

Taming the Duck: Can Stochastic Programming Help?

1. Formulations

1.1. Decision Variables and Parameters

Let \mathcal{G} represent the set of generators, and \mathcal{T} denote the set of discretized time periods. The set of buses is denoted with \mathcal{B} and the set of lines is denoted with \mathcal{L} . The generators that are located at a particular bus j are represented with the set \mathcal{G}_j . The following variables are defined for each generator $g \in \mathcal{G}$, and period $t \in \mathcal{T}$:

x_{gt} : 1 if generator g is operational at t , 0 otherwise,
 s_{gt} : 1 if generator g becomes online at t , 0 otherwise,
 z_{gt} : 1 if generator g becomes offline at t , 0 otherwise,
 G_{gt} : Production level of generator g at t ,

For each bus $j \in \mathcal{B}$ and period $t \in \mathcal{T}$, we define the following variables:

θ_{jt} : Voltage angle at bus j at t ,
 L_{jt} : Amount of load shed at bus j at t .
 O_{jt} : Amount of over-generation at bus j at t .

The latter two variables are penalized in the objective with the coefficients $L^{penalty}$ and $O^{penalty}$, respectively.

Each generator $g \in \mathcal{G}$ is characterized by the following parameters:

G_g^{\max} : Maximum generation capacity,
 G_g^{\min} : Minimum generation requirement when the generator is online,
 ΔG_g^{\max} : Ramp up limit,
 ΔG_g^{\min} : Ramp down limit,
 UT_g : Minimum required uptime before the generator can become offline,
 DT_g : Minimum required downtime before the generator can become online,
 c_g^{gen} : Generation cost,
 c_g^{start} : Start up cost,
 c_g^{noload} : No load cost.

The problem is defined over a topology that is defined by the following characteristics:

$F_{ij,t}^{\min} / F_{ij,t}^{\max}$: Flow lower / upper limits over line $(i, j) \in \mathcal{L}$ at t ,
 B_{ij} : Susceptance of line $(i, j) \in \mathcal{L}$.

Finally, each bus $j \in \mathcal{B}$ faces a demand D_{jt} at period $t \in \mathcal{T}$.

1.2. Unit Commitment

$$\min \sum_{g \in \mathcal{G}} \sum_{t \in \mathcal{T}} \left(c_g^{\text{gen}} G_{gt} + c_g^{\text{start}} s_{gt} + c_g^{\text{noload}} x_{gt} \right) + \sum_{j \in \mathcal{B}} \sum_{t \in \mathcal{T}} \left(L^{\text{penalty}} L_{jt} + O^{\text{penalty}} O_{jt} \right) \quad (1a)$$

$$\text{s.t.} \quad \text{State Equations} \quad \begin{cases} x_{gt} - x_{gt-1} = s_{gt} - z_{gt}, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \end{cases} \quad (1b)$$

$$\begin{aligned} \text{Minimum} & \\ \text{Up/Downtime} & \\ \text{Restrictions} & \begin{cases} \sum_{j=t-UT_g+1}^{t-1} s_{gt} \leq x_{gt}, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \\ \sum_{j=t-DT_g}^t s_{gt} \leq 1 - x_{gt}, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \end{cases} \end{aligned} \quad (1c)$$

$$\text{Generation Limits} \quad \begin{cases} G_g^{\min} x_{gt} \leq G_{gt} \leq G_g^{\max} x_{gt}, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \end{cases} \quad (1d)$$

$$\text{Ramping Limits} \quad \begin{cases} \Delta G_g^{\min} \leq G_{gt} - G_{gt-1} \leq \Delta G_g^{\max}, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \end{cases} \quad (1e)$$

$$\text{Flow Limits} \quad \begin{cases} F_{ij,t}^{\min} \leq F_{ij,t} \leq F_{ij,t}^{\max}, & \forall (i,j) \in \mathcal{L}, t \in \mathcal{T}, \end{cases} \quad (1f)$$

$$\text{Flow Balance} \quad \begin{cases} \sum_{i \in \mathcal{B}: (i,j) \in \mathcal{L}} F_{ij,t} - \sum_{i \in \mathcal{B}: (j,i) \in \mathcal{L}} F_{ji,t} + \sum_{g \in \mathcal{G}_j} G_{gt} + L_{jt} - O_{jt} = D_{jt}, & j \in \mathcal{B}, t \in \mathcal{T}, \end{cases} \quad (1g)$$

$$\text{Power Flow} \\ \text{Approximation} \quad \begin{cases} F_{ij,t} = B_{ij}(\theta_{it} - \theta_{jt}), & \forall (i,j) \in \mathcal{L}, t \in \mathcal{T}, \end{cases} \quad (1h)$$

$$\begin{aligned} \text{Integrality} & \\ \text{Restrictions \& Variable} & \\ \text{Bounds} & \begin{cases} x_{gt}, s_{gt}, z_{gt} \in \{0, 1\}, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \\ G_{gt} \geq 0, & \forall g \in \mathcal{G}, t \in \mathcal{T}, \\ -\pi \leq \theta_{jt} \leq \pi, L_{jt} \geq 0, O_{jt} \geq 0 & \forall j \in \mathcal{B}, t \in \mathcal{T}, \end{cases} \end{aligned} \quad (1i)$$

In deterministic UC formulations, alternative representations of the above constraints were used to gain computational performance. For brevity, such representations are omitted from this paper.

1.3. Economic Dispatch