
1. Like Divides, but more complicated

(a)

Let a, b, c be positive arbitrary integers.

Define the following relation R on the positive integers. $(x, y) \in R$ if and only if x/y is an even integer.

We prove R is a transitive relation which means S is transitive if and only if $\forall a \forall b \forall c [(a, b) \in R \wedge (b, c) \in R \rightarrow (a, c) \in R]$

Suppose $(a, b) \in R \wedge (b, c) \in R$.

By definition of R , we could know that there is an even integer z equals to a/b which means $a = b \cdot z$ where $z = 2 \cdot x$ for some integer x . By definition of R again, we could know there is an even integer w equals to b/c which means $b = c \cdot w$ where $w = 2 \cdot y$ for some integer y . By substituting b in $a = b \cdot z$ with $b = c \cdot w$, we have $a = c \cdot w \cdot z$ which is $a = c \cdot 2 \cdot y \cdot 2 \cdot x$. By algebra, we have $a = c \cdot 4 \cdot y \cdot x$. This means there is an even integer $4 \cdot y \cdot x$ equals to a/c by algebra. Since x, y are integers, $4 \cdot y \cdot x$ is also an even integer. Thus, we know $(a, c) \in R$.

Since a, b, c are arbitrary, we can conclude that if $(a, b) \in R \wedge (b, c) \in R$, then $(a, c) \in R$. Therefore, we prove R is a transitive relation.

(b)

We will prove R is not a symmetric relation by counter example.

R is symmetric if and only if $\forall a \forall b [(a, b) \in R \rightarrow (b, a) \in R]$.

Let $(4, 2) \in R$. We know that $(4, 2) \in R$ because $4/2 = 2 = 2 \cdot 1$ by the definition of even and R . But $(2, 4) \notin R$ since $2/4 = 1/2$ which is not an even integer. Thus, not all element satisfies $\forall a \forall b [(a, b) \in R \rightarrow (b, a) \in R]$.

Thus, R is not a symmetric relation.

2. Com-pair-isons

(a)

Let $X = \{x_1, x_2\}, Y = \{y_1, y_2\}$ with $x_1 < x_2$ and $y_1 < y_2$ be arbitrary elements that satisfy $X \in A \wedge Y \in A$. Suppose $X \preceq Y$ and $Y \preceq X$ satisfy the relation \preceq defined in the problem specification. We prove the relation \preceq is antisymmetry by showing for all $a, b \in A, [(a, b) \in R \wedge (b, a) \in R \rightarrow a = b]$.

Because of $X \preceq Y$, we know that $x_1 \leq y_1$ and $x_2 \leq y_2$ by the definition of \preceq . Because of $Y \preceq X$, we could know $y_1 \leq x_1$ and $y_2 \leq x_2$ by the definition of \preceq . Thus, since we know $x_1 \leq y_1$ and $x_2 \leq y_2$ and $y_1 \leq x_1$ and $y_2 \leq x_2$, we have $x_1 = y_1$ and $x_2 = y_2$ by algebra. Therefore, we get $X = Y$.

Since X, Y are arbitrary elements in A , we have for all $a, b \in A, [(a, b) \in R \wedge (b, a) \in R \rightarrow a = b]$. Thus, we prove the relation \preceq is antisymmetry.

(b)

No.

We prove it by counter example.

Consider $X = \{1, 5\} \in A, Y = \{2, 3\} \in A$. We know that $\{1, 5\} \not\preceq \{2, 3\}$ because $5 > 3$ and $\{2, 3\} \not\preceq \{1, 5\}$ because $2 > 1$.

Thus, \preceq is not a total order on A .

(c)

We prove it by counter example.

Consider $X = \{3, 9\} \in A, Y = \{4, 7\} \in A, Z = \{2, 8\} \in A$. $X \preceq Y$ because $3 \leq 4$ and $Y \preceq Z$ because $7 \leq 8$. However, $X \not\preceq Z$ since $3 \not\leq 2$ and $9 \not\leq 8$. Thus, we know not all elements in A satisfy $\forall X \forall Y \forall Z [X \preceq Y \wedge Y \preceq Z] \rightarrow X \preceq Z$.

