RASPBERRY PLASSEMBLER

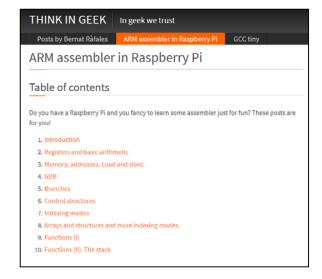
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Chapter 4: Raspberry Pi Assembler "Raspberry Pi Assembler" by R. Ferrer and W. Pervin

https://thinkingeek.com/2013/01/12/arm-assembler-raspberry-pi-chapter-4/





S. Winberg

- What is debugging?
 - Debugging refers to the process used to identify errors in your program, since it is rare to write bug-free code all at once.
- Example to explain the concept of debugging
 - You were tasked to develop a program called SortVector to sort a sequence of numbers in descending order
 - Let VectorUnsorted = [5 3 2 15 8]
 - After developing the program, the sorted vector VectorSorted was obtained

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VectorSorted = SortVector( VectorUnsorted )
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 Thus, the program needs to be debugged. Thereafter, the error in the program should be corrected.



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- Thus, the program needs to be debugged. Thereafter, the error in the program should be corrected.
- What does the debugging process involve? First, we predict the values of variables in the program. Then we step through the code and identify errors by comparing the difference between the predicted and the observed values.

GNU Debugger



Raspberry Pi Assembler GNU Debugger

- What is GNU Debugger?
 - GNU Debugger (GDB) is an open-source software used to debug other programs
- Basic functionality offered by GDB
 - User can step through the code
 - Print the contents of CPU registers and variables stored in the .data section, to the terminal
 - Disassemble machine instructions, ie. go from machine instructions to assembly code
 - Insert breakpoints in the code
 - ... many more
- Let's work with GDB and learn commands as we go along



GNU Debugger and store2 program



- Let's debug the store02 program developed at the end of Chapter 3
 - Why are we debugging a program that works correctly?

```
1 /* -- store02.s */
2 .data
 3 myvar1: .word 0
4 myvar2: .word 0
5 .text
6 .global main
7 main:
  ldr r1, =myvar1 @ r1 <- &myvar1
9 mov r3, #3 @ r3 <- 3
10 str r3, [r1] @ *r1 <- r3
11 ldr r2, =myvar2 @ r2 <- &myvar2
12 mov r3, #4 @ r3 <- 4
13 str r3, [r2] @ *r2 <- r3
14 ldr r1, =myvar1 @ r1 <- &myvar1
15 ldr r1, [r1] 0 r1 <- *r1
  ldr r2, =myvar2 @ r2 <- &myvar2
17 ldr r2, [r2] @ r2 <- *r2
   add r0, r1, r2
19 bx lr
```

- Let's debug the store02 program developed at the end of Chapter 3
 - Why are we debugging a program that works correctly?
 - To demonstration how to use GDB. Let's check the contents of CPU registers and variables in memory after certain lines of code have executed

1 /* -- store02.s */ 2 .data 3 myvar1: .word 0 4 myvar2: .word 0 5 .text After line 9 has executed 6 .global main 7 main: Check contents of r3 ldr r1, =myvar1 @ r1 <- &myvar1 ▲ 9 mov r3, #3 @ r3 <- 3 →10 str r3, [r1] 0 *r1 <- r3 After line 10 has executed ldr r2, =myvar2 @ r2 <- &myvar2 Check address of myvar1 and contents of r1. These 12 mov r3, #4 @ r3 <- 4 two should be the same. 13 str r3, [r2] @ *r2 <- r3 Check the contents of ldr r1, =myvar1 @ r1 <- &myvar1 myvar1 ldr r1, [r1] @ r1 <- *r1 15 ldr r2, =myvar2 @ r2 <- &myvar2 ldr r2, [r2] 0 r2 <- *r2 17 add r0, r1, r2 18

bx lr

- Let's debug the store02 program developed at the end of Chapter 3
- 1. Start GDB and specify the program to debug

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- 1. Start GDB and specify the program to debug
- 2. Type **start** to run the program up to and including the first instruction of main

(gdb) start

Temporary breakpoint 1 at 0x8394 : file store02.s, line 14. Starting program: /home/RPiA/Chapter03/store02

- Let's debug the store02 program developed at the end of Chapter 3
- 1. Start GDB and specify the program to debug
- 2. Type **start** to run the program up to and including the first instruction of main
- Disassemble the machine instructions into human readable assembly instructions and to observe the next instruction to be executed

