



# Introduction to Quantum Information and Quantum Machine Learning

Laboratory - class 4

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# Quantum Oracle - Preparation for Grover's Algorithm



**Grover's algorithm is a quantum algorithm** designed to run on a quantum computer, presented by Lov K. Grover in 1996 [1] and published in 2001 [2].

The algorithm involves searching a database consisting of  $N$  elements in order to find a highlighted element in it. This is a similar problem to "reverse" searching a phone book. In a book containing  $N$  data, we want to find the name of the owner of a given number.

### Computational complexity

While the number of steps required to solve the problem using the classical algorithm is of the order of  $O(N)$  Grover's quantum algorithm only needs about  $O(\sqrt{N})$  steps, so it allows for a quadratic acceleration of the program execution time.

The algorithm involves searching for a given element in an unsorted  $N$ -element set. The search problem comes down to determining, by means of unitary transformations, the appropriate index defining a given element in the set.

[1] L. K. Grover. *A fast quantum mechanical algorithm for database search*. In STOC '96: Proceedings p. 212-219, New York, NY, USA (1996)

[2] Grover L.K.: *From Schrödinger's equation to quantum search algorithm*, American Journal of Physics, 69(7): 769-777 (2001)



# 1. XOR Oracle and Phase Oracle

# XOR Oracle

The problem of N data in an unstructured database

Classical maps (function)

$$f: \{0,1\}^n \rightarrow \{0,1\}$$
$$f(x) = \begin{cases} 1, & x = a \\ 0, & x \neq a \end{cases}$$

$a = ?$

$f(x) = 1?$



Illustration of a searching process in a database with N entries.

A function  $f(x)$  is used to determine if the target is located at entry  $x$ .

## Computational Complexity

Classic search

→ Average  $\mathcal{O}\left(\frac{N}{2}\right) \sim \mathcal{O}(N)$

Grover's quantum algorithm

→ Average  $\mathcal{O}(\sqrt{N})$

Grover's algorithm provides a quadratic speedup over the classical one

# XOR Oracle

$$U_f (|y\rangle \otimes |x\rangle) = |y \oplus f(x)\rangle \otimes |x\rangle$$

$$a = 10$$

$ i\rangle$	$ y\rangle x\rangle$	$U_f ( y\rangle \otimes  x\rangle)$
$ 0\rangle$	$ 0\ 00\rangle$	$ 0\ 00\rangle$
$ 1\rangle$	$ 0\ 01\rangle$	...
$ 2\rangle$	$ 0\ \mathbf{10}\rangle$	$ 1\ \mathbf{10}\rangle$
$ 3\rangle$	$ 0\ 11\rangle$	...
...	...	...
$ n-1\rangle$	$ 1\ 11\rangle$	$ 1\ 11\rangle$

# Phase Oracle

$$U_f |x\rangle = (-1)^{f(x)} |x\rangle$$

$$a = 010$$

$ i\rangle$	$ x\rangle$	$U_f  x\rangle$
$ 0\rangle$	$ 0\ 00\rangle$	$ 0\ 00\rangle$
$ 1\rangle$	$ 0\ 01\rangle$	...
$ 2\rangle$	$ \mathbf{0\ 10}\rangle$	$-\mathbf{ 0\ 10\rangle}$
$ 3\rangle$	$ 0\ 11\rangle$	...
...	...	...
$ n-1\rangle$	$ 1\ 11\rangle$	$ 1\ 11\rangle$

# Create the Matrix of $U_f$ operation

XOR Oracle

Phase Oracle

	<i>j column</i>							
	0	1	2	3	4	5	6	7
<i>i row</i>	0	1	2	3	4	5	6	7
	1							
	2							
	3							
	4							
	5							
	6							
	7							



*quantum gate*

	<i>j column</i>							
	0	1	2	3	4	5	6	7
<i>i row</i>	0	1	2	3	4	5	6	7
	1							
	2							
	3							
	4							
	5							
	6							
	7							



