	Student information	Date	Number of session
	UO: 302165	27/02/2025	3
Algorithmics	Surname: Uña Garcia		✓ Escuela de



Ingeniería

Name: Lucas

Activity 1. [D&C BY SUBSTRACTION]

Substraction1:

a = 1, b = 1, k = 0

Complexity:

$$O(n^{k+1}) \Rightarrow O(n)$$

I cannot prove if the results fit or not with the theoretical complexity because the times are not accurate (below 50ms) and then the stack overflow exception rises.

Substraction2:

a = 1, b = 1, k = 1

Complexity:

 $O(n^{k+1}) => O(n^2)$

The results obtained fits with the theoretical complexity. Although for values of n < 2048are not accurate, for higher values until the stack overflow exception we can appreciate that the time grows by a factor of two, which is the square of the increment of the size (2²).

Substraction3:

a = 2, b = 1, k = 0

Complexity:

 $O(a^{n/b}) => O(2^n)$

The results obtained are the ones expected according to the theoretical complexity, because as n increases one by one, the time increases exponentially (the new time is the previously obtained times 2, as $2^{n+1} = 2^n * 2$).

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The operation is aborted not because the time but the memory occupied. As for each iteration a variable is created, the stack has less and less space until it is full occupied.

- N1=20
$$\rightarrow$$
 t1 = 41
N2=80 \rightarrow t2 = ?
t2 = $2^{N2}/2^{N1}$ * t1 = 2^{N2-N1} * t1 = 2^{60} * t1 = 47.269.781.688.880.726.016ms * 1s/1000ms * 1h/3600s * 1day/24h * 1year/365.25days = 1.497.888.993,1 years

Substraction4			
n time(ms)			
100	1,425		
200	10,65		
400	82,15		
800	650		
1600	5013		
3200	40532		

Substraction5			
n	time(ms)		
32	1098		
34	3291		
36	9873		
38	29285		

N1=32 → t1 = 1098
N2=80 → t2 = ?
t2 =
$$3^{N2/2}/3^{N1/2}$$
 * t1 = $3^{N2/2-N1/2}$ * t1 = 3^{24} * t1 = 310.107.631.056.138ms* 1s/1000ms
* 1h/3600s * 1day/24h * 1year/365.25days = 9.826,72years

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Activity 2. [D&C BY DIVISION]

Division1:

$$a = 1, b = 3, k = 1$$

Complexity:

$$O(n^k) \Rightarrow O(n)$$

The results obtained fits with the theoretical complexity because as the size increases by a factor of 2, same happens with the time, so the time increases linearly.

Division2:

$$a = 2, b = 2, k = 1$$

Complexity:

$$O(n^{k*}log(n)) \Rightarrow O(n*log(n))$$

According to the theoretical complexity, the results obtained are the ones expected. The size increments by a factor of 2 but the time is increased by a bit more than 2 which reflects that multiplication by log(n) of the theoretical calculus.

Division3:

$$a = 2, b = 2, k = 0$$

Complexity:

$$O(n^{\log_b a}) => O(n)$$

We are in the same case that with Division1 and the results obtained also fits the constraints required (has a linear time increment according to the size increment).

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Division4:

Division4			
n	time(ms)		
1000	7,9		
2000	29,2		
4000	116		
8000	465,2		
16000	1820,5		
32000	7293		
64000	28918		
128000	Oot		

Division5:

Division5			
n time(ms)			
1000	23,4		
2000	92		
4000	369,5		
8000	1502,5		
16000	5997		
32000	23637		
64000	Oot		

Activity 3. [TWO BASIC PRINCIPLES]

Fibonacci1:

Complexity:

O(n)

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Fibonacci2:

Complexity:

O(n)

Fibonacci3:

a = 1, b = 1, k = 0

Complexity:

 $O(n^{k+1}) \implies O(n)$

Fibonacci4:

a = 2, b = 2 or 1, k = 0

Complexity:

 $O(a^{n/b}) => < O(2^n)$

n	timeF1	timeF2	timeF3	timeF4
10	0,000088	0,000114	0,000191	0,00243
11	0,000092	0,000119	0,000222	0,00381
12	0,000093	0,000125	0,000236	0,00614
13	0,000097	0,000129	0,000253	0,00986
14	0,000101	0,000134	0,000268	0,01592
15	0,000105	0,000141	0,000284	0,02575
16	0,000109	0,000148	0,000312	0,04192
17	0,000112	0,000155	0,000318	0,06852
18	0,000116	0,000161	0,000335	0,105
19	0,000124	0,000167	0,000348	0,171
20	0,000125	0,000175	0,000364	0,276
21	0,000131	0,000183	0,000381	0,444
22	0,000134	0,000191	0,000397	0,745
23	0,00014	0,000201	0,000412	1,188
24	0,00014	0,000209	0,000434	1,909
25	0,000147	0,000213	0,000449	3,071

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As can be seen all the times fits with its theoretical complexity. As can be seen the first method, even if they have the same complexity, it's a bit better due to the simplicity of its implementation.

VectorSum1:

Complexity:

O(n)

VectorSum2:

a = 1, b = 1, k = 0

Complexity:

 $O(n^{k+1}) => O(n)$

VectorSum3:

a = 2, b = 2, k = 0

Complexity:

 $O(n^{\log_b a}) => O(n)$

n	timeSum1	timeSum2	timeSum3
3	0,00005	0,000074	0,000085
6	0,000067	0,000124	0,000163
12	0,000093	0,000235	0,000345
24	0,000141	0,000428	0,000709
48	0,000237	0,000841	0,001456
96	0,00043	0,001624	0,002876
192	0,000793	0,00318	0,00581
384	0,001497	0,006361	0,011583
768	0,002956	0,01274	0,02324
1536	0,00589	0,02513	0,0465
3072	0,01169	0,04912	0,09336

Here happens same as in the previous example with the Fibonacci series.

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Activity 4. [PETANQUE CHAMPIONSHIP]

Algorithm:

a = 2, b = 2, k = 1

Complexity:

 $O(n^k log n) \Rightarrow O(n log n)$

n	tCalendar
2	0,00067
4	0,00305
8	0,0118
16	0,0346
32	0,1435
64	0,54
128	2,22
256	8,9
512	37,04

The results obtained does not match the theoretical complexity so I think I have made an error in the calculation. The results corresponds with a quadratic complexity.