The next instruction to be executed is denoted by the => symbol

```
Type disassemble to convert the code
                         from machine instructions to assembly
                         instructions
(gdb) disassemble
Dump of assembler code for function main:
   0x00008390 <+0>: ldr r1, [pc, #40]; 0x83c0 <main+48>
=> 0x00008394 <+4>: mov r3, #3
   0x00008398 <+8>: str r3, [r1]
   0x0000839c <+12>: ldr r2, [pc, #32] ; 0x83c4 <main+52>
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   0x000083b4 <+36>: ldr r2, [r2]
   0x000083b8 < +40>: add r0, r1, r2
   0x000083bc < +44>: bx lr
   0x000083c0 < +48>: andeq r0,r1,r4,ror #10
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End of assembler dump.
```

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Let's make a few observations of the assembler code

The first instruction of **main** can be found at address 0x00008390

- The assembler has written code to store the real addresses of myvar1 and myvar2.
- However, the disassembler has simply translated these machine instructions into presumed assembler code

Address of myvar1 has been assigned to 0x83c0, which is 48 bytes from where **main** begins

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The next instruction to be executed is displayed to the user (gdb) stepi (executed is displayed to the user (gdb) stepi

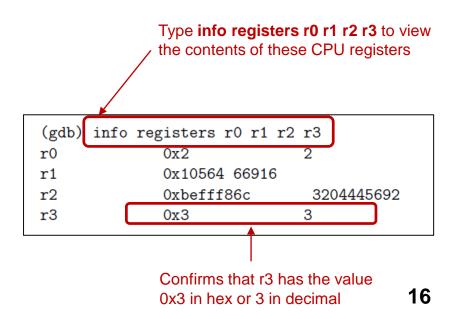
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- 5. Disassemble to confirm the last instruction that was executed

The last instruction that was executed was mov r3, #3 because it is the line of code before the => symbol

The next instruction to be executed

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End of assembler dump.
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If we disassembled the code, it would look like ...

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End of assembler dump.
```

The next instruction to be executed is displayed to the user

```
(gdb) stepi

→11 ldr r2, =myvar2 @ r2 <- &myvar2
```

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End of assembler dump.
```

```
Type p to print to terminal

Observation: counter for the printed result
```

```
contents of r1 is 0x10564 in hex
and 66919 in decimal
Great, it has changed. In fact this is the address of myvar1. Let's check that using its symbolic name and C syntax.

(gdb) p kmyvar1

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(cdata variable, no debug info> *) 0x10564

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(contents of r1 is 0x10564 in hex
and 66919 in decimal
address of myvar1

is 0x10564 in hex
```

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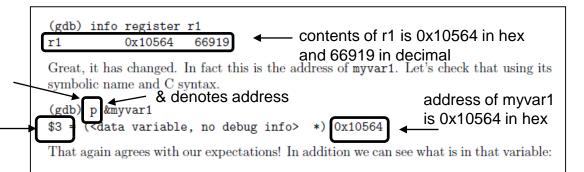
.. r1 is equal to the address of myvar1

Type p to print to terminal

Observation: counter for the printed result

If we disassembled the code, it would look like ...

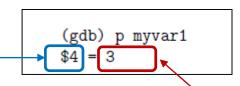
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End of assembler dump.
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- 1. Start GDB and specify the program to debug
- 2. Type **start** to run the program up to and including the first instruction of main
- Disassemble the machine instructions into human readable assembly instructions and to observe the next instruction to be executed
- 4. Type **stepi** to step through the code, ie. only execute the next instruction
- 5. Disassemble to confirm the last instruction that was executed
- 6. View the contents of registers r0, r1, r2, r3 to confirm that r3 was changed to the value of 3
- 7. Type **stepi** to step through the code, ie. only execute the next instruction
- 8. Check that r1 has the address of myvar1
 - View contents of r1
 - View address of myvar1
- 9. Lastly, check that the variable myvar1 has been assigned the value of r3

```
If we disassembled the code, it would look like ...
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End of assembler dump.
```

Observation: counter for the printed result



confirms that the contents of the variable myvar1 is 3, which is equal to register r3

GNU Debugger: commands



GNU Debugger: commands

GDB cheatsheet - page 1

Running

gdb --args program> <args...>
Start GDB and pass arguments

gdb --pid <pid>
Start GDB and attach to process.

set args <args...>

Set arguments to pass to program to be debugged.

run

Run the program to be debugged.

kill

Kill the running program.