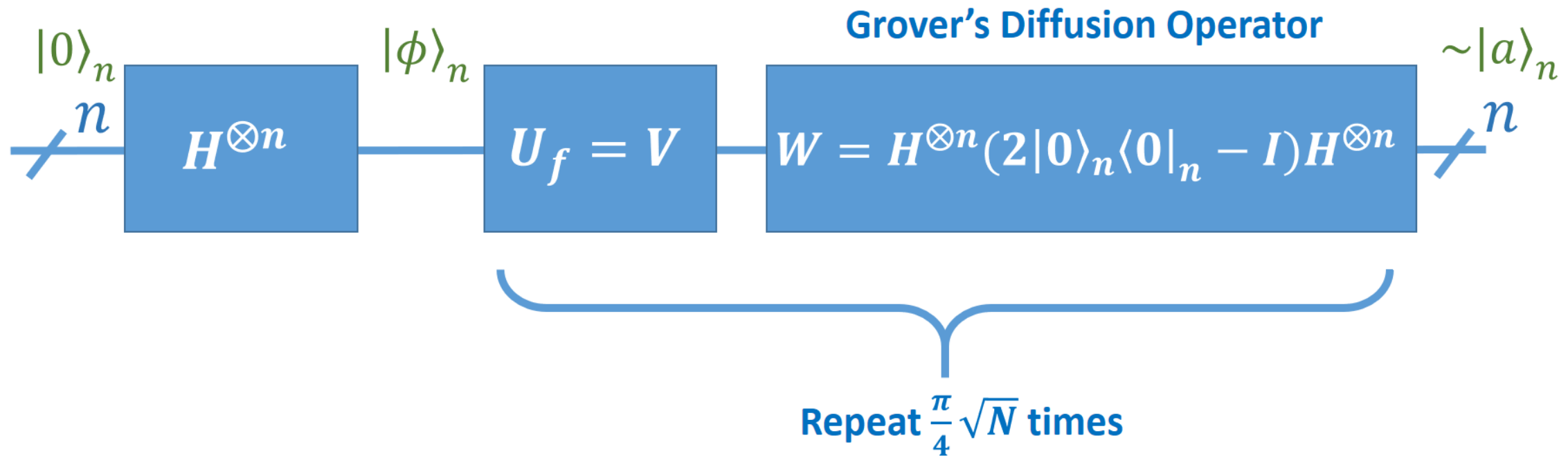
*quantum gate*



## 2. Grover's Diffusion operator



# Implementation of the Grover's diffusion operator



# Implementation of the Grover's diffusion operator

- $|0\rangle_n {}_n\langle 0| = ( \quad )$
- $\hat{\mathbb{I}} = \begin{pmatrix} 1 & 0 & \cdots & 0 \\ 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 1 \end{pmatrix}$
- $2 |0\rangle_n {}_n\langle 0| - \hat{\mathbb{I}} = 2( \quad ) - \begin{pmatrix} 1 & 0 & \cdots & 0 \\ 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 1 \end{pmatrix} = ( \quad )$

# Implementation of the Grover's diffusion operator

- $$\hat{X}^{\otimes n} (2 |0\rangle_{nn} \langle 0| - \hat{\mathbb{I}}) \hat{X}^{\otimes n} = 2 \hat{X}^{\otimes n} |0\rangle_{nn} \langle 0| \hat{X}^{\otimes n} - \hat{X}^{\otimes n} \hat{\mathbb{I}} \hat{X}^{\otimes n} = 2 |1\rangle_{nn} \langle 1| - \underbrace{\hat{X}^{\otimes n} \hat{X}^{\otimes n}}_{=\hat{\mathbb{I}}} =$$

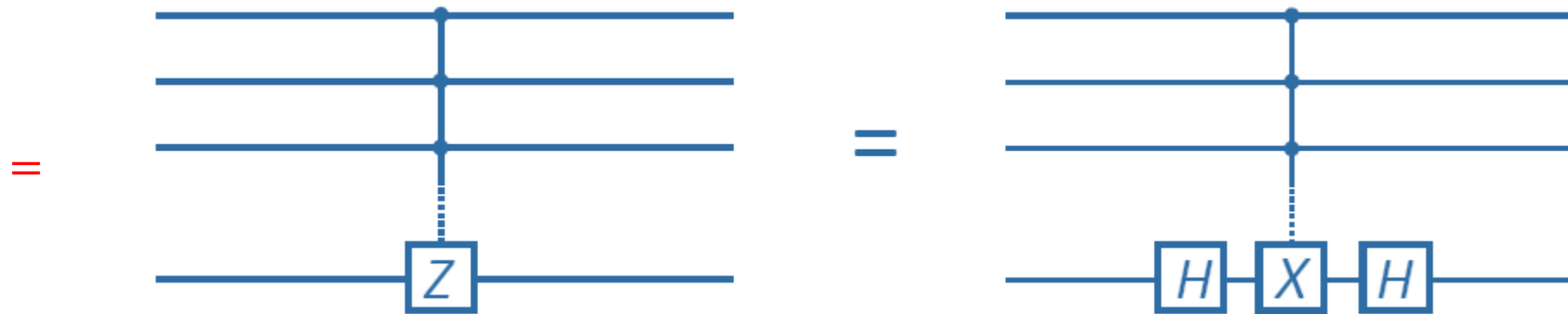
$$= 2 |1\rangle_{nn} \langle 1| - \hat{\mathbb{I}}$$

gdzie  $|1\rangle_n = \underbrace{|11 \dots 1\rangle}_n$

- $$2 |1\rangle_{nn} \langle 1| - \hat{\mathbb{I}} = 2 \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix} - \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix} = -(\hat{\mathbb{I}} - 2 |1\rangle_{nn} \langle 1|)$$

