

Lecture 9

- Memory Protection Unit
- Memory Management Unit



ARM: MPU, RISC-V: PMP





for this section

Joseph Yiu, The Definitive Guide to ARM® Cortex®-M0 and Cortex-M0+ Processors, 2nd Edition

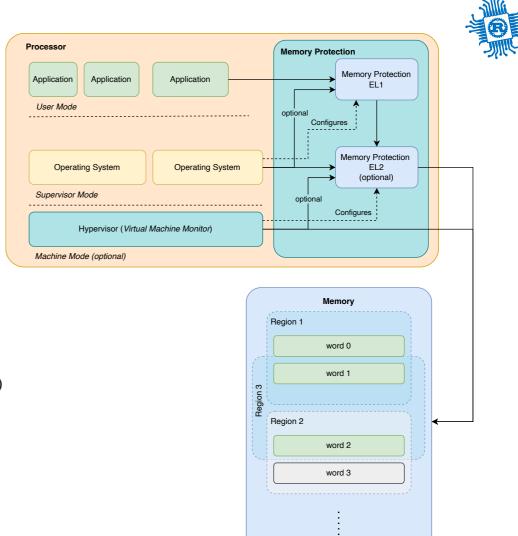
■ Chapter 12 - *Memory Protection Unit*

memory access defined region by region

- restricts access to physical memory
- uses physical addresses

The processor works in three modes:

- machine mode (optional) used at boot, allows access to everything
- supervisor mode restricts access to some registers and accesses memory through Memory Protection EL2 (if machine mode exists)
- user mode allows only ALU and memory access through Memory Protection



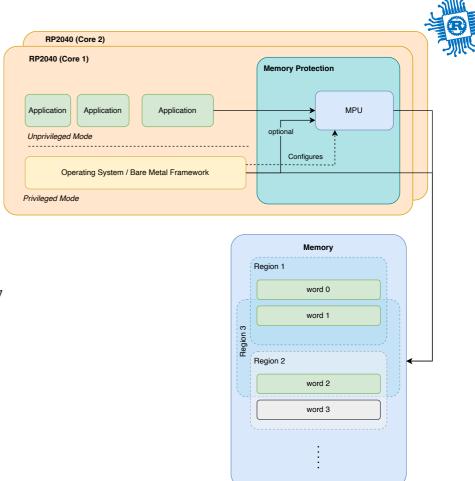
MPU for RP2040

Cortex-M0+ works in three modes

- handler mode no restrictions used while executing ISRs and Exception Handlers
- thread mode
 - privileged no restrictions usually used for the operating system
 - unprivileged mode allows only ALU and memory access through Memory Protection - used for applications

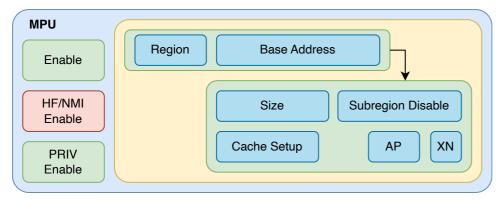
MPU allows 8 regions

- each region has up to 8 subregions
- permissions R W X



Memory Protection Unit

Cortex-M MPU

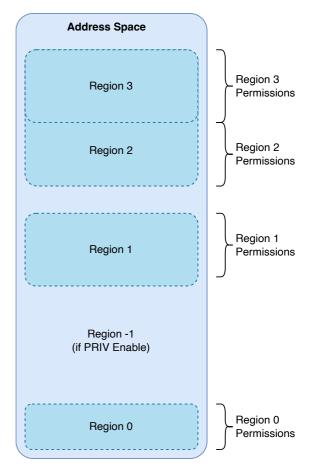


- allows the definition of memory regions
- regions can overlap, highest region number takes priority
- regions have access permissions (similar to rwx)

