

Stencil Pattern

Parallel Computing CIS 410/510

Department of Computer and Information Science



Table of Contents

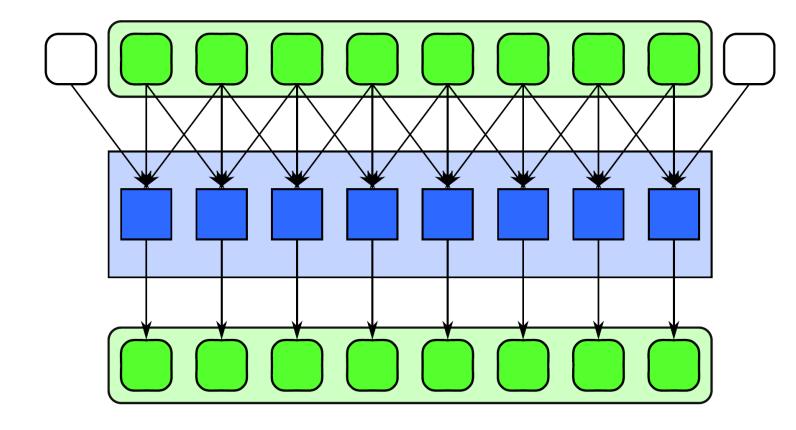
- □ What is the stencil pattern?
- □ Implementing stencil with shift
- □ Stencil and cache optimizations
- Stencil and communication optimizations
- □ Recurrence

Table of Contents

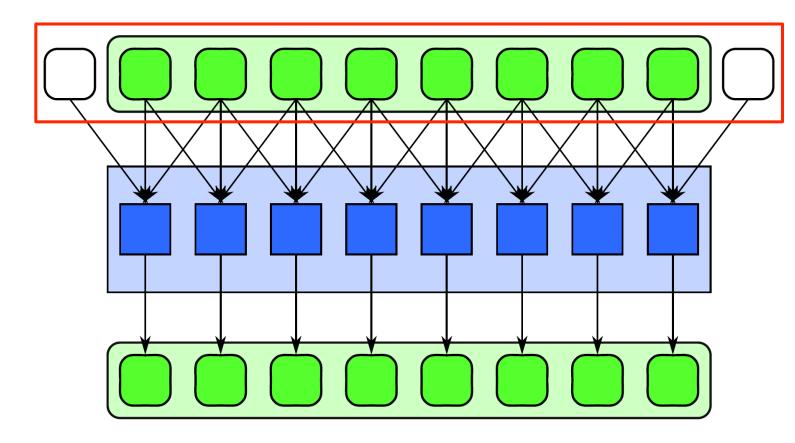
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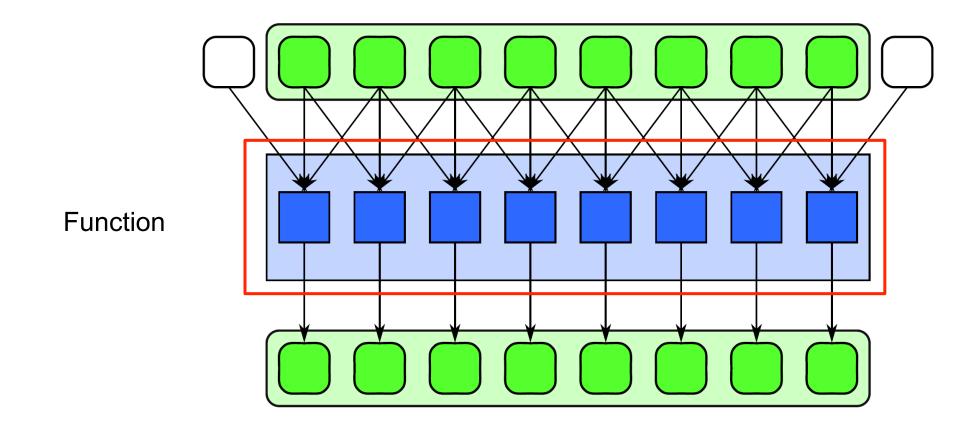
□ **Stencil**: A map where each output depends on a "neighborhood" of inputs. These inputs are a set of fixed offsets relative to the output position.

