# Optimization in Industry Assignment 2 Mathematical Model for Optimal ATM Refill Problem

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### Sets

- 1.  $\mathcal{I}$ : Set of ATMs, indexed by i (i.e.,  $i = 1, ..., |\mathcal{I}|$ ).
- 2.  $\mathcal{V}$ : Set of vehicles, indexed by v (i.e.,  $v = 1, \ldots, |\mathcal{V}|$ ).
- 3.  $\mathcal{D}$ : Set of days in the planning horizon, indexed by d (e.g.,  $d = 1, \ldots, 7$ ).
- 4.  $\mathcal{K}$ : Set of denominations, indexed by k.

# **Parameters**

- 1. Demand $_{i,d}$ : Cash demand at ATM i on day d.
- 2.  $I_i^0$ : Initial cash inventory at ATM i (at start of day 1).
- 3.  $C_i$ : Maximum cash capacity of ATM i.
- 4.  $L_i$ : Minimum cash level required at ATM i.
- 5.  $Q_{\min}$ : Minimum cash deposit per refill (e.g., 50K).
- 6.  $Q_{\text{unit}}$ : Cash deposit unit (e.g., 10K), so that any nonzero refill is a multiple of  $Q_{\text{unit}}$ .
- 7.  $Q_{\text{unit},k}$ : Cash deposit unit (e.g., 10K) for denomination k, so that any nonzero refill is a multiple of  $Q_{\text{unit},k}$ .
- 8.  $\operatorname{cap}_v$ : Maximum cash carrying capacity of vehicle v (security limited).
- 9.  $cost_v$ : Cost associated with using vehicle v for a day.
- 10.  $\max \operatorname{Visit}_{v,d}$ : Maximum number of ATMs vehicle v can service on day d (e.g., 20 on weekdays, 30 on weekends).
- 11. S: Security gap in days. If vehicle v visits ATM i on day d, it cannot visit the same ATM again in the next S-1 days.

- 12.  $\Delta_x$ : Maximum allowed difference in the number of ATMs visited by any two vehicles in the planning horizon.
- 13.  $\Delta_y$ : Maximum allowed difference in the cash deposited by any two vehicles in the planning horizon.
- 14.  $r_{i,d,k}$ : Denomination requirement for ATM i on day d for denomination k (soft constraint).
- 15.  $w_1$ ,  $w_2$ ,  $w_3$ ,  $w_4$ ,  $w_5$ : Weights for the objective function corresponding to vehicles used, ATM inventory holding cost, and number of visits, respectively.
- 16.  $M = 1 + \max_{v \in \mathcal{V}} \text{cap}_v$ : A sufficiently large constant (Big-M constant).

## **Decision Variables**

- 1.  $x_{i,v,d} \in \{0,1\}$ : Equals 1 if vehicle v refills ATM i on day d; 0 otherwise.
- 2.  $q_{i,v,d} \in \mathbb{Z}_{\geq 0}$ : Number of cash deposit units  $Q_{\text{unit}}$ , delivered to ATM i by vehicle v on day d.
- 3.  $y_{i,v,d} \geq 0$ : Amount of cash delivered to ATM i by vehicle v on day d.

$$y_{i,v,d} = Q_{\text{unit}} \cdot q_{i,v,d}, \quad q_{i,v,d} \in \mathbb{Z}_{>0}$$

with the additional condition that if  $q_{i,v,d} > 0$  then  $q_{i,v,d} \ge \frac{Q_{\min}}{Q_{\min}}$ .

- 4.  $z_{v,d} \in \{0,1\}$ : Equals 1 if vehicle v is used on day d; 0 otherwise.
- 5.  $I_{i,d}$ : Cash inventory at ATM i at the end of day d.
- 6.  $q_{i,v,d,k} \in \mathbb{Z}_{\geq 0}$ : Number of cash deposit units  $Q_{\text{unit},k}$  of currency k, delivered to ATM i by vehicle v on day d.
- 7.  $y_{i,v,d,k} \ge 0$ : Cash delivered in denomination k for ATM i by vehicle v on day d.

$$y_{i,v,d,k} = Q_{\text{unit},k} \cdot q_{i,v,d,k}, \quad q_{i,v,d,k} \in \mathbb{Z}_{>0}$$

8. penalty<sub>i,d,k</sub>  $\geq$  0: Unmet demand for cash delivered in denomination k for ATM i on day d.

