

Lenguajes y Sistemas Informáticos para la resolución de problemas complejos



Procesadores de Lenguajes

Casiano Rodríguez León

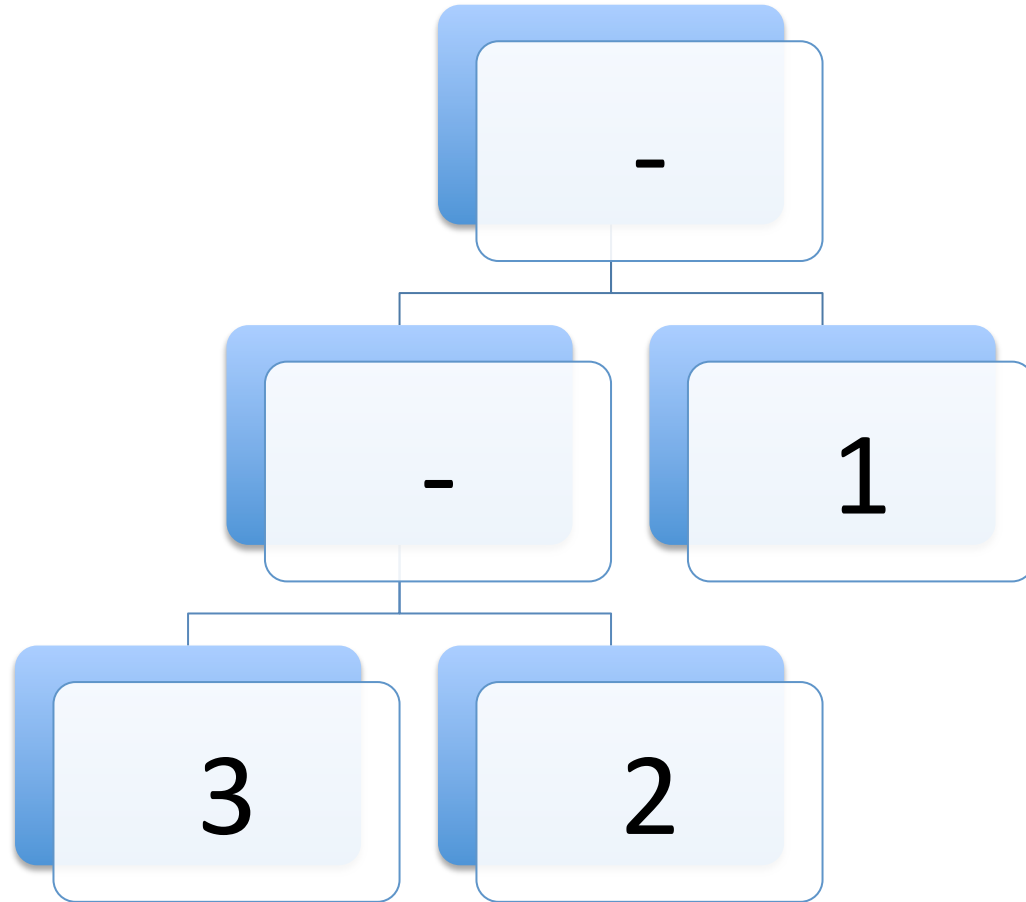
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- Optimización de Código

