

# Inter-networking

COS 460 & 540



# CAPTAIN JAMES T. KIRK

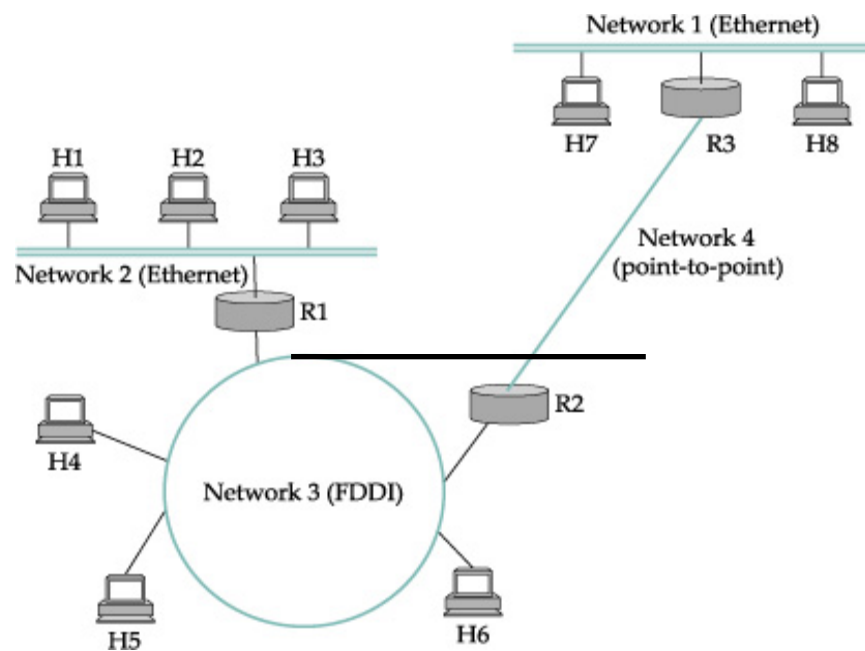
I'M SORRY, I CAN'T HEAR YOU OVER THE  
SOUND OF HOW AWESOME I AM.

# Problem

LAN's are great... but...

We want to **connect** them together

...across the world



# Inter-networking

- Internet Protocol (IP)
- Routing
- The Internet
- Multicast\*
- Multi-protocol Label Switching\*

# Internet Protocol (IP)

- What is an Internetwork
- Service Model
  - Datagrams, Packet Format, ...
- Addressing
- Datagram Forwarding
- ARP, DHCP, ICMP, ...

# What?

- “internet” vs “Internet”
- “network” vs “subnetwork”
- physical vs logical networks

# Bridge, Router, Switch

- repeater & hub - physical
- bridge - link
- switch - network
- router - internet
- gateway - application

# Alternatives?

- Novell's IPX



# Internet Protocol (IP)

✓ What is an Internetwork

- Service Model
  - Datagrams, Packet Format, ...
- Addressing
- Datagram Forwarding
- ARP, DHCP, ICMP, ...

# Service Model

- Datagram Delivery
- Packet Format
- Fragmentation and Reassembly

# Datagram Delivery

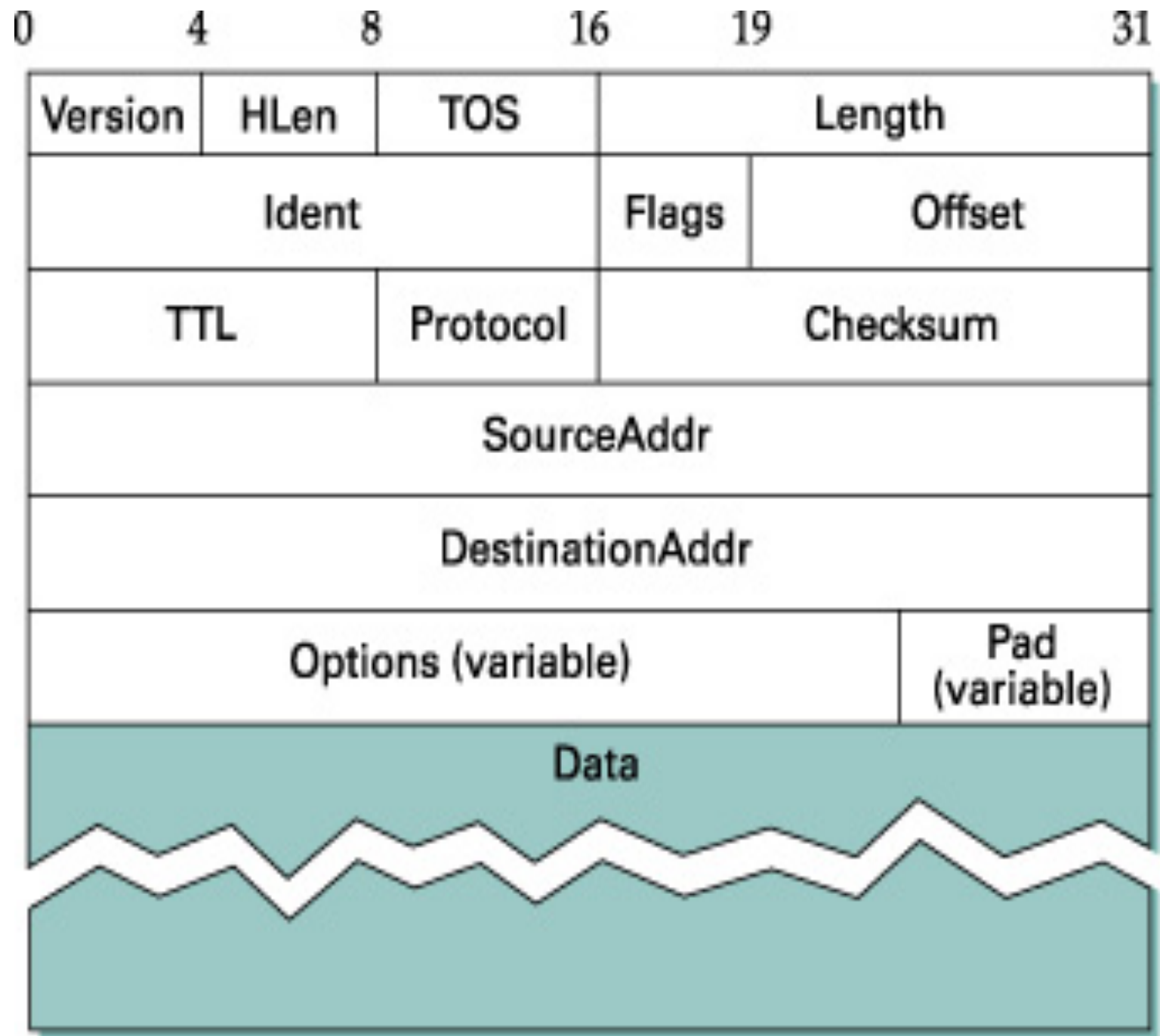
Underlying network is..

- reliable or unreliable
- connection-oriented or connectionless
- small or large packet/frame sized
- physical, logical, wireless, ...

# Datagram Delivery

- IP **Datagram** is basis of protocol
- Provides a “**best-effort**” or unreliable service
- May be **out of order**
- **Connectionless**

