CSC 411

Computer Organization (Spring 2024) Lecture 10: RISC-V load/store

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So far ...

Addition / subtraction

add rd, rs1, rs2 sub rd, rs1, rs2

Add immediate

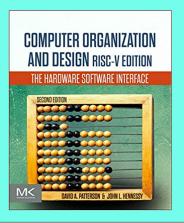
addi rd, rs1, imm

Disclaimer

Some figures and slides are adapted from:

Computer Organization and Design (Patterson and Hennessy)

The Hardware/Software Interface



Data transfer instructions

- RV32I is a load-store architecture
 - only load and store instructions access memory
 - arithmetic instructions only operate on CPU registers
- Memory is byte-addressed
 - · each address identifies an 8-bit byte
 - RISC-V is Little Endian
 - RISC-V and x86 do not require words to be aligned in memory (alignment strongly encouraged)
 - alignment restriction forces words to start at addresses that are multiples of the word size

Loading and storing data

Values can be loaded from memory



Values can be stored in memory



Loading words

The lw instruction loads a 32-bit value from memory into rd

destination register gets memory word from address (base + offset)

Storing words

► The sw instruction stores a 32-bit value from rs2 to memory

memory at address (base + offset) gets the value in source register

Practice

- C code
 - assume all are int variables

$$g = h + A[8];$$

- ► RISC-V code
 - index 8 requires offset of 32 (4 bytes)
 - assume g in x1, h in x2, and base address of A in x3

Practice

- C code
 - assume all are int variables

$$A[7] = h + A[5];$$

- ► RISC-V code
 - assume h in x1 and base address of A in x2
 - can use an additional register x3

Loading and storing bytes

- Instructions for transferring individual bytes from/to memory
- Use same format as lw and sw
 - load byte: lb
 - loads a 8-bit value from memory, then sign-extends to 32-bits before storing in rd
 - load unsigned byte: lbu">lbu
 - loads a 8-bit value from memory but then zero extends to 32-bits before storing in rd
 - store byte: sb
 - stores 8-bit value from the low bits of register rs2 to memory

Practice

- What are the resulting values in x3 and x4 respectively?
 - express the values using 32-bits (in hexadecimal notation)

```
addi x1, x0, 0x80AB

sw x1, 0(x2)

lb x3, 1(x2)

lbu x4, 0(x2)
```

Practice

- Translate the following C code into RISC-V
 - can use an additional register x3

```
// assume x, y are integer pointers
// stored in x1 and x2 respectively
*x = *y;
```

Registers vs memory

- Operating on memory data requires loads and stores
 - · consider memory latency
 - memory is byte-addressable but lw and sw can transfer entire words
- Registers are a fast read/write memory right on the CPU that can hold values
- Compiler must use registers for variables as much as possible
 - · only spill to memory for less frequently used variables
 - register optimization is important!