ECE-374-B: Lecture 1 - Logistics and Strings/Languages

Lecturer: Nickvash Kani

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University of Illinois at Urbana Champaign

Course Administration

Instructional Staff

- Instructor:
 - Nickvash Kani
- Teaching Assistants:
 - Sung Woo Jeon
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- Sindhu Vydana
- Sandhya Perumenki
- Office hours: TBD, See course webpage
- Contacting us: Use <u>private notes</u> on Piazza to reach course staff. Direct email only for sensitive or confidential information.

Section A vs B

This semester, the two sections will be run completely independently.

- Different lectures.
- Different homeworks, quizzes, exams.
- Different grading policies.

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Section B will be in-person only. Recordings will be attempted but not guaranteed.

Online resources

- Webpage: General information, announcements, homeworks, quizzes, course policies https://ecealgo.com
- Submission(Gradescope): Written homework submission and grading, regrade requests. Exams will be uploaded there as well.
- Communication(Piazza): Announcements, online questions and discussion, contacting course staff (via private notes)
- Gradebook (Canvas): Announcements, online questions and discussion, contacting course staff (via private notes)

See course webpage for links

Important: check Piazza/course web page at least once each day

Grading Policy

Grading Policy: Overview

- Homeworks: 24%
- There will be approximately 9 HWs with 3 questions each.
- Hoemworks need to be submitted on Gradescope.
- Only the top 21 question grades will be considered for final grade calculation

Grading Policy: Overview

- Homeworks: 24%
- Midterm/Final exams: 75% (3 × 25%)

Exam dates:

- Midterm 1: Thurs, Feb 16, 12:30pm-1:45pm
- Midterm 2: Tues, Apr 04, 12:30pm–1:45pm
- Midterm 3: Thurs, Apr 27, 12:30pm-1:45pm
- Final: TBD

One exam will be dropped Drop policies should eliminate need for conflict exams.

Grading Policy: Overview

■ Homeworks: 24%

■ Midterm/Final exams: 75% (3 × 25%)

■ ICES Form: 1%

Piazza is crucial for improving the course experience but only 20-35% participate.

Will ask to upload a screenshot that confirms your completion of the ICES eval before grade estimates.

Discussion Sessions/Labs

- 50min problem solving session led by TAs
- Two times a week
- Go to your assigned discussion section
- Bring pen and paper!

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Discussion sections will have questions that appear on the homework. If, you skip, you're just making more work for yourself later.

Advice

- Attend lectures, please ask plenty of questions.
- Attend discussion sessions.
- Don't skip homework and don't copy homework solutions.
 Each of you should think about <u>all</u> the problems on the home work do not divide and conquer.
- Start homework early! Your mind needs time to think.
- Study regularly and keep up with the course.
- This is a course on problem solving. Solve as many as you can! Books/notes have plenty.
- This is also a course on providing rigourous proofs of correctness. Refresh your 173 background on proofs.
- Ask for help promptly. Make use of office hours/Piazza.

Miscellaneous

Please contact instructors if you need special accommodations.

Lectures are being taped (hopefully). The issue is that these recording systems are prone to failure. While I make no promises, I will try my best to record the lectures.

See course webpage for additional information.

Over-arching course questions

High-Level Questions

This course introduces three distinct fields of computer science research:

- Computational complexity.
 - Given infinite time and a certain machine, is it possible to solve a given problem.
- Algorithms
 - Given a deterministic Turing machine, how fast can we solve certain problems.
- Limits of computation.
 - Are there tasks that our computers cannot do and how do we identify these problems?

Why not just focus on Algorithms?

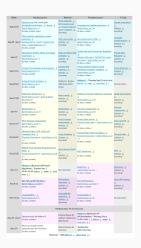
When someone asks you, "How fast can you compute problem X", they are actually asking:

- Is X solvable using the deterministic Turing machines we have at our disposal?
- If it is solvable, can we find the solution efficiently (in poly-time)?
- If it is solvable but we don't have a poly time solution, what problem(s) is it most similar too?

Course Structure

Course divided into three parts:

- Basic automata theory: finite state machines, regular languages, hint of context free languages/grammars, Turing Machines
- Algorithms and algorithm design techniques
- Undecidability and NP-Completeness, reductions to prove intractability of problems



Goals

- Algorithmic thinking
- Learn/remember some basic tricks, algorithms, problems, ideas
- Understand/appreciate limits of computation (intractability)
- Appreciate the importance of algorithms in computer science and beyond (engineering, mathematics, natural sciences, social sciences, ...)

(The Blue Weeks!)

Formal languages and complexity

Why Languages?

First 5 weeks devoted to language theory.

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But why study languages?

Multiplying Numbers

Consider the following problem:

Problem Given two *n*-digit numbers *x* and *y*, compute their product.

