# 16-网络传输机制(TCP)实验三报告

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### 1、实验内容

- (1) 在上个实验的基础上,实现基于超时重传的 TCP 可靠数据传输,使得节点之间在有丢包网络中能够建立连接并正确传输数据。
- (2) 在节点 h1 上运行 TCP server,在 h2 上运行 TCP client,在丢包率为 2%的链路上向 h1 发送一个大小为几 MB 的文件,比较 h1 接收的文件与 h2 发送的文件是否相同。然后,在一端用修改后的 tcp stack.py 替换 tcp stack 执行,测试另一端。

# 2、实验流程

- (1) 完成代码编写并编译。
- (2) 测试正确性。

首先通过 mininet 环境(tcp topo.py)构建网络拓扑如下图:



在节点 h1 上运行 TCP server,在 h2 上运行 TCP client,进行抓包测试:

#### 执行:

bash create\_randfile.sh

sudo python tcp topo.py

mininet> xterm h1 h2

hl# wireshark

hl# ./tcp stack server 10001

or h1# python tcp stack2.py server 10001

h2# ./tcp stack client 10.0.0.1 10001

or h2# python tcp stack2.py client 10.0.0.1 10001

mininet> quit

diff server-output.dat client-input.dat

### 3、分析与实现

#### (1) 超时重传实现。

① 发送一个新的数据包时(携带数据或 SYN | FIN): 将该数据包放进发送队列末尾: 如果定时器未启动,则启动定时器。 void tcp\_send\_packet(struct tcp\_sock \*tsk, char \*packet, int len) tcp\_send\_buf\_append(tsk,packet,len); tcp\_set\_retrans\_timer(tsk); void tcp\_send\_control\_packet(struct tcp\_sock \*tsk, u8 flags) if (flags & (TCP\_SYN|TCP\_FIN)){ tsk->snd nxt += 1; tcp\_send\_buf\_append(tsk,packet,pkt\_size); tcp\_set\_retrans\_timer(tsk); } . . . 每个 sock 有一个链表实现的发送队列,定时器和收发包线程同时访问,需要用锁维护。 发送队列的每一项包括 packet,len,retrans\_times。 void tcp\_send\_buf\_append(struct tcp\_sock \*tsk, char\* packet, int len) tcp\_send\_buf\_entry\_t\* entry = malloc(sizeof(tcp\_send\_buf\_entry\_t)); entry->packet = malloc(len); entry->len = len; entry->retrans\_times = 0; memcpy(entry->packet,packet,len); pthread\_mutex\_lock(&tsk->send\_buf\_lock); list add tail(&entry->list,&tsk->send buf); pthread\_mutex\_unlock(&tsk->send\_buf\_lock); } 启动定时器时,通过检查定时器的 enable 域判断是否已经启动,如果已经启动,则无须 其它操作。 定时器链表可能有收发线程进行 unset 操作,定时器线程进行扫描操作,访问时需要加 锁。 void tcp\_set\_retrans\_timer(struct tcp\_sock \*tsk)