$$region_size = min(256, 2^{size})$$

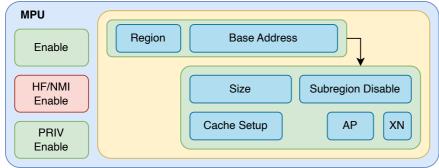
$$base_address = region_size \times N$$





Memory Protection Unit

Access Protection



AP Access Protection

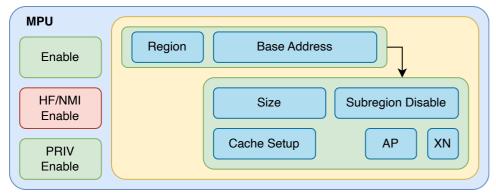
XN eXecute Never

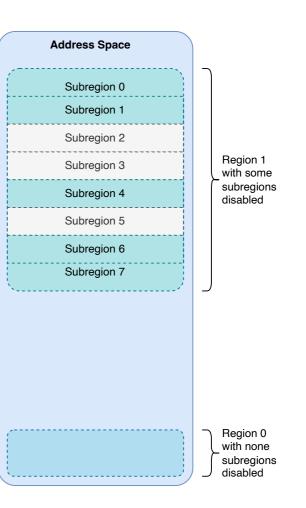
 faults if MCU has to read the next instruction from an XN region

AP	Privileged Mode	Unprivileged Mode
000	No Access	No Access
001	Read/Write	No Access
010	Read/Write	Read only
011	Read/Write	Read/Write
100	Do not use	Do not use
101	Read only	No Access
110	Read only	Read only
111	Read/Write	Read only

Subregions

- each region is divided in 8 subregion
- each bit in Subregion Disable disables a subregion
- a disabled subregion triggers a fault if accessed



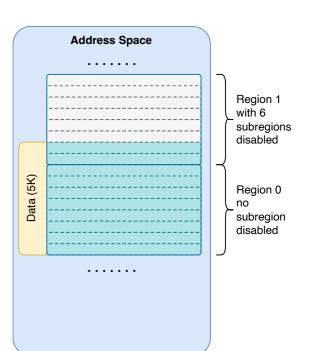


Subregions' Usage

improve granularity

$$egin{aligned} region_size &= min(256, 2^{size}) \ base_address &= region_size imes N \ subregion_size &= rac{region_size}{8} \end{aligned}$$

- a 5K region is not allowed (5K is not a power of 2)
- use two 4K regions back to back
- disable 6 of the subregions (subregion is 512B)



Memory Layout

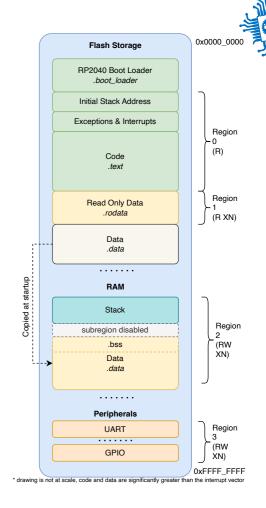
protection

Flash

- **Code** read and execute
- .rodata constants read only
- .data in flash initialized global variables
 - is copied to RAM at startup by the init function
 - should not be accessed after startup

RAM

- **stack** read and write
 - usually protected by unaccessible memory before and after
- .data in RAM global variables read and write
- .bss global variables (not initialized or initialized to 0) read and write





Memory Management

MMU

Bibliography

A BENTALON

for this section

- 1. Andrew Tanenbaum, Modern Operating Systems (4th edition)
 - Chapter 3 *Memory Management*
 - Subchapter 3.3 Virtual Memory
- 2. **Philipp Oppermann**, Writing an OS in Rust
 - Introduction to Paging
 - Paging Implementation

