

Constraints as a Denotational Semantics for Object Calculus

Midterm presentation

Thomas Sixuan Lou

February 3, 2017

Project Overview & Motivations

Motivation: We want to solve *local equations* and *local behaviors* over arbitrary algebras.

Project Overview & Motivations

Motivation: We want to solve *local equations* and *local behaviors* over arbitrary algebras.

- local equations (constraints) \sim classes.

Project Overview & Motivations

Motivation: We want to solve *local equations* and *local behaviors* over arbitrary algebras.

- local equations (constraints) \sim classes.
- local solutions \sim objects.

Project Overview & Motivations

Motivation: We want to solve *local equations* and *local behaviors* over arbitrary algebras.

- local equations (constraints) \sim classes.
- local solutions \sim objects.
- constraint resolution \sim semantics for object calculus.

Object Calculus in a Nutshell

- Values (“objects” / “instances”) of a type are classified into “classes”.

Object Calculus in a Nutshell

- Values (“objects” / “instances”) of a type are classified into “classes”.
- Functions (“methods”) have different behaviors depending on the “class” of its arguments.

Object Calculus in a Nutshell

- Values (“objects” / “instances”) of a type are classified into “classes”.
- Functions (“methods”) have different behaviors depending on the “class” of its arguments.
- Each “class” enforces different internal structures and different behaviors when passed to methods.

Object Calculus in a Nutshell

- Values (“objects” / “instances”) of a type are classified into “classes”.
- Functions (“methods”) have different behaviors depending on the “class” of its arguments.
- Each “class” enforces different internal structures and different behaviors when passed to methods.
- We represent each class by a record of field and function types.

Object Calculus in a Nutshell

- Values (“objects” / “instances”) of a type are classified into “classes”.
- Functions (“methods”) have different behaviors depending on the “class” of its arguments.
- Each “class” enforces different internal structures and different behaviors when passed to methods.
- We represent each class by a record of field and function types.
- We represent the behavior of each method with a “dispatch matrix”

$$\prod_{c \in \mathcal{C}} \prod_{d \in \mathcal{D}} (\tau^c \multimap \rho_d)$$

Object Calculus: examples

$$\tau^{cart} = \text{Class } \{x : \mathbb{R}, y : \mathbb{R}\}$$

$$\tau^{polar} = \text{Class } \{r : \mathbb{R}, \theta : \mathbb{R}\}$$

Object Calculus: examples

$$\tau^{cart} = \text{Class } \{x : \mathbb{R}, y : \mathbb{R}\}$$

$$\tau^{polar} = \text{Class } \{r : \mathbb{R}, \theta : \mathbb{R}\}$$

$$p = \text{Obj } \{x \mapsto 1, y \mapsto 1\} : \tau^{cart}$$

$$q = \text{Obj } \{r \mapsto 1, \theta \mapsto \frac{\pi}{4}\} : \tau^{polar}$$

Object Calculus: examples

$$\tau^{cart} = \text{Class } \{x : \mathbb{R}, y : \mathbb{R}\}$$

$$\tau^{polar} = \text{Class } \{r : \mathbb{R}, \theta : \mathbb{R}\}$$

$$p = \text{Obj } \{x \mapsto 1, y \mapsto 1\} : \tau^{cart}$$

$$q = \text{Obj } \{r \mapsto 1, \theta \mapsto \frac{\pi}{4}\} : \tau^{polar}$$

$$e : \left[\begin{array}{cc} \tau^{cart} \rightarrow \rho_{distSq} & \tau^{cart} \rightarrow \rho_{xCoord} \\ \tau^{polar} \rightarrow \rho_{distSq} & \tau^{polar} \rightarrow \rho_{xCoord} \end{array} \right]$$

$$e = \left[\begin{array}{cc} \lambda p : \tau^{cart}.(p.x^2 + p.y^2) & \lambda p : \tau^{cart}.p.x \\ \lambda p : \tau^{polar}.(p.r^2) & \lambda p : \tau^{polar}.(p.r \times \cos(p.\theta)) \end{array} \right]$$

Adding Constraints in Classes

Now, we allow the declaration of *local equations* in objects.

Adding Constraints in Classes

Now, we allow the declaration of *local equations* in objects.

$$\text{Class } \{s : A, t : B, s = t\}$$

Adding Constraints in Classes

Now, we allow the declaration of *local equations* in objects.

$$\text{Class } \{s : A, t : B, s = t\}$$

An instance solves constraints declared by its class.

Adding Constraints in Classes

Now, we allow the declaration of *local equations* in objects.

$$\text{Class } \{s : A, t : B, s = t\}$$

An instance solves constraints declared by its class.

$$\text{Obj } \{s \mapsto a, t \mapsto b\} : \text{Class } \{s : A, t : B, s = t\}$$

Adding Constraints in Classes

Now, we allow the declaration of *local equations* in objects.

$$\text{Class } \{s : A, t : B, s = t\}$$

An instance solves constraints declared by its class.

$$\text{Obj } \{s \mapsto a, t \mapsto b\} : \text{Class } \{s : A, t : B, s = t\}$$

and also first-order constraints!

Adding Constraints in Classes

Now, we allow the declaration of *local equations* in objects.

$$\text{Class } \{s : A, t : B, s = t\}$$

An instance solves constraints declared by its class.

$$\text{Obj } \{s \mapsto a, t \mapsto b\} : \text{Class } \{s : A, t : B, s = t\}$$

and also first-order constraints!

$$\text{Class } \{s = t, \neg\phi, \phi_1 \wedge \phi_2, \phi_1 \vee \phi_2, \exists x.\phi\}$$

Example: classes as problems

An arithmetic problem class:

$$\textit{Class } \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\}$$

Example: classes as problems

An arithmetic problem class:

Class $\{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\}$

evaluate!