if(!tsk->retrans\_timer.enable){
 tsk->retrans\_timer.enable = 1;

```
tsk->retrans_timer.type = 0;
   pthread mutex lock(&retrans timer lock);
   list_add_tail(&tsk->retrans_timer.list, &retrans_timer_list);
   pthread_mutex_unlock(&retrans_timer_lock);
   tsk->retrans_timer.timeout = TCP_RETRANS_TIME;
 }
}
   ② 定时器超时时
   扫描当前 sock 的发送队列,发出第一个数据包。如果该数据包的重传次数超过3,返回
失败并发送 RST。否则进行重传,重传时,需要更新 ACK 和重传次数,定时器时间翻倍。
int tcp send buf retrans(struct tcp sock *tsk)
 pthread_mutex_lock(&tsk->send_buf_lock);
 tcp_send_buf_entry_t *entry;
 int times = 1;
 if(list empty(&tsk->send buf)) return 1;
 entry = list_entry(tsk->send_buf.next,tcp_send_buf_entry_t,list);
 if(entry->retrans_times>3){
   pthread_mutex_unlock(&tsk->send_buf_lock);
   return 0;
 }
 int len = entry->len;
 times = ++entry->retrans_times;
 char *packet = malloc(len);
 memcpy(packet,entry->packet,len);
 struct tcphdr *tcp = packet_to_tcp_hdr(packet);
 struct iphdr *ip = packet_to_ip_hdr(packet);
 tcp->ack = htonl(tsk->rcv_nxt);
 tcp->checksum = tcp_checksum(ip, tcp);
 ip_send_packet(packet,len);
 pthread_mutex_unlock(&tsk->send_buf_lock);
 return times;
}
   ③ 收到新的 ACK 时
   重置当前定时器时间。(如果定时器未开启,则无须操作)
   扫描发送队列,移除 seq end 小于当前 ACK 的项。若发送队列清空,则关闭定时器。
int tcp_send_buf_clear(struct tcp_sock *tsk, u32 ack)
 pthread_mutex_lock(&tsk->send_buf_lock);
 tcp_send_buf_entry_t *entry,*q;
 list_for_each_entry_safe(entry,q,&tsk->send_buf,list){
   struct tcphdr *tcp = packet_to_tcp_hdr(entry->packet);
```

```
struct iphdr *ip = packet_to_ip_hdr(entry->packet);
   u32 len = ntohs(ip->tot_len) - IP_HDR_SIZE(ip) - TCP_HDR_SIZE(tcp);
   u32 seq_end = ntohl(tcp->seq) + len + ((tcp->flags & (TCP_SYN|TCP_FIN)) ? 1 : \theta
);
   if(less or equal 32b(seq end,ack)){
     list_delete_entry(&entry->list);
     free(entry->packet);
     free(entry);
   }
   else break;
 }
 int ret = list_empty(&tsk->send_buf);
 pthread_mutex_unlock(&tsk->send_buf_lock);
 return ret;
}
   关闭定时器时需要获取定时器链表的锁,注意锁的获取要在扫描发送队列之前,以免与扫
描定时器线程发生死锁。
static void handle_ack(struct tcp_sock *tsk, struct tcp_cb *cb)
{
 if(cb->flags & TCP ACK)
   if(less_than_32b(tsk->snd_una,cb->ack)){
     pthread_mutex_lock(&retrans_timer_lock);
     tsk->snd_una = cb->ack;
     tcp_reset_retrans_timer(tsk);
     if(tcp_send_buf_clear(tsk,tsk->snd_una)){ /
       tcp_unset_retrans_timer(tsk);
     pthread_mutex_unlock(&retrans_timer_lock);
   }
 tcp_update_window(tsk,cb);
 if(tsk->state != TCP_LISTEN && less_than_32b(cb->seq,cb->seq_end))
   tcp_send_control_packet(tsk,TCP_ACK);
}
   ④ 收到旧的包时
   可能是不必要的重传,也可能是 ACK 丢失。如果包含数据,则回复 ACK。
 if(!less or equal 32b(tsk->rcv nxt, cb->seq end)||
     (tsk->rcv_nxt==cb->seq_end && less_than_32b(cb->seq,cb->seq_end))){
   log(ERROR, "Dup pkt.DROP.");
   tcp_send_control_packet(tsk,TCP_ACK);
   return 0;
 }
```

#### (2) 处理乱序的数据包

如果收到的数据包是有效的(位于接收窗口内),但为乱序,则放入 ofo\_buf。否则,将数据中新的部分写入 rcv\_buf,然后检查是否可以读取 ofo\_buf。如果乱序的数据包含 FIN,处理较为繁琐,直接丢弃。

```
static void handle_tcp_data(struct tcp_sock *tsk, struct tcp_cb *cb)
{
  int len = (int)(cb->seq end-tsk->rcv nxt) - (cb->flags&(TCP SYN|TCP FIN)?1:0);
  if(!cb->pl_len || len<=0) {update_tsk(tsk,cb); return; }</pre>
  pthread_mutex_lock(&tsk->wait_recv->lock);
  if(ring_buffer_free(tsk->rcv_buf) < cb->pl_len)
    log(DEBUG, "RECV BUFFER FULL. DROPED.");
  else{
    if(less_or_equal_32b(cb->seq, tsk->rcv_nxt)){
      write ring buffer(tsk->rcv buf,cb->payload+cb->pl len-len,len);
      update_tsk(tsk,cb);
      tcp_check_ofo_buf(tsk);
      wake_with_lock(tsk->wait_recv);
      tsk->rcv_wnd = max(ring_buffer_free(tsk->rcv_buf),1);
    }
    else
      tcp_ofo_buf_append(tsk,cb);
  }
  pthread_mutex_unlock(&tsk->wait_recv->lock);
}
   写入 ofo_buf 时,按 seq_end 组成优先队列。
void tcp_ofo_buf_append(struct tcp_sock *tsk, struct tcp_cb *cb)
  assert((cb->flags&TCP_FIN)==0);
  tcp_ofo_buf_entry_t *entry,*q, *new_entry;
  list_for_each_entry_safe(entry,q,&tsk->rcv_ofo_buf,list){
    if(less or equal 32b(cb->seq end,entry->seq + entry->len)){
      if(greater_or_equal_32b(cb->seq,entry->seq))
        return;
      break;
    }
  }
  new_entry = malloc(sizeof(tcp_ofo_buf_entry_t));
  new entry->seq = cb->seq;
  new_entry->len = cb->pl_len;
  new_entry->data = malloc(cb->pl_len);
  memcpy(new entry->data,cb->payload,new entry->len);
  list_add_tail(&new_entry->list,entry->list.prev);
}
```