3 - 2 - 1

Árbol Sintáctico Abstracto

$(3-2)-1$



Semántica 3 - 2 - 1

$$0 = 1 - 1$$

-

$$1 = 3 - 2$$

-

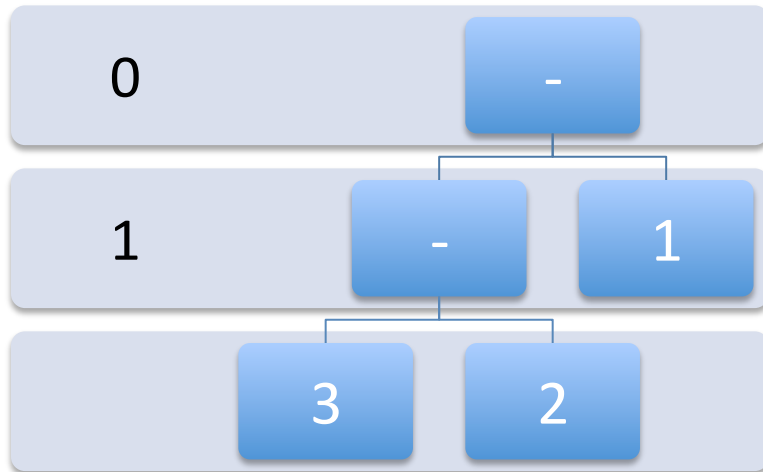
1

3

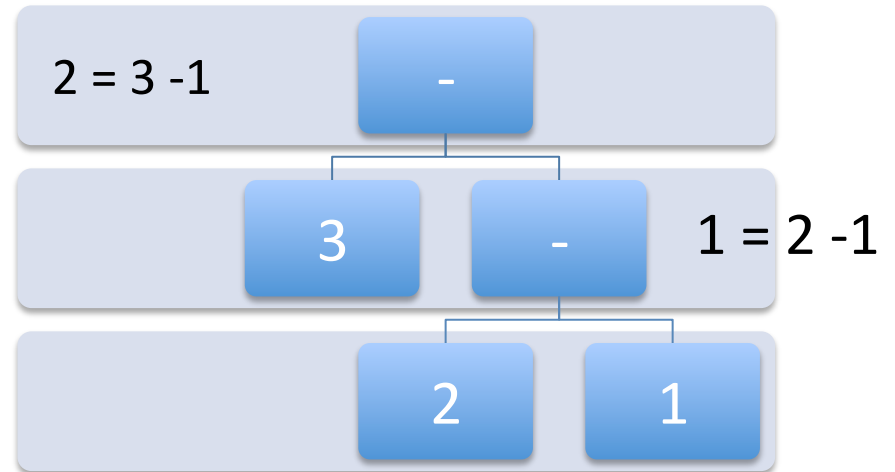
2

Semántica y Ambigüedad

$(3-2)-1$



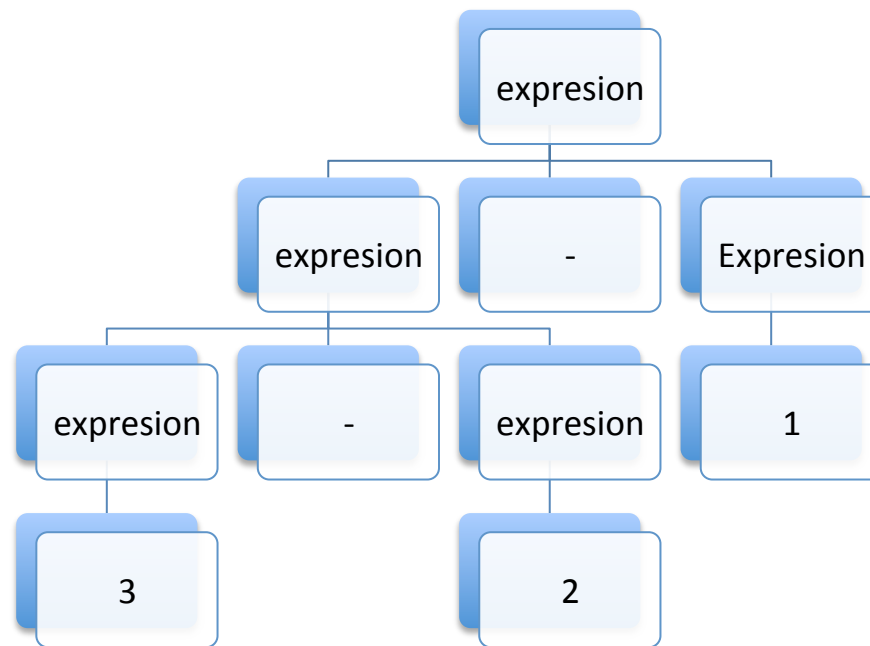
$3-(2-1)$



Gramática Independiente del Contexto

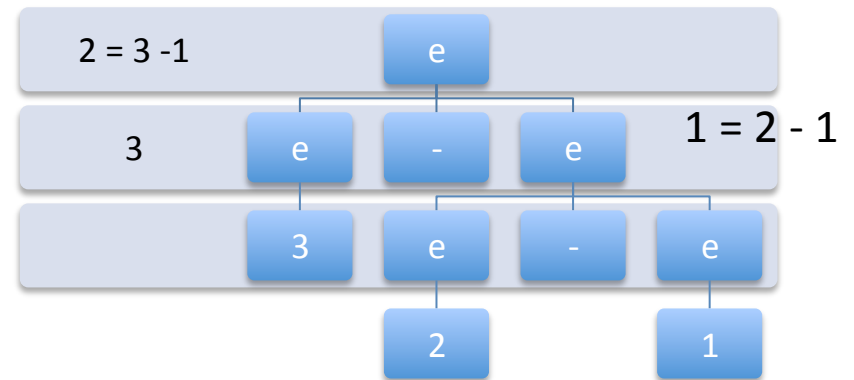
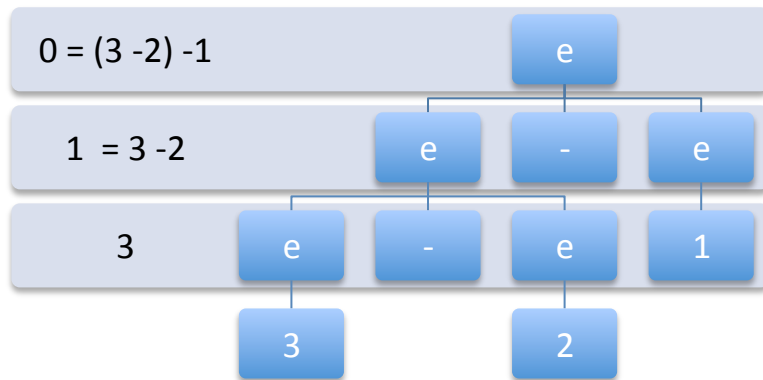
- $\text{expresion} \rightarrow \text{expresion} \text{ '-' expresion}$
- $\text{expresion} \rightarrow \text{NUMERO}$

(3-2)-1



Gramática Ambigua

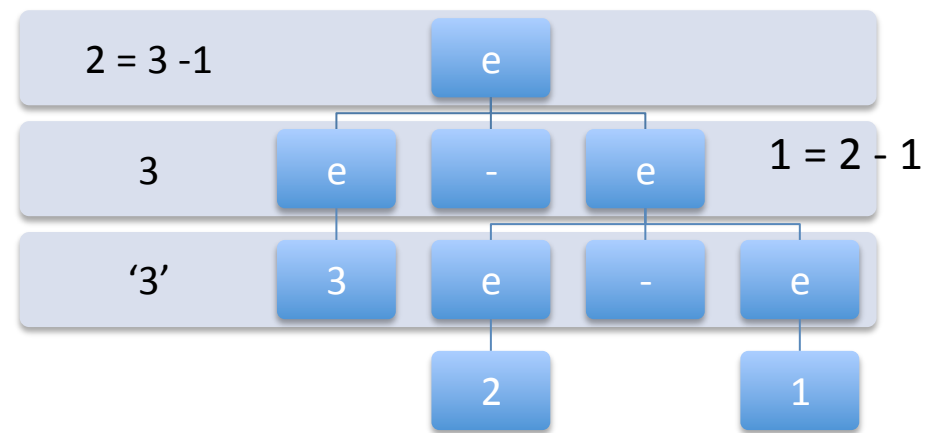
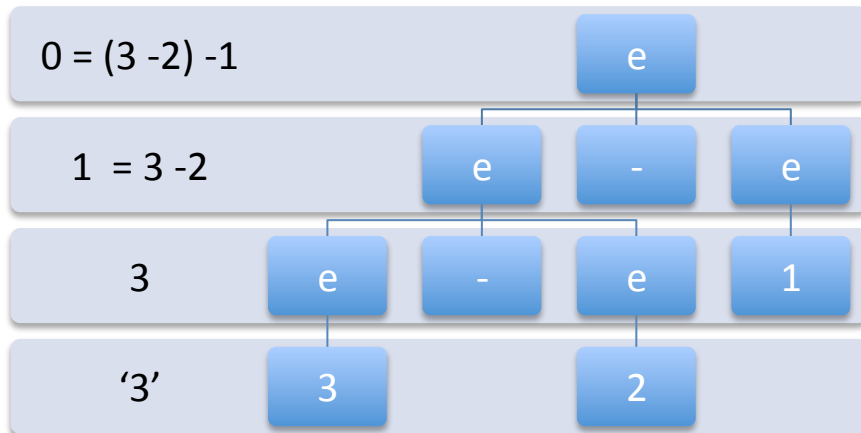
- $\text{expresion} \rightarrow \text{expresion} \text{ '-' } \text{expresion}$
- $\text{expresion} \rightarrow \text{NUMERO}$



Esquema de Traducción (yacc)

$e \rightarrow e \text{ '-' } e \quad \{ \$\$ = \$1 - \$3; \}$
 $e \rightarrow \text{NUM} \quad \{ \$\$ = \text{Number}(\$1); \}$