5. Extra-Arden-ary Translations

(a) $\forall A \forall B [\neg \text{Matches}(\varepsilon, B) \rightarrow \exists C ((C = A \cdot C \vee B) \wedge \forall D ((D = A \cdot D \vee B) \rightarrow (C = D)))]$

(b)

$$\begin{aligned}
& \neg \forall A \forall B [\neg \text{Matches}(\varepsilon, B) \rightarrow \exists C ((C = A \cdot C \cup B) \wedge \forall D ((D = A \cdot D \cup B) \rightarrow (C = D)))] \\
& \neg \forall A \forall B [\neg \neg \text{Matches}(\varepsilon, B) \vee \exists C ((C = A \cdot C \cup B) \wedge \forall D ((D = A \cdot D \cup B) \rightarrow (C = D)))] \\
& \neg \forall A \forall B [\text{Matches}(\varepsilon, B) \vee \exists C ((C = A \cdot C \cup B) \wedge \forall D ((D = A \cdot D \cup B) \rightarrow (C = D)))] \\
& \neg \forall A \forall B [\text{Matches}(\varepsilon, B) \vee \exists C ((C = A \cdot C \cup B) \wedge \forall D (\neg(D = A \cdot D \cup B) \vee (C = D)))] \\
& \exists A \neg \forall B [\text{Matches}(\varepsilon, B) \vee \exists C ((C = A \cdot C \cup B) \wedge \forall D (\neg(D = A \cdot D \cup B) \vee (C = D)))] \\
& \exists A \exists B \neg [\text{Matches}(\varepsilon, B) \vee \exists C ((C = A \cdot C \cup B) \wedge \forall D (\neg(D = A \cdot D \cup B) \vee (C = D)))] \\
& \exists A \exists B [\neg \text{Matches}(\varepsilon, B) \wedge \neg \exists C ((C = A \cdot C \cup B) \wedge \forall D (\neg(D = A \cdot D \cup B) \vee (C = D)))] \\
& \exists A \exists B [\neg \text{Matches}(\varepsilon, B) \wedge \forall C \neg ((C = A \cdot C \cup B) \wedge \forall D (\neg(D = A \cdot D \cup B) \vee (C = D)))] \\
& \exists A \exists B [\neg \text{Matches}(\varepsilon, B) \wedge \forall C (\neg(C = A \cdot C \cup B) \vee \neg \forall D (\neg(D = A \cdot D \cup B) \vee (C = D)))] \\
& \exists A \exists B [\neg \text{Matches}(\varepsilon, B) \wedge \forall C (\neg(C = A \cdot C \cup B) \vee \exists D (\neg \neg(D = A \cdot D \cup B) \wedge \neg(C = D)))] \\
& \exists A \exists B [\neg \text{Matches}(\varepsilon, B) \wedge \forall C ((C \neq A \cdot C \cup B) \vee \exists D ((D = A \cdot D \cup B) \wedge (C \neq D)))]
\end{aligned}$$

(c) There exists some regular expressions A and B such that the empty string does not match B and for all regular expression C, C does not satisfy $C = A \cdot C \cup B$ or C is not unique.

6. Proof By Contradiction

(a) $\forall x \forall y ((\text{Rational}(x) \wedge \neg \text{Rational}(y)) \rightarrow \neg \text{Rational}(x + y))$

(b) $\exists x \exists y (\text{Rational}(x) \wedge \neg \text{Rational}(y) \wedge \text{Rational}(x + y))$

(c)

I prove this claim by contradiction.

Denote r to be a rational number and i to be an irrational number

Suppose, for the sake of contradiction, the sum of a rational number and an irrational number, s , is rational which is $r + i = s$. By the definition of rational number, r and s can be made by dividing two integers, so $r = \frac{a}{b}$ and $s = \frac{m}{n}$ for some integers a, b, m , and n . Then, subtracting both sides of equation by r , we have $i = s - r$ which equals to $i = \frac{m}{n} - \frac{a}{b}$. By algebra, we have $i = \frac{mb - na}{nb}$. Since a, b, m , and n are integers, $mb - na$ and nb are also integers. So, we conclude i is a rational number which is a contradiction. Thus, our assumption that s is a rational number must be false. Therefore, I prove the claim.

7. Buggy

(a) The biggest bug is step 2. The negation is not correct, while when we do proof by contradiction, we need to assume the negation and start from that.

8. feedback

(a) 9 hours

(b) question 4

(c) No