

Algorithmics	Student information	Date	Number of session
	UO: 275725		3.1
	Surname: Gómez Menéndez		
	Name: Laura		

Activity 1. [Basic recursive models.]

A brief explanation for each of the given classes indicating how you calculated the complexity of that class.

- **Subtraction1**

```
public class Subtraction1{
    public static long rec1(int n) {
        long cont = 0;
        if (n<=0)
            cont++;
        else {
            cont++; // O(1)=O(n^0)
            rec1(n-1);
        }
        return cont;
    }
}
```

In this method we only have one subproblem ($\text{rec}(n-1)$), we can say that $a = 1$.

We know that $b = 1$ because the number that subtracts n in the subproblem is 1.

As there is no loops, $k = 0$.

In conclusion, as $a = 1$, the complexity is $O(n^{(k+1)})$ that is **$O(n)$** .

- **Subtraction2**

```
public static long rec2(int n) {
    long cont = 0;
    if (n<=0)
        cont++;
    else {
        for (int i=0;i<n;i++)
            cont++; // O(n)
        rec2(n-1);
    }
    return cont;
}
```

In this method we only have one subproblem ($\text{rec2}(n-1)$), we can say that $a = 1$.

We know that $b = 1$ because the number that subtracts n in the subproblem is 1.

As there is one loop, $k = 1$.

In conclusion, as $a = 1$, the complexity is $O(n^{(k+1)})$ that is **$O(n^2)$** .

Algorithmics	Student information	Date	Number of session
	UO: 275725		3.1
	Surname: Gómez Menéndez		
	Name: Laura		

- **Subtraction3**

```
public static long rec3(int n) {
    long cont = 0;
    if (n<=0)
        cont++;
    else {
        cont++; //O(1)
        rec3(n-1);
        rec3(n-1);
    }
    return cont;
}
```

In this method we have two subproblems ($\text{rec3}(n-1)$), we can say that $a = 2$.

We know that $b = 1$ because the number that subtracts n in the subproblem is 1.

As there is no loops, $k = 0$.

In conclusion, as $a > 1$, the complexity is $O(a^{(n/b)})$ that is **$O(a^n)$** .

- **Division1**

```
public static long rec1 (int n) {
    long cont = 0;
    if (n<=0) cont++;
    else {
        for (int i=1;i<n;i++)
            cont++ ; //O(n)
        rec1(n/3);
    }
    return cont;
}
```

In this method we have one subproblem ($\text{rec1}(n/3)$), we can say that $a = 1$.

We know that $b = 3$ because the number that divides n in the subproblem is 3.

As there is one loop, $k = 1$.

In conclusion, as $a < b^k$, the complexity is $O(n^k)$ that is **$O(n)$** .

Algorithmics	Student information	Date	Number of session
	UO: 275725		3.1
	Surname: Gómez Menéndez		
	Name: Laura		

- **Division2**

```
public static long rec2 (int n) {
    long cont = 0;
    if (n<=0) cont++;
    else {
        for (int i=1;i<n;i++)
            cont++ ; //O(n)
        rec2(n/2);
        rec2(n/2);
    }
    return cont;
}
```

In this method we have two subproblems ($\text{rec2}(n/2)$), we can say that $a = 2$.

We know that $b = 2$ because the number that divides n in the subproblem is 2.

As there is one loop, $k = 1$.

In conclusion, as $a = b^k$, the complexity is $O((n^k) \cdot \log n)$ that is **$O(n \log n)$** .

- **Division3**

```
public static long rec3 (int n) {
    long cont = 0;
    if (n<=0)
        cont++;
    else {
        cont++ ; // O(1)
        rec3(n/2);
        rec3(n/2);
    }
    return cont;
}
```

In this method we have two subproblems ($\text{rec3}(n/2)$), we can say that $a = 2$.

We know that $b = 2$ because the number that divides n in the subproblem is 2.

As there no loops, $k = 0$.

In conclusion, as $a > b^k$, the complexity is $O(n^{(\log_b a)})$ that is **$O(n^{\log_2 2})$** .

Algorithmics	Student information	Date	Number of session
	UO: 275725		3.1
	Surname: Gómez Menéndez		
	Name: Laura		

A brief explanation for each of the two new classes indicating how you calculate the complexity to get the requested one.

- **Subtraction4**

I was asked to write a recursive method by subtraction with a complexity $O(3^{(n/2)})$.

With that information I can know that $a > 1$, and $a = 3$, so I had to write three subproblems.

I can also know that $b = 2$, so the number that divides n in the subproblems is 2. The variable k does not matter at all, so I did not write loops and $k = 0$ in my class.

```
public static long rec1(int n) {
    long cont = 0;
    if (n<=0)
        cont++;
    else {
        cont++;
        rec1(n-2);
        rec1(n-2);
        rec1(n-2);
    }
    return cont;
}
```

- **Division4**

This time I was asked to create a new recursive method by division with a complexity $O(n^2)$ and a number of subproblems = 4. The number of subproblems is the same as the variable a , so $a = 4$. With that information I can know that $a < b^k$, and that $k = 2$ what means that I had to write two loops.

I had to look for a number whose square was bigger than 4, so I chose $b = 3$ ($4 < 9$), what means that the number that divides n in the subproblems is 3.

```
public static long rec2 (int n) {
    long cont = 0;
    if (n<=0) cont++;
    else {
        for (int i=1;i<n;i++)
            for(int j=i;j<n;j++)
                cont++ ; //O(n)

        rec2(n/3);
        rec2(n/3);
        rec2(n/3);
        rec2(n/3);
    }
    return cont;}
}
```

Algorithmics	Student information	Date	Number of session
	UO: 275725		3.1
	Surname: Gómez Menéndez		
	Name: Laura		