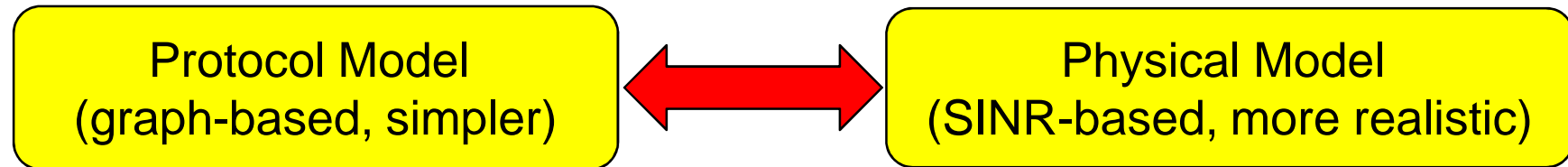


Wireless Networks Models



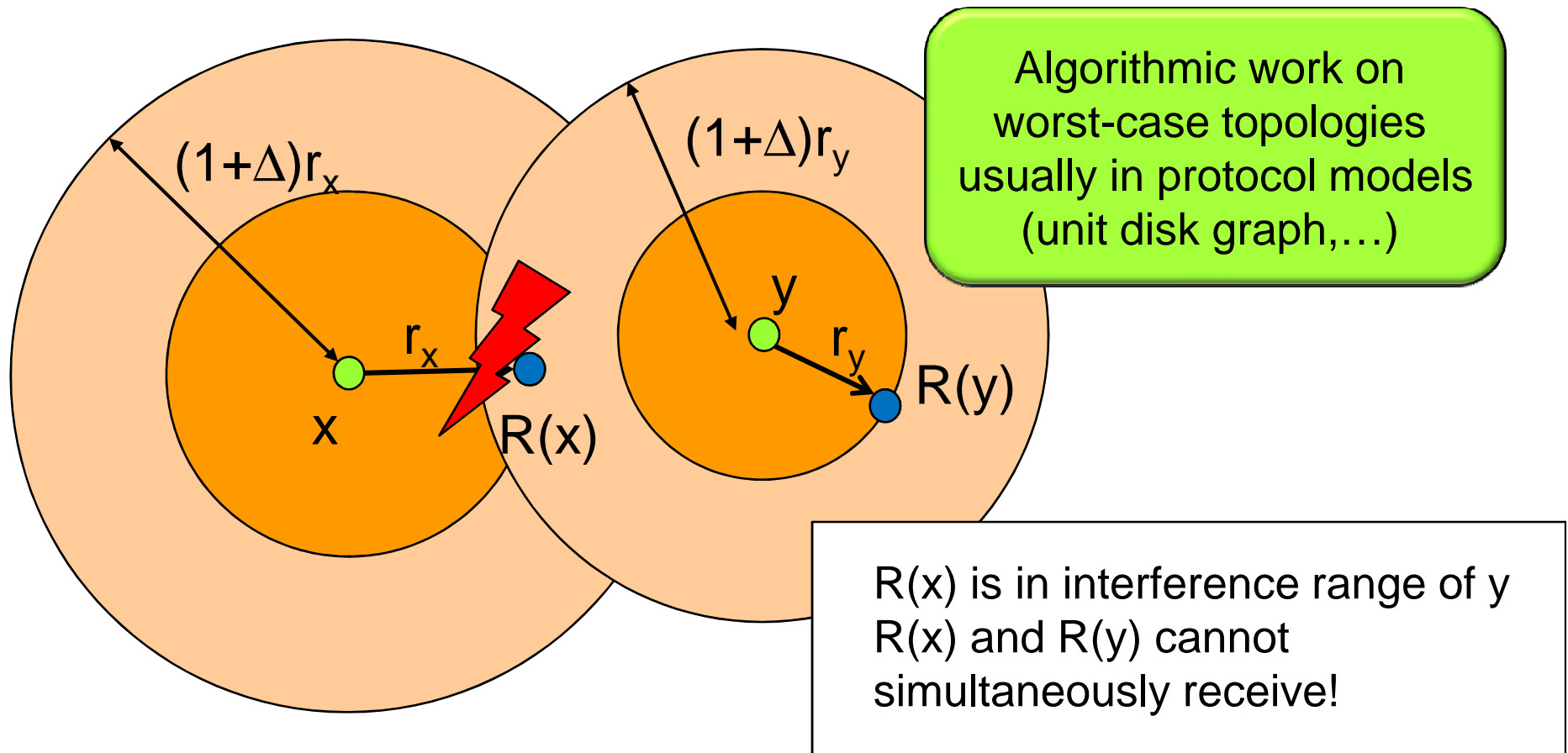
Models

- Two standard models in wireless networking



Protocol Model

- Based on **graph-based** notion of interference
- **Transmission range** and **interference range**



Models

Algorithmic models often inspired by

- “Connections” => **Graph Theory**
- Transmission ranges, interference, ... => **Geometry**

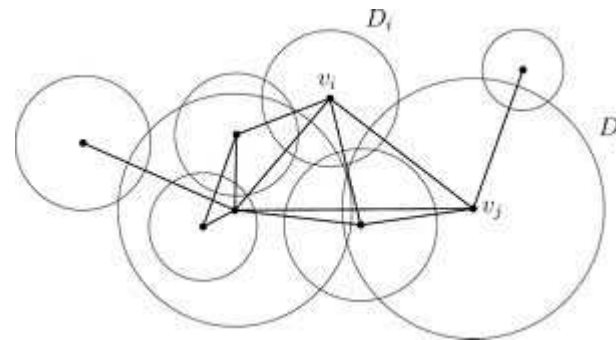
“Higher level abstractions“

Connectivity: Disk Graph

Which nodes are adjacent to a given node v ?

Disk Graph (DG)

- Each node has a communication range
- $\{u,v\} \in E \Leftrightarrow |u,v| < \text{Range}$



Connectivity: Unit Disk Graph

Which nodes are adjacent to a given node v ?

Example: **Unit Disk Graph (UDG)**

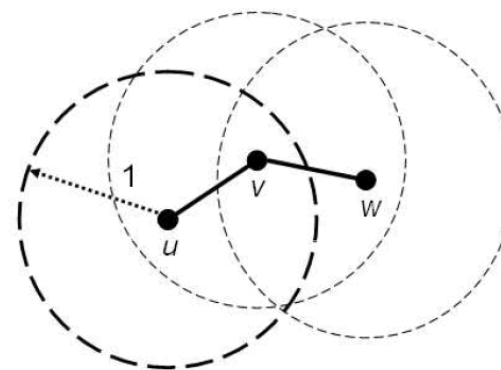
- Classic Model from computational geometry
- $\{u, v\} \in E \Leftrightarrow |u, v| \leq 1$

