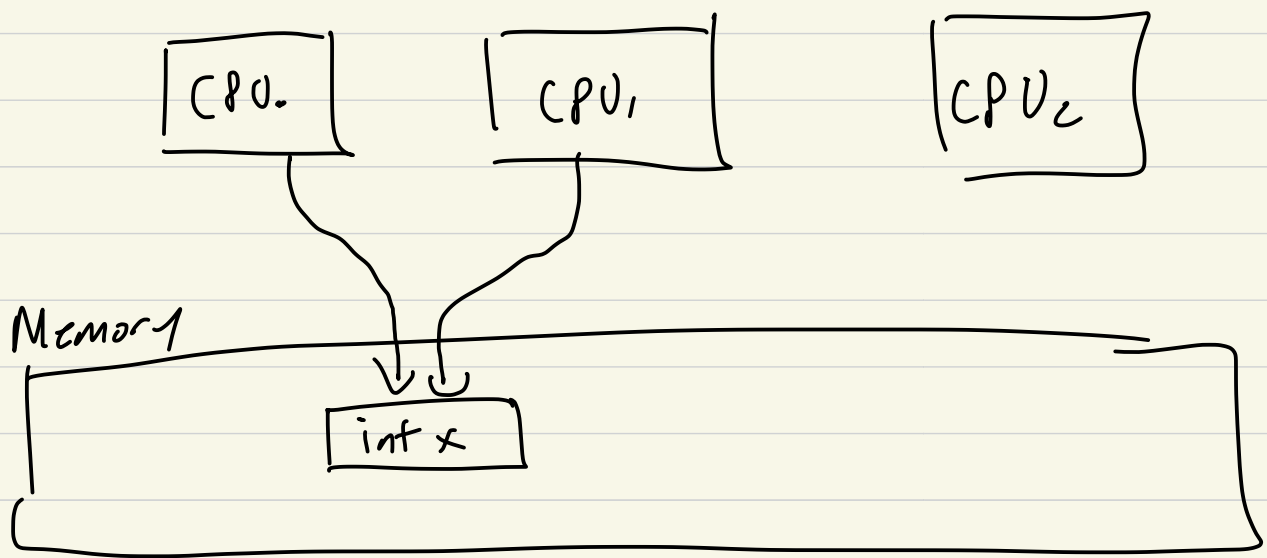


CS 326-01 Concurrency and Locks

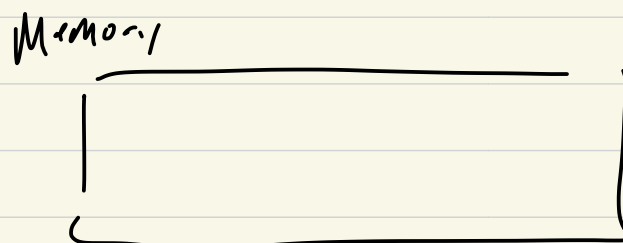
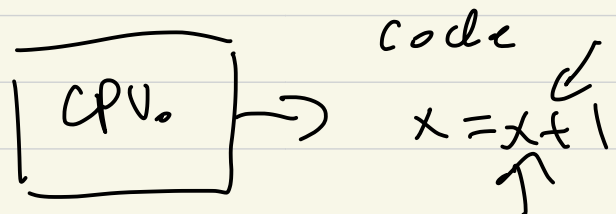
Project 03 \rightarrow Queue/List based scheduling