# **Objective Function**

We consider a weighted-sum objective function:

#### Minimize:

$$Z = w_1 \cdot \sum_{d \in \mathcal{D}} \sum_{v \in \mathcal{V}} z_{v,d} + w_2 \cdot \sum_{d \in \mathcal{D}} \sum_{i \in \mathcal{I}} I_{i,d} + w_3 \cdot \sum_{d \in \mathcal{D}} \sum_{v \in \mathcal{V}} \sum_{i \in \mathcal{I}} x_{i,v,d}$$
$$+ w_4 \cdot \sum_{i \in \mathcal{I}} \sum_{d \in \mathcal{D}} \sum_{k \in \mathcal{K}} \text{penalty}_{i,d,k} + w_5 \cdot \sum_{v \in \mathcal{V}} \sum_{d \in \mathcal{D}} \text{cost}_v \cdot z_{v,d}$$

#### **Explanation:**

- The first term represents the number of vehicles used.
- The second term represents the total ATM cash inventory holding cost.
- The third term represents the total number of ATM visits.
- The fourth term represents the unmet denomination preferences at each of the ATM.
- The fifth term represents the cost using the vehicles during the planning horizon.

## Constraints

#### A. ATM Inventory Balance and Capacity

(A1) Inventory Update: For each  $i \in \mathcal{I}$  and  $d \in \mathcal{D}$ ,

$$I_{i,d} = \begin{cases} I_i^0 + \sum_{v \in \mathcal{V}} y_{i,v,1} - \operatorname{Demand}_{i,1}, & d = 1, \\ I_{i,d-1} + \sum_{v \in \mathcal{V}} y_{i,v,d} - \operatorname{Demand}_{i,d}, & d \geq 2. \end{cases}$$

(A2) Inventory Limits: For all  $i \in \mathcal{I}$  and  $d \in \mathcal{D}$ ,

$$L_i \leq I_{i,d}$$

$$I_{i,d} \leq C_i$$

#### **Explanation:**

- (A1) Keeps track of the cash in the ATM across the planning horizon.
- (A2) Ensures the cash in the inventory does not go beyond the lower and upper limits.

- B. Service (Assignment) Constraints
- (B1) One Refill per ATM per Day: For every  $i \in \mathcal{I}$  and  $d \in \mathcal{D}$ ,

$$\sum_{v \in \mathcal{V}} x_{i,v,d} \le 1$$

**(B2)** Linking x and y: For all  $i \in \mathcal{I}$ ,  $v \in \mathcal{V}$ , and  $d \in \mathcal{D}$ ,

$$y_{i,v,d} \ge Q_{\min} x_{i,v,d}$$

$$y_{i,v,d} \leq M \cdot x_{i,v,d}$$

#### **Explanation:**

- (B1) Ensures that each ATM is filled atmost once in a day.
- (B2) Ensures that cash is either not deposited or if deposited at least  $Q_{\min}$  of cash is deposited.
- C. Vehicle Capacity and Workload
- (C1) Vehicle Loading Limit: For each  $v \in \mathcal{V}$  and  $d \in \mathcal{D}$ ,

$$\sum_{i \in \mathcal{T}} y_{i,v,d} \le \operatorname{cap}_v \cdot z_{v,d}$$

(C2) Maximum Visits per Vehicle per Day: For all  $v \in \mathcal{V}$  and  $d \in \mathcal{D}$ ,

$$\sum_{i \in \mathcal{T}} x_{i,v,d} \le \max_{i \in \mathcal{T}} \text{Visit}_{v,d} \cdot z_{v,d}$$

#### **Explanation:**

- (C1) Ensures that total cash to be delivered by vehicle v on a day d does not exceed its capacity.
- (C2) Ensures a vehicle v does not visit more than  $\max Visit_{v,d}$  ATMs on day d.

## D. Security and Scheduling Constraints

(D1) Spread Out ATM Visits: For all  $v \in \mathcal{V}$ , for each  $i \in \mathcal{I}$  and for each day  $d \in \mathcal{D}$  such that  $d + S - 1 \leq \max(\mathcal{D})$ ,

$$x_{i,v,d} + \sum_{d'=d+1}^{\min(d+S-1,\max(\mathcal{D}))} x_{i,v,d'} \le 1.$$

#### **Explanation:**

• (D1) Ensures if vehicle v visits ATM i on day d, it cannot visit the same ATM again in the next S-1 days.