3 - 2 - 1



Parsing: Construcción del Árbol

$e \rightarrow e \text{ '-' } e \quad \{ \$\$ = \$1 - \$3; \}$

$e \rightarrow \text{NUM} \quad \{ \$\$ = \text{Number}(\$1); \}$

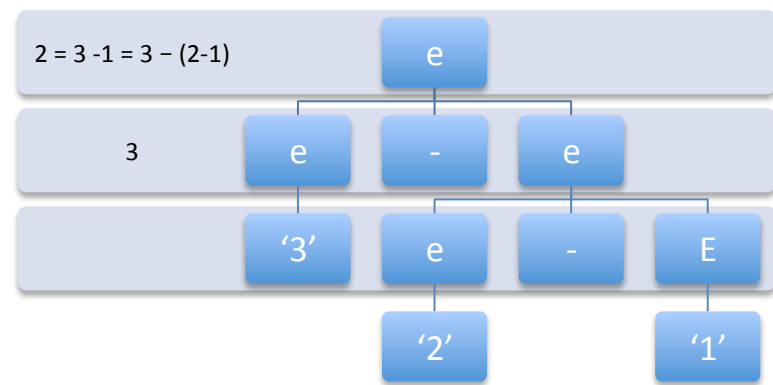
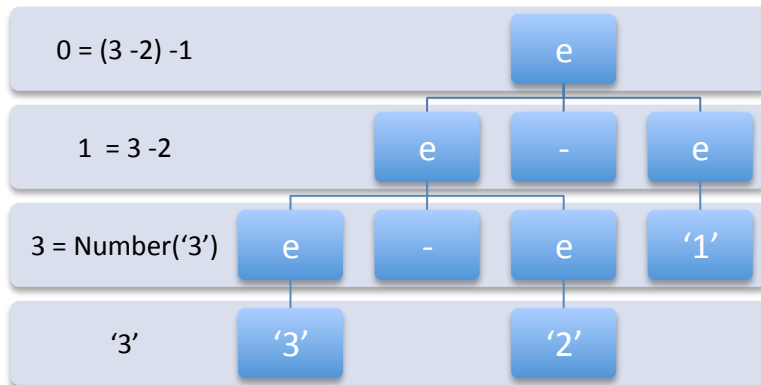
Análisis Sintáctico Ascendente:

$.3 - 2 - 1 \leq e. - 2 - 1 \leq = e -. 2 - 1 \leq = e - 2. - 1 \leq e - e. - 1$

¿Qué hacer?

1. $\leq e. - 1 \leq e -. 1 \leq e - 1. \leq e - e. \leq e.$

2. $\leq e - e -. 1 \leq e - e - 1. \leq e - e - e. \leq e - e. \leq e.$



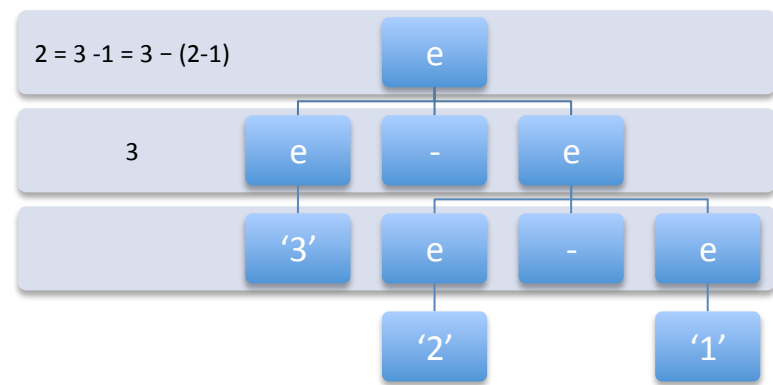
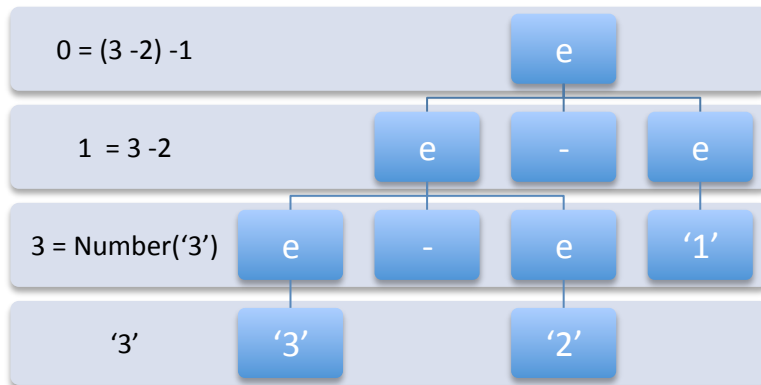
Conflicto Shift/Reduce

.3 - 2 - 1 <= e. - 2 - 1 <= = e -. 2 - 1 <= = e - 2. - 1 <= e - e. - 1

¿Qué hacer?

1. <= e. - 1 <= e - . 1 <= e - 1. <= e - e. <= e.
2. <= e - e - . 1 <= e - e - 1. <= e - e - e. <= e - e. <= e.

El conflicto puede verse como una lucha entre la regla $e \rightarrow e \text{ '-'}$ e y el terminal/token '-'



Un programa Yacc

%left ' _ ' ← En la lucha entre la regla $e \rightarrow e \text{ ' - '}$ e y el terminal/token ' - ' debe “ganar” la regla

%%

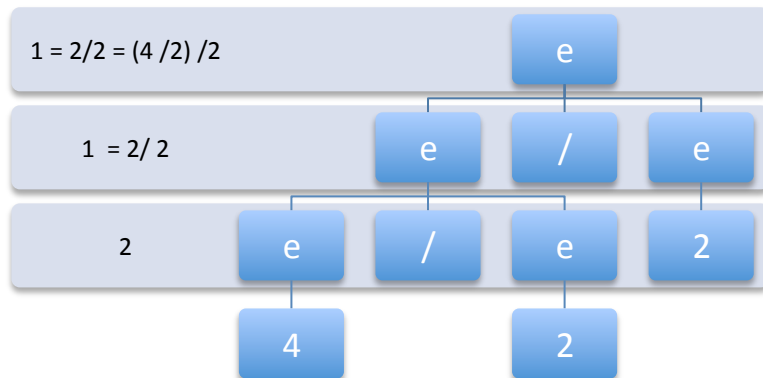
```
s : e      { return $1; }  
;
```

```
e : e ' _ ' e { $$ = $1 - $3; }  
  | NUM      { $$ = Number($1); }  
;
```