Breakpoints

break <where>

Set a new breakpoint.

delete <bre>delete <bre>delete <bre>delete <bre>delete
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Remove a breakpoint.

clear

Delete all breakpoints.

enable

breakpoint#>
Enable a disabled breakpoint.

disable <bre> <bre>breakpoint#>
Disable a breakpoint.

Watchpoints

watch <where>

Set a new watchpoint.

delete/enable/disable <watchpoint#>
 Like breakpoints.

<where>

function name

Break/watch the named function.

line number

Break/watch the line number in the current source file.

file:line number

Break/watch the line number in the named source file.

Conditions

break/watch <where> if <condition>
Break/watch at the given location if the
condition is met.

Conditions may be almost any C expression that evaluate to true or false.

condition

Set/change the condition of an existing break- or watchpoint.

Examining the stack

backtrace

where

Show call stack.

backtrace full

where full

Show call stack, also print the local variables in each frame.

frame <frame#>

Select the stack frame to operate on.

Stepping

step

Go to next instruction (source line), diving into function. next

Go to next instruction (source line) but don't dive into functions.

finish

Continue until the current function returns.

continue

Continue normal execution.

Variables and memory

print/format <what>

Print content of variable/memory location/register.

display/format <what>

Like "print", but print the information after each stepping instruction.

undisplay <display#>

Remove the "display" with the given number.

enable display <display#>

disable display <display#>

En- or disable the "display" with the given number.

x/nfu <address>

Print memory.

n: How many units to print (default 1).

f: Format character (like "print").

u: Unit.

Unit is one of:

b: Byte,

h: Half-word (two bytes)

w: Word (four bytes)

g: Giant word (eight bytes)).

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GNU Debugger: commands

GDB cheatsheet - page 2

Format a Pointer. c Read as integer, print as character. d Integer, signed decimal. f Floating point number. o Integer, print as octal. s Try to treat as C string. t Integer, print as binary (t = "two"). Integer, unsigned decimal. x Integer, print as hexadecimal.

<what>

expression

Almost any C expression, including function calls (must be prefixed with a cast to tell GDB the return value type).

file name::variable name

Content of the variable defined in the named file (static variables).

function::variable name

Content of the variable defined in the named function (if on the stack).

{type}address

Content at address, interpreted as being of the C type type.

\$register

Content of named register. Interesting registers are \$esp (stack pointer), \$ebp (frame pointer) and \$eip (instruction pointer).

Threads

thread <thread#>

Chose thread to operate on.

Manipulating the program

set var <variable name>=<value>

Change the content of a variable to the given value.

return <expression>

Force the current function to return immediately, passing the given value.

Sources

directory <directory>

Add *directory* to the list of directories that is searched for sources.

list

list <filename>:<function>

list <filename>:<line number>

list <first>, <last>

Shows the current or given source context. The *filename* may be omitted. If *last* is omitted the context starting at *start* is printed instead of centered around it.

set listsize <count>

Set how many lines to show in "list".

Signals

handle <signal> <options>

Set how to handle signles. Options are:

(no)print: (Don't) print a message when signals occurs.

(no)stop: (Don't) stop the program when signals occurs.

(no)pass: (Don't) pass the signal to the program.

Informations

disassemble

disassemble <where>

Disassemble the current function or given location.

info args

Print the arguments to the function of the current stack frame.

info breakpoints

Print informations about the break- and watchpoints.

info display

Print informations about the "displays".

info locals

Print the local variables in the currently selected stack frame.

info sharedlibrary

List loaded shared libraries

info signals

List all signals and how they are currently handled.

info threads

List all threads.

show directories

Print all directories in which GDB searches for source files.

show listsize

Print how many are shown in the "list" command.

whatis variable name

Print type of named variable.

Raspberry Pi Assembler Branching

RASPBERRY PLASSEMBLER

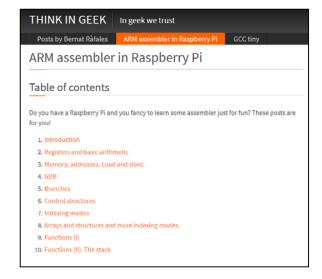
Roger Ferrer Ibáñez Cambridge, Cambridgeshire, U.K.