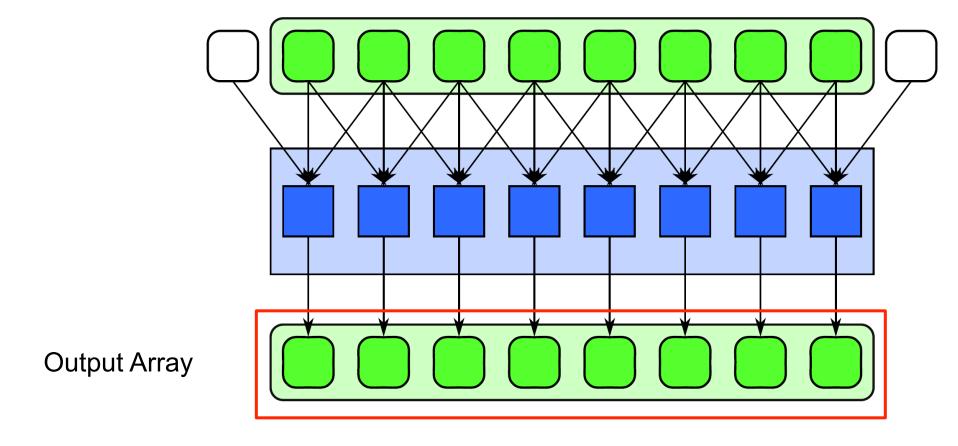
□ A stencil output is a function of a "neighborhood" of elements in an input collection



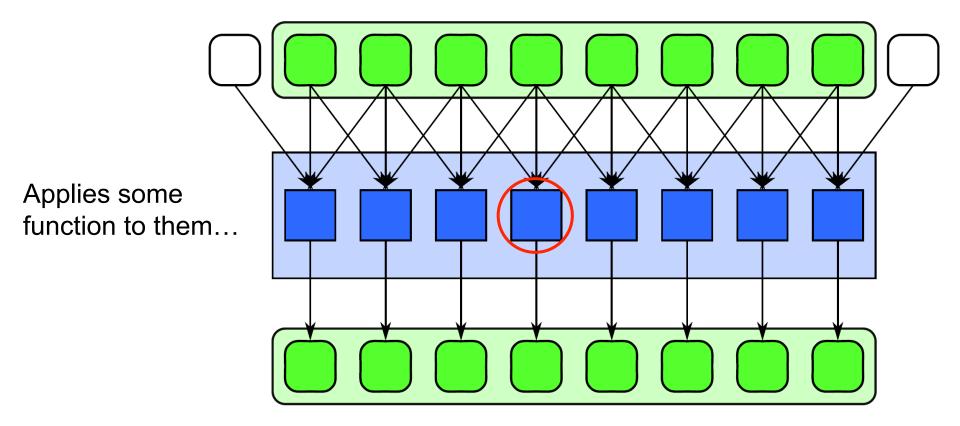
Input array

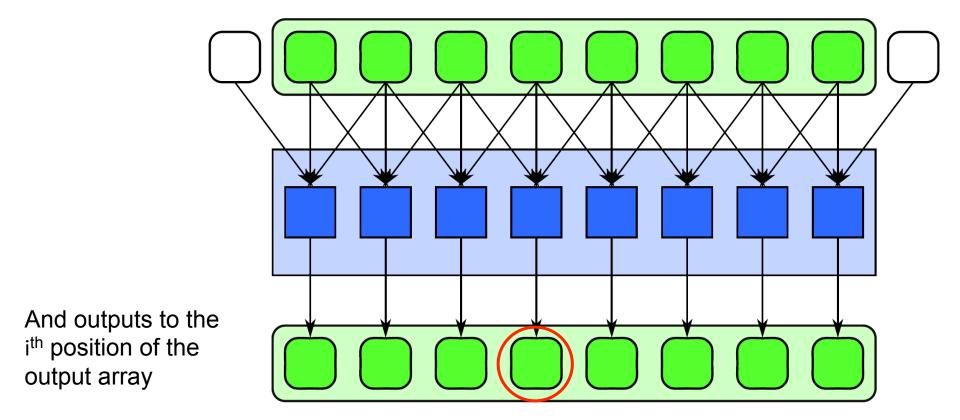






i+1This stencil takes the i-1, i, i+1 elements...

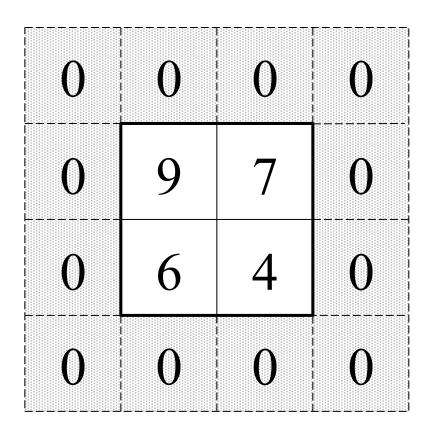




- □ Stencils can operate on one dimensional and multidimensional data
- □ Stencil neighborhoods can range from compact to sparse, square to cube, and anything else!

