Last name(s)	Name	ID

Midterm EDA exam

Length: 2.5 hours

05/11/2018

- The exam has 4 sheets, 8 sides and 4 problems.
- Write your full name and ID on every sheet.
- Write your answers to all problems in the exam sheets within the reserved space.

Problem 1 (2 points)

For each of the following statements, mark with an X the corresponding cell depending on whether it is true or false.

*Note:* Each right answer will add 0.2 points; each wrong answer will subtract 0.2 points, except in the case that there are more wrong answers than right ones, in which the grade of the exercise will be 0.

	(1)	(2)	(3)	(4)	(5)	(6)	(7)	(8)	(9)	(10)
TRUE										
FALSE										

- (1) The cost in time of quicksort in the worst case is  $\Theta(n^2)$ .
- (2) The cost in time of quicksort in the average case is  $\Theta(n^2)$ .
- (3) Any algorithm that computes the sum x + y of two naturals x, y of n bits each, must have cost in time  $\Omega(n)$ .
- (4) There exists an algorithm that computes the sum x + y of two naturals x, y of n bits each, in time O(n).
- (5) Any algorithm that computes the product  $x \cdot y$  of two naturals x, y of n bits each, must have cost in time  $\Omega(n^2)$ .
- (6) There exists an algorithm that computes the product  $x \cdot y$  of two naturals x, y of n bits each, in time  $O(n^2)$ .
- (7) Any algorithm that, given an integer  $k \ge 0$  and a matrix A of  $n \times n$  integer numbers, computes the matrix  $A^k$  must do  $\Omega(k)$  products of matrices.
- (8)  $2^{2n} \in O(2^n)$ .
- (9)  $2n \in O(n)$ .
- (10)  $\log(2n) \in O(\log n)$ .

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Problem 2 (3 points)

Given two matrices of Booleans A and B of size  $n \times n$ , we define their *logical product*  $P = A \cdot B$  as the  $n \times n$  matrix which at the coefficient of the i-th row and j-th column  $(0 \le i, j < n)$  contains:

$$p_{ij} = \bigvee_{k=0}^{n-1} (a_{ik} \wedge b_{kj})$$

(a) (0.5 pts.) Consider the following function for computing the logical product:

```
vector < vector < \mathbf{bool} \gg product1(\mathbf{const}\ vector < vector < \mathbf{bool} \gg A,
\mathbf{const}\ vector < vector < \mathbf{bool} \gg B) \  \{
\mathbf{int}\ n = A.size\  ();
vector < vector < \mathbf{bool} \gg P(n, vector < \mathbf{bool} > (n, \mathbf{false}));
\mathbf{for}\  (\mathbf{int}\ i = 0;\ i < n;\ ++i)
\mathbf{for}\  (\mathbf{int}\ j = 0;\ j < n;\ ++j)
\mathbf{for}\  (\mathbf{int}\ k = 0;\ k < n;\ ++k)
P[i][j] = P[i][j]\  \, \mathbf{or}\  (A[i][k]\  \, \mathbf{and}\  \, B[k][j]);
\mathbf{return}\  \, P;
```

(b) (1 pt.) Consider the following alternative for computing the logical product:

*vector* < *vector* < *bool* >> *product2*(*const vector* < *vector* < *bool* >> & A,

```
\mathbf{const}\ vector < vector < \mathbf{bool} \gg B)\ \{
\mathbf{int}\ n = A.size\ ();
vector < vector < \mathbf{bool} \gg P(n, vector < \mathbf{bool} > (n, \mathbf{false}));
\mathbf{for}\ (\mathbf{int}\ i = 0;\ i < n;\ ++i)
\mathbf{for}\ (\mathbf{int}\ j = 0;\ j < n;\ ++j)
\mathbf{for}\ (\mathbf{int}\ k = 0;\ k < n\ \mathbf{and}\ \mathbf{not}\ P[i][j];\ ++k)
\mathbf{if}\ (A[i][k]\ \mathbf{and}\ B[k][j])\ P[i][j] = \mathbf{true};
\mathbf{return}\ P;
```

Wha	at is the cost in time in the best case in terms of $n$ ? The cost is $\Theta($
Give	e an example of best case.
_	
Wha	at is the cost in time in the worst case in terms of $n$ ? The cost is $\Theta($
Give	e an example of worst case.
_	
(1 5	nt \ Evalain at a high level how to implement a function for computing the
(1.5 logi	pt.) Explain at a high level how to implement a function for computing the cal product of matrices that is more efficient asymptotically in the worst
	than those proposed at sections (a) and (b). What is its cost in time in the
	st case?
(	

Last name(s)	Name	ID	
Problem 3		(2 points)	)
Consider the following function:			
<pre>int mystery_rec(int n, int l, in   if (r == l+1) return l;   int m = (l+r)/2;   if (m*m ≤ n) return mystery_re   else</pre>	ec(n, m, r);		
<pre>int mystery(int n) {    return mystery_rec(n, 0, n+1); }</pre>			
a) (1 pt.) Given an integer $n \ge 0$ , wh	nat does function myst	tery compute?	
b) (1 pt.) What is the cost in time of $m$ Justify your answer.	ystery(n) in terms of $n$	? The cost is $\Theta($	

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Problem 4 (3 points)

Consider the problem of, given a vector of integers (possibly with repetitions), sorting it in increasing order. The following costs are in time and are asked in terms of n, the size of the vector.

- (a) (0.25 pts.) What is the cost of the insertion sort algorithm in the best case? The cost is  $\Theta($   $\bigcirc$   $\bigcirc$   $\bigcirc$   $\bigcirc$   $\bigcirc$   $\bigcirc$
- (b) (0.25 pts.) What is the cost of the insertion sort algorithm in the worst case? The cost is  $\Theta($   $\bigcirc$   $\bigcirc$   $\bigcirc$   $\bigcirc$   $\bigcirc$   $\bigcirc$
- (c) (0.25 pts.) What is the cost of the mergesort algorithm in the best case? The cost is  $\Theta(\bigcirc)$  ).
- (d) (0.25 pts.) What is the cost of the mergesort algorithm in the worst case? The cost is  $\Theta($
- (e) (0.75 pts.) Fill the gaps of the function *my\_sort* defined below, so that given a vector of integers *v*, it sorts it increasingly:

```
void my_sort(vector <int>& v) {
  int n = v. size ();
  double lim = n * log(n);
  int c = 0;
  for (int i = 1; i < n; ++i) {
    int x = v[i];
    int j;
    for (j = i; j > 0 and ; --j) {
      v[j] = ;
      ++c;
    }
  v[j] = ;
  if (c > lim) {
      merge_sort(v);
      return;
    }
}
```

The auxiliary function  $merge\_sort$  is an implementation of the mergesort algorithm, and function log computes the neperian logarithm (that is, in base e).

ustify your answers and give one example of best case and one of worst case.						