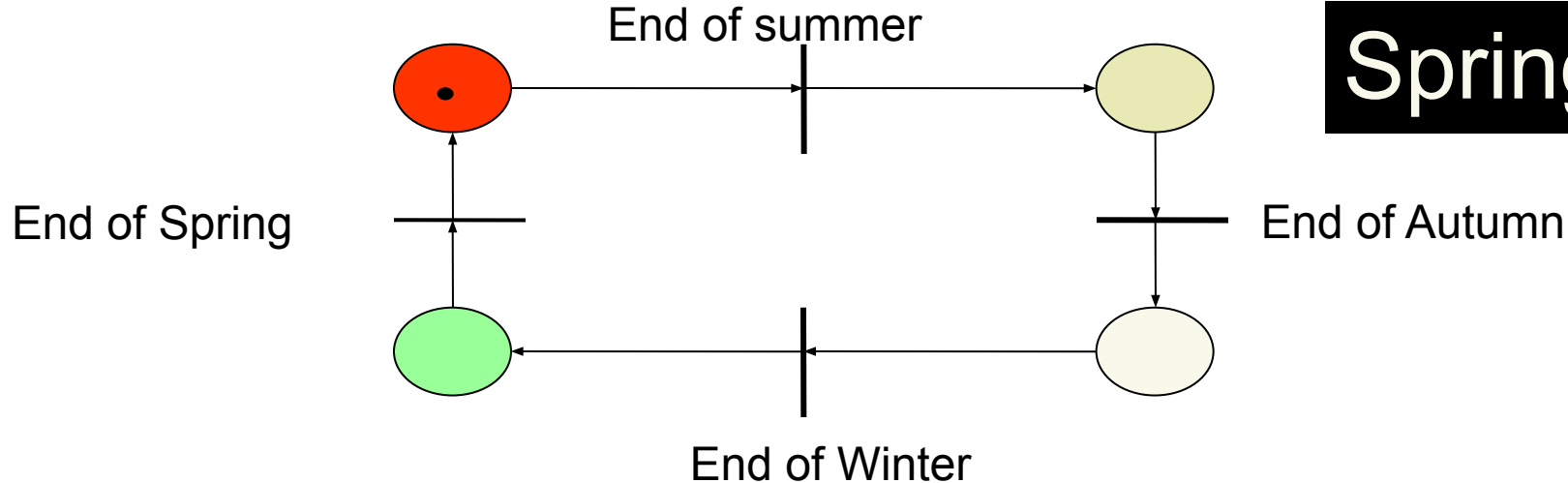
Dynamic modeling



One transition is triggered at a time

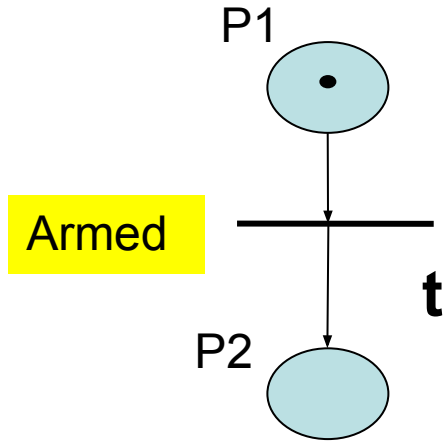
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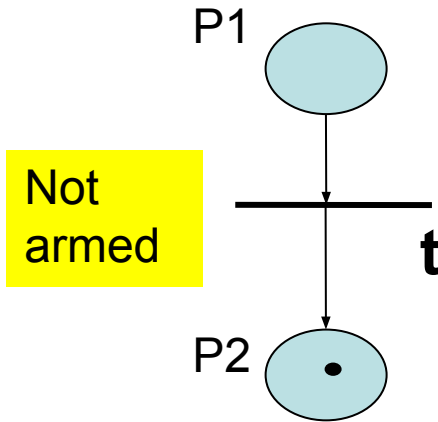