# Packet Format



# Fragmentation

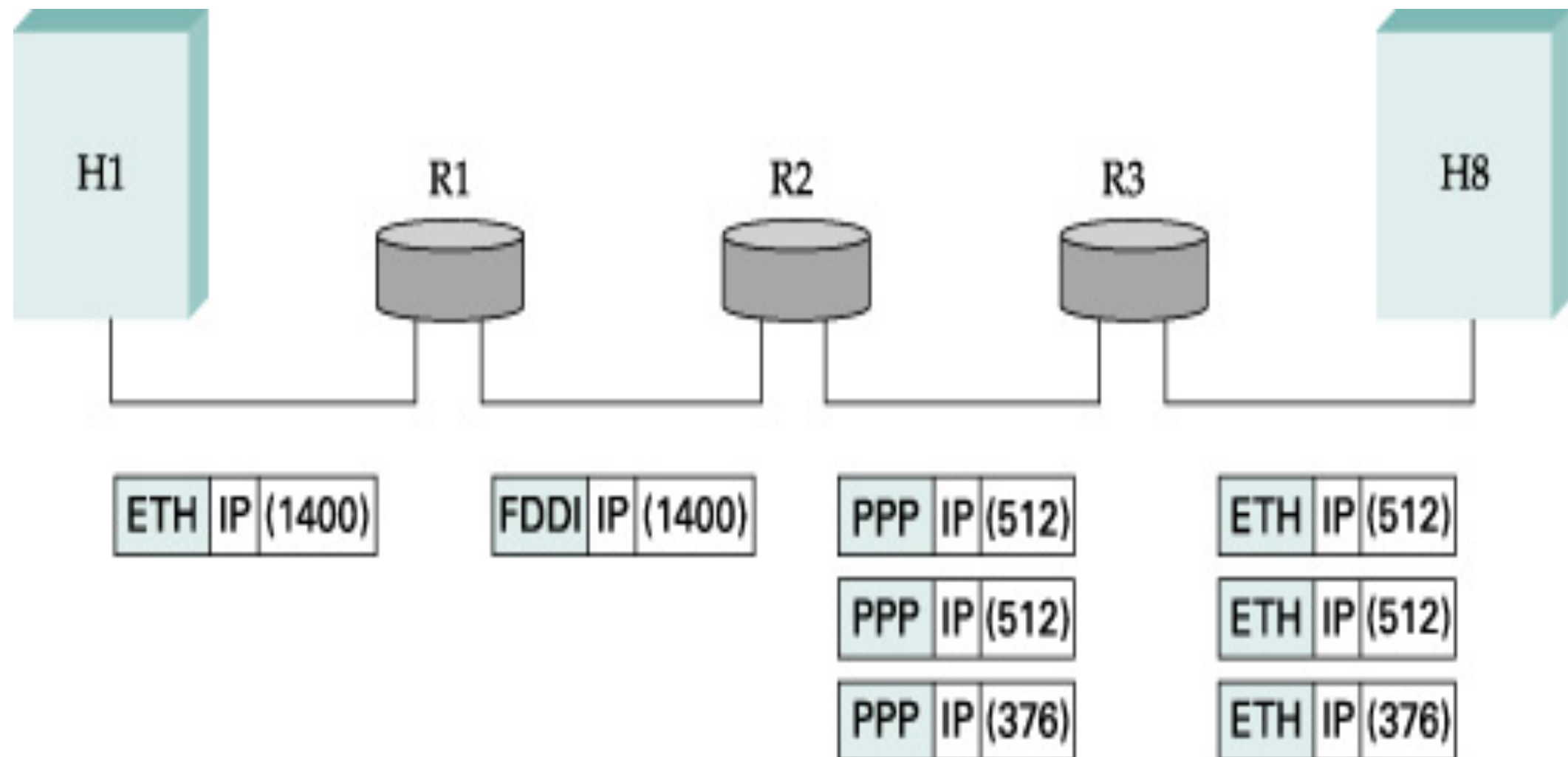
## **Problem:**

- IP datagrams can be 64 kB
- The underlying network may have 512 Byte MTU

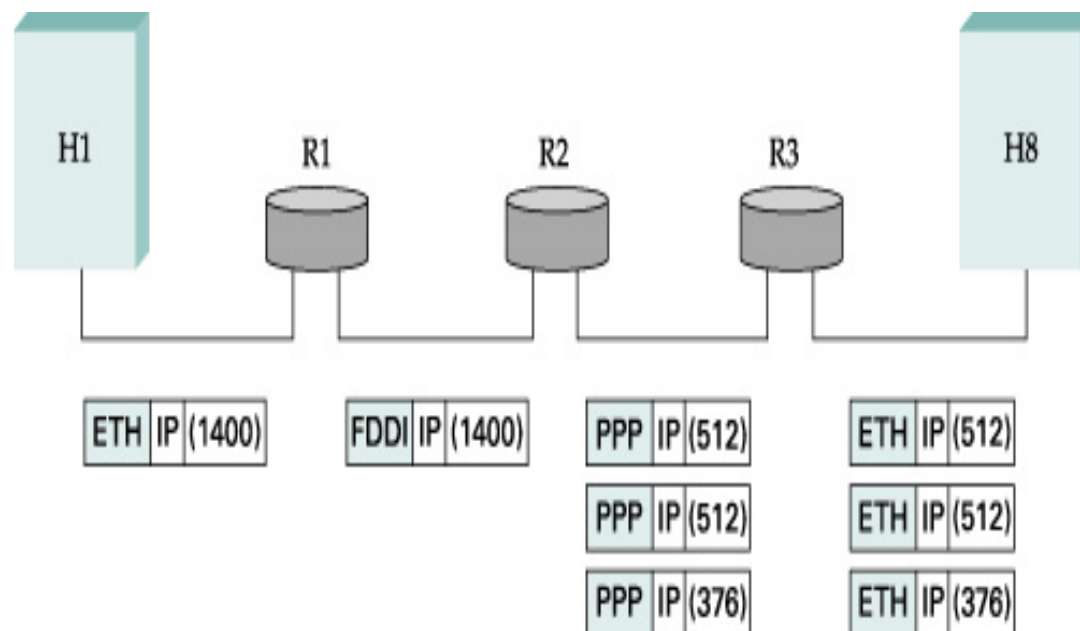
## **Solution:**

- Each fragment is an IP datagram
- Each IP datagram is rebuilt for each network

# Fragmentation



# Fragmentation



(a)

Start of header				
Ident = x			0	Offset = 0
Rest of header				
1400 data bytes				

(b)

Start of header				
Ident = x			1	Offset = 0
Rest of header				
512 data bytes				

Start of header				
Ident = x			1	Offset = 64
Rest of header				
512 data bytes				

Start of header				
Ident = x			0	Offset = 128
Rest of header				
376 data bytes				



# Internet Protocol (IP)

- ✓ What is an Internetwork
- ✓ Service Model
  - ✓ Datagrams, Packet Format, ...
- Addressing
- Datagram Forwarding
- ARP, DHCP, ICMP, ...

