Program Structures and Algorithms Spring 2023(SEC -01) Assignment-5

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Task:

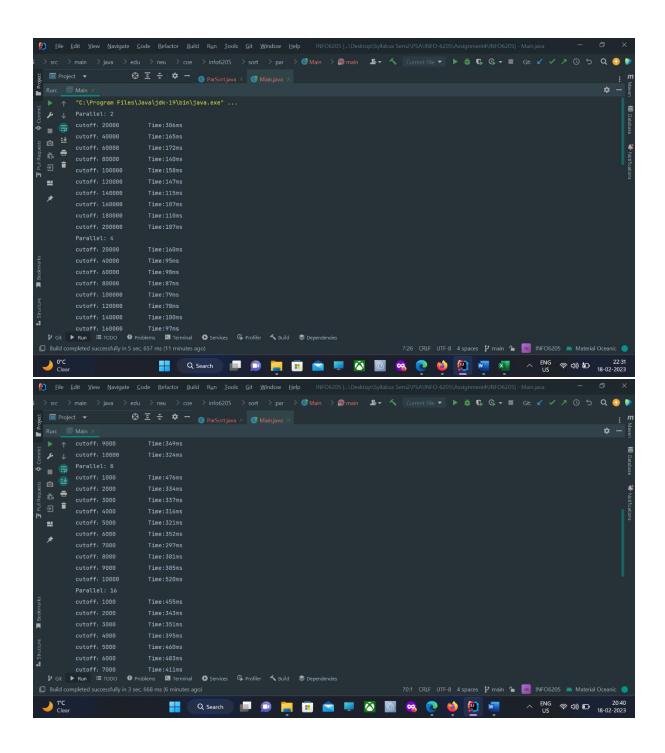
Please see the presentation on *Assignment on Parallel Sorting* under the *Exams. etc.* module. Your task is to implement a parallel sorting algorithm such that each partition of the array is sorted in parallel. You will consider two different schemes for deciding whether to sort in parallel.

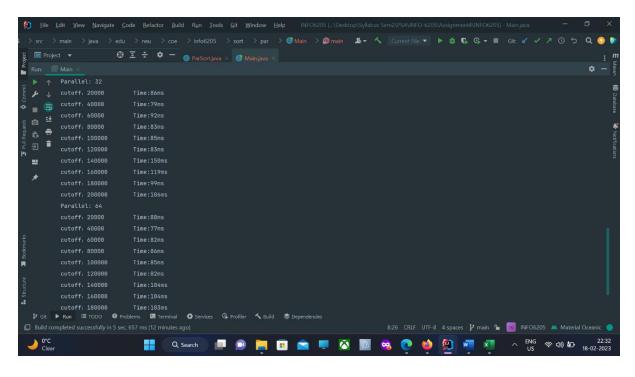
- 1. A cutoff (defaults to, say, 1000) which you will update according to the first argument in the command line when running. It's your job to experiment and come up with a good value for this cutoff. If there are fewer elements to sort than the cutoff, then you should use the system sort instead.
- 2. Recursion depth or the number of available threads. Using this determination, you might decide on an ideal number (*t*) of separate threads (stick to powers of 2) and arrange for that number of partitions to be parallelized (by preventing recursion after the depth of *lg t* is reached).
- 3. An appropriate combination of these.

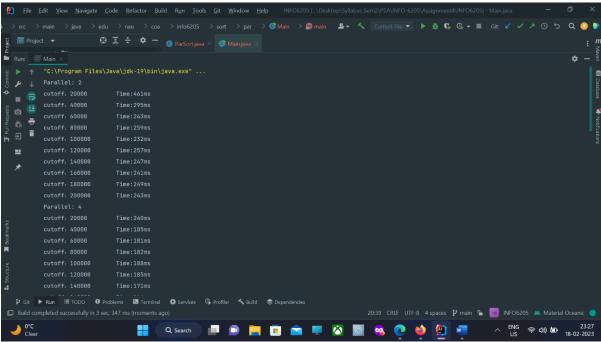
There is a *Main* class and the *ParSort* class in the *sort.par* package of the INFO6205 repository. The *Main* class can be used as is but the *ParSort* class needs to be implemented where you see "TODO..." [it turns out that these TODOs are already implemented]. Unless you have a good reason not to, you should just go along with the Java8-style future implementations provided for you in the class repository.