Examples

Before

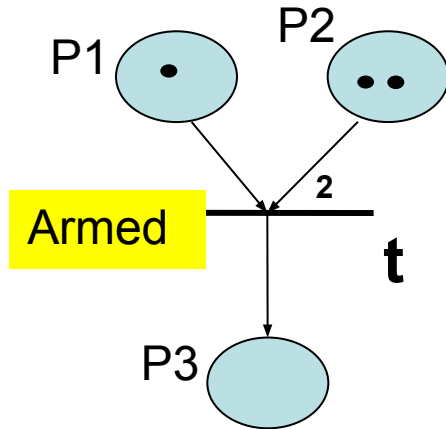


After

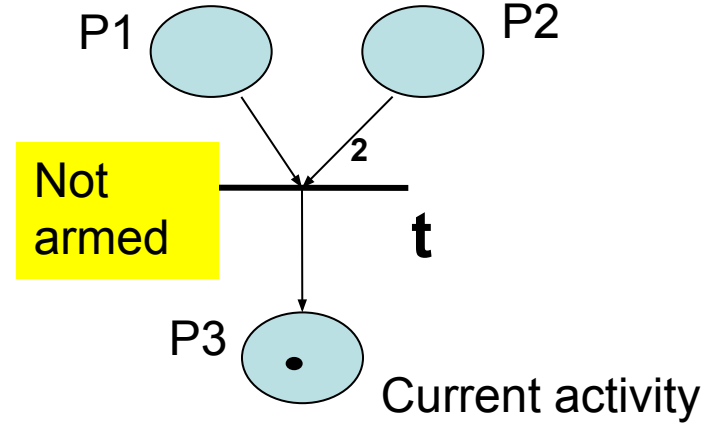


Examples: multiple pre-conditions

Before

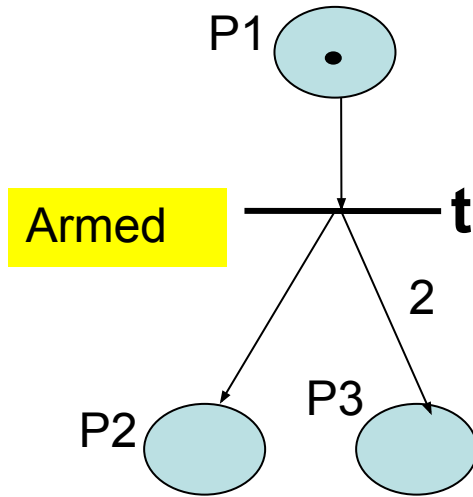


After

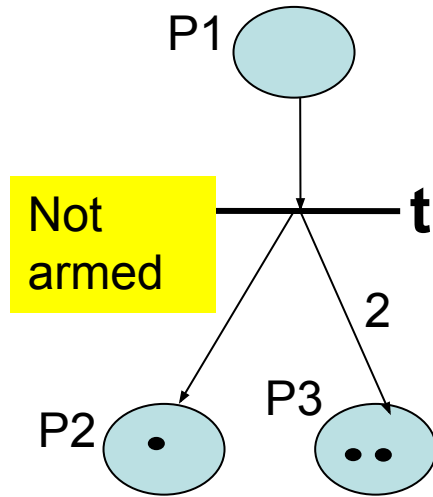


Examples: Diffusion

Before

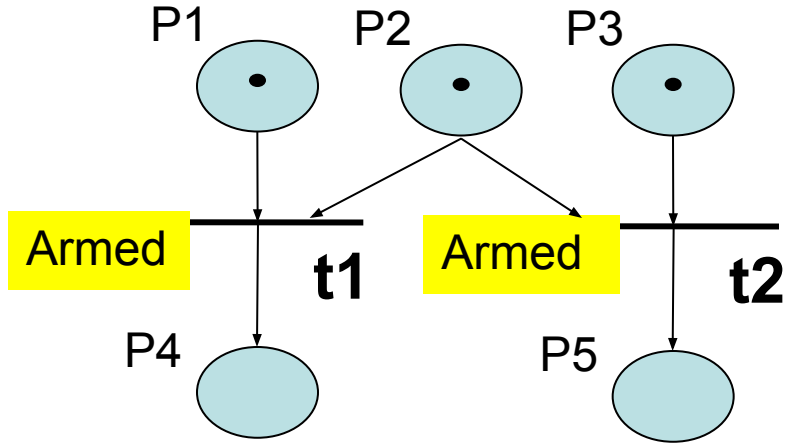