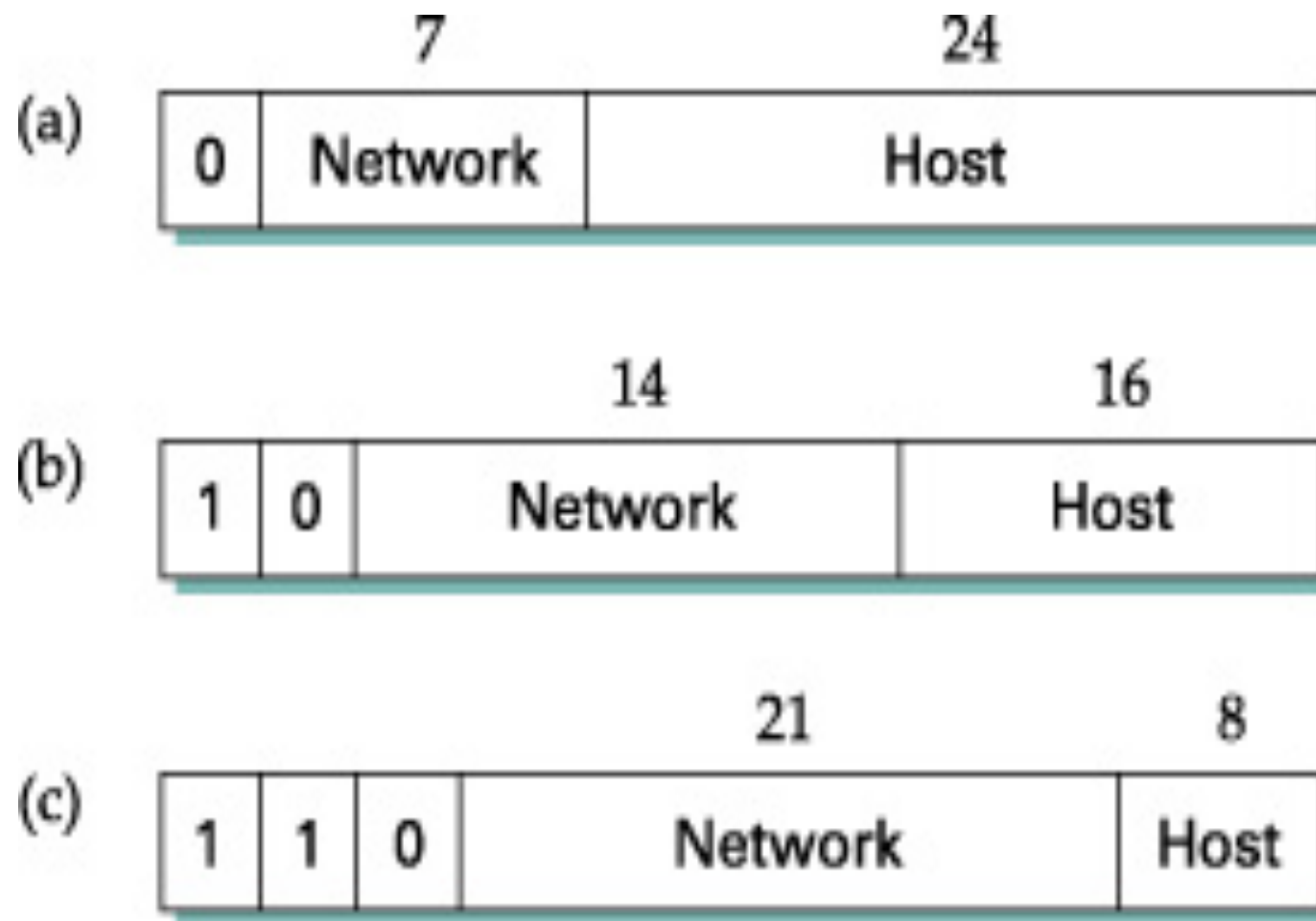
# IP Addressing

- Need to identify any host on the network
- Globally **unique**
- **Hierarchal** (not flat like Ethernet)
  - for easier routing
  - represents a network of networks

# IP Addressing

- 32 bit addresses
- “Dotted Decimal” format
  - 130.111.135.26
- Contains two parts
  - **Network**
  - **Host**

# IP Addressing



# IP Addressing

- Network portion
  - specifies a unique physical\* network
  - used for routing
- Host portion
  - specifies a unique host on the network
  - local delivery

# Internet Protocol (IP)

- ✓ What is an Internetwork
- ✓ Service Model
  - ✓ Datagrams, Packet Format, ...
- ✓ Addressing
  - Datagram Forwarding
  - ARP, DHCP, ICMP, ...

# Datagram Forwarding



**How do datagrams traverse the internetwork**

# Forwarding vs Routing

- ***forwarding*** is taking an input packet and sending it out the appropriate port
- ***routing*** is the process of building forwarding tables.



# Datagram Forwarding

- Every datagram has **destination IP**
- Network part **uniquely identifies** a physical network
  - All **hosts/routers** on network can **communicate** with all others
- Every network has a **router** on the net

# Datagram Forwarding

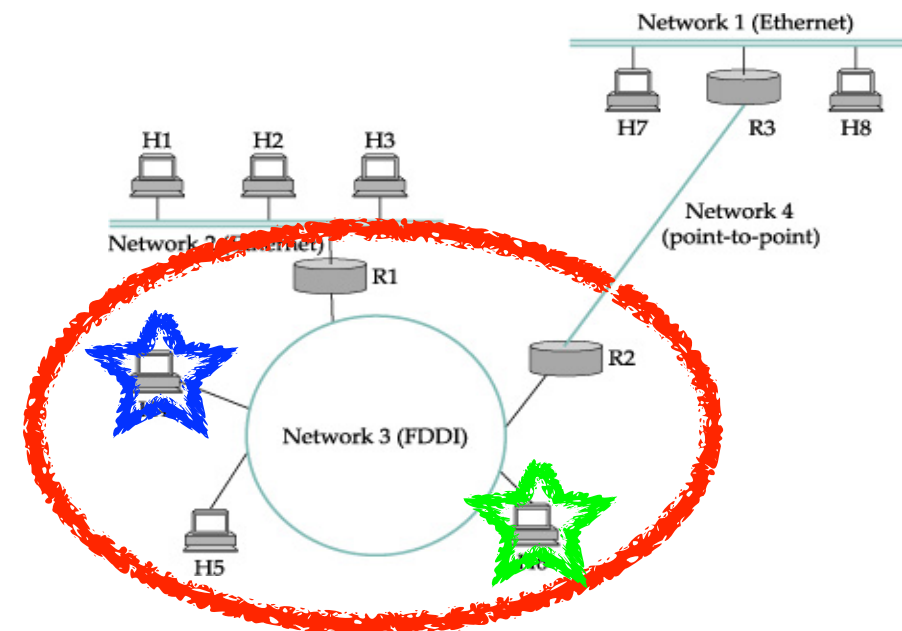
Each host and router maintains a forwarding table

<u>Network</u>	Next-hop
1 130.111	R1 130.111.32.1
2 141.114	R2 141.114.1.1
default	R2 76.5.4.3

# Datagram Forwarding

If destination network == source network

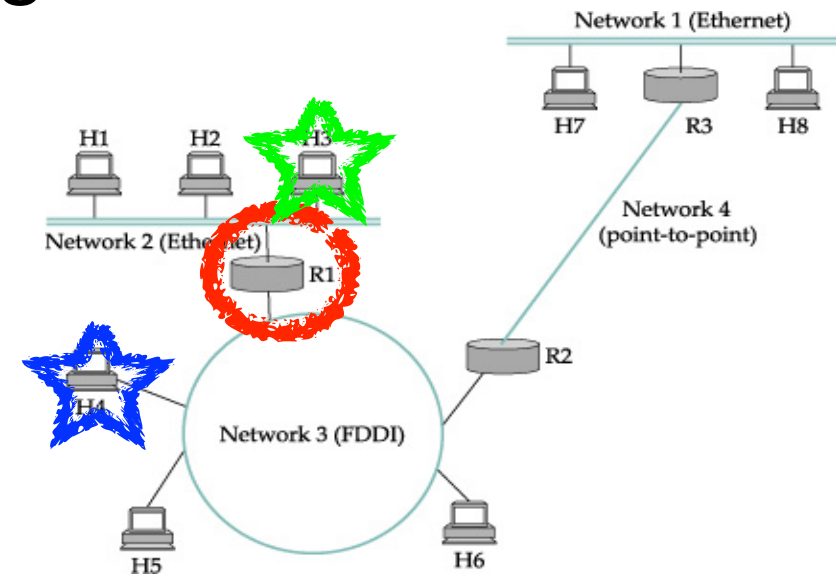
**deliver locally**



# Datagram Forwarding

if destination network in forwarding table

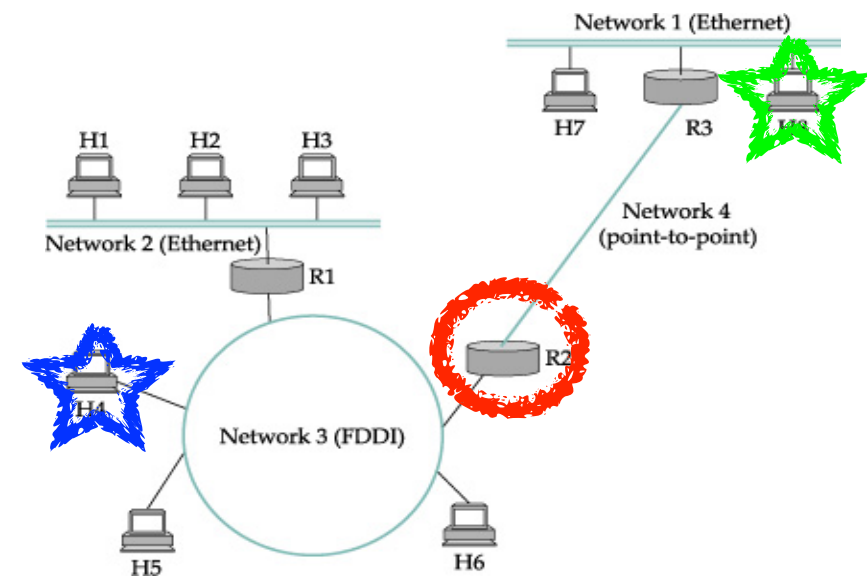
**deliver to next-hop**



# Datagram Forwarding

if destination network unknown

**deliver to default**



# Datagram Forwarding

- Local delivery for local datagrams
- Routers forward datagrams towards the physical network

# Internet Protocol (IP)

- ✓ What is an Internetwork
- ✓ Service Model
  - ✓ Datagrams, Packet Format, ...
- ✓ Addressing
- ✓ Datagram Forwarding
  - ARP, DHCP, ICMP, ...