Pro

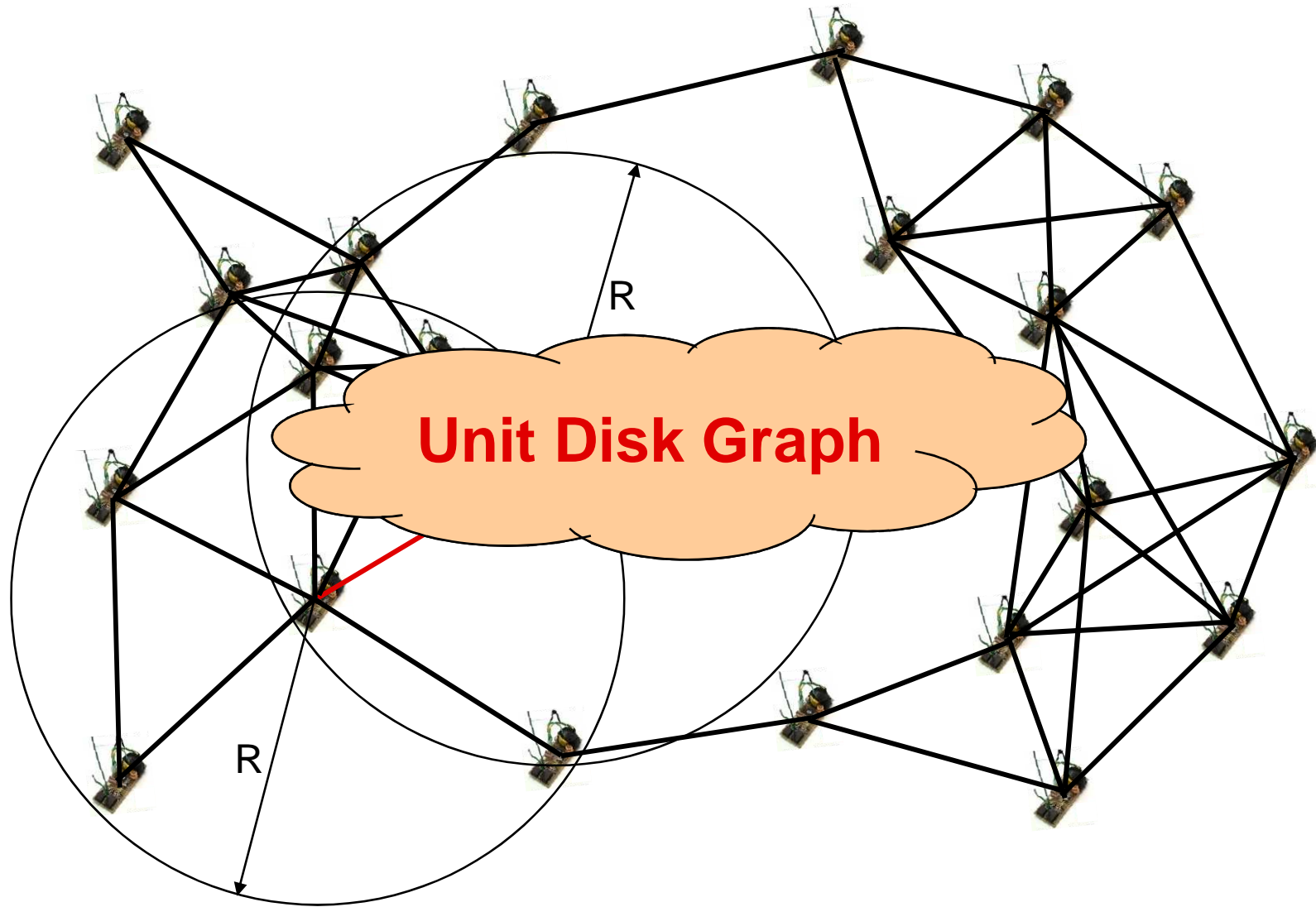
- Very simple
- Analytically tractable
- Realistic for unobstructed environments

Contra

- Too simple
- Not realistic for inner-city networks with many buildings etc.