Practice!

Here is our array, A



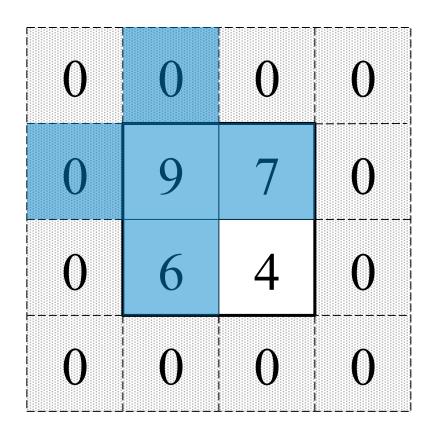
Here is our array, A

• Apply a stencil operation to the inner square of the form:

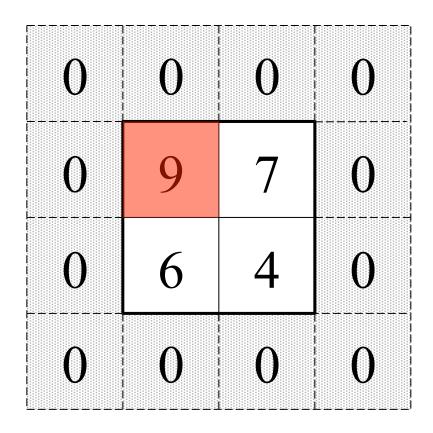
$$(i,j) = avg((i,j), (i-1,j), (i+1,j), (i,j-1), (i,j+1))$$

0	0	0	0
0	9	7	0
0	6	4	0
0	0	0	0

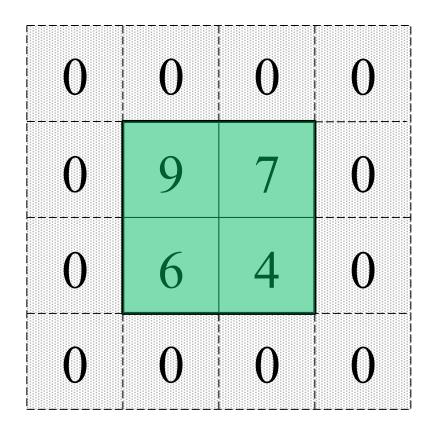
1) Average all blue squares



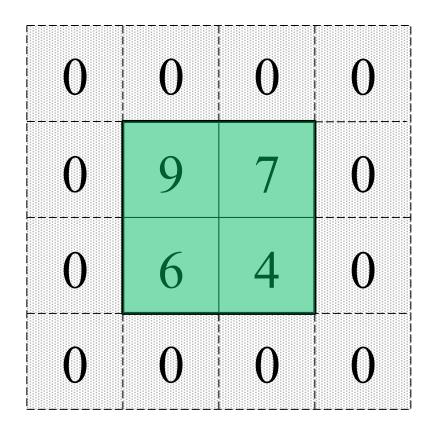
- 1) Average all blue squares
- 2) Store result in red square



- 1) Average all blue squares
- 2) Store result in red square
- 3) Repeat 1 and 2 for all green squares



- 1) Average all blue squares
- 2) Store result in red square
- 3) Repeat 1 and 2 for all green squares



Output!

• How would we parallelize this?

4.4	4	
3.8	3.4	

Serial Stencil Example (part 1)

```
template<
       int NumOff, // number of offsets
       typename In, // type of input locations
       typename Out, // type of output locations
       typename F // type of function/functor
    void stencil(
       int n, // number of elements in data collection
8
       const In a[], //input data collection (n elements)
       Out r[], // output data collection (n elements)
10
    In b, // boundary value
11
       F func, // function/functor from neighborhood inputs to output
12
       const int offsets[] // offsets (NumOffsets elements)
13
14
```

Serial Stencil Example (part 2)

```
// array to hold neighbors
15
       In neighborhood[NumOff];
16
       // loop over all output locations
17
       for (int i = 0; i < n; ++i) {
18
           // loop over all offsets and gather neighborhood
19
           for (int j = 0; j < NumOff; ++j) {
20
               // get index of jth input location
21
               int k = i + offsets[j];
22
               if (0 \le k \& \& k \le n) {
23
                   // read input location
24
                   neighborhood[j] = a[k];
25
               } else {
26
                   // handle boundary case
27
                   neighborhood[j] = b;
28
29
30
           // compute output value from input neighborhood
31
           r[i] = func(neighborhood);
32
33
34
```