# Address Resolution Protocol (ARP)

**Local delivery**

**it's not that simple**

**IP Address != Ethernet Address**



# Address Resolution

- Finds **local** or **link-level** address for an **IP** address
  - Both hosts on the same IP Network
  - Discover / Dynamic
- Uses *broadcast* feature of link-level

# Address Resolution

1. Do we have the IP-MAC addresses cached?
2. Send out broadcast query
3. Look for response, and fill in cache

# ARP Packet

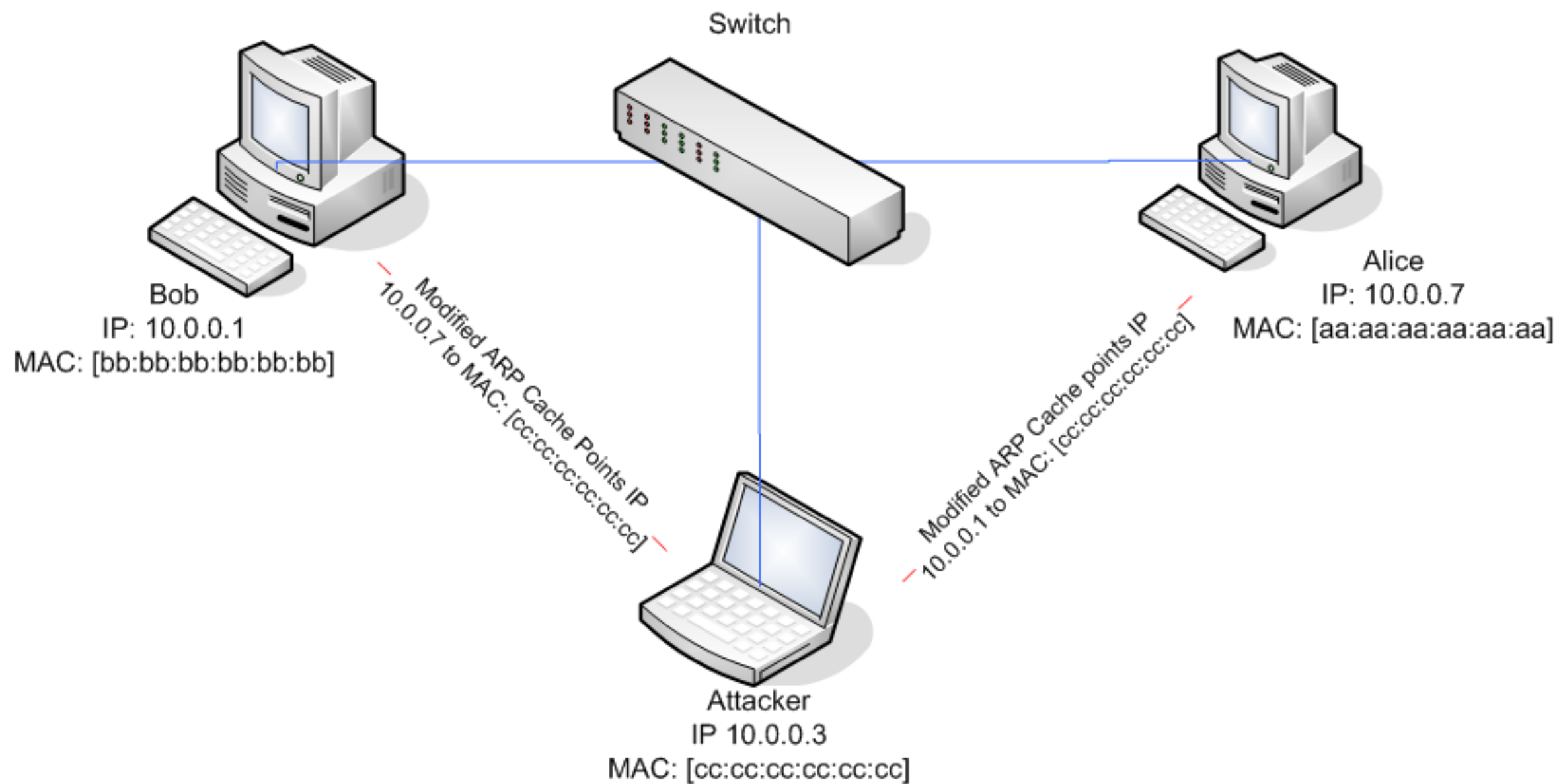
0	8	16	31
Hardware type = 1		ProtocolType = 0x0800	
HLen = 48	PLen = 32	Operation	
SourceHardwareAddr (bytes 0–3)			
SourceHardwareAddr (bytes 4–5)		SourceProtocolAddr (bytes 0–1)	
SourceProtocolAddr (bytes 2–3)		TargetHardwareAddr (bytes 0–1)	
TargetHardwareAddr (bytes 2–5)			
TargetProtocolAddr (bytes 0–3)			

Request  
or  
Response

# ARP

- Problems?
- ARP Spoofing  
([http://en.wikipedia.org/wiki/ARP\\_spoofing](http://en.wikipedia.org/wiki/ARP_spoofing))

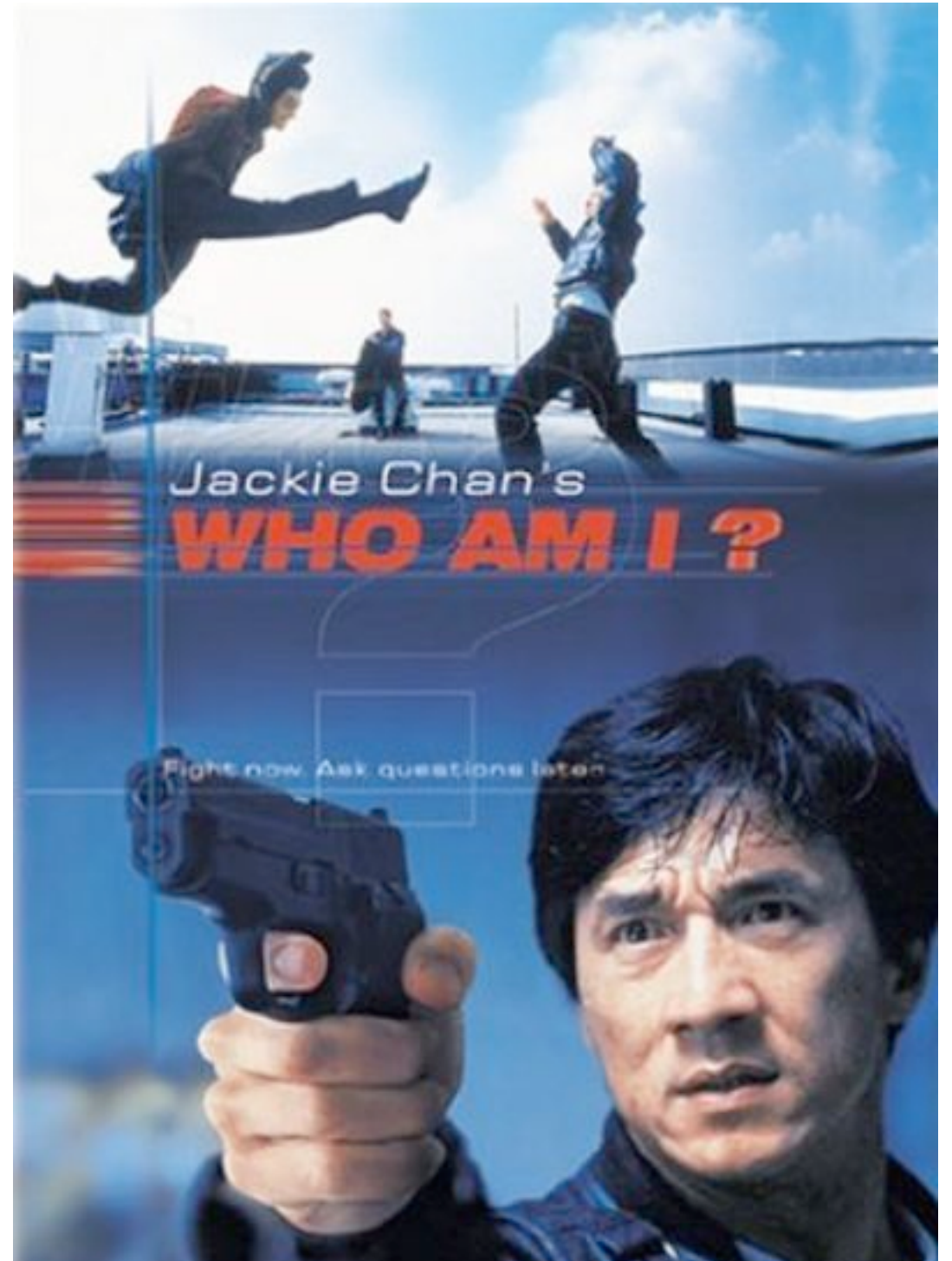
# ARP Spoofing



# DHCP

**Dynamic Host  
Configuration Protocol**

**Who Am I?**



# DHCP

- Static configuration of hosts
  - not flexible or adaptable to changes
  - cumbersome
- Dynamic configuration
  - provide host with an IP address
  - additional information: router, boot info...

