

## CS 188: Artificial Intelligence

### Constraint Satisfaction Problems I



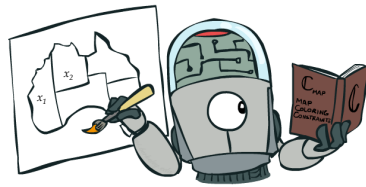
Instructor: Marco Alvarez  
University of Rhode Island

[These slides were created by Dan Klein and Pieter Abbeel for CS188 Intro to AI at UC Berkeley. All CS188 materials are available at <http://ai.berkeley.edu/>.]

1

## Constraint Satisfaction Problems

$N$  variables  
domain  $D$   
constraints



states  
partial assignment

goal test  
complete; satisfies constraints

successor function  
assign an unassigned variable

2

## What is Search For?

- Assumptions about the world: a single agent, deterministic actions, fully observed state, discrete state space

- Planning: sequences of actions
  - The path to the goal is the important thing
  - Paths have various costs, depths
  - Heuristics give problem-specific guidance

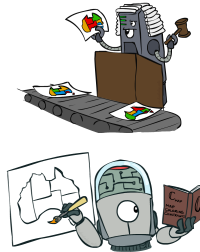
- Identification: assignments to variables
  - The goal itself is important, not the path
  - All paths at the same depth (for some formulations)
  - CSPs are specialized for identification problems



3

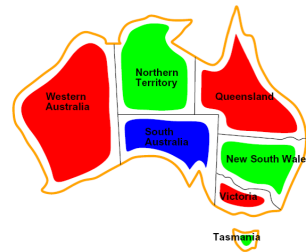
## Constraint Satisfaction Problems

- Standard search problems:
  - State is a "black box": arbitrary data structure
  - Goal test can be any function over states
  - Successor function can also be anything
- Constraint satisfaction problems (CSPs):
  - A special subset of search problems
  - State is defined by variables  $X_i$  with values from a domain  $D$  (sometimes  $D$  depends on  $i$ )
  - Goal test is a set of constraints specifying allowable combinations of values for subsets of variables
- Allows useful general-purpose algorithms with more power than standard search algorithms



4

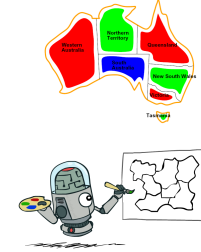
## CSP Examples



5

## Example: Map Coloring

- Variables: WA, NT, Q, NSW, V, SA, T
- Domains:  $D = \{\text{red, green, blue}\}$
- Constraints: adjacent regions must have different colors
  - Implicit:  $WA \neq NT$
  - Explicit:  $(WA, NT) \in \{(\text{red, green}), (\text{red, blue}), \dots\}$
- Solutions are assignments satisfying all constraints, e.g.:
  - $\{WA=\text{red}, NT=\text{green}, Q=\text{red}, NSW=\text{green}, V=\text{red}, SA=\text{blue}, T=\text{green}\}$



6

## Example: N-Queens

- Formulation 1:
  - Variables:  $X_{ij}$
  - Domains:  $\{0, 1\}$
  - Constraints



$$\begin{aligned} \forall i, j, k \quad (X_{ij}, X_{ik}) &\in \{(0, 0), (0, 1), (1, 0)\} \\ \forall i, j, k \quad (X_{ij}, X_{kj}) &\in \{(0, 0), (0, 1), (1, 0)\} \\ \forall i, j, k \quad (X_{ij}, X_{i+k, j+k}) &\in \{(0, 0), (0, 1), (1, 0)\} \\ \forall i, j, k \quad (X_{ij}, X_{i+k, j-k}) &\in \{(0, 0), (0, 1), (1, 0)\} \end{aligned}$$

$$\sum_{i,j} X_{ij} = N$$

7

## Example: N-Queens

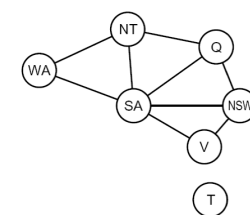
- Formulation 2:
  - Variables:  $Q_k$
  - Domains:  $\{1, 2, 3, \dots, N\}$
  - Constraints:



$$\begin{aligned} \text{Implicit: } \forall i, j \quad &\text{non-threatening}(Q_i, Q_j) \\ \text{Explicit: } (Q_1, Q_2) &\in \{(1, 3), (1, 4), \dots\} \\ &\dots \end{aligned}$$

8

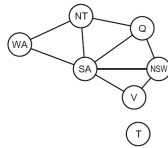
## Constraint Graphs



9

## Constraint Graphs

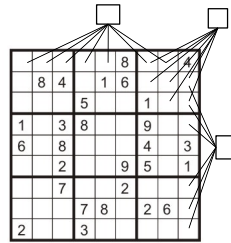
- Binary CSP: each constraint relates (at most) two variables
- Binary constraint graph: nodes are variables, arcs show constraints
- General-purpose CSP algorithms use the graph structure to speed up search. E.g., Tasmania is an independent subproblem!