After



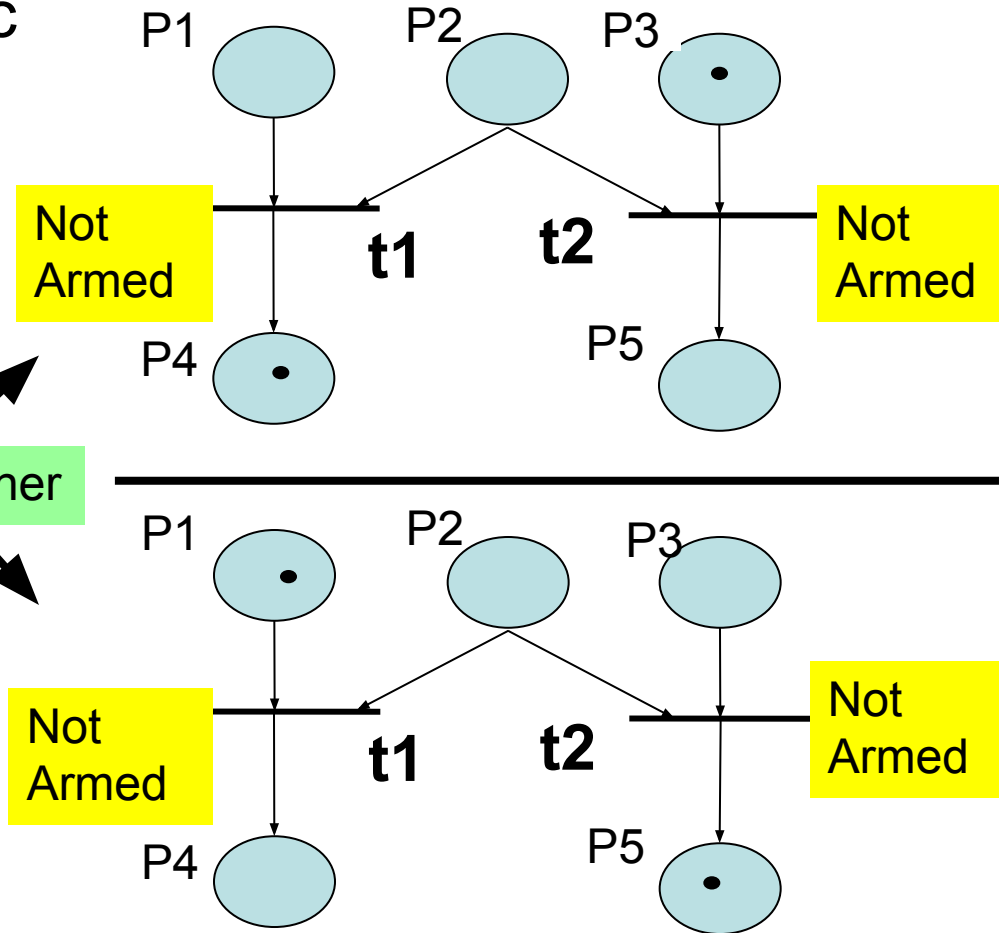
Examples: Non-deterministic

Before

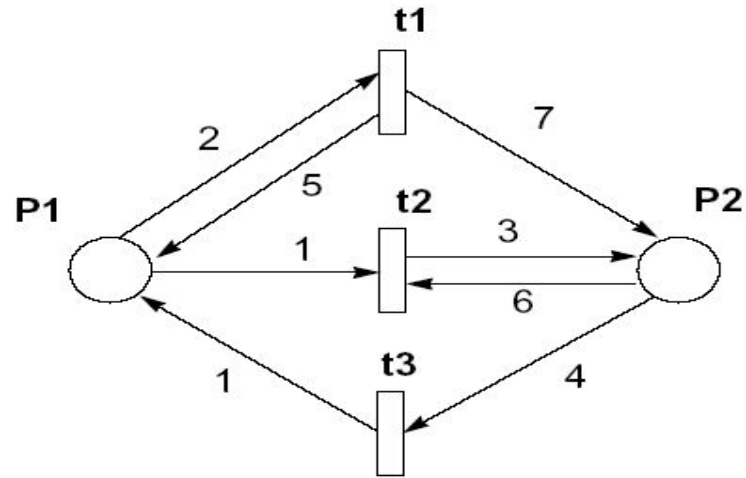


Either

After



Matrix representation



Input

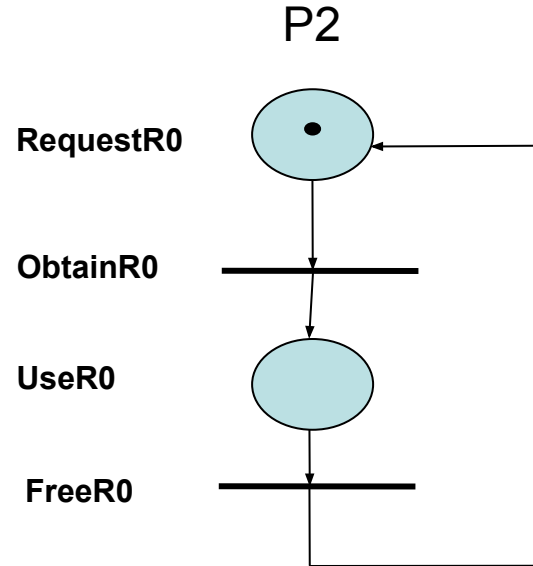
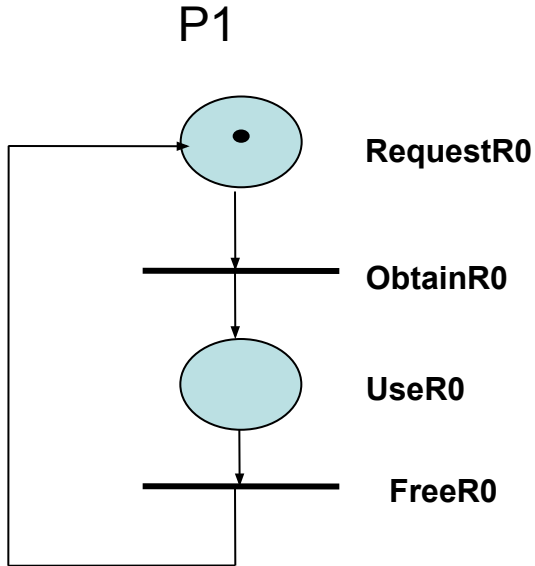
	$t1$	$t2$	$t3$
$P1$	2	1	0
$P2$	0	6	4