#### E. Workload Balance Across Vehicles

(E1) Number of ATMs visited by vehicles: For each  $v, v' \in \mathcal{V}$ ,

$$\sum_{i \in \mathcal{I}} \sum_{d \in \mathcal{D}} \left( x_{i,v,d} - x_{i,v',d} \right) \le \Delta_x$$

(E2) Total Cash Deposited in ATMs by vehicles: For each  $v, v' \in \mathcal{V}$ ,

$$\sum_{i \in \mathcal{T}} \sum_{d \in \mathcal{D}} \left( \frac{y_{i,v,d}}{\operatorname{cap}_v} - \frac{y_{i,v',d}}{\operatorname{cap}_{v'}} \right) \le \Delta_y$$

#### **Explanation:**

- (E1) Ensures the difference between the number of ATM visits by any two vehicles v and v' across the planning horizon does not exceed  $\Delta_x$ .
- (E2) Ensures the weighted difference between the total cash deposited by any two vehicles v and v' across the planning horizon does not exceed  $\Delta_y$ . Here we have taken the Weights to be the inverse of the vehicle capacity.

#### F. Denomination Constraints (Soft)

**(F1)** For each  $i \in \mathcal{I}$ ,  $d \in \mathcal{D}$ , and  $k \in \mathcal{K}$ ,

$$\sum_{v \in \mathcal{V}} y_{i,v,d,k} + \text{penalty}_{i,d,k} \ge r_{i,d,k}$$

**(F2)** For each  $i \in \mathcal{I}$ ,  $d \in \mathcal{D}$ 

$$\sum_{k \in \mathcal{K}} y_{i,v,d,k} = y_{i,v,d}$$

#### **Explanation:**

- (F1) Penalizes unmet denomination demands for denomination k on day d at ATM i.
- (F2) Ensures total cash deposited is the sum of the total cash deposited in each denomination.

#### G. Deposit Size Discreteness

(G1) To ensure that if cash is either not delivered or at least  $Q_{\min}$  (in multiples of  $Q_{\text{unit}}$ ):

$$y_{i,v,d} = Q_{\text{unit}} \cdot q_{i,v,d}$$

$$q_{i,v,d} \ge \frac{Q_{\min}}{Q_{\text{unit}}} \cdot x_{i,v,d}$$

(G2) To ensure that cash in denomination k is delivered, is in multiples of  $Q_{\text{unit},k}$ :

$$y_{i,v,d,k} = Q_{\text{unit},k} \cdot q_{i,v,d,k}$$

#### **Explanation:**

- (G1) Ensures that for ATM i, cash is either not deposited by a vehicle v on day d or at least  $Q_{\min}$  is deposited in multiples of  $Q_{\min}$
- (G2) Ensures that cash in each denomination k is deposited in non-zero multiples of  $Q_{\min,k}$ .

## Assumptions

- The planning horizon is one week (7 days).
- Daily demands Demand<sub>i,d</sub> and initial inventories  $I_i^0$  are known.
- ATMs must always have cash between the levels  $L_i$  and  $C_i$ .
- Vehicles are loaded only at the beginning of the day (no mid-day refilling).
- Vehicles' cash capacities are limited by security considerations, which may be lower than their physical capacities.
- A refill at an ATM is done by at most one vehicle per day.
- For security, if a vehicle visits an ATM on day d, it cannot visit the same ATM again for the next S-1 days.
- Refills must be in multiples of  $Q_{\text{unit}}$  (e.g., 10K), with any nonzero refill being at least  $Q_{\text{min}}$  (e.g., 50K).
- The denomination requirement is a soft constraint; unmet denomination preferences are penalized but do not prevent meeting the overall cash demand.

# Output

The solution to this Mixed-Integer Programming (MIP) model will determine:

- Which vehicles are used on each day (via  $z_{v,d}$ ).
- Which ATMs are refilled by which vehicle on each day (via  $x_{i,v,d}$ ).
- The amount of cash delivered to each ATM on each day (via  $y_{i,v,d}$ ).

- The amount of cash in a denomination delivered to each ATM on each day (via  $y_{i,v,d,k}$ ).
- $\bullet\,$  The evolution of ATM cash inventories  $I_{i,d}$  over the planning horizon.