\Downarrow

$\llbracket \textit{Class } \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\} \rrbracket_{AE}$

Example: classes as problems

An arithmetic problem class:

Class $\{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\}$

evaluate!

\Downarrow

$\llbracket \text{Class } \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\} \rrbracket_{AE}$

\Downarrow

Class $\{x : \mathbb{N}, y : \mathbb{N}, \llbracket (3 + x) \cdot 2 = y \rrbracket_{AE} \wedge \llbracket y = 2 \cdot 3 + 2 \rrbracket_{AE}\}$

Example: classes as problems

An arithmetic problem class:

Class $\{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\}$

evaluate!

\Downarrow

$\llbracket \text{Class } \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 2 \cdot 3 + 2\} \rrbracket_{AE}$

\Downarrow

Class $\{x : \mathbb{N}, y : \mathbb{N}, \llbracket (3 + x) \cdot 2 = y \rrbracket_{AE} \wedge \llbracket y = 2 \cdot 3 + 2 \rrbracket_{AE}\}$

\Downarrow

Class $\{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 8\}$

Example: objects as solutions

Possible solutions to a problem given by the semantics of equations

$$\mathit{Obj} \{x \mapsto 1, y \mapsto 8\} : C$$

We denote the “solving” process into the CLP land:

$$\text{let } \theta = \{x \mapsto 1, y \mapsto 8\}$$

$$\llbracket \mathit{Class} \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 8\} \rrbracket(\theta)$$

Example: objects as solutions

Possible solutions to a problem given by the semantics of equations

$$\text{Obj } \{x \mapsto 1, y \mapsto 8\} : C$$

We denote the “solving” process into the CLP land:

$$\text{let } \theta = \{x \mapsto 1, y \mapsto 8\}$$

$$\llbracket \text{Class } \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 8\} \rrbracket(\theta)$$

$$\Downarrow$$

$$8 = \llbracket (3 + x) \cdot 2 \rrbracket \theta \equiv \llbracket y \rrbracket \theta$$

$$8 = \llbracket y \rrbracket \theta = \llbracket 8 \rrbracket \theta$$

Example: objects as solutions

Possible solutions to a problem given by the semantics of equations

$$\text{Obj } \{x \mapsto 1, y \mapsto 8\} : C$$

We denote the “solving” process into the CLP land:

$$\text{let } \theta = \{x \mapsto 1, y \mapsto 8\}$$

$$\llbracket \text{Class } \{x : \mathbb{N}, y : \mathbb{N}, (3 + x) \cdot 2 = y \wedge y = 8\} \rrbracket(\theta)$$

$$\Downarrow$$

$$8 = \llbracket (3 + x) \cdot 2 \rrbracket \theta \equiv \llbracket y \rrbracket \theta$$

$$8 = \llbracket y \rrbracket \theta = \llbracket 8 \rrbracket \theta$$

$$\Downarrow$$

$$\{\theta\}$$

As in ordinary object calculus, there are inheritance to create class hierarchies.

$$\tau^{cart} = \text{Class } \{x : \mathbb{R}, y : \mathbb{R}\}$$

$$\tilde{\tau}^{cart} <: \tau^{cart} = \text{Class } \{color : colorT, x : \mathbb{R}, y : \mathbb{R}\}$$

same for constraints

$$\tilde{\tau}^{cart} <: \tau^{cart} = \text{Class } \{x + y = 1, x : \mathbb{R}, y : \mathbb{R}\}$$

$$\hat{\tau}^{cart} <: \tilde{\tau}^{cart} = \text{Class } \{x = 1 \wedge x + y = 1, x : \mathbb{R}, y : \mathbb{R}\}$$

Recall methods are (internally) represented as a dispatch matrix, when adding method to a class, we are adding a column vector into the class.

$$\begin{aligned}\tilde{\tau}^{cart} <: \tau^{cart} &= \text{Class} \left\{ \text{dist} : \prod_{c \in C} \tau^c \rightarrow \mathbb{R} \right\} \\ \hat{\tau}^{cart} <: \tilde{\tau}^{cart} &= \text{Class} \left\{ \text{dist}(\text{self}) = 1 \right\}\end{aligned}$$

Recall methods are (internally) represented as a dispatch matrix, when adding method to a class, we are adding a column vector into the class.

$$\tilde{\tau}^{cart} <: \tau^{cart} = \text{Class} \left\{ \text{dist} : \prod_{c \in C} \tau^c \rightarrow \mathbb{R} \right\}$$

$$\hat{\tau}^{cart} <: \tilde{\tau}^{cart} = \text{Class} \left\{ \text{dist}(\text{self}) = 1 \right\}$$

$$\text{Obj} \{ x \mapsto 1, y \mapsto 0 \} : \hat{\tau}^{cart}$$

Dispatch methods

Let $\theta = \{x \mapsto 1, y \mapsto 0\}$,

$$\begin{aligned} & \llbracket \text{Class}\{x : \mathbb{R}, y : \mathbb{R}, \text{dist} : \prod_{c \in \mathcal{C}} \tau^c \rightarrow \mathbb{R}\}, \text{dist}(\text{self}) = 1 \rrbracket \theta \\ & \quad \Downarrow \\ & \llbracket \text{dist}(\text{self}) \rrbracket \theta \stackrel{\text{dispatch}}{=} \llbracket (\lambda p : \tau^{\text{cart}}. (p.x^2 + p.y^2)) \text{self} \rrbracket \theta \\ & \quad = \llbracket x^2 + y^2 \rrbracket \theta = 1 = \llbracket 1 \rrbracket \theta \\ & \quad \Downarrow \\ & \quad \{\theta\} \end{aligned}$$