# Implementation of the Grover's diffusion operator

- $(\hat{\mathbb{I}} - 2|1\rangle_n\langle 1|)$



$$C_{n-1}\hat{Z} = (\hat{H} \otimes \hat{\mathbb{I}})C_{n-1}\hat{X}(\hat{H} \otimes \hat{\mathbb{I}})$$

$$\hat{Z} = \hat{H} \hat{X} \hat{H}, \hat{X}^{\otimes n} \hat{X}^{\otimes n} = \hat{\mathbb{I}}$$

$$\hat{W} = -\hat{H}^{\otimes n} \hat{X}^{\otimes n} C_{n-1} \hat{Z} = \hat{X}^{\otimes n} (\hat{H} \otimes \hat{\mathbb{I}}) C_{n-1} \hat{X} (\hat{H} \otimes \hat{\mathbb{I}}) \hat{X}^{\otimes n} \hat{H}^{\otimes n}$$

## Tasks:

$$f(x) = \begin{cases} 1, & x = a \\ 0, & x \neq a \end{cases}$$

1. Create the Matrix of  $U_f$  operation:

1.1 as XOR Oracle (**Group @9:45 take  $a=00$ , Group @11:45 take  $a=01$** )

1.2 as Phase Oracle (**Group @9:45 take  $a=011$ , Group @11:45  $a=100$** )

2. From the Matrix of  $U_f$  operation create quantum gate

**For  $n=3$**

3. Calculate the explicit matrix form of operator (lecture 6, frame 35):  $|0\rangle_{nn}\langle 0|$

4. Calculate the explicit matrix form of operator (lecture 6, frame 35-36):  $2|0\rangle_{nn}\langle 0| - \hat{\mathbb{I}}$

5. Prove that (lecture 6, frame 34):  $\hat{X}^{\otimes n}(2|0\rangle_{nn}\langle 0| - \hat{\mathbb{I}})\hat{X}^{\otimes n} = 2|1\rangle_{nn}\langle 1| - \hat{\mathbb{I}}$

6. Calculate the explicit matrix form of operator: **W** Grover's diffusion operator (lecture 6, frame 33-43)

7. Using the Qiskit template (file lab4\_template.ipynb) create quantum gates representing  $U_f$  and **W**. (Tip use X, H and CNOT gates to create **W** )



# A1 Appendix

# XOR Oracle

$$U_f (|y\rangle \otimes |x\rangle) = |y \oplus f(x)\rangle \otimes |x\rangle \quad \mathbf{a = 10}$$

$ i\rangle$	$ y\rangle x\rangle$	$U_f ( y\rangle \otimes  x\rangle)$
$ 0\rangle$	$ 0\ 00\rangle$	$ 0\ 00\rangle$
$ 1\rangle$	$ 0\ 01\rangle$	$ 0\ 01\rangle$
$ 2\rangle$	$ 0\ \mathbf{10}\rangle$	$ 1\ \mathbf{10}\rangle$
$ 3\rangle$	$ 0\ 11\rangle$	$ 0\ 11\rangle$
$ 4\rangle$	$ 1\ 00\rangle$	$ 1\ 00\rangle$
$ 5\rangle$	$ 1\ 01\rangle$	$ 1\ 01\rangle$
$ 6\rangle$	$ 1\ \mathbf{10}\rangle$	$ 0\ \mathbf{10}\rangle$
$ 7\rangle$	$ 1\ 11\rangle$	$ 1\ 11\rangle$

# Phase Oracle

$$U_f |x\rangle = (-1)^{f(x)} |x\rangle \quad \mathbf{a = 010}$$

$ i\rangle$	$ x\rangle$	$U_f  x\rangle$
$ 0\rangle$	$ 0\ 00\rangle$	$ 0\ 00\rangle$
$ 1\rangle$	$ 0\ 01\rangle$	$ 0\ 01\rangle$
$ 2\rangle$	$ \mathbf{0\ 10}\rangle$	$-\mathbf{ 010\rangle}$
$ 3\rangle$	$ 0\ 11\rangle$	$ 0\ 11\rangle$
$ 4\rangle$	$ 1\ 00\rangle$	$ 1\ 00\rangle$
$ 5\rangle$	$ 1\ 01\rangle$	$ 1\ 01\rangle$
$ 6\rangle$	$ 1\ 10\rangle$	$ 1\ 10\rangle$
$ 7\rangle$	$ 1\ 11\rangle$	$ 1\ 11\rangle$