# DHCP

1. Broadcast request: DHCPDISCOVER
2. Look for responses: DHCPOFFERs
3. Pick one and DHCPREQUEST
4. Wait for DHCPACK



# DHCP

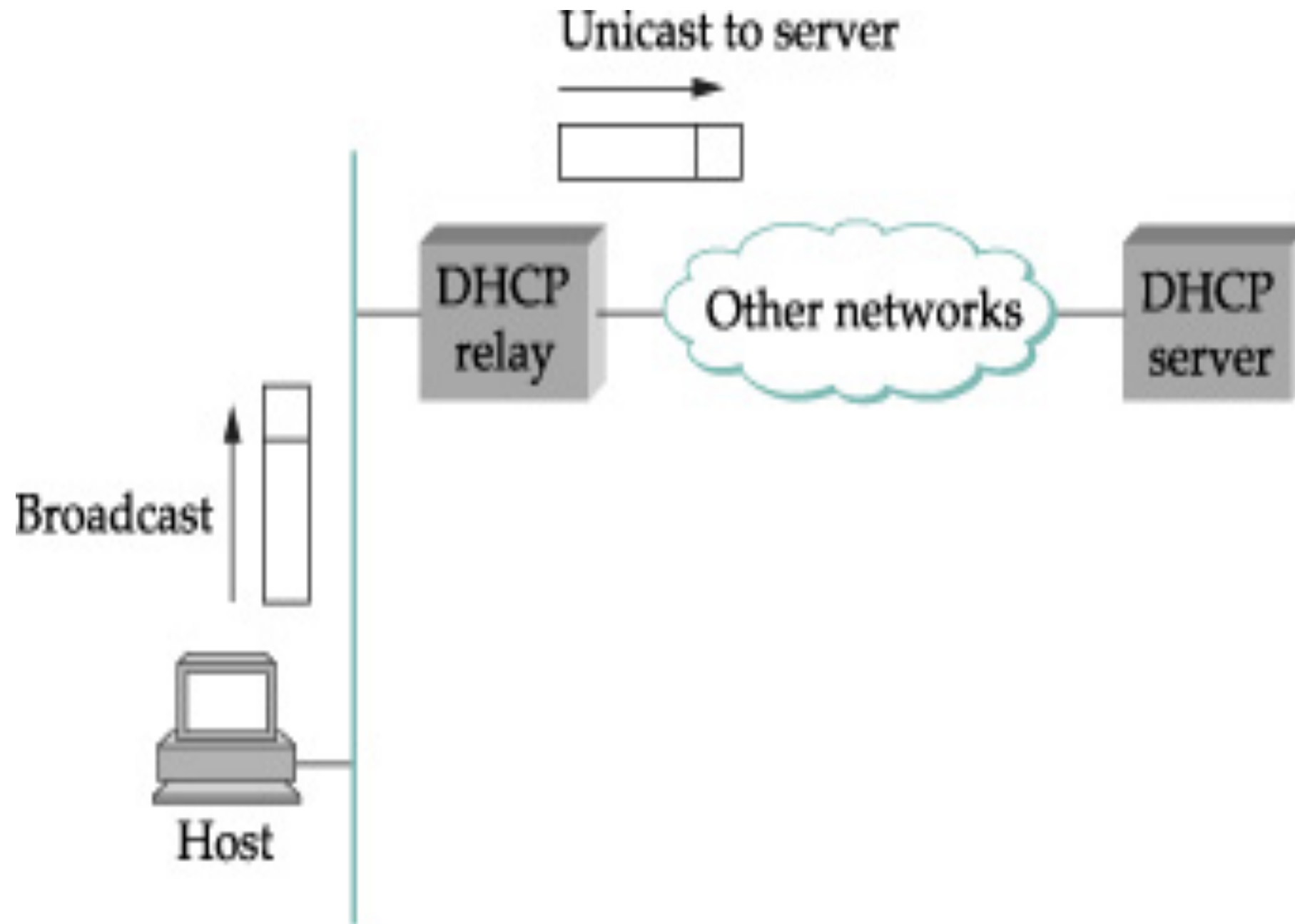
## **Problem:**

**Broadcast means you need a DHCP server on every network.**

## **Solution:**

**DHCP forwarding**

# DHCP Forwarding



# DHCP

- IP Address is ***leased*** for a given amount of time.
- Host must ***renew*** the lease with the server.
- Server can deny renewal
  - Can request a new lease

# ICMP

When things don't go as planned.



# ICMP

**Remember: Best-effort service**

- Diagnostic purposes (ping)
- TTL (hops) reaches 0 at a router
- Host not reachable (network error)
- Network redirection

# ARP, DHCP, ICMP

**These protocols are at the edge between the Network and IP layers**

- Address Resolution Protocol (ARP)
- Dynamic Host Configuration Protocol (DHCP)
- Internet Control Message Protocol (ICMP)

# Internet Protocol (IP)

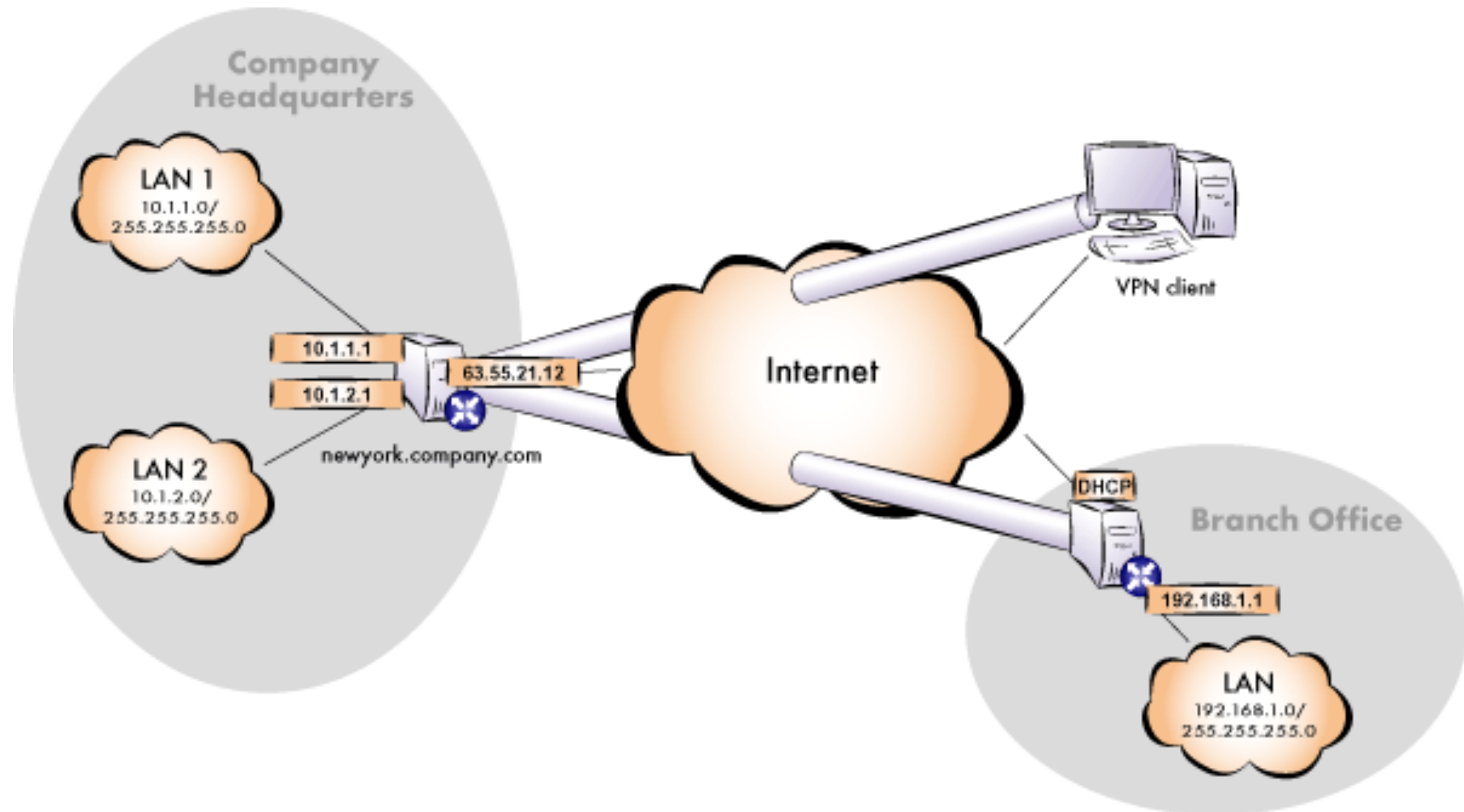
- ✓ What is an Internetwork
- ✓ Service Model
  - ✓ Datagrams, Packet Format, ...
- ✓ Addressing
- ✓ Datagram Forwarding
- ✓ ARP, DHCP, ICMP, ...