Connectivity: Unit Disk Graph



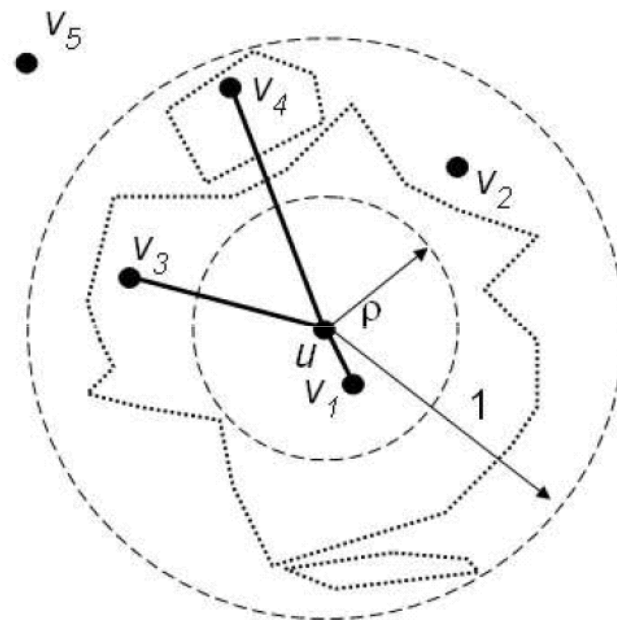
Connectivity: Quasi Unit Disk Graph

More realistic: *Quasi* UDG (QUDG)

- two radii
- $\{u,v\} \in E \Leftrightarrow |u,v| \leq \rho$
- $\{u,v\} \notin E \Leftrightarrow |u,v| > 1$
- otherwise: **It depends!**

It depends...

- ... on an adversary,
- ... on probabilistic model,
- etc.!

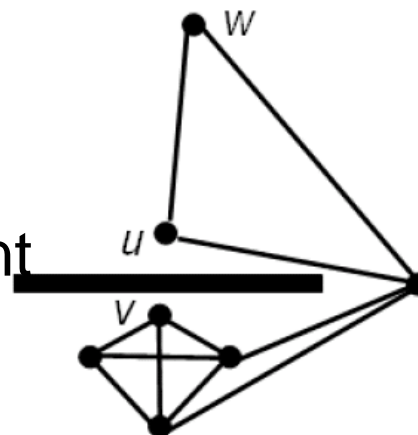


Advantage: More flexible and realistic than UDG!

Connectivity: Drawbacks of QUDG

How realistic is QUDG?

- if there is a **wall**...
- ... u and v can be close but not adjacent
- => QUDG model requires **very small ρ**



However, although if there are walls, connectivity typically still adheres to certain geometric constraints!