Output

	$t1$	$t2$	$t3$
$P1$	5	0	1
$P2$	7	3	0

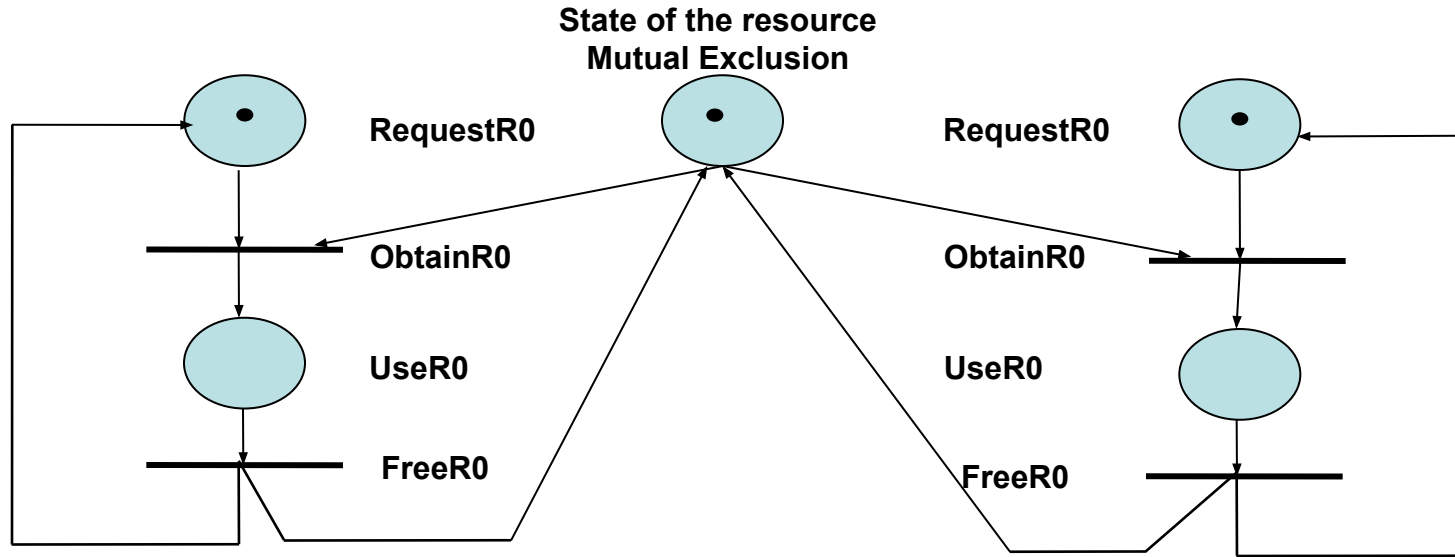
Example: Model two cyclic processes that share a resource in mutual exclusion

User process modeling. What about R0?