# Internet Protocol (IP)

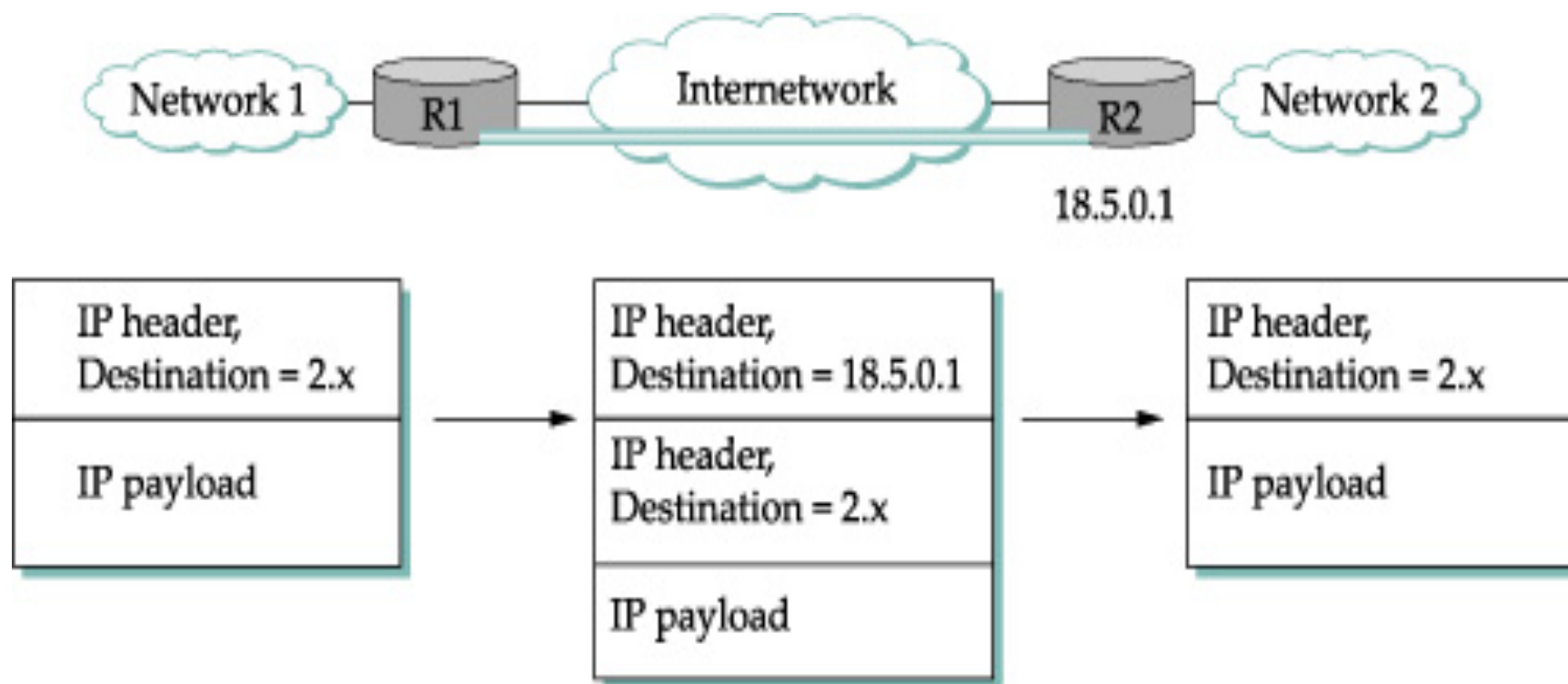
- ✓ What is an Internetwork
- ✓ Service Model
  - ✓ Datagrams, Packet Format, ...
- ✓ Addressing
- ✓ Datagram Forwarding
- ✓ ARP, DHCP, ICMP, ...



# Virtual Networks (VPN)



# VPN



# Inter-networking

✓ Internet Protocol (IP)

★ Routing

- The Internet
- Multicast\*
- Multi-protocol Label Switching\*

# Routing

- Network as a graph
- Distance Vector (RIP)
- Link State
  - Open Shortest Path First (OSPF)
- Mobile Routing

# Forwarding vs Routing

- ***forwarding*** is taking an input packet and sending it out the appropriate port
- ***routing*** is the process of building forwarding tables.