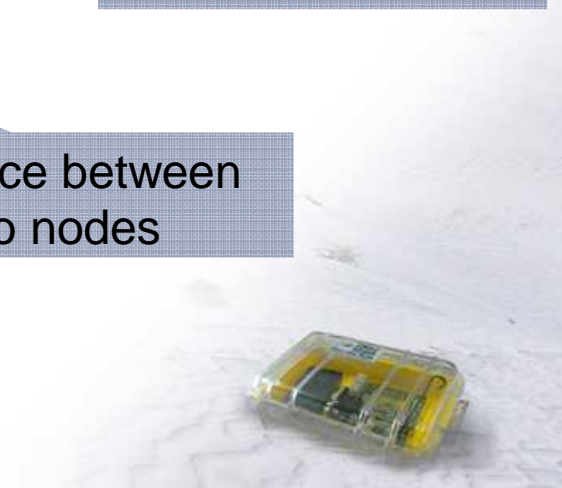
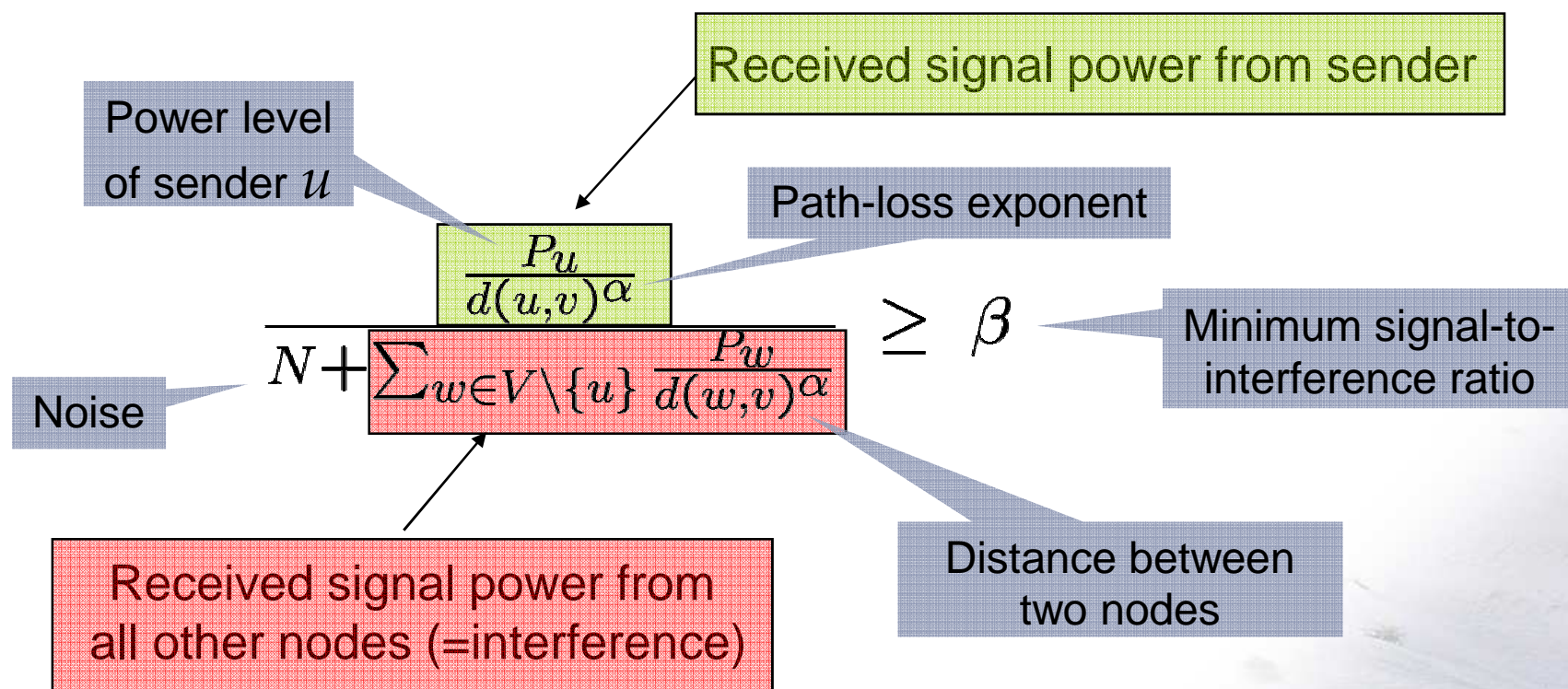
- Resort to general connectivity graphs too pessimistic!

Observation: Even in **complex environments**, the neighbors of a node are often also neighboring (cf wall example)

- Motivation for **Bounded Independence Graph**!

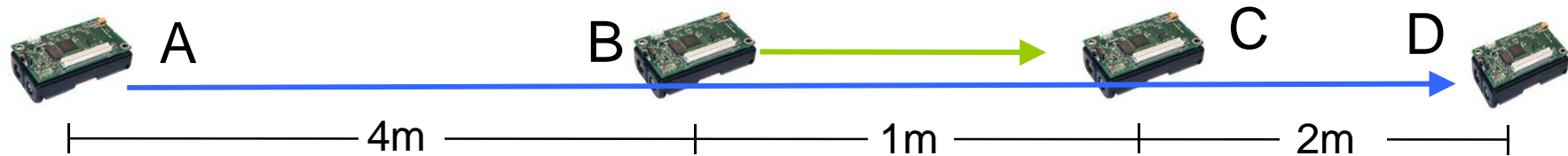
Physical Model

- Based on **signal-to-noise-plus-interference (SINR)**
- Simplest case:
 - packets can be decoded if SINR is larger than β at receiver