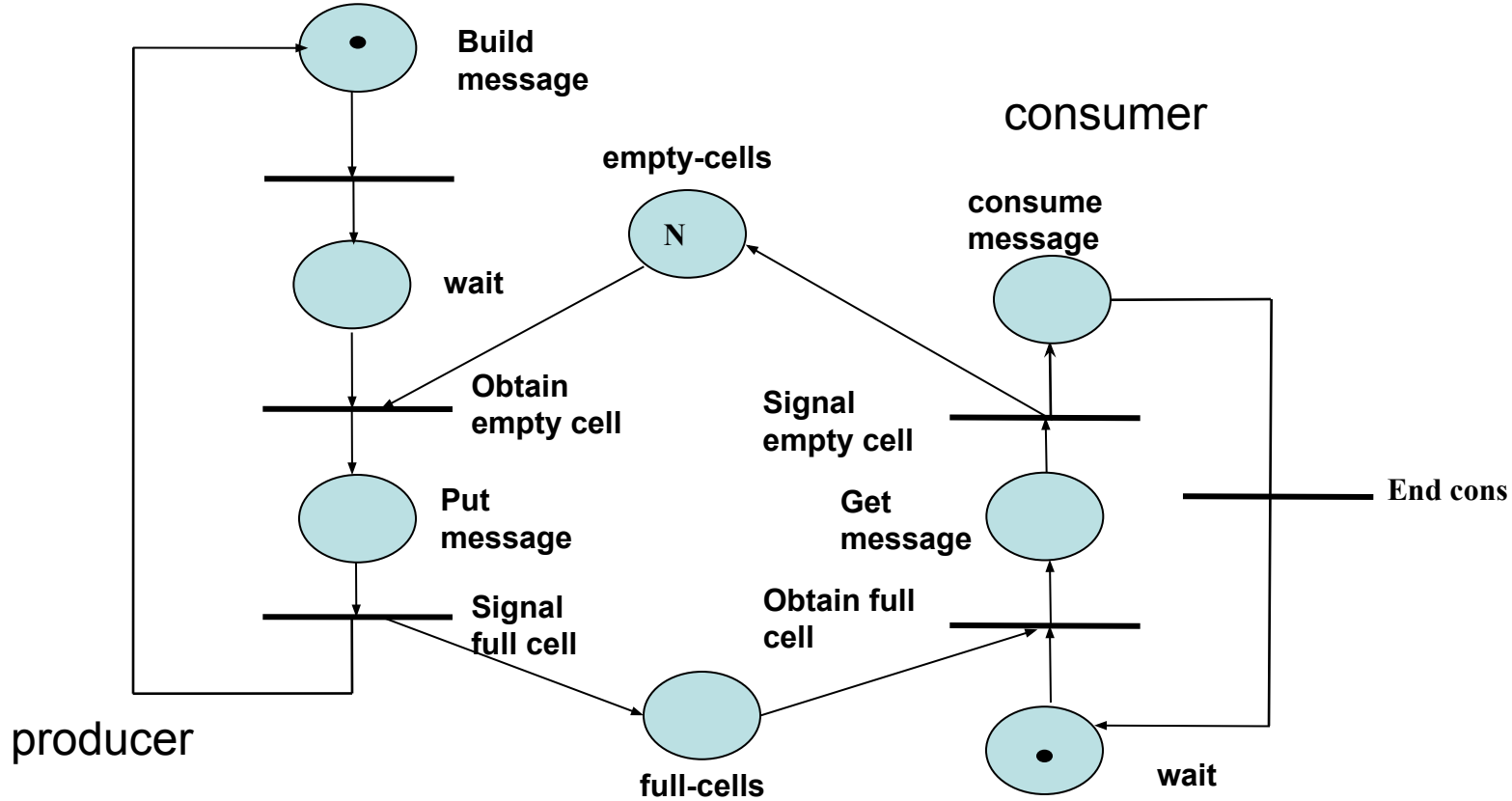


Example: Model two cyclic processes that share a resource in mutual exclusion

Ensure mutual exclusion



Example: Producer-consumer



Graph of markings

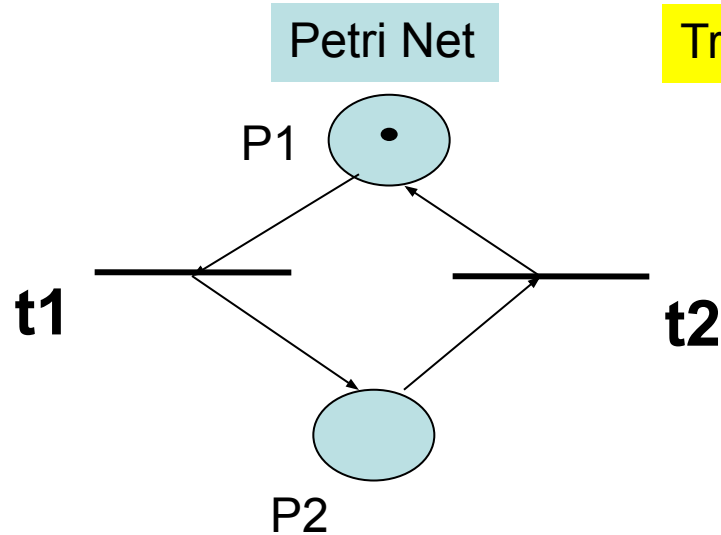
Detection of absence of global or partial blocking.