Grade School Multiplication

Compute "partial product" by multiplying each digit of y with x and adding the partial products.

 $\begin{array}{r}
 3141 \\
 \times 2718 \\
 \hline
 25128 \\
 3141 \\
 21987 \\
 \hline
 6282 \\
 8537238
 \end{array}$

Time analysis of grade school multiplication

- Each partial product: $\Theta(n)$ time
- Number of partial products: $\leq n$
- Adding partial products: n additions each $\Theta(n)$ (Why?)
- Total time: $\Theta(n^2)$
- Is there a faster way?

Fast Multiplication

- $O(n^{1.58})$ time [Karatsuba 1960] disproving Kolmogorov's belief that $\Omega(n^2)$ is best possible
- O(n log n log log n) [Schonhage-Strassen 1971].
 Conjecture: O(n log n) time possible
- $O(n \log n \cdot 2^{O(\log^* n)})$ time [Furer 2008]
- $O(n \log n)$ [Harvey-van der Hoeven 2019]

Can we achieve O(n)? No lower bound beyond trivial one!

Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort?

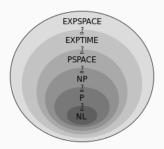
Equivalent Complexity

Does this mean multiplication is as complex as another problem that has a $O(n \log n)$ algorithm like sorting/QuickSort? How do we compare? The two problems have:

- Different inputs (two numbers vs n-element array)
- Different outputs (a number vs n-element array)
- Different entropy characteristics (from a information theory perspective)

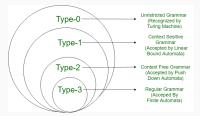
Languages, Problems and Algorithms ... oh my! II

An algorithm has a runtime complexity.



Languages, Problems and Algorithms ... oh my! III

A problem has a complexity class!



Problems do not have run-time since a problem \neq the algorithm used to solve it. *Complexity classes are defined differently.* How do we compare problems? What if we just want to know if a problem is "computable".

Algorithms, Problems and Languages ... oh my! I

Definition

- 1. An algorithm is a step-by-step way to solve a problem.
- A problem is some question that we'd like answered given some input. It should be a decision problem of the form "Does a given input fulfill property X."
- 3. A Language is a set of strings. Given a alphabet, Σ a language is a subset of Σ^*

Algorithms, Problems and Languages ... oh my! I

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- 3. A Language is a set of strings. Given a alphabet, Σ a language is a subset of Σ^* A language is a formal realization of this problem. For problem X, the corresponding language is:

 $L = \{w \mid w \text{ is the encoding of an input y to problem } X \text{ and}$ the answer to input y for a problem X is "YES" } A decision problem X is "YES" is the string is in the language.

Language of multiplication

How do we define the multiplication problem as a language?

Define L as language where inputs are separated by comma and output is separated by \mid .

Machine accepts a $x^*y=z$ if " $x^*y|z$ " is in L. Rejects otherwise.

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Define L as language where inputs are separated by comma and output is separated by |.

Machine accepts a $x^*y=z$ if " $x^*y|z$ " is in L. Rejects otherwise.

$$L_{MULT2} = \begin{cases} 1 \times 1 | 1, & 1 \times 2 | 2, & 1 \times 3 | 3, \dots \\ 2 \times 1 | 2, & 2 \times 2 | 4, & 2 \times 3 | 6, \dots \\ \vdots & \vdots & \vdots \\ n \times 1 | n, & n \times 2 | 2n, & n \times 3 | 3n, \dots \end{cases}$$
(1)

Language of sorting

We do the same thing for sorting.

Define L as language where inputs are separated by comma and output is separated by \mid .

Machine accepts a $[i_1, i_2, \ldots] = sort(\{i_1, i_2, \ldots\})$ if "x[]|z[]" is in L. Rejects otherwise.

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$$(2)$$

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$$(2)$$

If the same type of machine can recognize both languages, then that gives us an upperbound top their hardness. How do we formulate languages?

Strings

Alphabet

An alphabet is a **finite** set of symbols.

Examples of alphabets:

- $\Sigma = \{0, 1\},$
- ASCII.
- UTF8.
- $\Sigma = \{\langle \text{moveforward} \rangle, \langle \text{moveback} \rangle, \langle \text{moveleft} \rangle, \langle \text{moveright} \rangle \}$

String Definition

Definition

- 1. A string/word over Σ is a **finite sequence** of symbols over Σ . For example, '0101001', 'string', ' $\langle \text{moveback} \rangle \langle \text{rotate} 90 \rangle$ '
- 2. $x \cdot y \equiv xy$ is the concatenation of two strings
- 3. The length of a string w (denoted by |w|) is the number of symbols in w. For example, |101|=3, $|\epsilon|=0$
- 4. For integer $n \geq 0$, Σ^n is set of all strings over Σ of length n. Σ^* is the set of all strings over Σ .
- 5. Σ^* set of all strings of all lengths including empty string.