> William J. Pervin Dallas, Texas, U.S.A.

Chapter 5: Raspberry Pi Assembler "Raspberry Pi Assembler" by R. Ferrer and W. Pervin

https://thinkingeek.com/2013/01/19/arm-assembler-raspberry-pi-chapter-5/

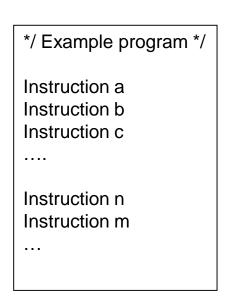


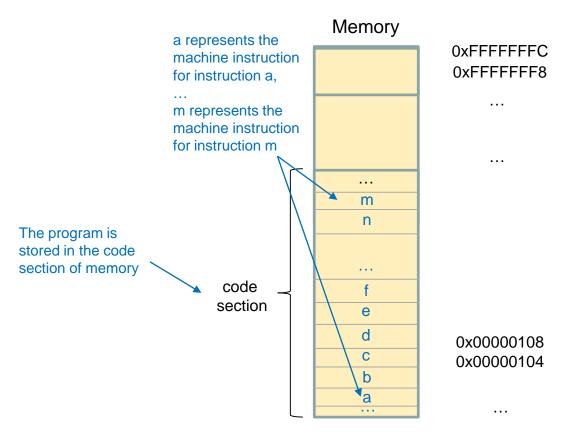


S. Winberg

Introduction: programs

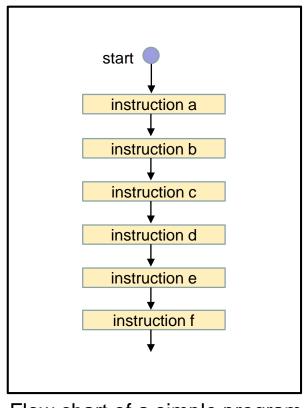
- Consider a program with many operations.
 - The first instruction is denoted by 'instruction a'. The second instruction is denoted by 'instruction b', ...
 - Furthermore, let the letter 'a' denote the 32-bit machine instruction related to instruction
 a, 'b' denote the machine instruction related to instruction b, ...



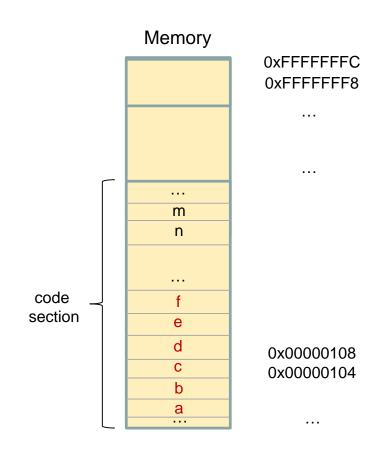


Programs without branches

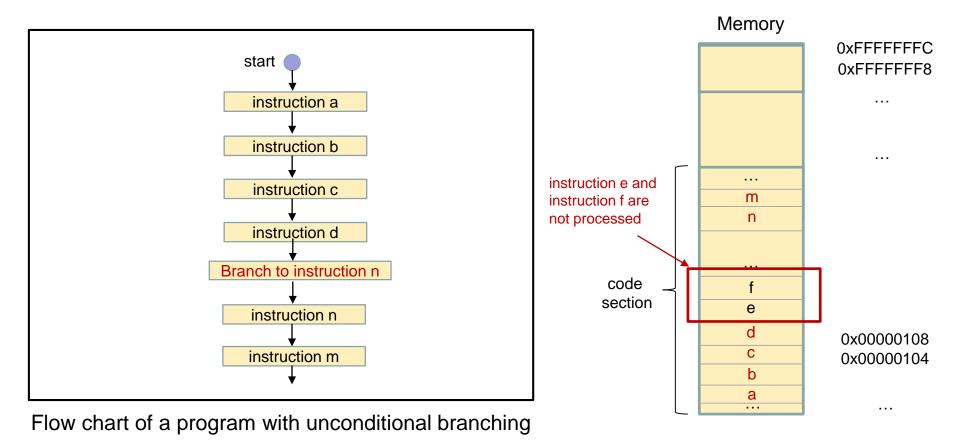
 Typically, in a simple program, after the next instruction is fetched from memory, the PC (R15) is incremented by 4. Thus, instructions are processed sequentially.