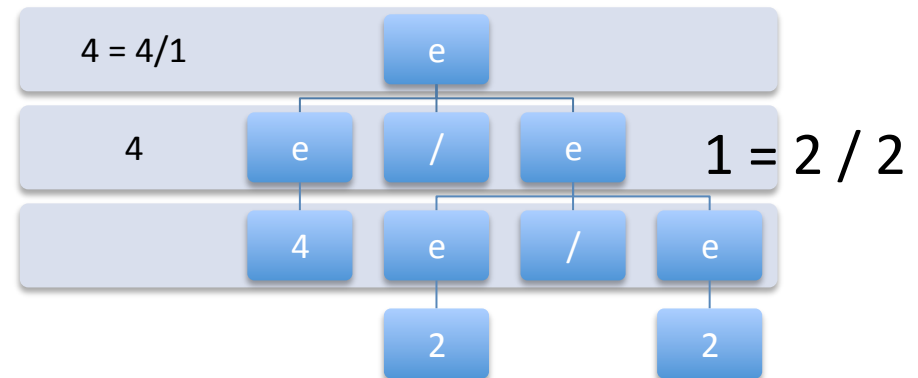
Ambigüedad: Asociatividad

$4/2/2$

$(4/2)/2$



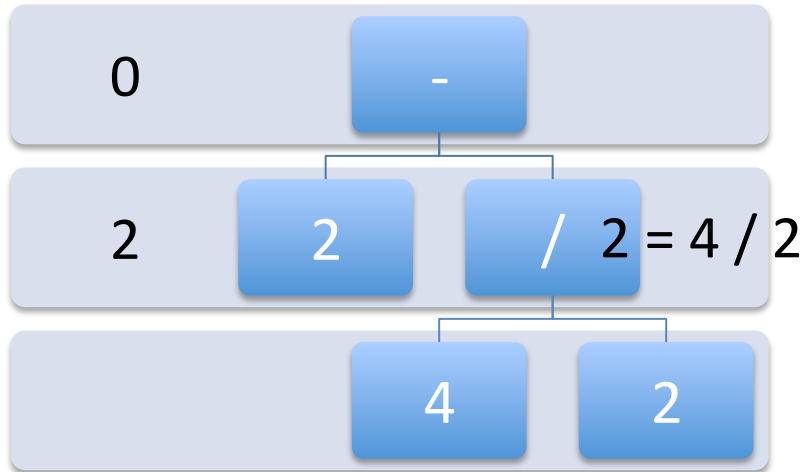
$4/(2/2)$



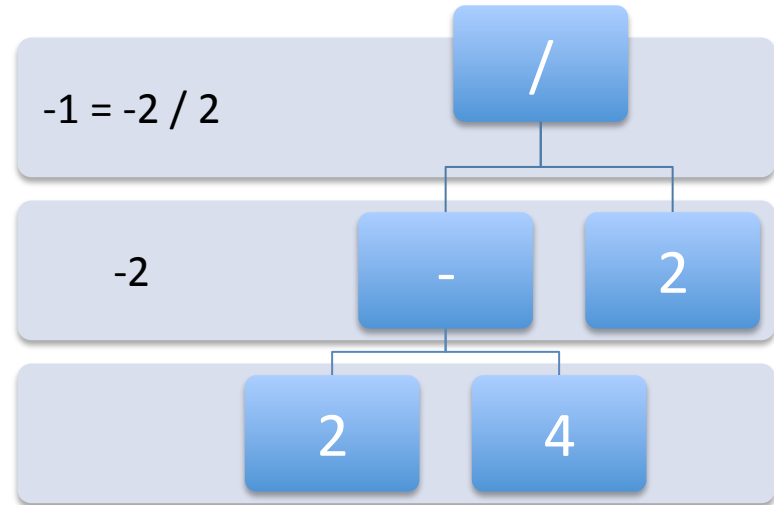
Ambigüedad: Prioridad

```
e : e ' _ ' e { $$ = $1 - $3; }  
  | e ' / ' e { $$ = $1 / $3; }  
  | NUM { $$ = Number($1); }  
  ;
```

2-(4/2)



(2-4)/2



Ambigüedad: Prioridad

2-4/2

2.-4/2 <= e.-4/2<=e-.4/2<=e-4./2<=e-e./2

¿Qué hacer?

1. <= e./2 <= e/. 2 <= e/e.<= e.
2. <= e-e/.2 <= e-e/2.<= e-e/e.<= e-e. <= e.

El conflicto es entre la regla $e \rightarrow e \text{ '-' }$ e y el terminal '/'

Ambigüedad: Prioridad



Mas prioridad

%left '-'

%left '/'

%%

e : e '-' e { \$\$ = \$1 - \$3; }

| e '/' e { \$\$ = \$1 / \$3; }

| NUM { \$\$ = Number(\$1); }

;

En la lucha entre reducir por la regla $e \rightarrow e '-' e$ y desplazar el terminal '/' debe "ganar" el token

Dynamic Resolution of Shift-Reduce Conflicts

Write a language that accepts lists of two kind of commands: *arithmetic expressions* like $4-2-1$ or one of two *commands*: *left* or *right*.

- When a *right* command is issued, the semantic of the '-' operator is changed to be right associative.
- When a *left* command is issued the semantic for '-' returns to left associative interpretation.

Dynamic Resolution of Shift-Reduce Conflicts

```
eyapp-examples — casiano@sanclemente-2:~/.../lsi-4-rpc-1819/casiano/eyapp-examples — -bash — 106x21
...vi .gitignore  ...les — -bash  ...on — -bash  ...ad — -bash  ...pp — -bash  ...as — -bash  ...to — -bash  ...ng — -bash  ...ng — -bash  ash  +

[~/.../lsi-4-rpc-1819/casiano/eyapp-examples(master)]$ cat input_for_dynamicgrammar.txt
2-1-1 # left: 0
RIGHT
2-1-1 # right: 2
LEFT
3-1-1 # left: 1
RIGHT
3-1-1 # right: 3
[~/.../lsi-4-rpc-1819/casiano/eyapp-examples(master)]$ eyapp -C dynamicgrammar.eyp
[~/.../lsi-4-rpc-1819/casiano/eyapp-examples(master)]$ ./dynamicgrammar.pm -f input_for_dynamicgrammar.txt

0
2
1
3
[~/.../lsi-4-rpc-1819/casiano/eyapp-examples(master)]$ █
```

