C Code Conversion to One Counter Automata for Reachability Analysis

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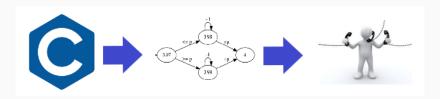
Overview

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- 2. Related work
- 3. Motivation
- 4. One counter automata
- 5. C code converter tool
- 6. Conclusion
- 7. Future work

Problem Overview

C code conversion

- · Analyse code reachability
- Convert C code to counter automaton
- · Code restrictions



Related Work

Related work I

- On parametric timed automata and one-counter machines

 Daniel Bundala, Joel Ouaknine, 2017

 Information and Computation
- Programs with Lists Are Counter Automata

 Ahmed Bouajjani, Marius Bozga, Peter Habermehl, Radu Iosif,
 Pierre Moro, TomÃąÅą Vojnar, 2006

 Computer Aided Verification
- Flat Counter Automata Almost Everywhere!

 JÃľrÃťme Leroux, GrÃľgoie Sutre, 2005

 Automated Technology for Verification and Analysis

Related work II



Motivation

Motivation

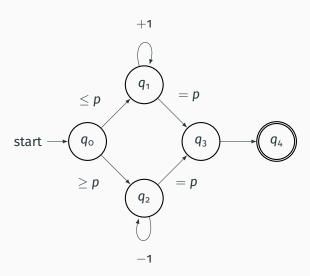
- · Resource constrained environments
- Code efficiency
- Development time

One Counter Automata

Counter automata I

- Non-Deterministic Finite Automata with counter(s)
- limited set of operations
 - increment the counter with parameter/constant
 - compare the counter to parameter/constant
- transition relation $\sigma \subseteq \mathbf{Q} \times \mathbf{op} \times \mathbf{Q}$

Counter automata II

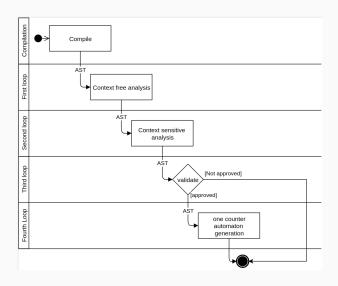


Counter automata III

- Advantages
 - · Simple conversion
 - · Known proven reachability solution
 - PSPACE-complete
 - 2NEXP time
- Disadvantages
 - · Enforced restrictions
 - Counters
 - Operations

C code converter Tool

Tool overview



Compilation



- ANTLR compilation
- · Generates AST
- · Boolean extension

First loop

- · Context free optimization
- · Remove redundant chains of nodes
- Unify node structures
- Enables improvements in further loops

Second loop

- · Main cleaner loop
- · Symbol table
- Constant folding/substitution
- · Dead code elimination
- · Counter tracking
- · No folding and/or substitution in conditional branches

Third loop

- validation loop
- · checks counter generation conditions
 - · Counter overlap
 - Counter type(s)
 - Return type
 - · Parameter types
 - · Parameter modification
 - · Used operations on counter
 - · Used conditional statements with counter

Fourth Loop

- · Counter automaton generation
- Simple loop
- One node per line
- Transitions
 - · Empty if no counter related operation occurs
 - · Condition if a conditional statement with counter occurs
 - Operation if an operation on the counter occurs

Conclusion

Conclusion

- · Improvement to current algorithms
- · Cost in execution time
- · Need for additional operation support
- · Suited for resource constrained environments
- (l+3)*m loops
 - m = number of nodes in AST
 - · l = number of lines in c code

Future work

Future work

- The reachability algorithm for the counter automata
- Additional supported operations
- Comparisons with other reachability algorithms