## CSP PRIMER<sup>2</sup>

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## 1. Preliminary Definitions and Notations

**a.** If A and B are sets, then the Cartesian product of A and B is denoted by  $A \times B$  and is defined to be the set of all ordered pairs (a, b) such that a belongs to A and b belongs to B; that is,

$$A \times B = \{(a, b) : a \in A, b \in B\}.$$

If A is a set and n is a natural number, then the n-th Cartesian power of A, denoted by  $A^n$ , is the set of n-tuples of elements of A; that is,

$$A^n = A \times A \times \cdots \times A = \{(a_1, a_2, \dots, a_n) : a_i \in A \text{ for } 1 \le i \le n\}.$$

If |A| denotes the number of elements in the set A, then  $A^n$  contains exactly  $n^{|A|}$  elements. We define the k-th projection on n-tuples to be the function  $\pi_k : A^n \to A$  given by

$$\pi_k(a_1, a_2, \dots, a_n) = a_k.$$

That is,  $\pi_k$  simply projects onto the k-th "coordinate" of  $A^n$  by picking out the k-th entry of  $(a_1, a_2, \ldots, a_n)$ .

It is important to note that each n-tuple  $(a_1, a_2, \ldots, a_n) \in A^n$  in the Cartesian power of A defines a function mapping the set  $\{1, 2, \ldots, n\}$  to the set A. Specifically,  $(a_1, a_2, \ldots, a_n)$  is the function  $a: \{1, 2, \ldots, n\} \to A$  given by  $a(k) = \pi_k(a_1, \ldots, a_n) = a_k$ . It may seem like we're making something out of nothing here, but this slight change of viewpoint (tuples as functions) can be very useful. Thus, the Cartesian power  $A^n$  "is" the set of all functions from  $\{1, 2, \ldots, n\}$  to A.

**Exercise 1** (easy). Use a counting principle to explain why there are  $n^{|A|}$  elements in  $A^n$ .

**b.** For two sets X and Y, we denote by  $Y^X$  the set of all functions  $f: X \to Y$  that map each element  $x \in X$  to some element  $f(x) \in Y$ . Take a moment to observe the analogy with tuples. Indeed, the function f is an |X|-tuple of elements of Y. This analogy is exact when X happens to be a countable set, in which case we can enumerate its elements,  $X = \{x_1, x_2, \dots\}$ . This allows us to represent f by the tuple consisting of its values:  $f = (f(x_1), f(x_2), \dots)$ , and applying f to an "index"  $x_k \in X$  gives the k-th element in the tuple f:

$$f(x_k) = \pi_k(f(x_1), f(x_2), \dots) = \pi_k f.$$

But the analogy with tuples is useful even when the domain is not countable and we may think of f as a "tuple"  $(f(x):x\in X)\in Y^X$ . There is no harm in thinking of  $Y^X$  as a Cartesian

Date: 12 May 2015.

power in this case as well, as long as you don't make the mistake of assuming we can always enumerate the "index set" X. In particular, if X is an uncountable set, then you shouldn't write  $(f(x_1), f(x_2), \ldots)$  instead of  $(f(x) : x \in X)$ . That is, you cannot write the values of f as an enumerated list of elements of Y. There are simply too many values!

**Exercise 2** (easy). If X has |X| = m elements and Y has |Y| = n elements, how many functions are there from X to Y? In other words, what is the cardinality of the set  $Y^X$ ? How many of these functions are one-to-one? (Hint: handle the cases  $m \le n$  and m > n separately.)

**c.** The k-th projection operation on n-tuples defined above has domain  $A^n$  and codomain A, so it belongs to the set  $A^{(A^n)}$ . Note that  $A^{(A^n)}$  has the form  $Y^X$  described above; in this case Y = A and  $X = A^n$ .

Unfortunately, we have to use a slightly uglier, but more precise, notation for projections. We let  $\pi_k^n: A^n \to A$  denote the k-th projection on  $A^n$ , since we will occasionally refer to projections of other arities, say,  $\pi_k^m: A^m \to A$  in the same context.

Let A be any set and let n be any natural number. Then  $A^{(A^n)}$  denotes the set of all n-ary functions on A, that is, the set of all functions  $f: A^n \to A$  taking an n-tuple  $(a_1, \ldots, a_n) \in A^n$  to some value  $f(a_1, \ldots, a_n) \in A$ . In symbols,

$$A^{(A^n)} = \{f : A^n \to A\}$$

(You are probably familiar with such "multivariable" functions from calculus.) We let Op(A) denote the set of all functions from  $A^n$  to A for all natural numbers n. In symbols,

$$\operatorname{Op}(A) = \bigcup_{n \in \mathbb{N}} A^{(A^n)}.$$

**d.** Let n and k be natural numbers, and suppose that  $f \in A^{(A^n)}$  and  $g_1, g_2, \ldots, g_n \in A^{(A^k)}$ . Then we define a new k-ary operation  $f[g_1, g_2, \ldots, g_n]$  by

$$(a_1, a_2, \dots, a_k) \mapsto f(g_1(a_1, \dots, a_k), \dots, g_n(a_1, \dots, a_k))$$

called the *generalized composition* of f with  $g_1, \ldots, g_n$ . Note that, unlike the ordinary composition of unary functions, the generalized composition exists only when the arities match up correctly.

Just as the set of unary operations forms a monoid<sup>1</sup> under the operation of composition, we can form an algebraic structure whose elements are members of Op(A) with the operation of generalized composition.

**Definition 1.** Let A be a nonempty set. A clone on A is a subset  $\mathscr{C}$  of  $\operatorname{Op}(A)$  that contains all projection operations and is closed under generalized composition.

**Exercise 3.** Show that the set Proj(A) of all projections  $\{\pi_k^n : n \in \mathbb{N}, k \in \mathbb{N}\}$  on the set A is a clone.

<sup>&</sup>lt;sup>1</sup>A monoid,  $\langle X, \circ, e \rangle$ , is a set X together with an associative binary operation  $\circ$  and an identity element e. Note that the set  $A^A$  of all unary functions  $f: A \to A$ , along with function composition  $f \circ g$  and the identity map  $\mathrm{id}_A$ , forms a monoid,  $\langle A^A, \circ, \mathrm{id}_A \rangle$ .

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**Exercise 4.** Show that the set  $\mathcal{E}(A)$  of all idempotent operations on A is a clone. An operation f is called idempotent if f(a, a, ..., a) = a for all  $a \in A$ .

Given a set  $F \subseteq \operatorname{Op}(A)$  of functions, we can consider the smallest clone that contains F. This is called the *clone generated by* F and is denoted by  $\operatorname{Clo}(F)$ . It is no too hard to prove that the clone  $\operatorname{Clo}(F)$  can be built recursively, as in the following theorem:

**Theorem 1.** Let A be a set and  $F \subseteq Op(A)$  a set of operations on A. Define

$$F_0 = \operatorname{Proj}(A)$$

$$F_{n+1} = F_n \cup \{f[g_1, \dots, g_k] : f \in F, k = \operatorname{arity}(f), g_1, \dots, g_k \in F_n \cap \operatorname{Op}(A)\},$$

$$Then \operatorname{Clo}(F) = \bigcup_n F_n.$$