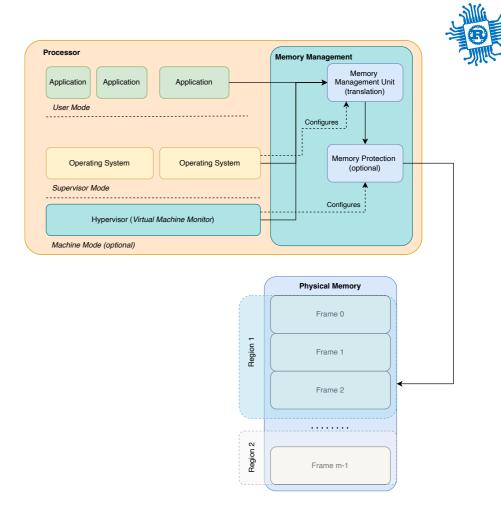
Memory Management

memory access defined page by page

- uses logical addresses
- translates to physical addresses

The processor works in at least two modes:

- supervisor mode
 - restricts access to some registers
 - accesses virtual addresses through Memory
 Protection (*if machine mode exists*)
- user mode
 - allows only ALU and memory load and store
 - accesses memory access through the Memory Management Unit (MMU)



Paging

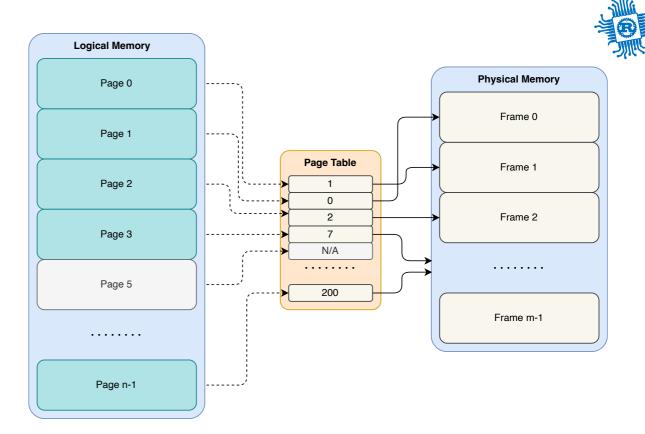
the memory unit is the page

- Physical Memory (RAM) is divided in frames
- Logical Memory is divided in pages
- *page* = *frame* = **4 KB** (usually)

logical addresses are translated to physical addresses using a page table

the **page table** is located in the **physical memory**

 each memory access requires at least memory 2 accesses



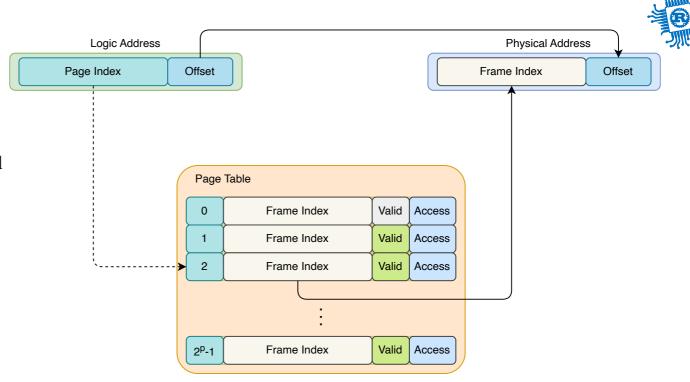
Address Translation

page to frame

the logic address is divided in two parts:

- page index
- offset within the page

the MMU translates every logic address into a physical address using a page table



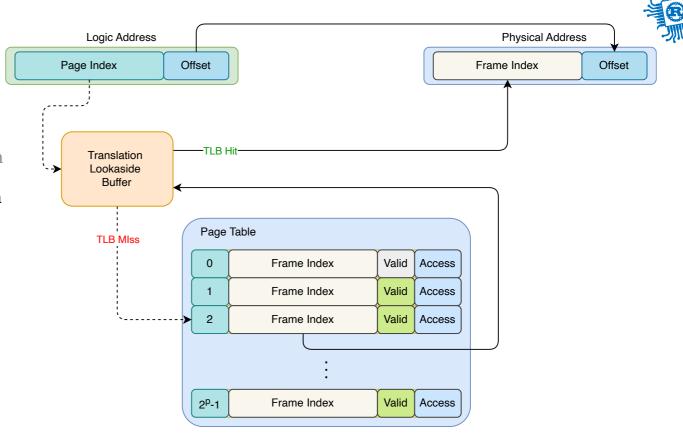
Translation Lookaside Buffer (TLB)