Question: $\{'a', 'c'\}^* =$

Emptiness

- ullet is a string containing no symbols. It is not a set
- $\{\epsilon\}$ is a set containing one string: the empty string. It is a set, not a string.
- \emptyset is the empty set. It contains no strings.

Question: What is $\{\emptyset\}$

Concatenation and properties

- If x and y are strings then xy denotes their concatenation.
- Concatenation defined recursively :
 - xy = y if $x = \epsilon$
 - xy = a(wy) if x = aw
- xy sometimes written as x y.
- concatenation is associative: (uv)w = u(vw) hence write $uvw \equiv (uv)w = u(vw)$
- not commutative: uv not necessarily equal to vu
- The identity element is the empty string ϵ :

$$\epsilon u = u\epsilon = u$$
.

Substrings, prefixes, Suffixes

Definition

v is substring of $w \iff$ there exist strings x, y such that w = xvy.

- If $x = \epsilon$ then v is a prefix of w
- If $y = \epsilon$ then v is a suffix of w

Subsequence

A subsequence of a string w[1...n] is either a subsequence of w[2...n] or w[1] followed by a subsequence of w[2...n].

Example

kapa is a supsequence of knapsack

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Question: How many sub-sequences are there in a string |w| = 5?

String exponent

Definition

If w is a string then w^n is defined inductively as follows:

$$w^n = \epsilon$$
 if $n = 0$
 $w^n = ww^{n-1}$ if $n > 0$

Question: $(blah)^3 =$.

Rapid-fire questions -strings

Answer the following questions taking $\Sigma = \{0, 1\}$.

- 1. What is Σ^0 ?
- 2. How many elements are there in Σ^n ?
- 3. If |u| = 2 and |v| = 3 then what is $|u \cdot v|$?
- 4. Let u be an arbitrary string in Σ^* . What is ϵu ? What is $u\epsilon$?

Languages

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A language L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

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Standard set operations apply to languages.

- For languages A, B the concatenation of A, B is $AB = \{xy \mid x \in A, y \in B\}.$
- For languages A, B, their union is $A \cup B$, intersection is $A \cap B$, and difference is $A \setminus B$ (also written as A B).
- For language $A \subseteq \Sigma^*$ the complement of A is $\bar{A} = \Sigma^* \setminus A$.

Set Concatenation

Definition

Given two sets X and Y of strings (over some common alphabet Σ) the concatenation of X and Y is

$$XY = \{xy \mid x \in X, y \in Y\} \tag{3}$$

Question:
$$X = \{fido, rover, spot\}, Y = \{fluffy, tabby\} \implies XY = .$$

Σ^* and languages

Definition

1. Σ^n is the set of all strings of length n. Defined inductively:

$$\Sigma^{n} = \{\epsilon\} \text{ if } n = 0$$

$$\Sigma^{n} = \Sigma \Sigma^{n-1} \text{ if } n > 0$$

- 2. $\Sigma^* = \cup_{n \geq 0} \Sigma^n$ is the set of all finite length strings
- 3. $\Sigma^+ = \bigcup_{n \geq 1} \Sigma^n$ is the set of non-empty strings.

Definition

A language L is a set of strings over Σ . In other words $L \subseteq \Sigma^*$.

Question: Does Σ^* have strings of infinite length?

Rapid-Fire questions - Languages

Problem

Consider languages over $\Sigma = \{0, 1\}$.

- 1. What is \emptyset^0 ?
- 2. If |L| = 2, then what is $|L^4|$?
- 3. What is \emptyset^* , $\{\epsilon\}^*$, ϵ^* ?
- 4. For what L is L* finite?
- 5. What is \emptyset^+ ?
- 6. What is $\{\epsilon\}^+$, ϵ^+ ?

Terminology Review

Let's review what we learned.

- A character(a, b, c, x) is a unit of information represented by a symbol: (letters, digits, whitespace)
- A alphabet(Σ) is a set of characters
- A string(w) is a sequence of characters
- A language(A, B, C, L) is a set of strings

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- A character(a, b, c, x) is a unit of information represented by a symbol: (letters, digits, whitespace)
- A $alphabet(\Sigma)$ is a set of characters
- A string(w) is a sequence of characters
- A language(A, B, C, L) is a set of strings
- A grammar(G) is a set of rules that defines the strings that belong to a language

Languages: easiest, easy, hard, really hard, really hard

- Regular languages.
 - Regular expressions.
 - DFA: Deterministic finite automata.
 - NFA: Non-deterministic finite automata.
 - Languages that are not regular.
- Context free languages (stack).
- Turing machines: Decidable languages.
- TM Undecidable/unrecognizable languages (halting theorem).