# Tables

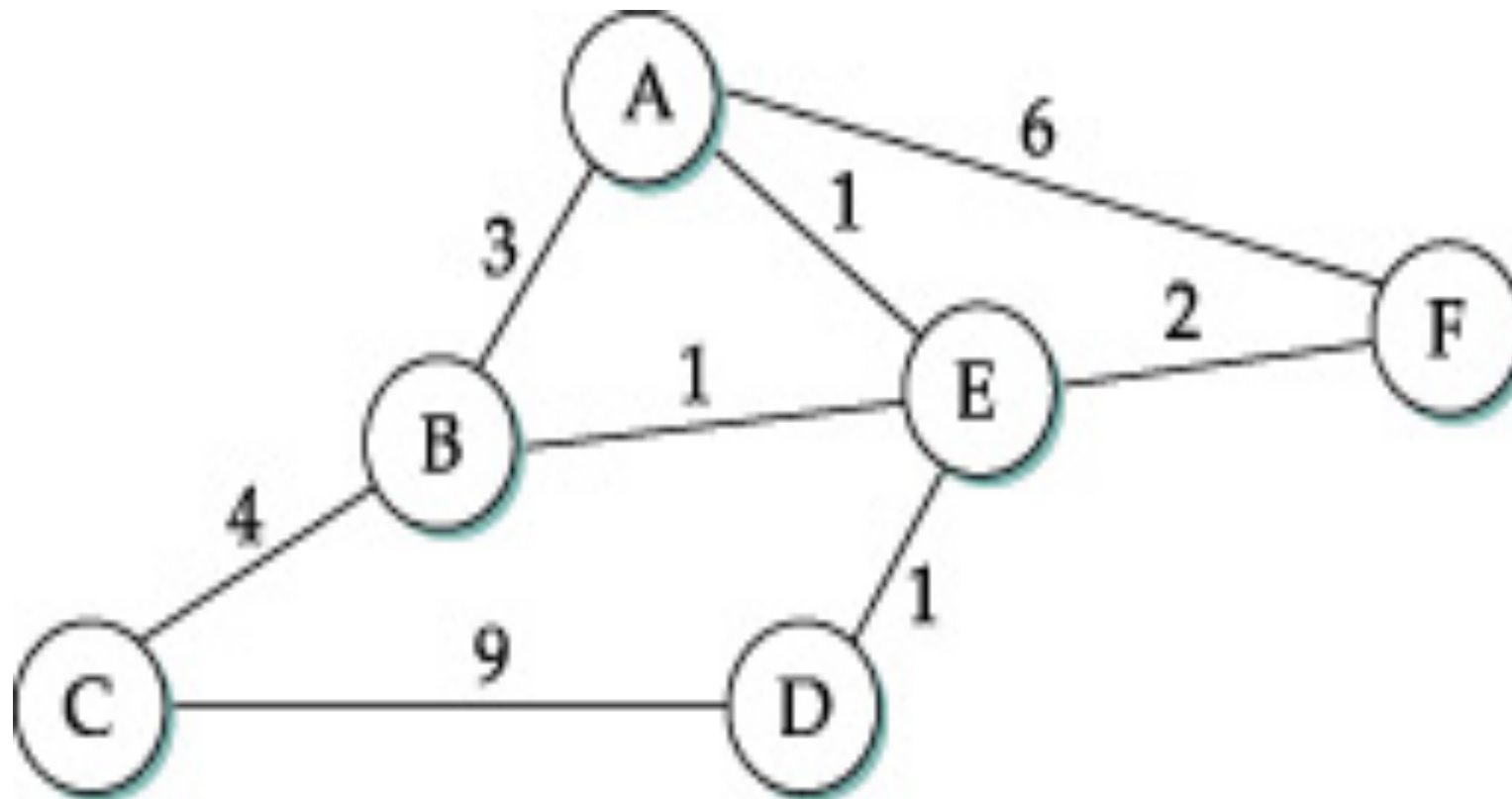
## Routing Table

Network	Next Hop
42	192.168.1.1

## Forwarding Table

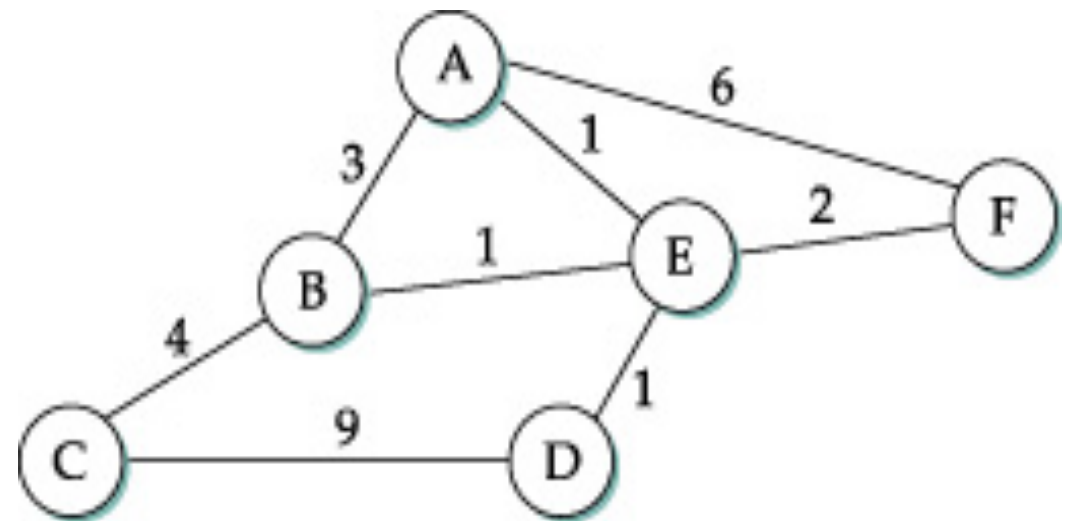
Network	Port	MAC
42	1	00:BC:D4:34:32:0B

# Network Graph



# Distance Vector (RIP)

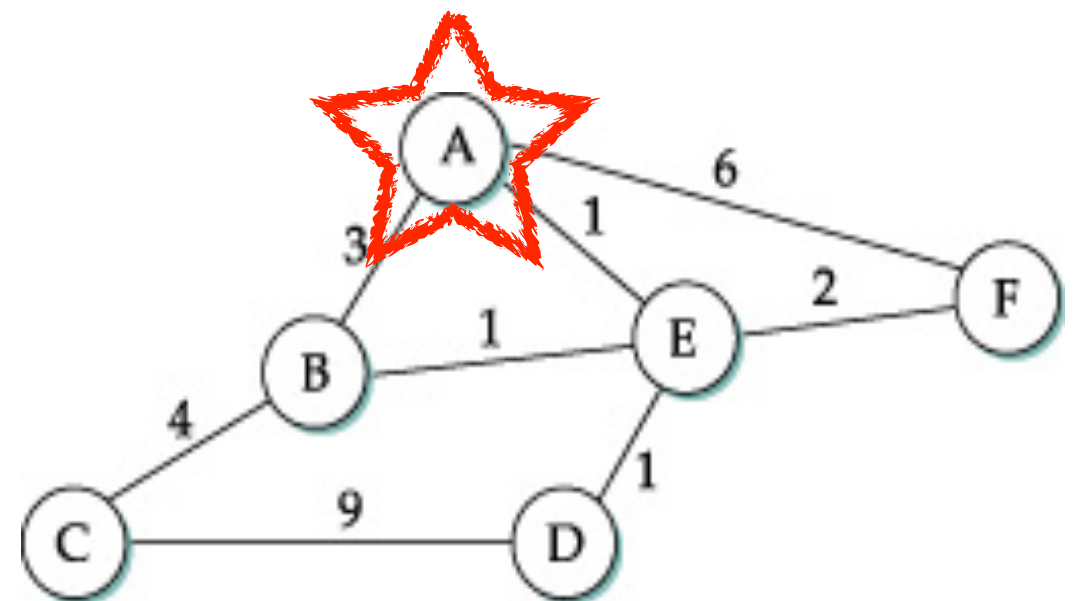
- Table of distance/cost to all nodes
- Distribute to immediate neighbors
- Link Down =  $\infty$
- periodic & triggered