caching address translation

the **page table** is **stored in RAM**

each memory access requires 2 accesses

- read the page table entry to translate the address
- 2. the requested access



Page Directory

caching address translation ->

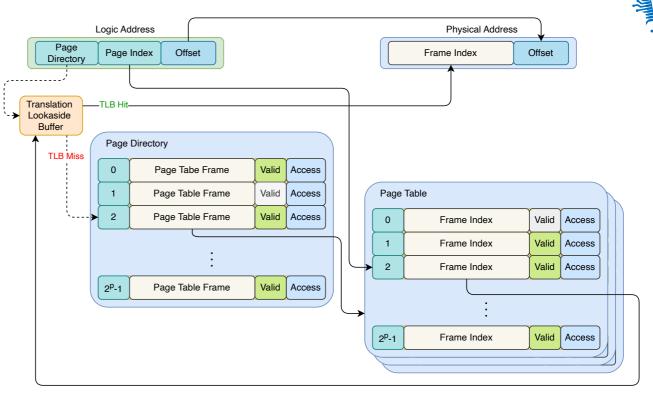
$$size_{table} = rac{size_{ram}}{size_{page}}$$

- each table entry is 4B
- the address space is 4GB (for 32 bits processors)

$$size_{table_32_bits} = rac{2^{32}}{4 imes 2^{10}}$$

$$size_{table 32 \ bits} = 4MB$$

RAM was counted in MB when paging started being used



two levels, page directory and table, usually used for 32 bits systems

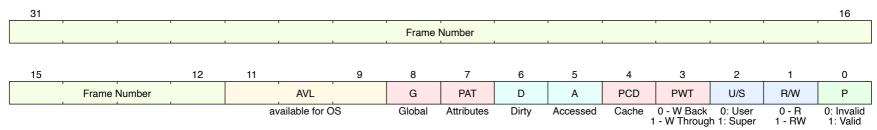
Page Table Entry



for x86 - 32 bits

this is one entry of the page table

- **P** is the page's frame present in RAM?
- **R/W** read only or read write access
- **U/S** can the page be accessed in user mode?
- **D** and **A** has this page been written since the OS has reset these bits?
- **AVL** bits available for the OS to use, ignored by MMU



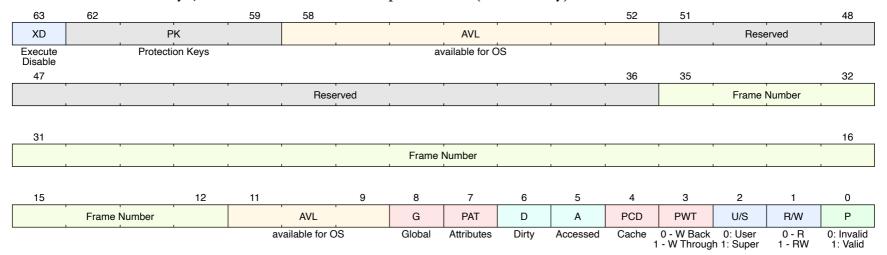
Page Table Entry



for x86 - 32 bits with PAE

this is one entry of the page table using Physical Address Extension (PAE)

- **XD** eXecute Disable (aka *DEP*), if set triggers a fault if an instruction is read from the page
- **PK** Protection Keys, allows user mode to set protection (64 bit only)



Microcontroller (MCU)

Integrated in embedded systems for certain tasks

- low operating frequency (MHz)
- a lot of I/O ports
- controls hardware
- does not require an Operating System
- costs \$0.1 \$25
- uses Memory Protection Unit



Microprocessor (CPU)

General purpose, for PC & workstations

- high operating frequency (GHz)
- limited number of I/O ports
- usually requires an Operating System
- costs \$75 \$500
- uses Memory Management Unit





Conclusion

we talked about

- Memory Protection Unit
- Memory Management Unit