xv6 Book: Chapters 6 and 7

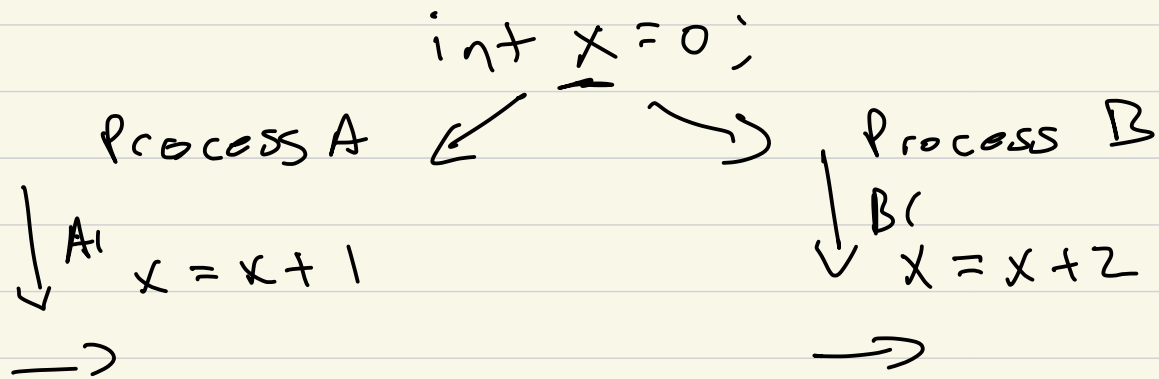
Concurrency



Shared Memory/
architecture



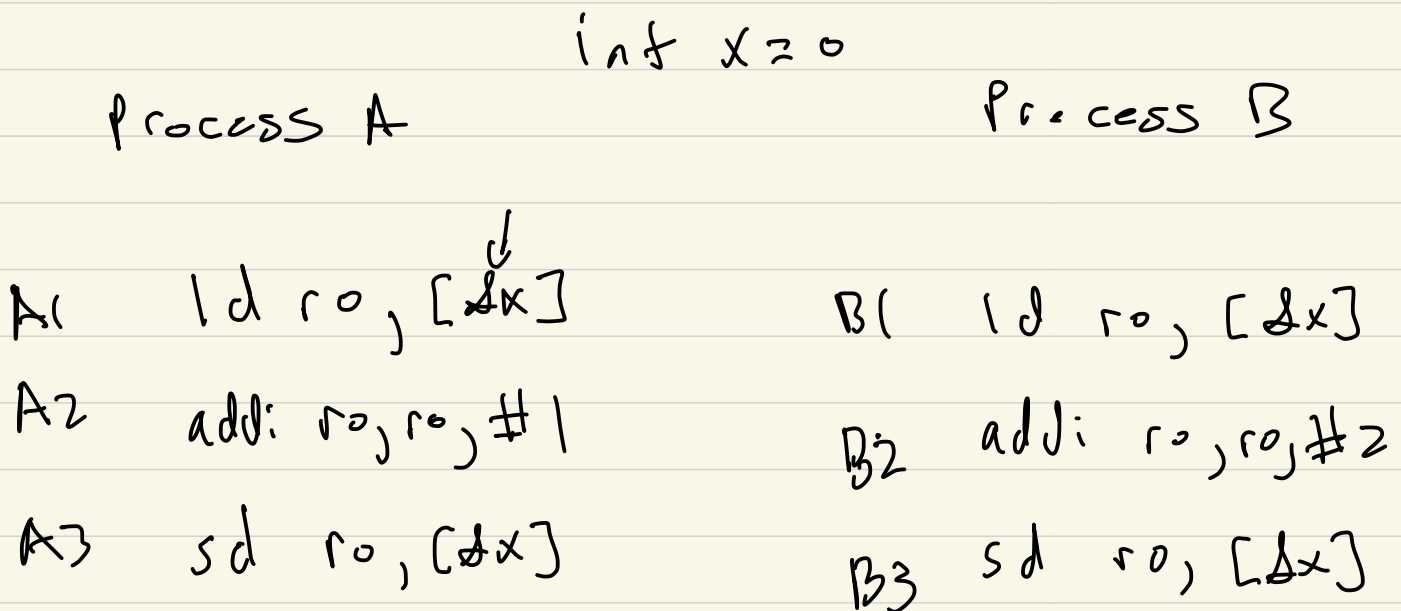
Race conditions



possible values
of x?