Dynamic Resolution of Shift-Reduce Conflicts

```
eyapp-examples — casiano@sanclemente-2:~/.../lsi-4-rpc-1819/casiano/eyapp-examples — -bash — 130x32
...vi .gitignore ...les — -bash ...on — -bash ...ad — -bash ...pp — -bash ...as — -bash ...to — -bash ...ng — -bash ...ng — -bash ...les — -bash ...20 — -bash bash +

%whites /\s*(?:#.|*)?\s*/
%token NUM = /\d+/

%conflict leftORright {
    if ($reduce) { $self->YYSetReduce('-', ':M') } else { $self->YYSetShift('-') }
}

%expect 1

%%
p: c * {} ;

c:
    $expr { print "$expr\n" }
    | RIGHT { $reduce = 0 }
    | LEFT { $reduce = 1 }

;

expr:
    '(' $expr ')' { $expr }
    | %name :M
      expr.left          %PREC leftORright
        '-' expr.right   %PREC leftORright
        { $left - $right }
    | NUM

;

%%
[~/.../lsi-4-rpc-1819/casiano/eyapp-examples(master)]$
```

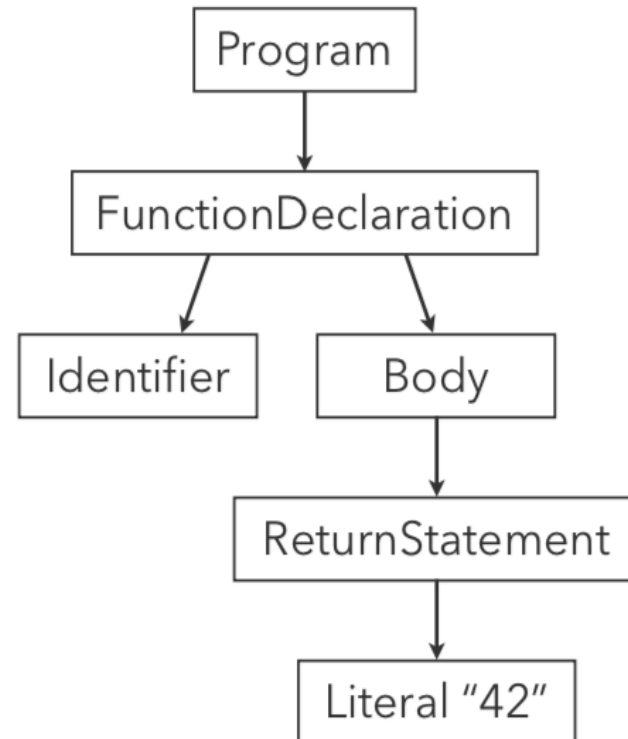
Parsing, Traversing and Code Generation

```
esprima-examples — casiano@sanclemente-2:~/campus-virtual/1819/lsi-4-rpc-1819/casiano/esprima-examples — -bash — 94x31
[~/campus-virtual/1819/lsi-4-rpc-1819/casiano/esprima-examples(master)]$ ls
ast-talk-codemotion-170406094223.pdf  esprima-pegjs-jsconfeu-talk  jsconfeu-parsing.pdf
checkstyle.js                        hello-ast.js                  node_modules
[~/campus-virtual/1819/lsi-4-rpc-1819/casiano/esprima-examples(master)]$ cat hello-ast.js
const util = require('util');
const esprima = require('esprima');
const ast = esprima.parse(`
function getAnswer() {
  return 42;
}
`);
console.log(util.inspect(ast, {depth: Math.Infinity}));
[~/campus-virtual/1819/lsi-4-rpc-1819/casiano/esprima-examples(master)]$ node hello-ast.js
Script {
  type: 'Program',
  body:
    [ FunctionDeclaration {
      type: 'FunctionDeclaration',
      id: Identifier { type: 'Identifier', name: 'getAnswer' },
      params: [],
      body:
        BlockStatement {
          type: 'BlockStatement',
          body:
            [ ReturnStatement {
              type: 'ReturnStatement',
              argument: Literal { type: 'Literal', value: 42, raw: '42' } } ] ],
      generator: false,
      expression: false,
      async: false } ],
  sourceType: 'script' }
```

Parsing, Traversing and Code Generation

```
1  
2 function getAnswer() {  
3   return 42;  
4 }
```

```
Script {  
  type: 'Program',  
  body:  
    [ FunctionDeclaration {  
      type: 'FunctionDeclaration',  
      id: Identifier { type: 'Ident  
      params: [],  
      body:  
        BlockStatement {  
          type: 'BlockStatement',  
          body:  
            [ ReturnStatement {  
              type: 'ReturnStatemen  
              argument: Literal { t  
            generator: false,  
            expression: false,  
            async: false } ],  
            sourceType: 'script' }  
          ]  
        }  
      ]  
    }  
  ]  
}
```



Parsing and Traversing Example: Logging function calls

```
esprima-examples — casiano@sanclemente-2:~/.../lsi-4-rpc-1819/casiano/esprima-examples — -bash — 110x28
[~/.../lsi-4-rpc-1819/casiano/esprima-examples(master)]$ ./logging-dibad.js prueba-logging-dibad.js
input:
function foo(a, b) {
  var x = 'blah';
  var y = (function (z) {
    return z+3;
  })(2);
}
foo(1, 'wut', 3);

—
output:
function foo(a, b) {
  console.log(`Entering foo(${ a },${ b })`);
  var x = 'blah';
  var y = function (z) {
    console.log(`Entering <anonymous function>(${ z })`);
    return z + 3;
  }(2);
}
foo(1, 'wut', 3);

[~/.../lsi-4-rpc-1819/casiano/esprima-examples(master)]$ node out-prueba-logging-dibad.js
Entering foo(1,wut)
Entering <anonymous function>(2)
[~/.../lsi-4-rpc-1819/casiano/esprima-examples(master)]$
[~/.../lsi-4-rpc-1819/casiano/esprima-examples(master)]$
```

<https://astexplorer.net/>

The screenshot shows the AST Explorer web application. The left pane contains the source code of a JavaScript function:

```
1 function foo(a, b) {  
2   var x = 'blah';  
3   var y = (function (z) {  
4     return z+3;  
5   })(2);  
6 }  
7 foo(1, 'wut', 3);
```

The right pane displays the Abstract Syntax Tree (AST) for this code, generated by the `esprima-4.0.1` parser. The tree structure is as follows:

- `- Program {`
 - `type: "Program"`
 - `- body: [`
 - `- FunctionDeclaration {`
 - `type: "FunctionDeclaration"`
 - `- id: Identifier {`
 - `type: "Identifier"`
 - `name: "foo"`
 - `+ range: [2 elements]`
 - `+ params: [2 elements]`
 - `- body: BlockStatement {`
 - `type: "BlockStatement"`
 - `- body: [`
 - `+ VariableDeclaration {type, declarations, kind, range}`
 - `+ VariableDeclaration {type, declarations, kind, range}`
 - `+ range: [2 elements]`
- `generator: false`
- `expression: false`

At the bottom of the interface, it states: "Built with [React](#), [Babel](#), [Font Awesome](#), [CodeMirror](#), [Express](#), and [webpack](#) | [GitHub](#)".

