


Algorithmics	Student information	Date	Number of session
	UO:UO277653	04/03/2021	5
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## Activity 1. Basic recursive models.

- **A brief explanation for each of the given classes indicating how you calculated the complexity of that class.**

We have to take into consideration that  $a$  is the number of subproblems,  $b$  is the constant by which we reduce the subproblem, and  $k$  is the complexity of the code when excluding the recursive calls, and also, for divide and conquer by subtraction, we consider the following:

### ➤ Complexity

- $O(n^k)$  if  $a < 1$  (never happens)
- $O(n^{k+1})$  if  $a = 1$
- $O(a^{n \div b})$  if  $a > 1$

With this information we can compute the complexity of the classes having  $a$ ,  $b$ , and  $k$ :

Subtraction 1:  $a=1$ ,  $b=1$  and  $k=0$ . So the complexity is  $O(1)$

Subtraction 2:  $a=1$ ,  $b=1$  and  $k=1$ . So the complexity is  $O(n^2)$

Subtraction 3:  $a=2$ ,  $b=1$  and  $k=0$ . So the complexity is  $O(2^n)$

For divide and conquer by division we consider this other one:

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## ➤ Complexity

- $O(n^k)$  if  $a < b^k$
- $O(n^k * \log n)$  if  $a = b^k$
- $O(n^{\log_b a})$  if  $a > b^k$

Division 1:  $a=1$  ,  $b=3$  and  $k=1$ . So the complexity is  $O(n)$

Division 2:  $a=2$  ,  $b=2$  and  $k=1$ . So the complexity is  $O(n \log n)$

Division 3:  $a=2$  ,  $b=2$  and  $k=0$ . So the complexity is  $O(n)$

- **A brief explanation for each of the two new classes indicating how you calculate the complexity to get the requested one.**

Using the same images as before:

For Subtraction4, for getting a complexity  $O(3^{n/2})$ , we see that the exponential complexity is obtained when  $a > 1$ , and looking at the arguments, we can deduce that  $b$  should be equal to 2 and  $a = 3$ . The value of  $k$  is not relevant. For getting  $b = 2$  we reduce the value of  $n$  by 2 in every recursive call, and to get  $a = 3$  put 3 recursive calls to the method.

For Division4, for getting a complexity  $O(n^2)$ , with a number of subproblems  $a = 4$ , we have to make sure that  $a < b^k$ , and that  $k = 2$ , for that purpose we can use for example  $b = 3$ , as it satisfies the condition.