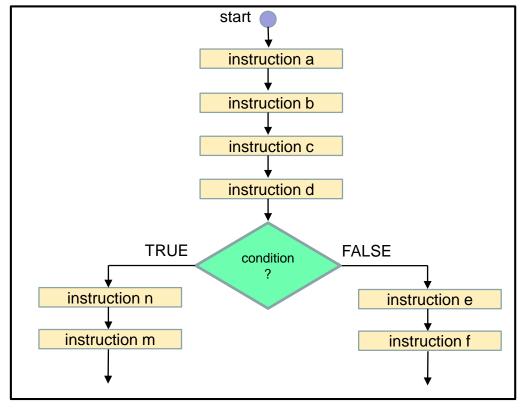
Flow chart of a simple program without branching

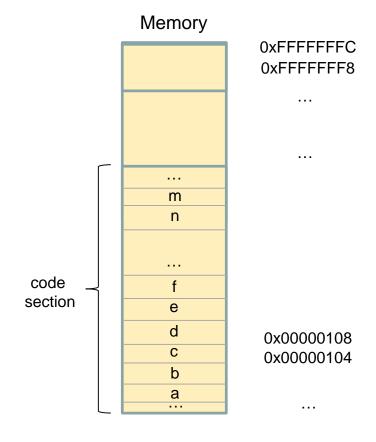


- What is branching?
 - Branching refers to changing the sequential processing of instructions
 - When an instruction modifies the value of the Program Counter (PC), then the new PC is used to address the next instruction. In this way, branching is implemented. There is unconditional branching



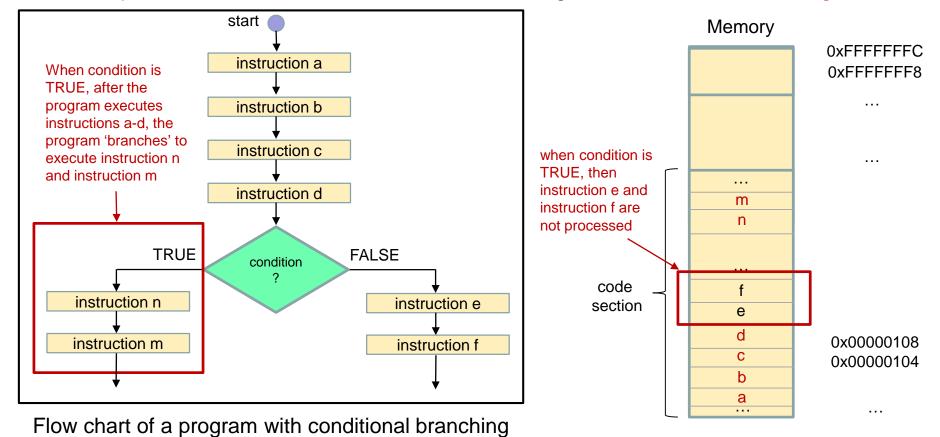
- What is branching?
 - Branching refers to changing the sequential processing of instructions
 - When an instruction modifies the value of the Program Counter (PC), then the new PC is used to address the next instruction. In this way, branching is implemented. There is unconditional branching and conditional branching



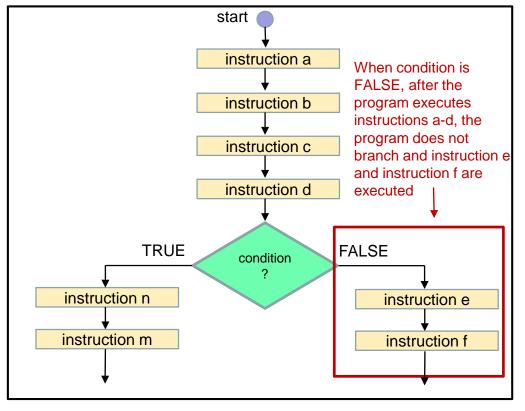


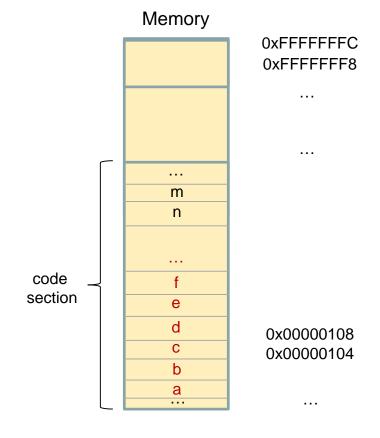
Flow chart of a program with conditional branching