<https://astexplorer.net/>

The screenshot shows the AST Explorer interface with the following components:

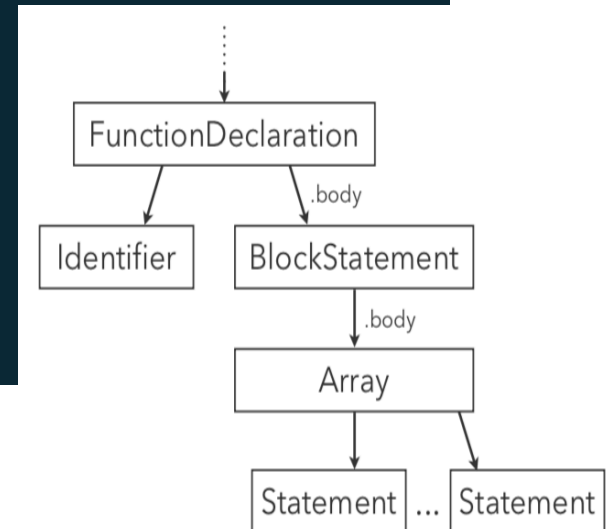
- Code Editor:** Contains the JavaScript code:

```
1 function foo(a, b) {  
2   var x = 'blah';  
3   var y = (function (z) {  
4     return z+3;  
5   })(2);  
6 }  
7 foo(1, 'wut', 3);
```
- AST View:** Shows the JSON representation of the AST. The root node is `Program`, which has a `body` array containing a `FunctionDeclaration` node. The `FunctionDeclaration` node has an `id` of `foo` and a `params` array of `[a, b]`. The `body` of the function is a `BlockStatement` containing a `VariableDeclaration` for `x` and an `ExpressionStatement` for the function call `foo(1, 'wut', 3)`.
- Diagram:** A tree diagram illustrating the structure of the AST. The root node is `FunctionDeclaration`, which has two children: `Identifier` (representing the function name `foo`) and `BlockStatement`. The `BlockStatement` node has a `.body` property that points to an `Array` node. This `Array` node contains two `Statement` nodes, representing the function's body.

Mostrar todas

Parsing, Traversing and Generating Code

```
function addLogging(code) {  
  var ast = esprima.parse(code);  
  estraverse.traverse(ast, {  
    enter: function(node, parent) {  
      if (node.type === 'FunctionDeclaration'  
        || node.type === 'FunctionExpression') {  
        addBeforeCode(node);  
      }  
    }  
  });  
  return escodegen.generate(ast);  
}
```



API de estraverse: <https://github.com/estools/estraverse>

Parsing, Traversing and Generating Code

```
function addLogging(code) {  
  var ast = esprima.parse(code);  
  estraverse.traverse(ast, {  
    enter: function(node, parent) {  
      if (node.type === 'FunctionDeclaration'  
        || node.type === 'FunctionExpression') {  
        addBeforeCode(node);  
      }  
    }  
  });  
}
```

AST Explorer

Snippet JavaScript </> esprima Transform default ? Parser: [esprima-4.0.1](#)

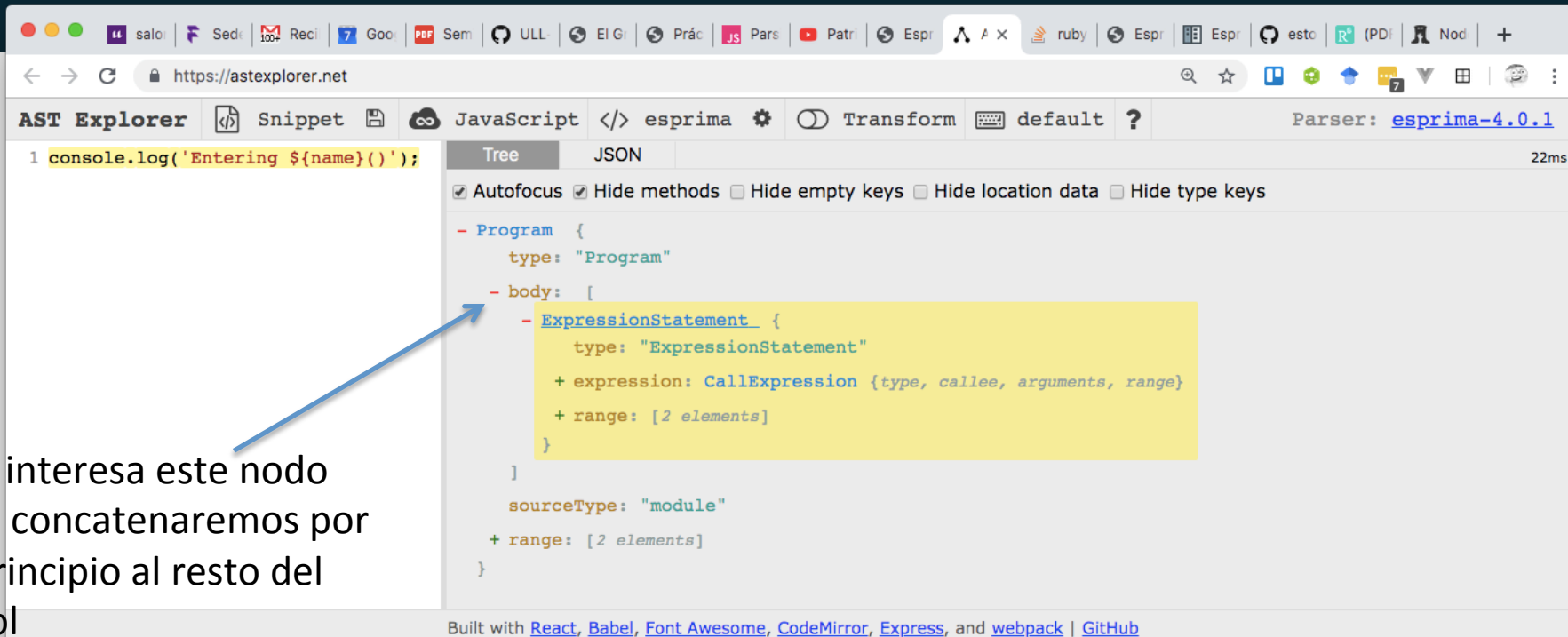
Tree JSON 19ms

☒ Autofocus ☒ Hide methods ☐ Hide empty keys ☐ Hide location data ☐ Hide type keys

```
- body: [  
  + VariableDeclaration {type, declarations, kind, range}  
  - VariableDeclaration {  
    type: "VariableDeclaration"  
    - declarations: [  
      - VariableDeclarator {  
        type: "VariableDeclarator"  
        + id: Identifier {type, name, range}  
        - init: CallExpression {  
          type: "CallExpression"  
          + callee: FunctionExpression {type, id, params, body, generator, ... +3}  
          + arguments: [1 element]  
          + range: [2 elements]  
        }  
        + range: [2 elements]  
      }  
    ]  
  }  
]
```

