

- (1) Let $L_1 = \{w : w \text{ has 3 or 4 } a\text{'s}\}$. $\Sigma = \{a, b\}$. Give a DFA for L_1 . Give a regular expression α_1 for L_1 .
- (2) Let $L_2 = \{w : w \text{ has 3 } a\text{'s or } w \text{ has } b\text{'s}\}$. $(\Sigma\{a, b\})$. Given an NFA M_2 s.t. $L(M_2) = L_2$. Give a regular expression α_2 s.t. $L(\alpha_2) = L_2$.
- (3) Prove or disprove: $\{a^i b^{i^2} : i \geq 0\}$ is context-free. (Either give a PDA or a CFG; or prove its not context-free.
- (4) Prove or disprove: $\{a^{i^2} : i \geq 0\}$ is context-free.
- (5) Prove or disprove: $\{a^i b^j : i \leq j \text{ or } i \text{ is a perfect square}\}$ is not regular.
- (6) Prove or disprove: $\{a^i b^j : 100 \leq i < j \leq 1000\}$ is context-free.
- (7) Prove or disprove: $\{a^i b^j : 100 \leq i < j \leq 1000\}$ is regular.
- (8) Let L be a CFL. Let k, K be as in the strong pumping lemma for L . Prove that if $L \neq \emptyset$ then L has a string of length $\leq K$.
- (9) Construct a PDA which accept $\{a^i b^j c^k d^\ell : i = k \text{ or } j = \ell\}$.
- (10) Give a CFG that generates this language $\{a^i b^j c^k d^\ell : i = k \text{ or } j = \ell\}$.
- *(11) Show $\{a^m b^n : m > n \text{ or } (m \text{ is prime and } n \leq m)\}$ is not context-free. **Hint:** for any $k > 0$, there are arbitrarily large primes p such that $p \pm i$ is composite for $1 \leq i \leq k$.