[Demo: CSP applet (made available by aispac.org) -- n-queens]

10

## Example: Sudoku

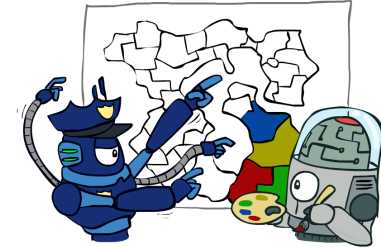


- Variables:
  - Each (open) square
- Domains:
  - {1,2,...,9}
- Constraints:

9-way alldiff for each column  
9-way alldiff for each row  
9-way alldiff for each region  
(or can have a bunch of pairwise inequality constraints)

11

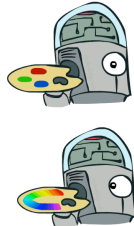
## Varieties of CSPs and Constraints



12

## Varieties of CSPs

- Discrete Variables
  - Finite domains
    - Size  $d$  means  $O(d^n)$  complete assignments
    - E.g., Boolean CSPs, including Boolean satisfiability (NP-complete)
  - Infinite domains (integers, strings, etc.)
    - E.g., job scheduling, variables are start/end times for each job
    - Linear constraints solvable, nonlinear undecidable
- Continuous variables
  - E.g., start/end times for Hubble Telescope observations
  - Linear constraints solvable in polynomial time by LP methods (see cs170 for a bit of this theory)



13

## Varieties of Constraints

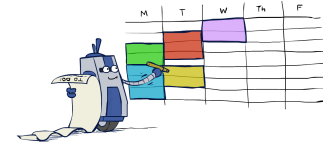
- Varieties of Constraints
  - Unary constraints involve a single variable (equivalent to reducing domains), e.g.:
    - $SA \neq \text{green}$
  - Binary constraints involve pairs of variables, e.g.:
    - $SA \neq WA$
  - Higher-order constraints involve 3 or more variables:
    - e.g., cryptarithmic column constraints
- Preferences (soft constraints):
  - E.g., red is better than green
  - Often representable by a cost for each variable assignment
  - Gives constrained optimization problems
  - (We'll ignore these until we get to Bayes' nets)



14

## Real-World CSPs

- Assignment problems: e.g., who teaches what class
- Timetabling problems: e.g., which class is offered when and where?
- Hardware configuration
- Transportation scheduling
- Factory scheduling
- Circuit layout
- Fault diagnosis
- ... lots more!



- Many real-world problems involve real-valued variables...

15

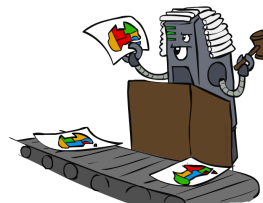
## Solving CSPs



16

## Standard Search Formulation

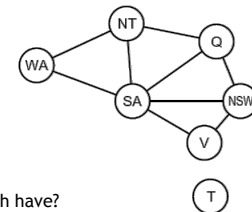
- Standard search formulation of CSPs
- States defined by the values assigned so far (partial assignments)
  - Initial state: the empty assignment, {}
  - Successor function: assign a value to an unassigned variable
  - Goal test: the current assignment is complete and satisfies all constraints
- We'll start with the straightforward, naive approach, then improve it



17

## Search Methods

- What would BFS do?
- What would DFS do?
- What problems does naive search have?



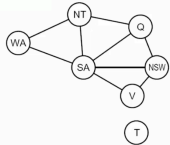
[Demo: coloring -- dfs]

18

## Video of Demo Coloring -- DFS

### Search Methods

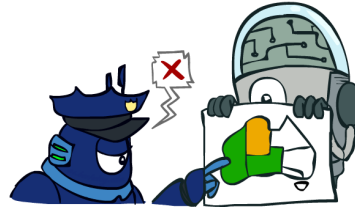
- What would BFS do?
- What would DFS do?



[demo: dfs]

19

## Backtracking Search



20

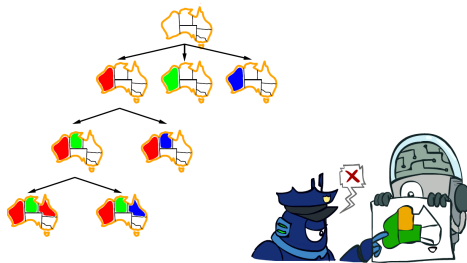
## Backtracking Search

- Backtracking search is the basic uninformed algorithm for solving CSPs
- Idea 1: One variable at a time
  - Variable assignments are commutative, so fix ordering
  - I.e., [WA = red then NT = green] same as [NT = green then WA = red]
  - Only need to consider assignments to a single variable at each step
- Idea 2: Check constraints as you go
  - I.e., consider only values which do not conflict previous assignments
  - Might have to do some computation to check the constraints
  - "Incremental goal test"
- Depth-first search with these two improvements is called *backtracking search* (not the best name)
- Can solve n-queens for  $n = 25$



21

## Backtracking Example



22

## Backtracking Search

```

function BACKTRACKING-SEARCH(csp) returns solution/failure
  return RECURSIVE-BACKTRACKING({}, csp)
function RECURSIVE-BACKTRACKING(assignment, csp) returns soln/failure
  if assignment is complete then return assignment
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
    if value is consistent with assignment given CONSTRAINTS[csp] then
      add {var = value} to assignment
      result ← RECURSIVE-BACKTRACKING(assignment, csp)
      if result ≠ failure then return result
      remove {var = value} from assignment
  return failure
  
```

- Backtracking = DFS + variable-ordering + fail-on-violation

[Demo: coloring -- backtracking]

23