

## Midterm S20

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Due Date .....Saturday Nov 19, 2022 4pm MT  
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Quiz Code (enter in Canvas to get access to the LaTeX template) .....**7X0vUY9sMh**

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### Instructions

- You may either type your work using this template, or you may handwrite your work and embed it as an image in this template. **If you choose to handwrite your work, the image must be legible, and oriented so that we do not have to rotate our screens to grade your work.** We have included some helpful LaTeX commands for including and rotating images commented out near the end of the LaTeX template.
- You should submit your work through the **class Gradescope page** only. Please submit one PDF file, compiled using this LaTeX template.
- You may not need a full page for your solutions; pagebreaks are there to help Gradescope automatically find where each problem is. Even if you do not attempt every problem, please submit this document with no fewer pages than the blank template (or Gradescope has issues with it).
- You **may not collaborate with other students. Copying from any source is an Honor Code violation. Furthermore, all submissions must be in your own words and reflect your understanding of the material.** If there is any confusion about this policy, it is your responsibility to clarify before the due date.
- Posting to **any** service including, but not limited to Chegg, Discord, Reddit, StackExchange, etc., for help on an assignment is a violation of the Honor Code.
- You **must** virtually sign the Honor Code. Failure to do so will result in your assignment not being graded.

## Honor Code (Make Sure to Virtually Sign)

- Problem HC.**
- My submission is in my own words and reflects my understanding of the material.
  - Any collaborations and external sources have been clearly cited in this document.
  - I have not posted to external services including, but not limited to Chegg, Reddit, StackExchange, etc.
  - I have neither copied nor provided others solutions they can copy.

*I agree to the above, Tyler Huynh.*

□

## 20 Standard 20: QuickSort

**Problem 20.** Consider a modified version of QUICKSORT, in which the PARTITION subroutine *always chooses the  $(2/3)n$ -th largest element of the list*. You may assume PARTITION still takes  $\Theta(n)$  time on lists of length  $n$ .

Do the following **three** parts of the question (note the third part on the next page).

1. **Write down** the recurrence relation for the runtime of this modified version of QUICKSORT. **Justify your recurrence relation in 1–2 sentences.**

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P20.1

1.)

The recurrence relation for the runtime of Quicksort.

$$T(n) = \begin{cases} \Theta(1) & : n \leq 3 \\ T(\frac{n}{3}) + T(\frac{2n}{3}) + \Theta(n) & : n > 3 \end{cases}$$

This is my recurrence relation for this problem because  $\Theta(1)$  represents the runtime of  $T(n)$ , when we reach the base case.

$T(\frac{n}{3}) + T(\frac{2n}{3}) + \Theta(n)$  represents the partition subroutine choosing the  $(\frac{2}{3})n^{\text{th}}$  largest element of the list, this subroutine will always take  $\Theta(n)$  time on lists of length  $n$ .

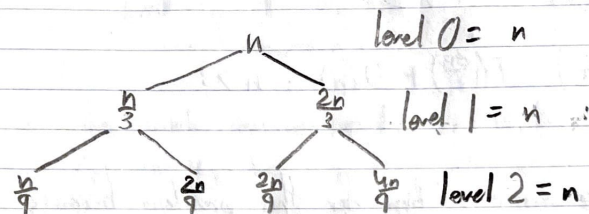
Answer. □

2. Solve your recurrence relation from part (1), by the tree method. **Show your work.** Find a function  $f(n)$  such that the runtime  $T(n)$  is  $T(n) = \Theta(f(n))$ .

2.)

I will solve my recurrence relation from part 1 by using the tree method.

The tree



Work done per level:  $n$

I will now find  $k$  where it will represent the amount of times it will take us to reach the base case:

I will compare  $n/3$  and  $2n/3$  to see which  $k$  value will terminate last:

$$T(n/3) = \frac{n}{3^{k_1}} \leq 3$$

$$n \leq 3(3^{k_1})$$

$$n \leq 3^{k_1+1}$$

$$\log_3 n \leq \log_3 (3^{k_1+1})$$

$$\log_3 n \leq k_1 + 1$$

$$k_1 > \log_3 n - 1$$

$$T(2n/3) = \left(\frac{2}{3}\right)^{k_2} n \leq 3$$

$$\left(\frac{2}{3}\right)^{k_2} \leq \frac{3}{n}$$

$$\log_{\frac{2}{3}} \left(\frac{2}{3}\right)^{k_2} \leq \log_{\frac{2}{3}} \frac{3}{n}$$

$$k_2 \leq \log_{\frac{2}{3}} \frac{3}{n}$$

This value will be greater than  $k_1$

Answer.

$$\begin{aligned}
 T(n) &= (\text{cost of base case}) (\# \text{ of times base case is reached}) + \sum_{i=0}^{\log_{\frac{2}{3}} n} n \\
 &= (\theta(1)) \left(1^{\log_{\frac{2}{3}} n}\right) + n \left(\log_{\frac{2}{3}} n\right) \\
 &= (\theta(1)) + n \left(\log_{\frac{2}{3}} n - \log_{\frac{2}{3}} n\right) \\
 &= (\theta(1)) + n \log_{\frac{2}{3}} n - n \log_{\frac{2}{3}} n
 \end{aligned}$$

From the above we can see that the highest element is  $n \log n$ , such that our runtime complexity for  $T(n)$  will be:

$$T(n) = \theta(n \log n)$$

3. Write down the following three functions and put them in order from smallest to largest. If two are asymptotically equal ( $\Theta$  of one another), you should clearly indicate this.

- the best-case running time of QUICKSORT (when it always chooses the best possible pivot),
- the worst-case running time of QUICKSORT (when it always chooses the worst possible pivot),
- your answer from part (2) above.

**You do not** need to recalculate the best- and worst-case runtimes, those you can get from class notes.

3.)

Best-case for quicksort: chooses median for the pivot, splitting the list in half, such that:

$$2T\left(\frac{n}{2}\right) + \Theta(n) = \Theta(n \log n)$$

Worst-case for quicksort: chooses the last or first element in a list as the pivot, such that:

$$T(n-1) + T(1) + \Theta(n) = \Theta(n^2)$$

Runtime for our recurrence relation, choosing  $\left(\frac{2}{3}\right)^{\text{th}}$  largest element in a list, such that:

$$T\left(\frac{n}{3}\right) + T\left(\frac{2n}{3}\right) + \Theta(n) = \Theta(n \log n)$$

Note the runtime for our recurrence and the best-case for quicksort are asymptotically equal to one another.

Ordering the functions from smallest to greatest:

$$(\Theta(n \log n), \Theta(n \log n)) = (\text{Best-case, our recurrence relation})$$

$$\Theta(n^2) = \text{worst-case}$$