The construction of the graph is done from the tree of markings.

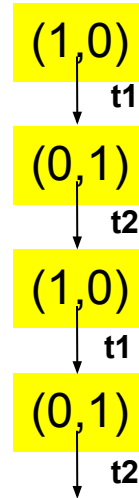
There are two cases:

- Finite number of markings: markings graph
- Infinite number of markings: cover graph

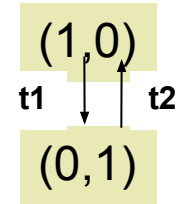
Graph of markings



Tree of markings



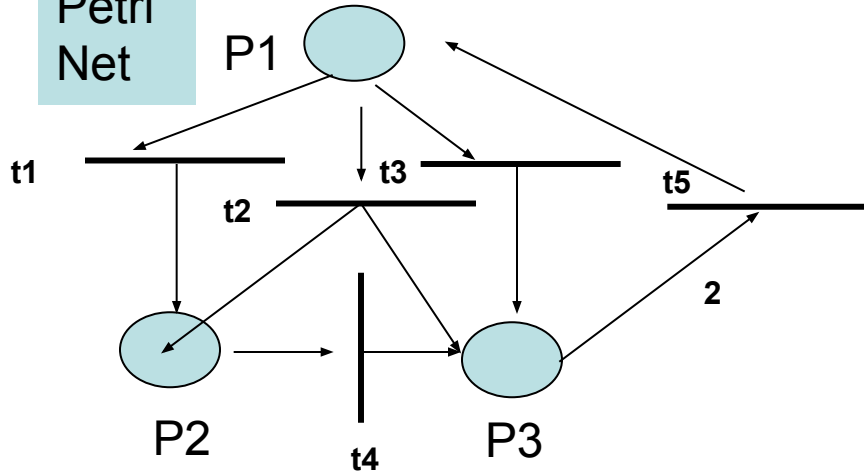
Graph of markings



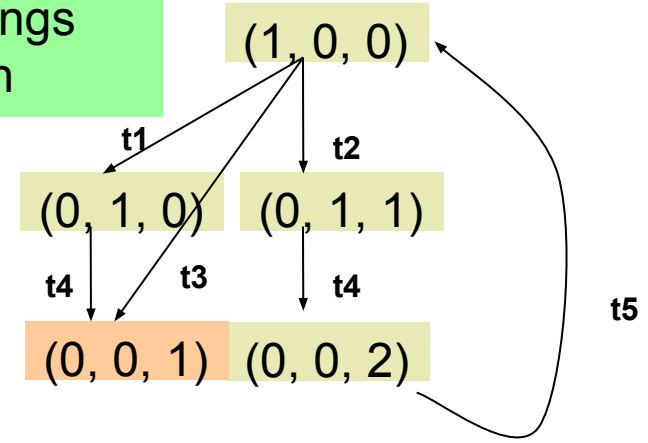
A marking is a vector of which each component represent the number of tokens in a place
It represents the state of a Petri Net at a given time