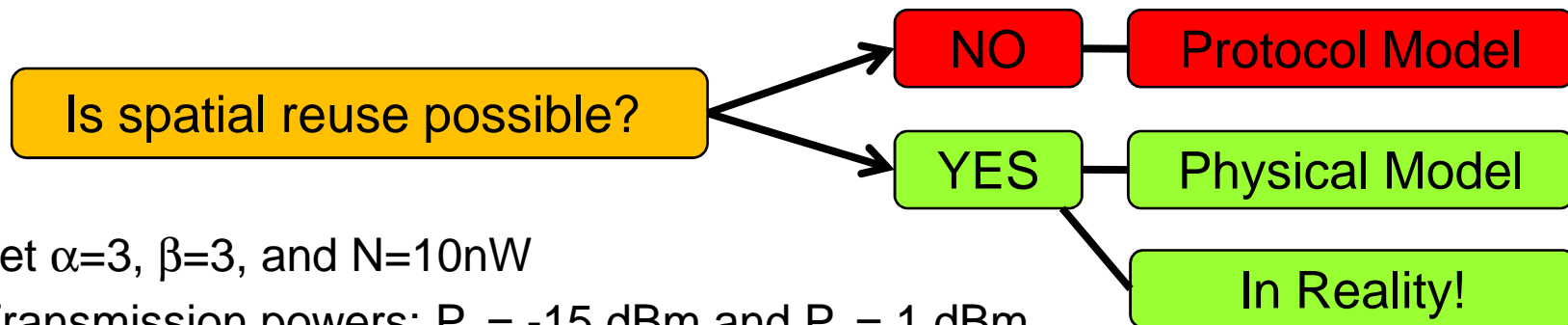


Example: Protocol vs. Physical Model

A sends to D, B sends to C





Assume a **single frequency** (and no fancy decoding techniques!)



Let $\alpha=3$, $\beta=3$, and $N=10\text{nW}$

Transmission powers: $P_B = -15\text{ dBm}$ and $P_A = 1\text{ dBm}$

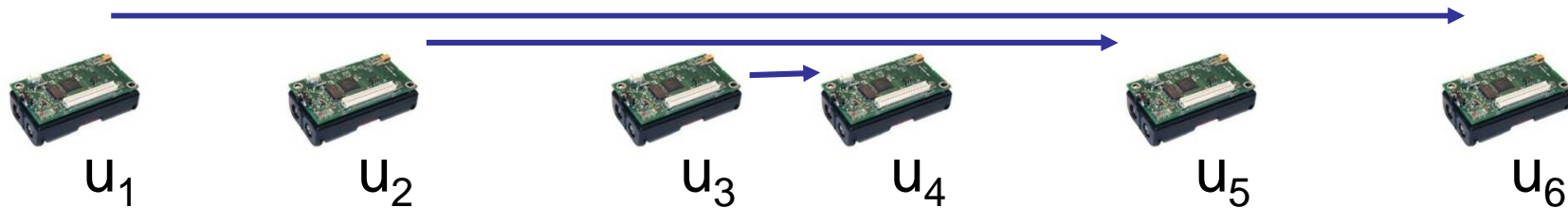
SINR of A at D: $\frac{1.26\text{mW}/(7\text{m})^3}{0.01\mu\text{W} + 31.6\mu\text{W}/(3\text{m})^3} \approx 3.11 \geq \beta$ 

SINR of B at C: $\frac{31.6\mu\text{W}/(1\text{m})^3}{0.01\mu\text{W} + 1.26\text{mW}/(5\text{m})^3} \approx 3.13 \geq \beta$ 



Real work in practice!

- Measurements using standard **mica2** nodes!
- Replaced standard MAC protocol by a (tailor-made) „**SINR-MAC**“
- Measured for instance the following deployment...



- Time for successfully transmitting 20'000 packets:

	Time required	
	standard MAC	“SINR-MAC”
Node u_1	721s	267s
Node u_2	778s	268s
Node u_3	780s	270s

	Messages received	
	standard MAC	“SINR-MAC”
Node u_4	19999	19773
Node u_5	18784	18488
Node u_6	16519	19498

Speed-up is almost a factor 3