

(a) Consider these three transactions:

- $T_1 : R_1(A), R_1(B), W_1(A), W_1(B), Co_1$
- $T_2 : R_2(B), W_2(B), R_2(C), W_2(C), Co_2$
- $T_3 : R_3(C), W_3(C), R_3(A), W_3(A), Co_3$

i. Schedule 1:

$R_2(B), W_2(B), R_3(C), W_3(C), R_3(A), W_3(A), Co_3, R_2(C), W_2(C), Co_2, R_1(A), R_1(B), W_1(A), W_1(B), Co_1$

Is this schedule conflict-serializable? If yes, indicate a serialization order.

ii. Schedule 2:

$R_2(B), W_2(B), R_3(C), W_3(C), R_1(A), R_1(B), W_1(A), W_1(B), Co_1, R_2(C), W_2(C), Co_2, R_3(A), W_3(A), Co_3$

Is this schedule conflict-serializable? If yes, indicate a serialization order.

(b) Consider the following three transactions:

- $T_1 : R_1(A), W_1(B), Co_1$
- $T_2 : R_2(B), W_2(C), Co_2$
- $T_3 : R_3(C), W_3(D), Co_3$

Given an example of a conflict-serializable schedule that has the following properties:
transaction T_1 commits before transaction T_3 starts, and the equivalent serial order is T_3, T_2, T_1 .