CSC 411

Computer Organization (Fall 2024) Lecture 7: Floating Point (part 1)

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Fractional binary numbers

- Used to represent fractional numbers
- Binary point
 - separates the integer and fractional parts (similar to the decimal point)
 - bits to the right of binary point represent fractional powers of 2

$$\sum_{k=-i}^{i} b_k 2^k$$

1 1 0 1 . 1 0 1 =

Practice

Convert fractional binary numbers to decimal

$$1.10 =$$
 $11.001 =$
 $100.01 =$
 $11.111 =$
 $0.010 =$

Practice

Convert 349.84375 to a fractional binary number

$$349/2 = 174 \quad R \quad 1$$

$$174/2 = 87 \quad R \quad 0$$

$$87/2 = 43 \quad R \quad 1$$

$$43/2 = 21 \quad R \quad 1$$

$$21/2 = 10 \quad R \quad 1$$

$$10/2 = 5 \quad R \quad 0$$

$$5/2 = 2 \quad R \quad 1$$

$$2/2 = 1 \quad R \quad 0$$

$$1/2 = 0 \quad R \quad 1$$

$$0.84375 * 2 = 1 6875$$

$$0.6875 * 2 = 1 375$$

$$0.375 * 2 = 0 75$$

$$0.75 * 2 = 1 5$$

$$0.5 * 2 = 1 0$$

$$0.0 * 2 = 0$$

101011101.110110

Observations

- Not all decimal fractions have exact binary equivalents
 - . can only represent numbers of the form $\frac{x}{2^k}$
 - irrational and some rational numbers can't be represented
 - e.g., 1/3 = 0.33333... or 0.1
 - both require an infinite series of 0s and 1s
 - large repeating decimals can only be approximated within a certain degree of accuracy
- Note that 0.111111... represents a decimal number just below 1.0

$$\frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{2^i} + \dots = 1.0 - \varepsilon$$

Practice

- Apply the 'multiply by 2' method to convert 0.1 from decimal to fractional binary
 - explain what is happening

- ► Type 0.1 + 0.2 on a python terminal
 - · explain the resulting value

Normalized scientific notation

- Fractional binary numbers in the form $1.xxx... \times 2^k$
 - · analogous to decimal scientific notation
 - always has a single 1 to the left of the binary point
 - · exponent adjusts to maintain the same value
- Examples
 - 101.11 => normalized 1.0111×2^2

Forms the basis of the IEEE 754 standard

• 0.00101 => normalized 1.01×2^{-3}

• 1000.101 => normalized 1.000101×2^3

Practice

- Convert to normalized scientific notation
 - single 1 to the left of the binary point

IEEE 754 standard

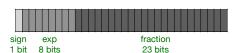
IEEE standard 754

- Defines a common format for representing real numbers in computers
 - · developed in response to divergence of representations
 - first standardized in 1985 and revised in 2008
 - supported by major CPUs (almost universally adopted)
 - provides different <u>precision levels</u> (half, single, double, quadruple) for various needs
- Standardizes the following layout :
 - more exp bits leads to a wider range of numbers
 - · more fraction bits leads to higher precision

sign exp fraction

Common precision levels

- ► Single-precision (32 bits) float
 - · good balance of performance and range



- ▶ Double-precision (64 bits) double
 - · higher precision for demanding calculations



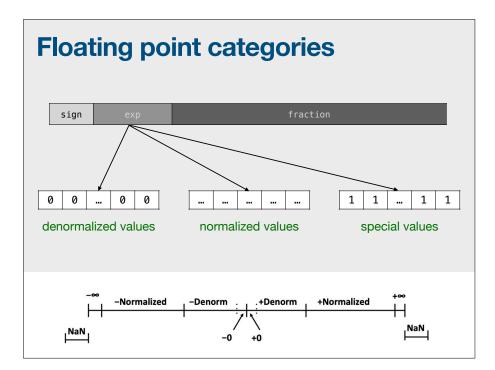
Floating point encoding/decoding

▶ sign bit S

 $(-1)^s M2^E$

- 0 for positive, 1 for negative
- exponent E
 - magnitude of the number, the term 2^E scales the mantissa
- mantissa M
 - a.k.a. significand, it captures the significant digits of the number
 - normally in range [1.0,2.0)

Floating point categories



Floating point categories

- Normalized values
 - represent the majority of values, using the full precision of the significand, and a wide range of magnitudes
 - exp!= 000...000 and exp!= 111...111
- Denormalized values
 - used for <u>very small numbers</u>, extending the range of representable values closer to zero
 - lower precision compared to normalized numbers
 - exp == 000...000
- Special numbers
 - represent special cases like NaN, infinity (positive and negative)
 - exp == 111...111

Normalized values

Normalized values

 $(-1)^s M2^E$

- ► E = exp bias
 - exp is the unsigned integer represented by the bits in the exp field
 - bias is $2^{k-1} 1$, where k is the length in bits of the exp field
 - single precision (bias = $2^{8-1} 1 = 127$)
 - double precision (bias = $2^{11-1} 1 = 1023$)
- M = 1.bb...bb
 - · bb...bb are the bits in the fraction field
 - M is the decimal that corresponds to 1.bb...bb
 - note that M has an implied leading 1

sign ex

fractio

Practice (encoding)

 $(-1)^s M 2^E$

- Assume a float F = 2024.0
 - convert to fractional binary
 - 11111101000
 - make it form 1.xxx
 - 1.1111101000 x 2¹⁰
 - · write M and frac
 - M = 1.1111101000
 - write exp using E = exp bias
 - exp = E + bias = 10 + 127 = 137 = 10001001
 - final binary sequence
 - 0 10001001 111110100000000000000000

0x44FD0000

Practice (encoding)

 $(-1)^{s}M2^{E}$

- ► Assume a float F = 0.5
 - convert to fractional binary
 - 0.
 - make it form 1.xxx
 - 1.0 x 2⁻¹
 - · write M and frac
 - M = 1.0
 - write exp using E = exp bias
 - $\exp = E + \text{bias} = -1 + 127 = 126 = 01111110$
 - final binary sequence
 - 0 01111110 000000000000000000000000

0x3F000000

Practice (encoding)

 $(-1)^{s}M2^{E}$

- Convert the following decimals to hexadecimal representation (single precision)
 - 1.3125
- 3.53125

Practice (decoding)

 $(-1)^s M2^E$

- ► Assume a float F = 0x43CDA000
 - write the binary
 - divide into s, exp, frac
 - **0** 10000111 100110110100000000000000
 - calculate M
 - M = 1.1001101101000000000000 = 1.6064453125
 - calculate E
 - $E = \exp bias = 135 127 = 8$
 - write final number
 - $(-1)^0 * 1.6064453125 * 2^8 = 411.25$

Practice (decoding)

 $(-1)^s M2^E$

- Convert the following values in hexadecimal representation (single precision) to decimal
 - 0x41850000
 - 0x43cdb400