

Lamret Kaur  
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KRGV3A

Test cases:

input 1:  $A = [1, 3, 5]$

output 1: 8

input 2:  $A = [2, 3]$

output 2: 2

We define  $f[x, y]$  as number of different corresponding bits in the binary representation of  $x$  &  $y$ . For example,  $f[2, 7] = 2$ , since the binary representation of 2 & 7 are 010 and 111, respectively. The first and the third bit differ, so  $f[2, 7] = 2$ .

You are given an array of  $N$  positive integers,  $A_1, A_2, \dots, A_N$ . Find sum of  $f(A_i, A_j)$  for all pairs  $(i, j)$  such that  $1 \leq i, j \leq N$ . Return the answer modulo  $10^9 + 7$ .

$\Rightarrow$  Given

$f(x, y) = \text{no. of different bits in binary of } x \text{ & } y$   
[Checking every bit  $\rightarrow TC = (N^2) \rightarrow TLE$ ]

To solve this, we can use the concept of bit manipulation - Try checking the bits of  $x \oplus y$ .

Let's look at a particular  $k^{th}$  bit

if  $k^{th}$  bit of  $x$  is 1 &  $k^{th}$  bit of  $y$  is 0.

Or vice versa

so ans = 2 for pair  $(x, y)$  &  $(y, x)$

$\hookrightarrow$  Let  $\text{cnt}_1 = \text{nos having bit } k=1$

$\text{cnt}_2 = \text{nos having bit } k=0$

contribution =  $2 \times \text{cnt}_1 \times \text{cnt}_0$   
for  $(x, y)$  and  $(y, x)$

→ Steps to solve :

- For every bit ( $0 \rightarrow 31$ )  
1) Count numbers having that bit set.  
2) Compute pairs.  
3) Add to answer.

→ Code :

```
#include <bits/stdc++.h>
using namespace std;
```

```
const long long MOD = 1e9 + 7;
```

```
long long totalPairs (vector<int> &A) {
    long long n = A.size();
    long long ans = 0;
```

```
    for (int bit = 0; bit < 32; bit++) {
        long long cnt1 = 0;
        for (int i = 0; i < n; i++) {
            if (A[i] & (1 << bit)) {
                cnt1++;
            }
        }
    }
}
```

```
long long contribution = (cnt1 * cnt0) % MOD;
contribution = (2 * contribution) % MOD;
```

```
ans = (ans + contribution) % MOD
```

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```
return ans;
```

Y