#1

A1	x = 1	B1	x = 2
B1	x = 3	A1	x = 3
	↑		



Possible interleavings

A1	$r_0 \leftarrow 0$	B1
A2	$r_0 \leftarrow r_0 + 1$	B2
A3	$x \leftarrow r_0 \quad x = 1$	B3
B1	$r_0 \leftarrow 1$	A1
B2	$r_0 \leftarrow r_0 + 2$	A2
B3	$x \leftarrow r_0 \quad x = 3$	A3

$x = 3$

A1 $r_0 \leftarrow 0$

B1 $r_0 \leftarrow 0$

A2 $r_0 \leftarrow r_0 + 1 \quad r_0 = 1$

A3 $x \leftarrow r_0 \quad x = 1$

B2 $r_0 \leftarrow r_0 + 1 \quad r_0 = 2$

B3 $x \leftarrow r_0 \quad x = 2$

race
condition

A1 $r_0 \leftarrow 0$

B1 $r_0 \leftarrow 0$

B2 $r_0 \leftarrow r_0 + 2 \quad r_0 = 2$

B3 $x \leftarrow r_0 \quad x = 2$

A2 $r_0 \leftarrow r_0 + 1 \quad r_0 = 1$

A3 $x \leftarrow r_0 \quad x = 1$

Locks

```
int x = 0;  
lock x_lock = 0;
```

Process A

```
→ acquire(&x_lock)  
↓  
[ x = x + 1; ]  
↑  
release(&x_lock)
```

atomically

critical section

Process B

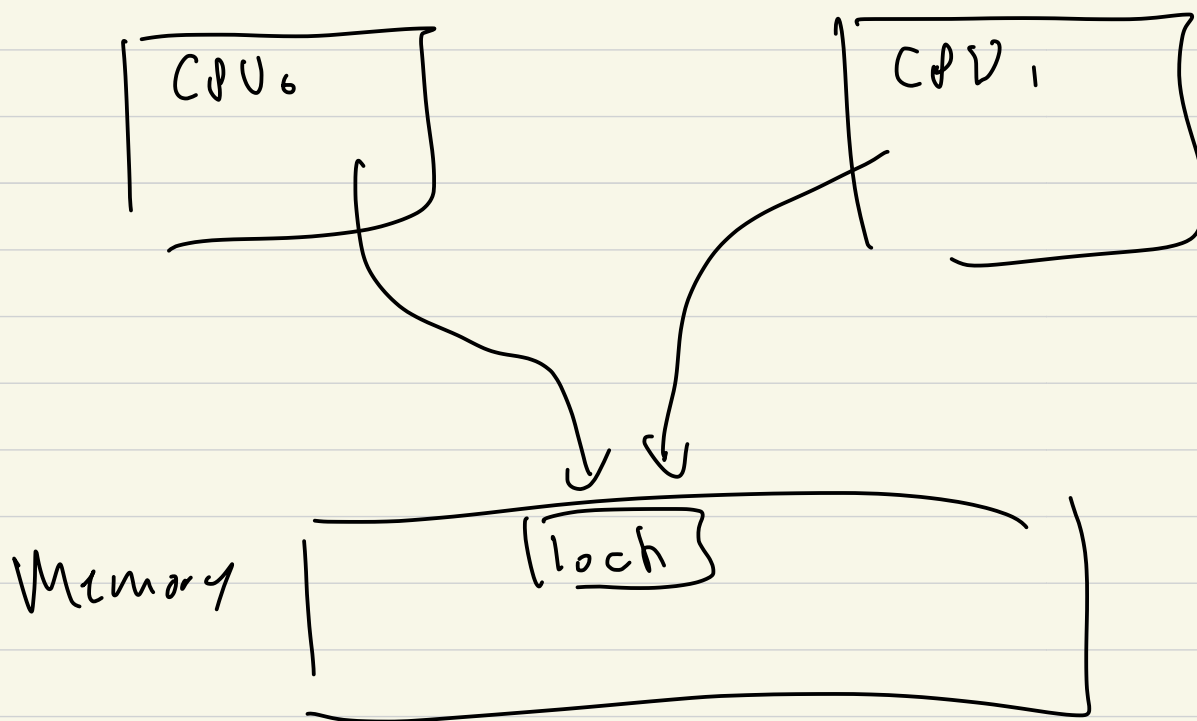
```
→ acquire(&x_lock)  
  
x = x + 2;  
release(&x_lock)
```

Spinlocks

acquire(int *lock) {

{ while (*lock == 1);
*lock = 1;
}

X
does
not
work



Atomic swap

$$\text{while } (\text{atomic_swap}(\text{lock}, 1) == 1);$$

Diagram illustrating the atomic swap function call: `atomic_swap(address, value)`. The parameter `lock` is identified as the *address*, and the parameter `1` is identified as the *value*.