Traversing and Modifying the AST

```
function addBeforeCode(node) {  
  var name = node.id ? node.id.name : '<anonymous function>';  
  var beforeCode = "console.log('Entering " + name + "()');";  
  var beforeNodes = esprima.parse(beforeCode).body;  
  node.body.body = beforeNodes.concat(node.body.body);  
}
```



The screenshot shows the AST Explorer interface. The input code is `1 console.log('Entering ${name}()');`. The AST tree view is displayed, showing the following structure:

```
{  
  type: "Program",  
  body: [  
    {  
      type: "ExpressionStatement",  
      expression: {  
        type: "CallExpression",  
        callee: {  
          type: "Identifier",  
          name: "console",  
          range: [1, 14]  
        },  
        arguments: [  
          {  
            type: "TemplateLiteral",  
            quasis: [  
              {  
                value: "Entering ",  
                range: [18, 30]  
              },  
              {  
                value: "${name}",  
                range: [32, 40]  
              },  
              {  
                value: "()",  
                range: [42, 47]  
              }  
            ],  
            range: [18, 47]  
          }  
        ],  
        range: [15, 47]  
      },  
      range: [1, 47]  
    }  
  ],  
  sourceType: "module",  
  range: [1, 47]  
}
```

A blue arrow points from the text "Nos interesa este nodo" to the `ExpressionStatement` node in the `body` array.

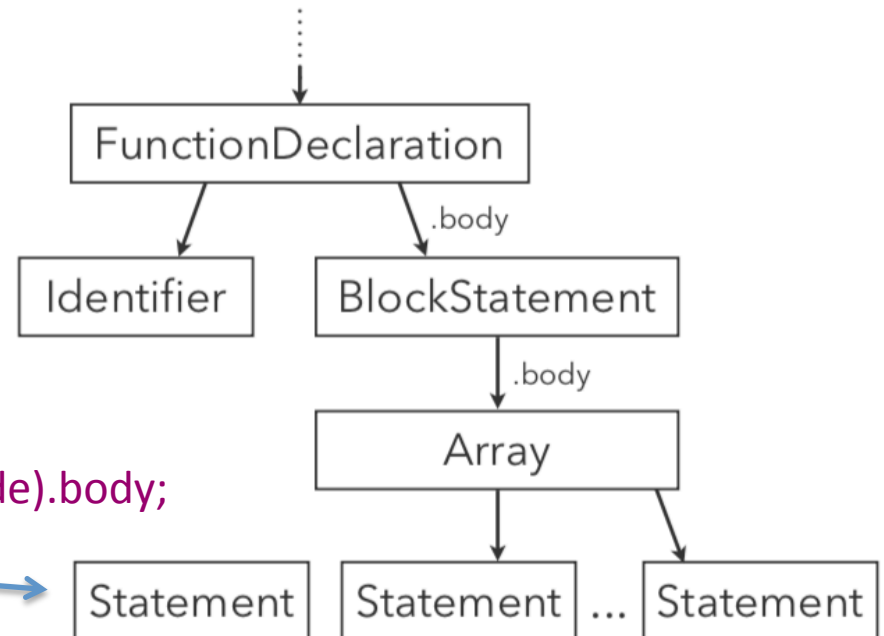
Nos interesa este nodo
Que concatenaremos por
el principio al resto del
árbol

Built with [React](#), [Babel](#), [Font Awesome](#), [CodeMirror](#), [Express](#), and [webpack](#) | [GitHub](#)

Traversing and Modifying the AST

```
function addBeforeCode(node) {  
  var name = node.id ? node.id.name : '<anonymous function>';  
  var beforeCode = "console.log('Entering " + name + "()');";  
  var beforeNodes = esprima.parse(beforeCode).body;  
  node.body.body = beforeNodes.concat(node.body.body);  
}
```

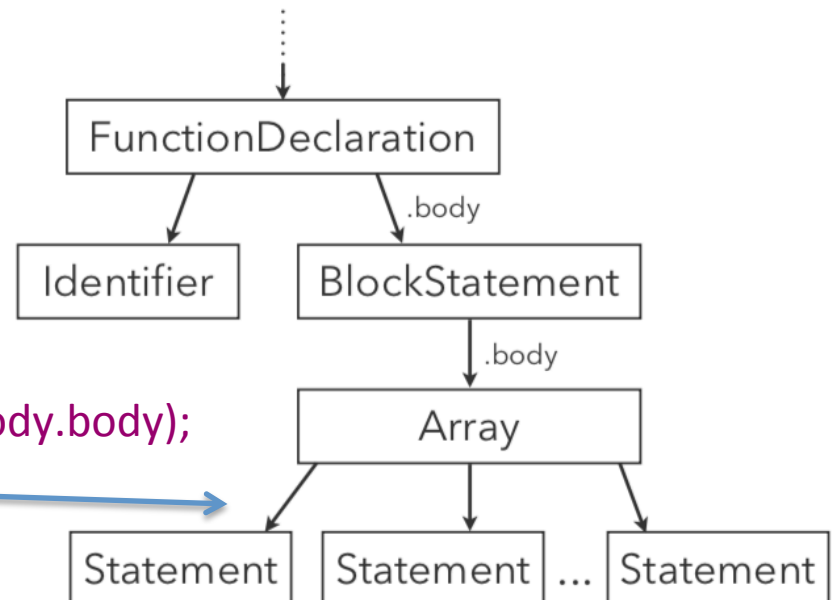
`var beforeNodes = esprima.parse(beforeCode).body;`



Parsing, Traversing and Modifying the AST

```
function addBeforeCode(node) {  
  var name = node.id ? node.id.name : '<anonymous function>';  
  var beforeCode = "console.log('Entering " + name + "()');";  
  var beforeNodes = esprima.parse(beforeCode).body;  
  node.body.body = beforeNodes.concat(node.body.body);  
}
```