$$\begin{matrix} & P1 & P2 \\ M0(& 1 & , & 0) \end{matrix}$$

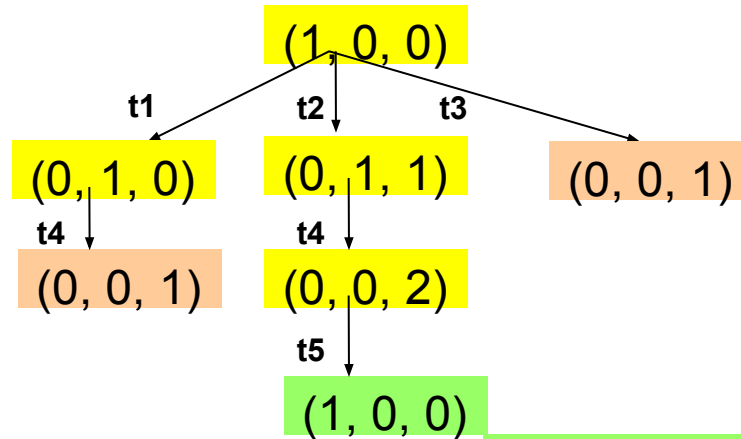
Petri Net



Markings Graph



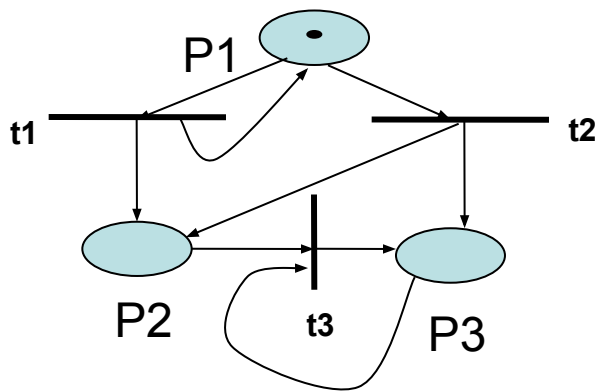
Markings Tree



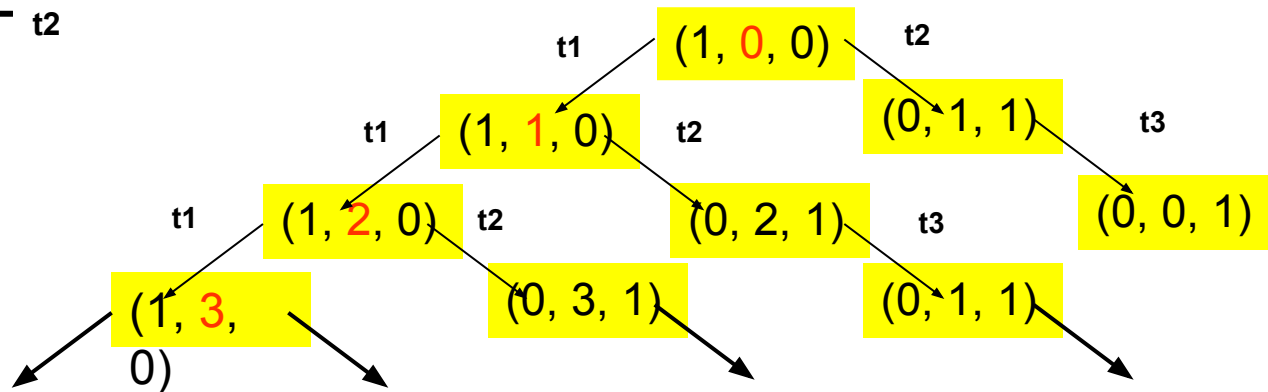
Terminal blocking marking.
All transitions from this marking onwards are disarmed.

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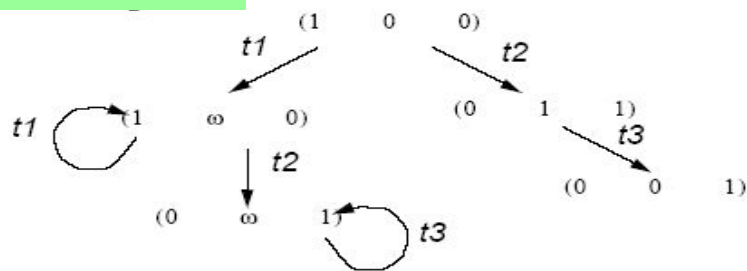
Terminal blocking marking. We loop



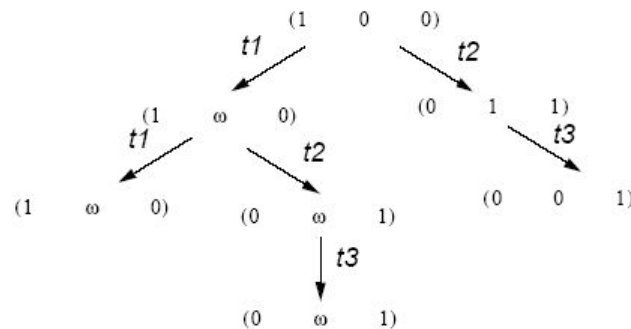
Infinite tree



Graph of marking



Finite tree
Usage of ω



ω symbol

ω symbol:

- Arbitrarily large quantity of tokens.
- Can be approximated as infinite: $\omega \notin \mathbb{N}$
- Properties (valid for all n in \mathbb{N}) :
 - $\omega + n = \omega$
 - $\omega - n = \omega$
 - $n < \omega$
 - $\omega \leq \omega$
- Used to construct the spanning tree in the case of an infinite markings graph

Build the graph

We start from the initial marking

Repeat

For each non-terminal leaf marking, we draw all the armed transitions

For each new marking obtained:

we label it as terminal:

{ if there is no armed transition

Or

if it is already encountered on the path from the root (parent)

}

if it is not labeled then if it is $>$ than a parent marking then we replace the $>$ component by ω

Until (any leaf is terminal)

$M_i > M_j$ if all components of M_i are \geq than the one of M_j except at least one which is $>$.