- What is branching?
 - Branching refers to changing the sequential processing of instructions
 - When an instruction modifies the value of the Program Counter (PC), then the new PC is used to address the next instruction. In this way, branching is implemented. There is unconditional branching and conditional branching



- What is branching?
 - Branching refers to changing the sequential processing of instructions
 - When an instruction modifies the value of the Program Counter (PC), then the new PC is used to address the next instruction. In this way, branching is implemented. There is unconditional branching and conditional branching





Flow chart of a program with conditional branching

Unconditional branching



Unconditional branches

- Let's look at an assembly program that uses the branch instruction
 - The unconditional branch command is denoted by b



Unconditional branches

- Let's look at an assembly program that uses the branch instruction
 - The unconditional branch command is denoted by b

Observations

```
/* -- branch01.s */
.text
.global main

main:

mov r0, #2 @ r0 <- 2
b end @ branch to 'end'
mov r0, #3 @ r0 <- 3
end:
bx lr
```



Unconditional branches

- Let's look at an assembly program that uses the branch instruction
 - The unconditional branch command is denoted by b

Stepping through code to understand the program flow

After this instruction has executed, the register r0 will have the value 2



Unconditional branches

- Let's look at an assembly program that uses the branch instruction
 - The unconditional branch command is denoted by b

Stepping through code to understand the program flow

After this instruction has executed, the Program Counter (PC) will be loaded with the value of the memory address of label **end**



Unconditional branches

- Let's look at an assembly program that uses the branch instruction
 - The unconditional branch command is denoted by b

Stepping through code to understand the program flow

program branches to the **end** label and program terminates



Unconditional branches

- Let's look at an assembly program that uses the branch instruction
 - The unconditional branch command is denoted by b

If you execute this program you will see that it returns an error code of 2.

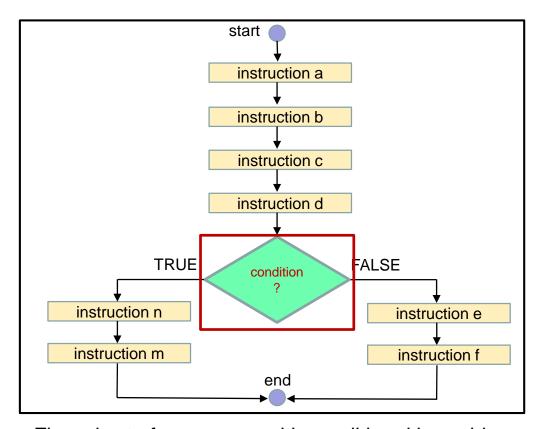
\$./branch01 ; echo \$?
2

Conditional branching



Conditional branches

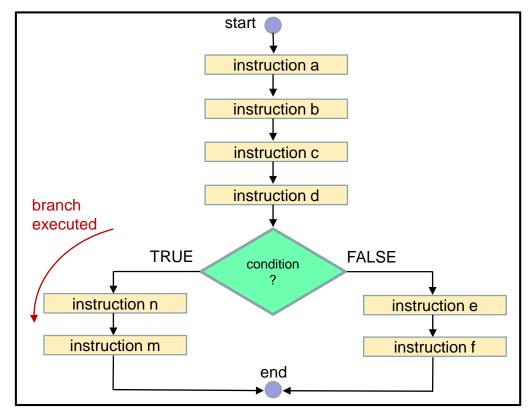
- Conditional branching involves two steps
 - 1. Evaluate a condition



Flow chart of a program with conditional branching

Conditional branches

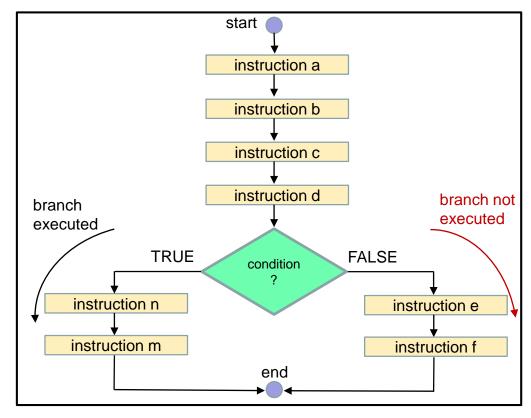
- Conditional branching involves two steps
 - Evaluate a condition
 - 2. If the condition is TRUE, then execute the branch instruction