`node.body.body = beforeNodes.concat(node.body.body);`



El método `concat()` se usa para unir dos o más arrays

Generating Code from the AST

```
const escodegen = require('escodegen');

let result = escodegen.generate({
  type: 'BinaryExpression',
  operator: '+',
  left: { type: 'Literal', value: 40 },
  right: { type: 'Literal', value: 2 }
});

console.log(result); //40 + 2
```

```
[~/.../lsi-4-rpc-1819/casiano/esprima-examples(master)]$ node escodegen-hello.js
40 + 2
```

API de escodegen: <https://github.com/estools/escodegen/wiki/API>

Generating Code from the AST

```
function addLogging(code) {  
  var ast = esprima.parse(code);  
  estraverse.traverse(ast, {  
    enter: function(node, parent) {  
      if (node.type === 'FunctionDeclaration'  
          || node.type === 'FunctionExpression') {  
        addBeforeCode(node);  
      }  
    }  
  });  
  return escodegen.generate(ast);  
}
```

Code Optimization

A Survey on Compiler Autotuning using Machine Learning

Accepted in ACM Computing Surveys 2018 (Received Nov 2016, Revised Aug 2017)

<https://doi.org/10.1145/3197978>

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Since the mid-1990s, researchers have been trying to use machine-learning based approaches to solve a number of different compiler optimization problems. These techniques primarily enhance the quality of the obtained results and, more importantly, make it feasible to tackle two main compiler optimization problems: optimization selection (choosing which optimizations to apply) and phase-ordering (choosing the order of applying optimizations). The compiler optimization space continues to grow due to the advancement of applications, increasing number of compiler optimizations, and new target architectures. Generic optimization passes in compilers cannot fully leverage newly introduced optimizations and, therefore, cannot keep up with the pace of increasing options. This survey summarizes and classifies the recent advances in using machine learning for the compiler optimization field, particularly on the two major problems of (1) selecting the best optimizations, and (2) the phase-ordering of optimizations. The survey highlights the approaches taken so far, the obtained results, the fine-grain classification among different approaches and finally, the influential papers of the field.

Additional Key Words and Phrases: Compilers, Autotuning, Machine Learning, Phase ordering, Optimizations, Application Characterization

Code Optimization

4.2.2 Evolutionary Algorithms. Evolutionary algorithms are inspired by biological evolution such as the process of natural selection and mutation.

space play the role of individuals in a population. A solution is using a fitness function such as an execution takes place after the repeated application of the fitness some of the more notable techniques used in the literature

Genetic Algorithm (GA) is a meta-heuristic algorithm any other machine learning technique or be used independently. NSGA-II (Non-dominated Sorting Genetic Algorithm II) multi-objective optimization problems and have had numerous architecture domain [168, 229]. NSGA-II is shown to be a classic GA algorithms.

Another interesting evolutionary model is Neuro Evolution. They proved to be a powerful model for learning complex the network topology and parameter weight to find the best-balanced fitness function. NEAT specifically has been used in many notable recent research work as well [66, 151, 152]. This section summarized a few notable research work that used evolutionary algorithms.

Cooper et al. [69, 70] addressed the code size issue of the generated binaries by using genetic algorithms. The results of this approach were compared to an iterative algorithm generating fixed optimization sequence and also at random frequency. Given the comparison, the authors concluded that by using their GAs they could develop new fixed optimization sequences that generally work well on reducing the code-size of the binary.

Knijnenburg et al. [142] proposed an iterative compilation approach to tackle the selection size of the tiling and the unrolling factors in an architecture independent manner. The authors evaluated their approach using a number of state-of-the-art iterative compilation techniques, e.g., simulated annealing and GAs, and a native Fortran77 or g77 compiler enabling optimizations for Matrix-Matrix Multiplication ($M \times M$), Matrix-Vector Multiplication ($M \times V$), and Forward Discrete Cosine Transform.

Juan Fumero

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CV



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About me

I am a research associate working as part of the Advanced Processor Technologies (APT) Research Group at The University of Manchester on Heterogeneous High-Level Languages Virtual Machines using GPUs and FPGAs. Currently, I am working as part of the E2Data European project for bringing automatic GPU and FPGA JIT compilation and execution for Java programs. I am interested in combining GPGPU computing and FPGA acceleration with interpreted programming languages such as R, Ruby and Java through automatic parallelisation, compilation and transparent execution.

I did my PhD at The University of Edinburgh on Accelerating Interpreted Programming Languages on GPUs with Just-In-Time and Runtime Optimisations. I extended the Graal JIT compiler and the Partial Evaluator to allow programmers to automatically execute Java, R and Ruby programs on GPUs via OpenCL. My PhD was fully funded by Oracle Labs. You can see more details in my PhD Thesis.

Before doing the PhD, I did my degree at The University of La Laguna (Tenerife, Spain) and I developed my final project on the accULL compiler, the first open-source OpenACC compiler under the supervision of Dr. Ruyman Reyes and Dr. Francisco de Sande.

Recursos

- [Repositorio GitHub con los recursos de la charla: https://github.com/ULL-LSI/campus-america-2019](https://github.com/ULL-LSI/campus-america-2019)
- [Apuntes de Procesadores de Lenguajes. Curso 2018/2019: https://ull-esit-pl-1819.github.io/introduccion/](https://ull-esit-pl-1819.github.io/introduccion/)
- [Rodriguez-Leon, Casiano & Garcia-Forte, L. \(2011\). Solving Difficult LR Parsing Conflicts by Postponing Them. Comput. Sci. Inf. Syst.. 8. 517-531. 10.2298/CSIS101116008R.](#)
- [Parse Eyapp](#) en CPAN
- [Parsing Strings and Trees with Parse::Eyapp](#) (An Introduction to Compiler Construction). 2010
- [Patrick Dubroy: Parsing, Compiling, and Static Metaprogramming](#)
- <https://astexplorer.net/>

Any(Questions)?

The screenshot shows the AST Explorer interface at <https://astexplorer.net>. The input code is `1 Any(questions)? 'yes' : 'no'`. The AST is displayed in the right pane, showing a conditional expression (CallExpression) with a callee of type Identifier (name: "Any") and an argument of type Identifier (name: "questions"). The consequent is a Literal value "yes" and the alternate is a Literal value "no".

```
- test: CallExpression {
  type: "CallExpression"
  start: 0
  end: 14
  - callee: Identifier {
    type: "Identifier"
    start: 0
    end: 3
    name: "Any"
  }
  - arguments: [
    - Identifier = $node {
      type: "Identifier"
      start: 4
      end: 13
      name: "questions"
    }
  ]
}
+ consequent: Literal {type, start, end, value, raw}
- alternate: Literal {
```

Parser: [acorn-6.1.1](#)

1ms

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campus-ameri....png ^

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