读取 ofo\_buf 时,从 seq\_end>rcv\_nxt 开始,一直读取到 seq>rcv\_nxt 时停止。每读取一个包,需要更新 rcv nxt。

```
void tcp_check_ofo_buf(struct tcp_sock *tsk)
  tcp_ofo_buf_entry_t *entry,*q;
  list_for_each_entry_safe(entry,q,&tsk->rcv_ofo_buf,list){
    if(less_or_equal_32b(entry->seq+entry->len,tsk->rcv_nxt)){
      list delete entry(&entry->list);
      free(entry->data);
      free(entry);
    }
    else if(greater_or_equal_32b(tsk->rcv_nxt,entry->seq)){
      int len = (int)(entry->seq+entry->len-tsk->rcv_nxt);
      write_ring_buffer(tsk->rcv_buf,entry->data+entry->len-len,len);
      tsk->rcv nxt += len;
      list_delete_entry(&entry->list);
      free(entry->data);
      free(entry);
    }
    else break;
  }
}
```

# (3) 考虑连接建立/断开的情况

不需要为连接建立/断开实现额外的重传机制,上面的实现已经能够处理包括 SYN | FIN 的数据包。需要考虑的是状态切换:作为发送方,发出包后状态随即切换,重传时不会导致状态切换。作为接收方,收到包后需要判断包的 ack 等于 snd\_nxt 才进行状态切换。这样,如果收到重复的 ACK,不会造成错误的切换。另外,在 CLOSE\_WAIT 状态,需要等到发送的数据对方都接收了才切换状态,判断条件是 snd nxt = snd una。

### 4、实验结果

网络拓扑如图 1。生成的文件大小为 4052632B。

#### (1) C 程序发送, C 程序接收

shell 输出结果如图 2。由 diff 命令输出空结果,我的程序结果正确。

```
'Node: h2'
send window full. sleep. 599021116 598957246 63870
Retrans: 1
send window full. sleep. 599022781 598958646 64135
Retrans: 1
send window full. sleep. 599024181 598961446 62735
Retrans: 1
send window full. sleep. 599026981 598965646 61335
Retrans: 1
send window full. sleep. 599031181 599026981 4200
send window full. sleep. 599096716 599031446 65270
send window full. sleep. 599124981 599060846 64135
send window full. sleep. 599189381 599124981 64400
send window full. sleep. 599190516 599139246 51270
Retrans: 1
send window full. sleep. 599204781 599140646 64135
Retrans: 1
send window full. sleep. 599206181 599204781 1400 send window full. sleep. 599270316 599206181 64135 send window full. sleep. 599308381 599245646 62735
DEBUG: Closed. Sock freed.
^C
root@ubuntu:~/netexp/16/16-tcp_stack# diff server-output.dat client-input.dat
root@ubuntu:~/netexp/16/16-tcp_stack#
root@ubuntu:~/netexp/16/16-tcp_stack# ./tcp_stack server 10001
DEBUG: find the following interfaces: h1-eth0.
Routing table of 1 entries has been loaded.
DEBUG: listen to port 10001.
DEBUG: accept a connection.
ERROR: Dup pkt.DROP.
tot:4052632
DEBUG: close this connection.
```