Examples:

- $M_i(0,0,0,1) > M_j(0,0,0,0)$
- $M_i(0,2,0,1) > M_j(0,1,0,0)$

Problem 1

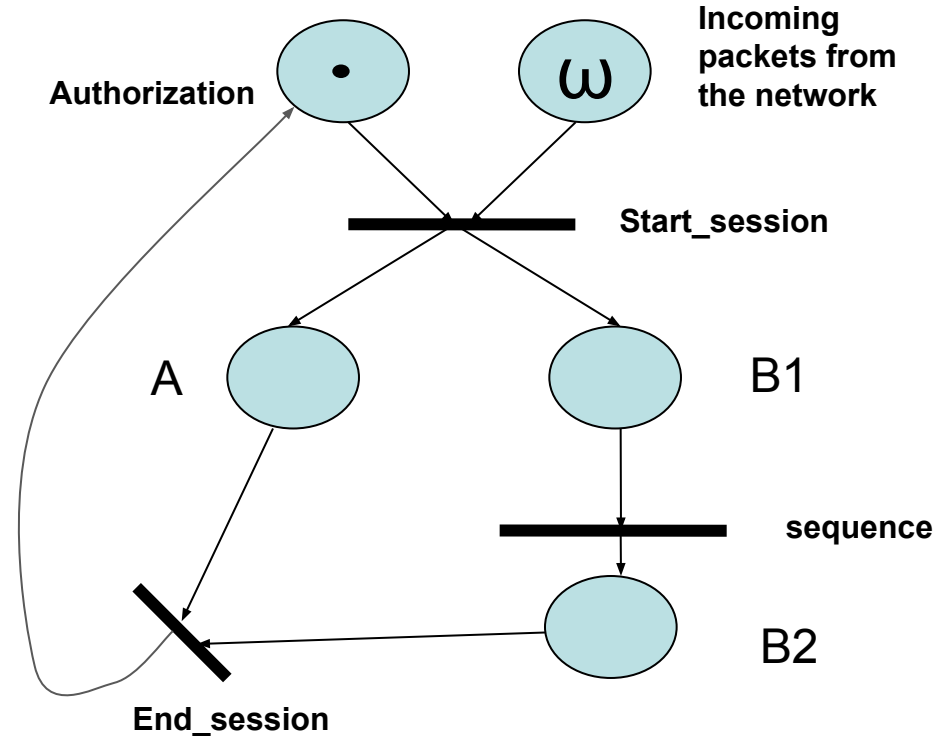
Model a system that receives a message from the network, launches two tasks in parallel:

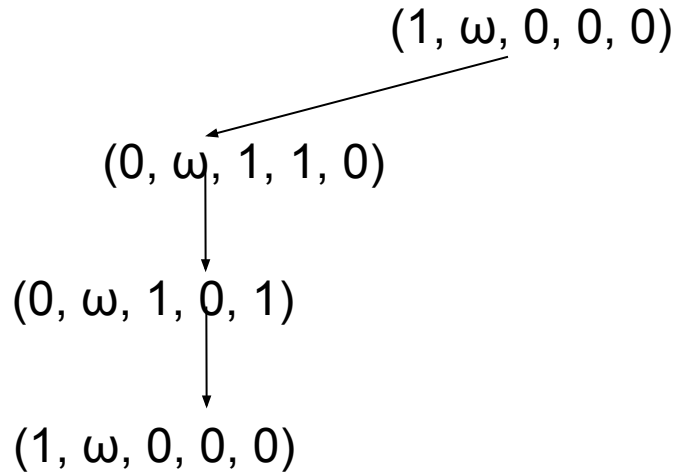
- The first one is composed by an action A
- The second is composed by a sequence of actions B1; B2.

A new session can only be started if the current processing is finished.

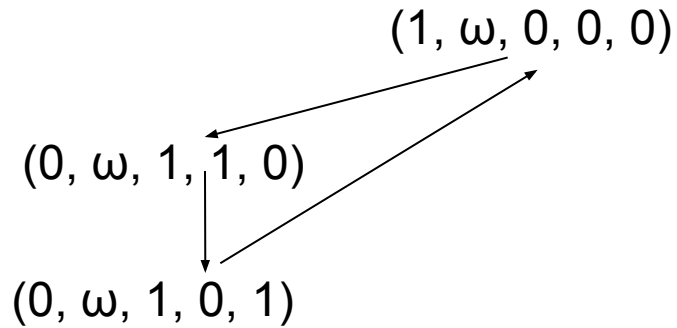
Model this system by a Petri Net

Construct the markings graph of the PN





Tree of markings



Graph of markings

Problem 2

Model a system that receives a message from the network, launches two tasks in parallel:

- The first one is composed by an action A
- The second is composed by N actions B

A new session can only be started if the current processing is finished.

Model this system by a Petri Net

Construct the markings graph of the PN