Flow chart of a program with conditional branching

Conditional branches

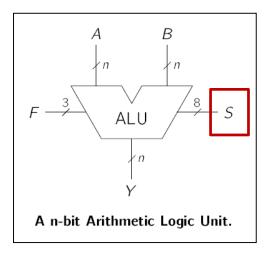
- Conditional branching involves two steps
 - Evaluate a condition
 - 2. If the condition is TRUE, then execute the branch instruction. If the condition is FALSE, then do not execute the branch instruction



Flow chart of a program with conditional branching

Raspberry Pi Assembler Conditional branches: condition evaluated

- What condition is evaluated?
 - Typically, one or more of the four flags in the Current Program Status Register (CPSR) is evaluated when a conditional branch instruction is used
- What is the CPSR?
 - The CPSR is a 32-bit register that is found in the CPU
 - This is similar to the status bits S from the output of a generic ALU



Conditional branches: condition evaluated

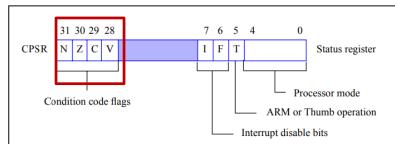
- Example: the ARM CPSR
 - There are four condition code flags which may be updated after the previously executed instruction

Assumes the operands and the result are both **unsigned** numbers

- Carry (C) : set to 1:
 - If the last operation was addition and there was a unsigned overflow
 - If the last operation was subtraction and a borrow was not needed
 In all other cases, it is cleared to zero
- Zero (Z): is set to 1 when the last result was zero. Otherwise, it is cleared to 0
- Negative (N): is set to 1 when the last result was negative, ie. the MSB is set to 1.
 Otherwise, it is cleared to 0.
- Overflow (V): is set to 1 when the last operation had a signed overflow. Otherwise, it
 is cleared to 0

Assumes the operands and the result are both **signed** numbers





Conditional branches: condition evaluated

- Examples of instructions that update the CPSR
 - adds r1, r1, r2
 - sub**s** r2, r1, r0
 - cmp r2, r1
 - ...



Conditional branches: condition evaluated

- Examples of instructions that update the CPSR
 - adds r1, r1, r2
 - sub**s** r2, r1, r0
 - cmp r2, r1
 - ...

The **s** suffix is used to denote that the CPSR will update after this instruction has executed



Raspberry Pi Assembler Conditional branches: condition evaluated

- How is branching done based on the four condition code flags?
 - There are many branch instructions that can be used
 - Example: BEQ LABEL
 - The BEQ instruction causes a branch to the location LABEL if the conditional code flag Z is equal to one, when this branch instruction was executed



Raspberry Pi Assembler Conditional branches: condition evaluated

- How is branching done based on the four condition code flags?
 - There are many branch instructions that can be used
 - Example: BEQ LABEL
 - The BEQ instruction causes a branch to the location LABEL if the conditional code flag Z is equal to one, when this branch instruction was executed
- Examples of branch suffixes and the related conditional code flag that are evaluated
 - Note: place the letter b in front of the suffix to create the full branch instruction.

Example: NE becomes BNE EQ becomes BEQ



Condition field encoding in ARM instructions		
Condition suffix	Condition name	Condition Code test
EQ	Equal (zero)	Z = 1
NE	Not equal (nonzero)	Z = 0
CS/HS	Carry set/Unsigned higher or same	C = 1
CC/LO	Carry clear/Unsigned lower	C = 0
MI	Minus (negative)	N = 1
PL	Plus (positive or zero)	N = 0
VS	Overflow	V = 1
VC	No overflow	V = 0
HI	Unsigned higher	$\overline{C} \vee Z = 0$
LS	Unsigned lower or same	$\overline{C} \vee Z = 1$
GE	Signed greater than or equal	$N \oplus V = 0$
LT	Signed less than	$N \oplus V = 1$
GT	Signed greater than	$Z \lor (N \oplus V) = 0$
LE	Signed less than or equal	$Z \lor (N \oplus V) = 1$
AL	Always	
	not used	

```
1 /* -- compare01.s */
 2 .text
 3 .global main
 4 main:
   mov r1, #2
                  0 r1 <- 2
                   0 r2 <- 2
   mov r2, #2
   cmp r1, r2
                   @ update cpsr condition codes with r1-r2
     beg case_equal
                      @ branch to case_equal only if Z = 1
9 case_different:
      mov r0, #2
                      0 \text{ r} 0 < -2
10
                      @ branch to end
11
      b end
12 case_equal:
13
      mov r0, #1 @ r0 <- 1
14 end:
15
      bx lr
```