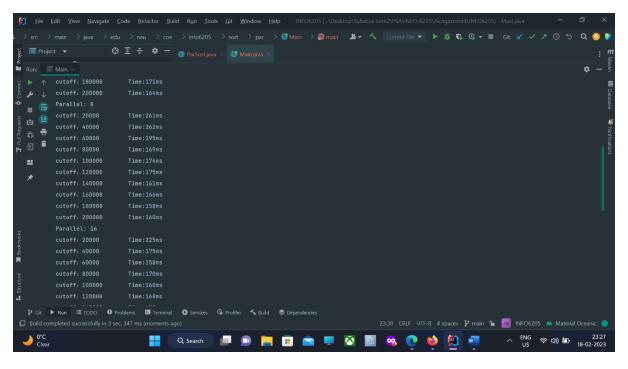
You must prepare a report that shows the results of your experiments and draws a conclusion (or more) about the efficacy of this method of parallelizing sort. Your experiments should involve sorting arrays of sufficient size for the parallel sort to make a difference. You should run with many different array sizes (they must be sufficiently large to make parallel sorting worthwhile, obviously) and different cutoff schemes.

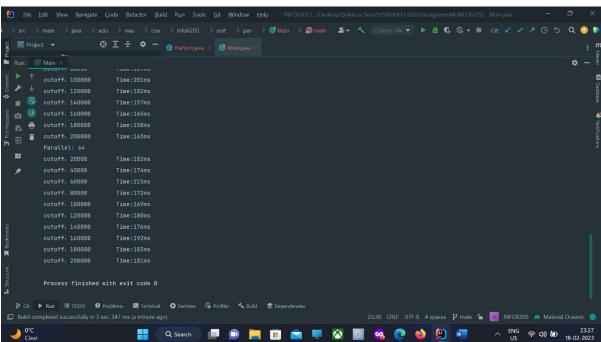
Outputs and Test cases:

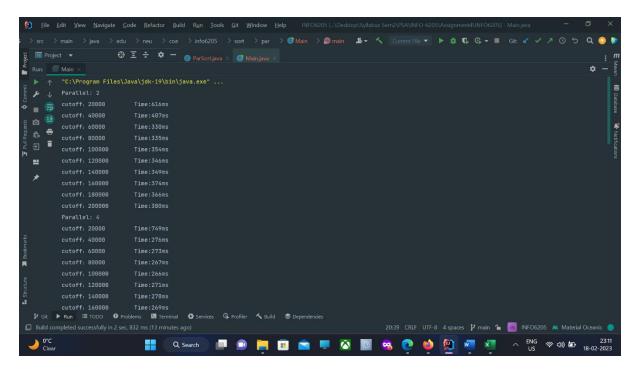


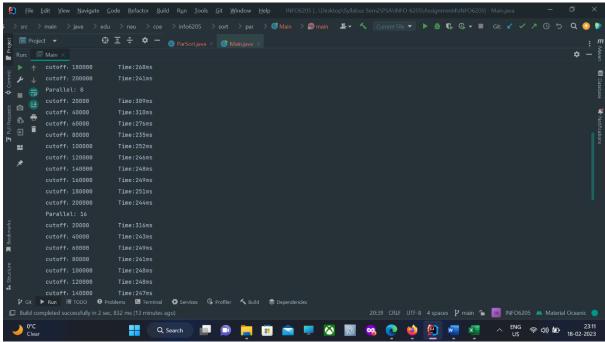


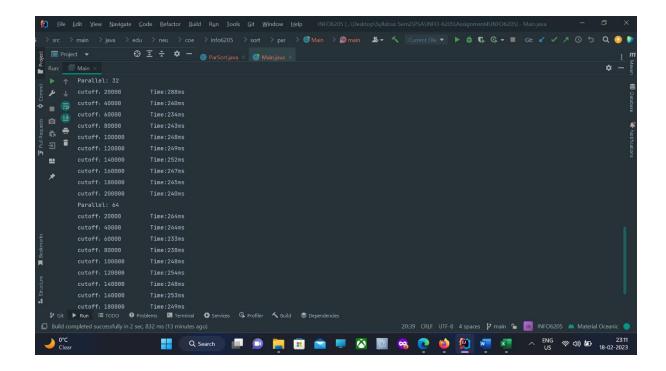












Relationship Conclusion/ Evidence Graph:

Cutoff values for array of size: 250,000

Cutoff	Thread:2	Thread:4	Thread:8	Thread:16	Thread:32	Thread:64
20000	64.2	52	13.2	12.7	13.1	14.5
40000	30.5	12.8	9.1	9.3	9.6	10.7
60000	25.7	12.1	9.1	10.5	10.4	9
80000	16.5	9.4	8.3	10	35.6	8.8
100000	17.6	8.6	8.5	8.8	12.5	8.9
120000	34.3	8.9	8.5	8.9	12.4	8.7
140000	17.5	9.4	8.8	8.9	9.4	8.7
160000	25.5	9.5	8.5	8.3	8.7	7.9
180000	19	8.7	8.7	8.5	8.5	7.9
200000	20.9	8.6	8.6	8.2	7.7	8.4

