

Commitment-Schemes

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XX.XX.XXXX

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Problem(s)

Commitment
Schemes

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Basics

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to do: What are the problems we need to adress

- A **commits** to B
- B keeps commitment, is unable to read or process it
- A reveals to B
- B can verify the commitment

TODO: Image

- ① **Binding:** The Values Alice put in the Commitment cannot be changed after B recieved it
- ② **XXX:** Bob cannot gain any information from the commitment itself

Additional for *real-life-applications*:

- ① Bob's are able compare commitments
- ② Commitments are *tradeable*

Lean-Login:

You're able to access your Youtube-Account and favorites without login, but if you want to change your credentials you need to authenticate
this is done storing the commitment of your login in a cookie

JSON-Web-Tokens (JWT):

A payload (e.g. some account details) are encrypted to a commitment and passed to a third party.

You can verify yourself at the third-party revealing the commitment

this is done *automatic* via session or system attributes

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Hash-Based Commitments

General Concept

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- 1 Alice produces $h = \text{Hash}(m)$ and sends Bob h and Hash
- 2 Bob keeps h
- 3 Alice reveals herself by sending Bob m
- 4 Bob checks if $\text{Hash}(m) \equiv h$

Hash-Based Commitments

Problem: unlimited range - limited domain

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Usually: Bob (and Eve) are not able to *guess* m from h and $Hash$

But: if the *plausible domain* of m is known, its possible for modern computers to brute force reveal your m

Example: Alice commits to Bob about the result of a soccer game Germany vs. Brazil.

Therefore she chooses a score of 0:7 and sends Bob $h = SHA_3(str(0 : 7))$ and the Hashfunction SHA_3

Eve catches the commitment and knows the context of the soccer game. she can know try reasonable combinations of results from 0:0 up to 20:20. She only needs to try $20 \cdot 20 = 400$ results

Improved Concept:

- Alice chooses a random value s
- Alice produces $h = \text{Hash}(m, s)$ and sends h and Hash to Bob
- Bob keeps h and Hash
- Alice reveals herself by sending bob m and s
- Bob checks if $\text{Hash}(m, s) \equiv h$

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Binary-Concept

Requirements and Definitions

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Discrete Logarithm

Requirements and Definitions

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