Parallel Stencil Example

Table of Contents

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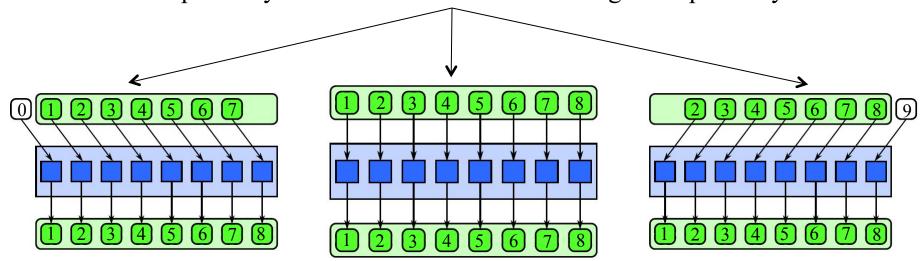
Implementing Stencil with Shift

□ One possible implementation of the stencil pattern includes shifting the input data

□ For each offset in the stencil, we gather a new input vector by **shifting** the original input by the offset amount

Implementing Stencil with Shift

All input arrays are derived from the same original input array



Implementing Stencil with Shift

□ This implementation is only beneficial for one dimensional stencils or the memory-contiguous dimension of a multidimensional stencil

□ Memory traffic to external memory is not reduced with shifts

□ But, shifts allow vectorization of the data reads, which may reduce the total number of instructions

Table of Contents

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- □ Assuming 2D array where rows are contiguous in memory...
 - Horizontally related data will tend to belong to the same cache line
 - Vertical offset accesses will most likely result in cache misses

- □ Assigning rows to cores:
 - Maximizes horizontal data locality
 - Assuming vertical offsets in stencil, this will create redundant reads of adjacent rows from each core
- □ Assigning columns to cores:
 - Redundantly read data from same cache line
 - o Create false sharing as cores write to same cache line

- □ Assigning "strips" to each core can be a better solution
- □ **Strip-mining**: an optimization in a stencil computation that groups elements in a way that avoids redundant memory accesses and aligns memory accesses with cache lines

- □ A strip's size is a multiple of a cache line in width, and the height of the 2D array
- □ Strip widths are in increments of the cache line size so as to avoid false sharing and redundant reads
- □ Each strip is processed serially from top to bottom within each core

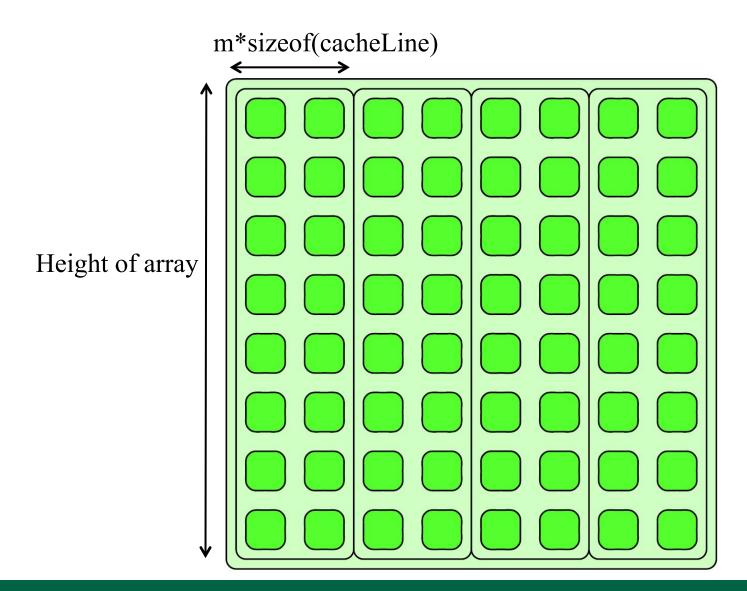
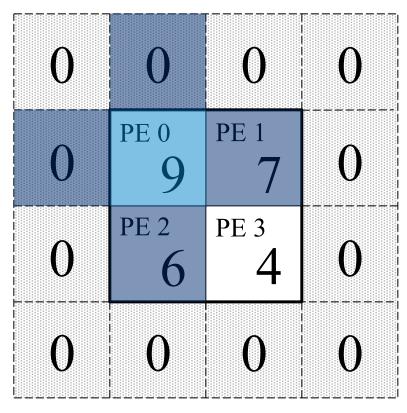