Cutoff values for array of size: 500,000

Cutoff	Thread:2	Thread:4	Thread:8	Thread:16	Thread:32	Thread:64
20000	42.7	25.5	24.7	22.9	16.4	17.2
40000	26.7	18.6	25.3	15.8	15.9	16.1
60000	22.8	18.5	18.1	15.6	16.4	15.9

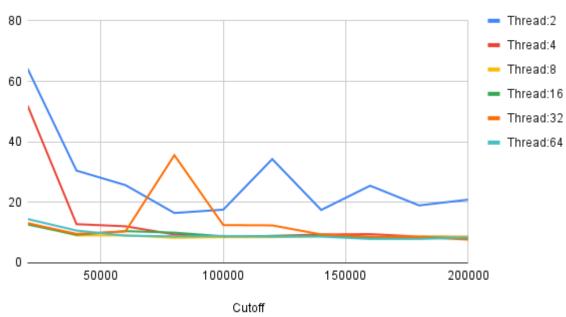
80000	23.3	18.8	17.5	17	17.5	17.5
100000	24	18.3	17.8	16.6	17.1	17.3
120000	24.3	18.4	17.8	16.7	16.9	16.3
140000	23.5	17.2	19.7	17	16.4	16.4
160000	25.2	17.4	19.4	16.5	16.4	16.8
180000	25.1	17.1	16.4	16.4	16.7	16.9
200000	25.6	16.8	16.3	16.5	16.4	16.6

Cutoff values for array of size: 750,000

Cutoff	Thread:2	Thread:4	Thread:8	Thread:16	Thread:32	Thread:64
20000	61.6	74.9	30.9	31.6	28.8	26.4
40000	40.7	27.6	31	24.3	24	24.4
60000	33	27.3	27.6	24.9	23.4	23.3
80000	33.5	26.7	23.5	24.1	24.3	23.8
100000	35.4	26.6	25.2	24.8	24.8	24.8
120000	34.6	27.1	24.6	24.8	24.9	25.4
140000	34.9	27.8	24.8	24.7	25.2	24.8
160000	37.4	26.9	24.9	25	24.7	25.3
180000	36.6	26.8	25.1	24.8	24.5	24.9
200000	38	24.1	24.4	24.6	24	24.4

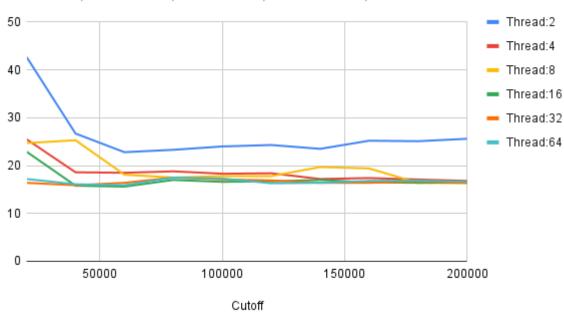
Graph for array size: 250,000

Thread:2, Thread:4, Thread:8, Thread:16, Thread:32...



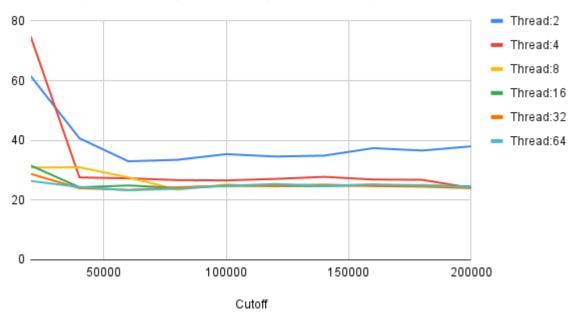
Graph for array size: 500,000

Thread:2, Thread:4, Thread:8, Thread:16, Thread:32...



Graph for array size: 750,000





From the graph above we can conclude that:

After plotting the above three graphs from the different values tested for their cut-off limits. It can be understood that the cut-off ratio, defined as the ratio of cut off to the array size helps us understand the relationship between both. The ratio tends towards a constant time when there are sufficiently large gaps between the parallel sorting parameters and array size. The values plotted for the largest array size shows constant behaviour in plotting for the graph. Thus, it could be understood to keep the values of array size to be sufficiently large in comparison to the parallel sorting values to yield a comparable and appropriate cut-off value. The above values shows least constant behaviour when the array size is closest to parallel sorting parameters, there the graph plotted is erratic and not of much linearity.