- In the code below
 - The CoMPare (cmp) instruction is executed before the branch instruction (beq)

```
1 /* -- compare01.s */
 2 .text
 3 .global main
 4 main:
      mov r1, #2
                        @ r1 <- 2
                                                                              This instruction
      mov r2, #2
                        0 r2 <- 2
       cmp r1, r2
                        @ update cpsr condition codes with r1-r2
                                                                             performs the
       beg case_equal
                        @ branch to case_equal only if Z = 1
                                                                             subtraction: r1 - r2 and
 9 case_different:
                                                                             then updates the
10
       mov r0, #2
                        0 r0 <- 2
                                                                             CPSR
                        @ branch to end
11
           end
12 case_equal:
13
      mov r0, #1
                        @ r0 <- 1
14 end:
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       bx lr
```

- In the code below
 - The CoMPare (cmp) instruction is executed before the branch instruction (beq).
 After this instruction has executed, the Z flag of the CPSR register is set to 1, since (r1 r2) = 0 and the value of the last result was zero

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 2 .text
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 4 main:
      mov r1, #2
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                                                                             This instruction
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       cmp r1, r2
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                                                                             performs the
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9 case_different:
                                                                             then updates the
10
      mov r0, #2
                        0 r0 <- 2
                                                                             CPSR
                        @ branch to end
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           end
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13
      mov r0, #1
14 end:
15
       bx lr
```

- In the code below
 - The CoMPare (cmp) instruction is executed before the branch instruction (beq).
 After this instruction has executed, the Z flag of the CPSR register is set to 1, since (r1 r2) = 0 and the value of the last result was zero
 - When the branch instruction (BEQ) is executed, the Z flag of the CPSR is evaluated and the program will branch to the label 'case_equal', since Z has a value of 1

```
1 /* -- compare01.s */
 2 .text
 3 .global main
 4 main:
      mov r1, #2
                        0 r1 <- 2
      mov r2, #2
                        @ r2 <- 2
     cmp r1, r2
                        @ update cpsr condition codes with r1-r2
                        @ branch to case_equal only if Z = 1
      beg case_equal
9 case_different:
10
      mov r0, #2
                        0 \text{ r} 0 < -2
                        @ branch to end
11
           end
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      mov r0, #1
                        @ r0 <- 1
13
14 end:
15
       bx lr
```

- In the code below
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 After this instruction has executed, the Z flag of the CPSR register is set to 1, since (r1 r2) = 0 and the value of the last result was zero
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 3 .global main
 4 main:
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      mov r2, #2
                    0 r2 <- 2
   cmp r1, r2
                   @ update cpsr condition codes with r1-r2
     beg case_equal
                      @ branch to case_equal only if Z = 1
9 case_different:
                       @ r0 <- 2
10
      mov r0, #2
                       @ branch to end
11
          end
12 case equal:
                       0 r0 <- 1
13
      mov r0, #1
14 end:
15
      bx lr
```

- In the code below
 - The CoMPare (cmp) instruction is executed before the branch instruction (beq).
 After this instruction has executed, the Z flag of the CPSR register is set to 1, since (r1 r2) = 0 and the value of the last result was zero
 - When the branch instruction (BEQ) is executed, the Z flag of the CPSR is evaluated and the program will branch to the label 'case_equal', since Z has a value of 1
 - When the program completes, r0 = 1 and the value 1 is displayed to the terminal

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1 /* -- compare01.s */
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 4 main:
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      mov r2, #2
                     0 r2 <- 2
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 9 case_different:
                       0 \text{ r} 0 < -2
10
      mov r0, #2
11
                       @ branch to end
          end
12 case_equal:
      mov r0, #1
13
                       0 r0 <- 1
14 end:
15
       bx lr
```

- In the code below
 - The CoMPare (cmp) instruction is executed before the branch instruction (beq). After this instruction has executed, the Z flag of the CPSR register is set to 1, since (r1 r2) = 0 and the value of the last result was zero
 - When the branch instruction (BEQ) is executed, the Z flag of the CPSR is evaluated and the program will branch to the label 'case_equal', since Z has a value of 1
 - When the program completes, r0 = 1 and the value 1 is displayed to the terminal