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Stencil and Communication Optimizations

- □ When data is distributed, overlapping regions (or "ghost cells") must be explicitly communicated between nodes between loop iterations
- Darker cells are PE 0's ghost cells
- After first iteration of stencil computation
 - PE 1 & PE 2 must send their stencil results to PE 0
 - PE 0 can perform another iteration of stencil



Stencil and Communication Optimizations

- □ Generally better to replicate ghost cells in each local memory and swap after each iteration than to share memory
 - Fine-grained sharing can lead to increased communication cost

Stencil and Communication Optimizations

- □ Halo: set of all ghost cells
- □ Halo must contain all neighbors needed for one iteration
- □ Larger halo (deep halo)
 - Less communications
 - More redundant local computation

□ Latency Hiding: Compute interior of stencil while waiting for ghost cell updates

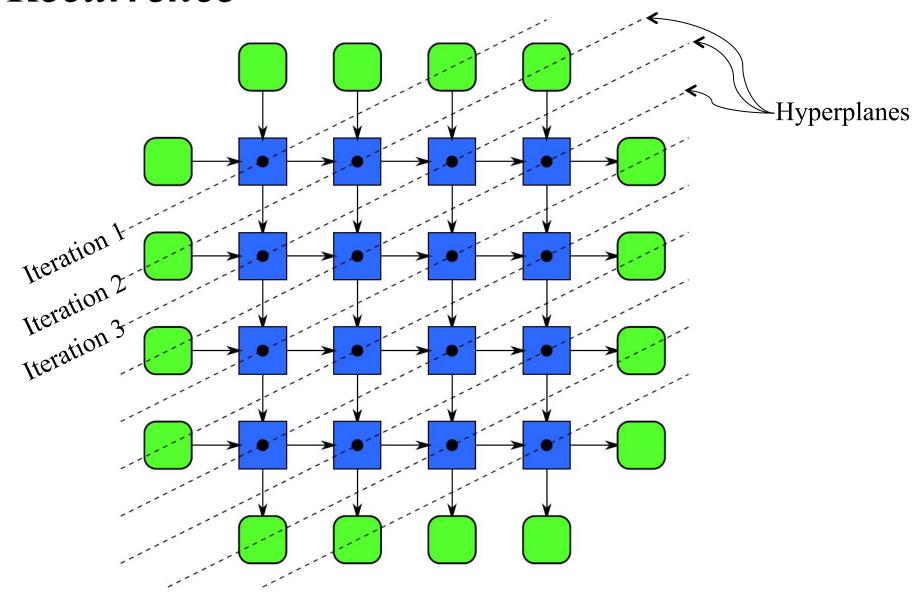
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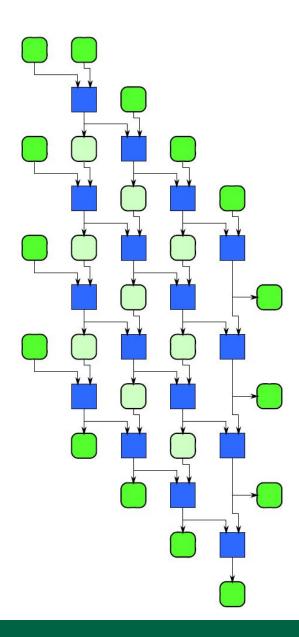
□ What if we have several nested loops with data dependencies between them when doing a stencil computation?

```
void my_recurrence(
       size_t v, // number of elements vertically
       size_t h, // number of elements horizontally
       const float a[v][h], //input 2D array
       float b[v][h] // output 2D array (boundaries already initialized )
5
       for (int i=1; i < v; ++i)
          for (int j=1; j < h; ++j)
              b[i][j] = f(b[i-1][j], b[i][j-1], a[i][j]);
9
10
                           Data dependencies between loops
```

- □ This can still be parallelized!
- □ Trick: find a plane that cuts through grid of intermediate results
 - Previously computed values on one side of plane
 - Values to still be computed on other side of plane
 - Computation proceeds perpendicular to plane through time (this is known as a sweep)
- □ This plane is called a *separating hyperplane*



- Same grid of intermediate results
- Each level corresponds to a loop iteration
- Computation proceeds downward



Example Implementation

Conclusion

- □ Examined the stencil and recurrence pattern
 - Both have a regular pattern of communication and data access
- □ In both patterns we can convert a set of offset memory accesses to shifts
- □ Stencils can use strip-mining to optimize cache use
- □ Ghost cells should be considered when stencil data is distributed across different memory spaces