# Construct the matrix for XOR Oracle $a = 10$

**base** =  $\{|0\rangle, |0\rangle, \dots, |n-1\rangle\}$   
*base vector numbers  $\Leftrightarrow$  decimal numbers*

$$\langle 0|0\rangle \Leftrightarrow \langle 0|0\rangle = 1$$

$|j\rangle$  – row,  $\langle i|$  – column

$$U_f = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j|U_f|i\rangle |j\rangle \langle i|$$

$$U_{f,j,i} = \langle j|U_f|i\rangle$$

$$U_{f20} = \langle 2|U_f|0\rangle = \langle 010|000\rangle = 0$$

$$U_{f21} = \langle 2|U_f|1\rangle = \langle 010|001\rangle = 0$$

$$U_{f22} = \langle 2|U_f|2\rangle = \langle 010|1\mathbf{10}\rangle = 0$$

$$U_{f26} = \langle 2|U_f|6\rangle = \langle 010|0\mathbf{10}\rangle = 1$$

$$U_{f62} = \langle 6|U_f|2\rangle = \langle 110|1\mathbf{10}\rangle = 1$$

$$\delta_{ji} = \begin{cases} 1, & j = i \\ 0, & j \neq i \end{cases}$$

$$\begin{aligned} \langle 2|U_f|6\rangle &= \langle 2|\sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j|U_f|i\rangle |j\rangle \langle i|6\rangle = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j|U_f|i\rangle \langle 2|j\rangle \langle i|6\rangle = \\ &= \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j|U_f|i\rangle \delta_{2j} \delta_{i6} = \sum_{i=0}^{N-1} \langle j|U_f|6\rangle \delta_{2j} = \langle 2|U_f|6\rangle = \langle 010|0\mathbf{10}\rangle = 1 \end{aligned}$$

$U_f(|y\rangle \otimes |x\rangle)$   
See table



# Create the Matrix of $U_f$ operation

XOR Oracle

*j column*

*i row*

	0	1	2	3	4	5	6	7
0	1	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0
2	0	0	0	0	0	0	1	0
3	0	0	0	1	0	0	0	0
4	0	0	0	0	1	0	0	0
5	0	0	0	0	0	1	0	0
6	0	0	1	0	0	0	0	0
7	0	0	0	0	0	0	0	1



quantum gate

Phase Oracle

*j column*

*i row*

	0	1	2	3	4	5	6	7
0	1	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0
2	0	0	-1	0	0	0	0	0
3	0	0	0	1	0	0	0	0
4	0	0	0	0	1	0	0	0
5	0	0	0	0	0	1	0	0
6	0	0	0	0	0	0	1	0
7	0	0	0	0	0	0	0	1



quantum gate

## Hint Task 3-5

$$|0\rangle_{n=3} = |000\rangle = |0\rangle \otimes |0\rangle \otimes |0\rangle$$

$${}_{n=3}\langle 0| = \langle 000| = \langle 0| \otimes \langle 0| \otimes \langle 0|$$

$$|0\rangle_{n=3} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \otimes (1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0) = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

## Construct the matrix for a 2 bits NOT-gate (N=2)

- $|j\rangle$  – row,  $\langle i|$  – column (22.6) ,  $\forall |x\rangle_n = (-1)^{f(x)} |x\rangle_n$  (23.9)

$$\begin{aligned} U_{NOT} &= \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j| U_{NOT} |i\rangle |j\rangle \langle i| = \\ &= \langle 0| U_{NOT} |0\rangle |0\rangle \langle 0| + \langle 1| U_{NOT} |0\rangle |1\rangle \langle 0| + \langle 0| U_{NOT} |1\rangle |0\rangle \langle 1| + \langle 1| U_{NOT} |1\rangle |1\rangle \langle 1| = \\ &= \langle 0||1\rangle |0\rangle \langle 0| + \langle 1||1\rangle |1\rangle \langle 0| + \langle 0||0\rangle |0\rangle \langle 1| + \langle 1||0\rangle |1\rangle \langle 1| = \\ &= 0 * |0\rangle \langle 0| + 1 * |1\rangle \langle 0| + 1 * |0\rangle \langle 1| + 0 * |1\rangle \langle 1| = |1\rangle \langle 0| + |0\rangle \langle 1| = \\ &= \begin{pmatrix} 0 \\ 1 \end{pmatrix} * (1 \ 0) + \begin{pmatrix} 1 \\ 0 \end{pmatrix} * (0 \ 1) = \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \end{aligned}$$



The End