图 2

### (2) Python 程序发送, C程序接收

抓包结果如图 3,可以看到 FIN 包进行了重传,连接最后正确关闭。 实验结果如图 4,结果正确。

15554 5.741395659 10.0.0.	2 10.0.0.1 TC	P 590	0 [TCP Previous segment not captured] 53002 - 10001 [ACK] Seq=4051024 Ack=1 Win=42340 Len=536
15555 5.741398256 10.0.0.			0 53002 - 10001 [PSH, ACK] Seq=4051560 Ack=1 Win=42340 Len=536
15556 5.741618934 10.0.0.	1 10.0.0.2 TC	P 54	4 10001 → 53002 [ACK] Seq=1 Ack=4049952 Win=64999 Len=0
15557 5.741631005 10.0.0.	1 10.0.0.2 TC	P 54	4 10001 → 53002 [ACK] Seq=1 Ack=4050488 Win=64463 Len=0
15558 5.741888719 10.0.0.	2 10.0.0.1 TC		0 53002 → 10001 [ACK] Seq=4052096 Ack=1 Win=42340 Len=536
15559 5.741889872 10.0.0.			5 53002 - 10001 [FIN, PSH, ACK] Seq=4052632 Ack=1 Win=42340 Len=1
15560 5.741981422 10.0.0.			4 [TCP Dup ACK 15557#1] 10001 53002 [ACK] Seq=1 Ack=4050488 Win=64463 Len=0
15561 5.741984691 10.0.0.			4 [TCP Dup ACK 15557#2] 10001 → 53002 [ACK] Seq=1 Ack=4050488 Win=64463 Len=0
15562 5.742153301 10.0.0.			4 [TCP Dup ACK 15557#3] 10001 - 53002 [ACK] Seq=1 Ack=4050488 Win=64463 Len=0
15563 5.948139279 10.0.0.			0 [TCP Retransmission] 53002 → 10001 [ACK] Seq=4050488 Ack=1 Win=42340 Len=536
15564 5.949082805 10.0.0.			4 10001 - 53002 [ACK] Seq=1 Ack=4051024 Win=64999 Len=0
15565 5.950233878 10.0.0.			0 [TCP Retransmission] 53002 - 10001 [ACK] Seq=4051024 Ack=1 Win=42340 Len=536
15566 5.950255907 10.0.0.			0 [TCP Retransmission] 53002 → 10001 [PSH, ACK] Seq=4051560 Ack=1 Win=42340 Len=536
15567 5.950778279 10.0.0.			4 10001 - 53002 [ACK] Seq=1 Ack=4052632 Win=63927 Len=0
15568 5.951095588 10.0.0.			4 [TCP Dup ACK 15567#1] 10001 → 53002 [ACK] Seq=1 Ack=4052632 Win=63927 Len=0
15569 5.951594377 10.0.0.			5 [TCP Retransmission] 53002 - 10001 [FIN, PSH, ACK] Seq=4052632 Ack=1 Win=42340 Len=1
15570 5.953294502 10.0.0.			4 10001 - 53002 [ACK] Seq=1 Ack=4052634 Win=65534 Len=0
15571 5.973494479 10.0.0.			4 10001 → 53002 [FIN, ACK] Seq=1 Ack=4052634 Win=65534 Len=0
L 15572 5.973795095 10.0.0.	2 10.0.0.1 TC	P 54	4 53002 10001 [ACK] Seq=4052634 Ack=2 Win=42340 Len=0

图 3

```
ERROR: Dup pkt.DROP.
ERROR: SYN|FIN out of order.
ERROR: Dup pkt.DROP.
tot:4052632
DEBUG: close this connection.
DEBUG: Closed. Sock freed.
                                     "Node: h2"
root@ubuntu:~/netexp/16/16-tcp_stack# python tcp_stack2.py client 10.0.0.1 1000
root@ubuntu:~/netexp/16/16-tcp_stack# python tcp_stack2.py client 10.0.0.1 1000
root@ubuntu:~/netexp/16/16-tcp_stack# diff server-output.dat client-input.dat
root@ubuntu:~/netexp/16/16-tcp_stack#
```

图 4

# (3) C程序发送,Python程序接收

```
send window full. sleep. 1548302270 1548236735 65535
send window full. sleep. 1548349870 1548284335 65535
Retrans: 1
send window full. sleep. 1548361070 1548295535 6<u>5535</u>
Retrans: 1
send window full. sleep. 1548372270 1548306735 65535
Retrans: 1
send window full. sleep. 1548382070 1548316535 6<u>5535</u>
Retrans: 1
send window full. sleep. 1548411470 1548345935 65535
Retrans: 1
send window full. sleep. 1548484270 1548418735 65535
Retrans: 1
send window full. sleep. 1548566870 1548501335 65535
send window full. sleep. 1548572470 1548506935 65535
Retrans: 1
Retrans: 2
send window full. sleep. 1548583670 1548518135 65535
Retrans: 1
send window full. sleep. 1548601870 1548536335 65535
Retrans: 1
DEBUG: Closed. Sock freed.
root@ubuntu:~/netexp/16/16-tcp_stack# python tcp_stack2.py server 10001 ('10.0.0.2', 12345)
('tot:', 4052632)
root@ubuntu:~/netexp/16/16-tcp_stack# diff server-output.dat client-input.dat
root@ubuntu:~/netexp/16/16-tcp_stack#
```

实验结果如图 5, 结果正确。

每次重传一个包时,传输 4MB 数据需要十几秒。若改为每次重传发送队列中的所有包,则 传输时间仅需不到 5 秒。

# 5、遇到的问题

(1) 处理含 FIN 的乱序包

为了简化问题,不把它加入乱序队列,而是直接丢弃。