# Distance Vector (RIP)

Dest	Cost	Hop
B	3	B
C	$\infty$	-
D	$\infty$	-
E	1	E
F	6	F

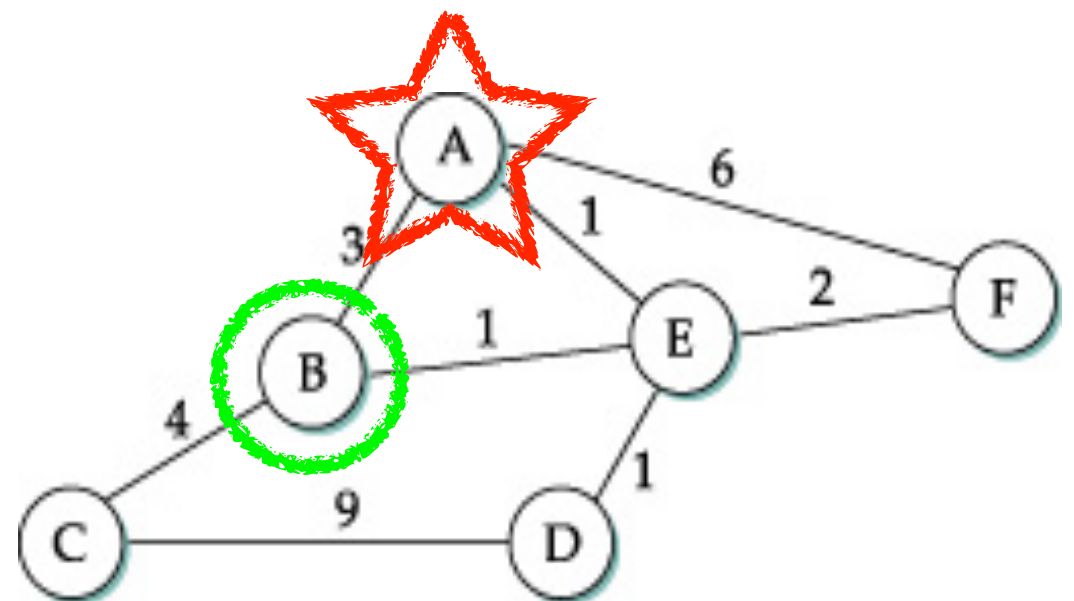


# Distance Vector (RIP)

Dest	Cost	Hop
B	3	B
C	7	B
D	$\infty$	-
E	1	E
F	6	F

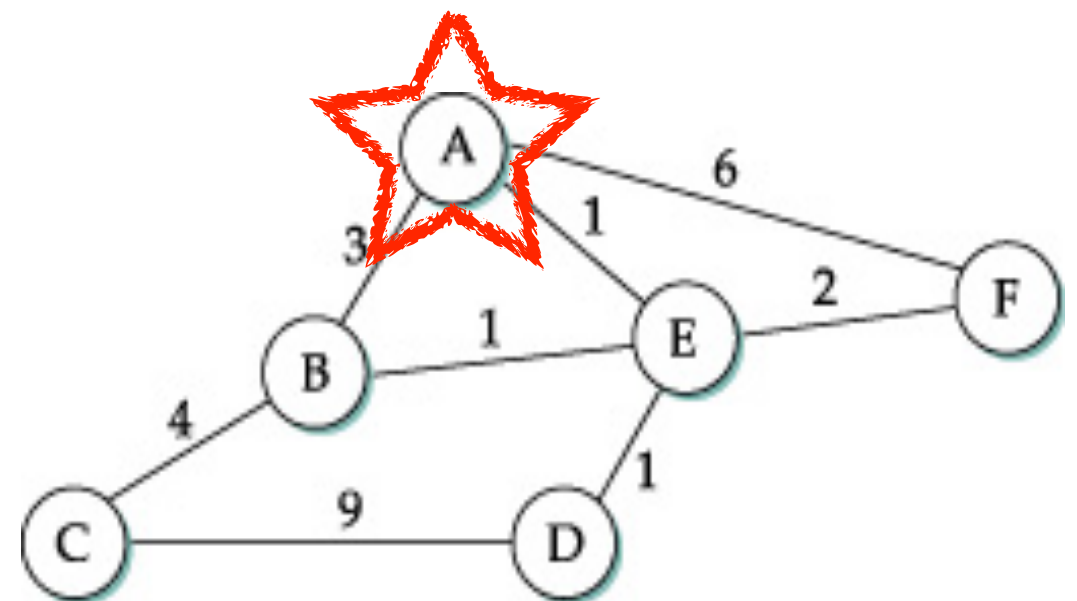
From: B

Dest	Cost	Hop
A	3	A
C	4	C
D	$\infty$	-
E	1	E
F	$\infty$	-



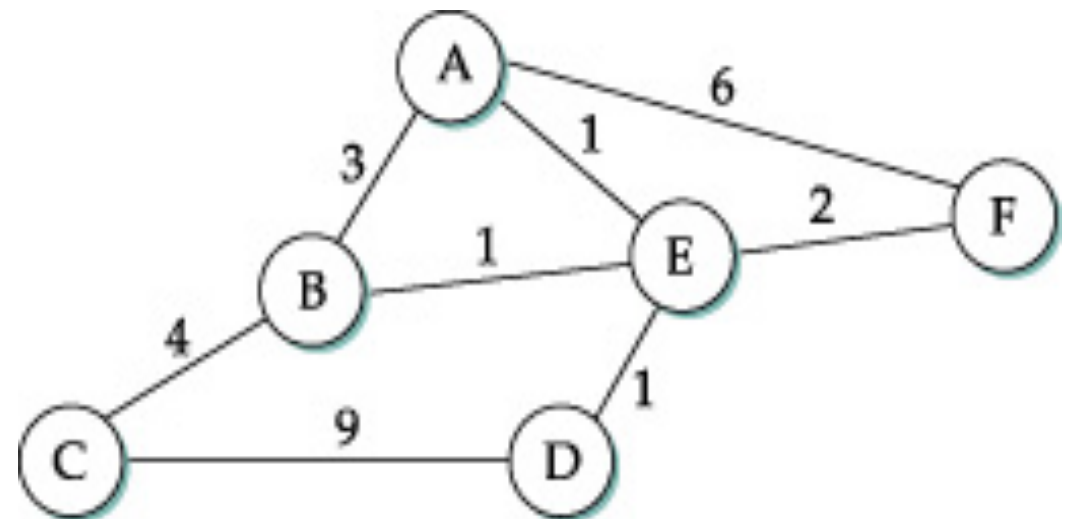
# Distance Vector (RIP)

Dest	Cost	Hop
B	2	E
C	6	E
D	2	E
E	1	E
F	3	E



# Distance Vector (RIP)

- Complete Map
- Exchange with immediate neighbors



# Distance Vector Game

1. Make “Routing Table” of network

Destination	Cost	Next Hop
-------------	------	----------

2. Send messages to neighbors with

From: \_\_\_\_\_

Destination
-------------

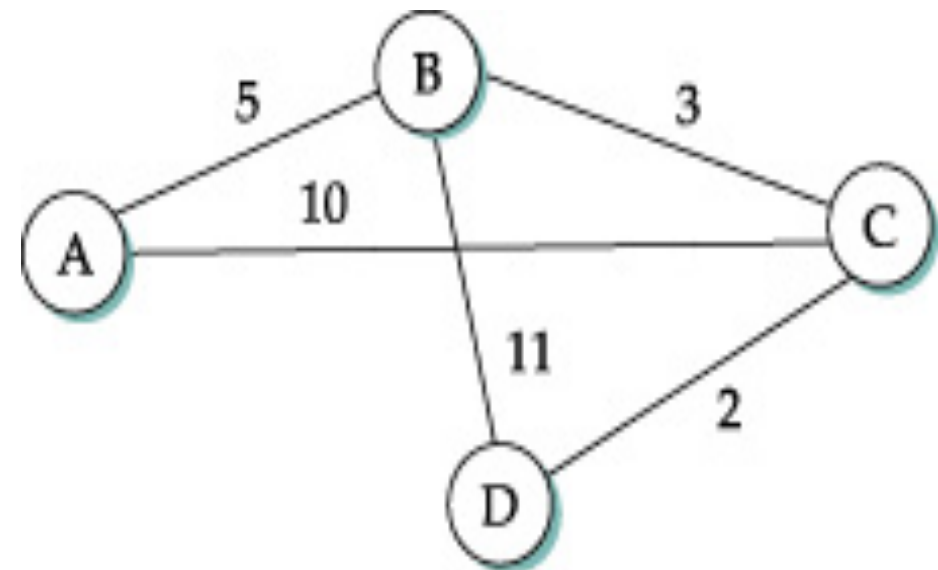
Cost
------

3. Use messages from neighbors and destroy

4. Repeat from step 2

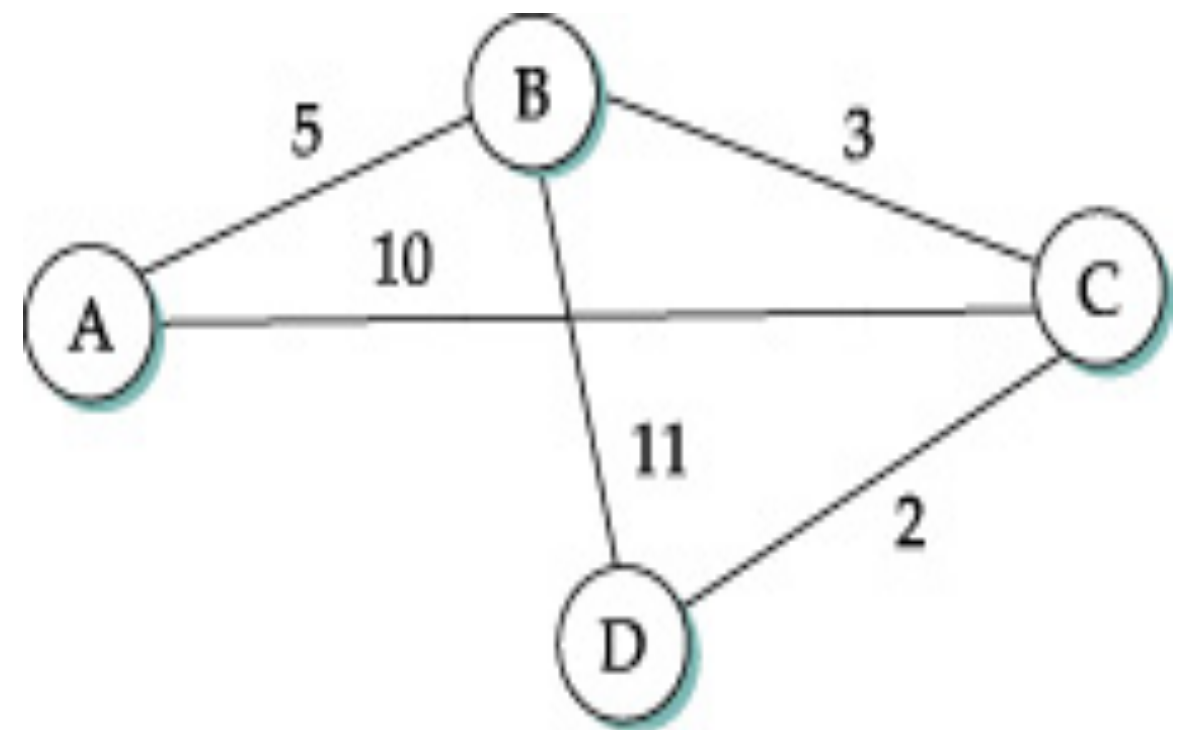
# Link State

- Link State Packet (LSP)
- Distance/Cost of neighbors
- Flood to all routers



# Link State

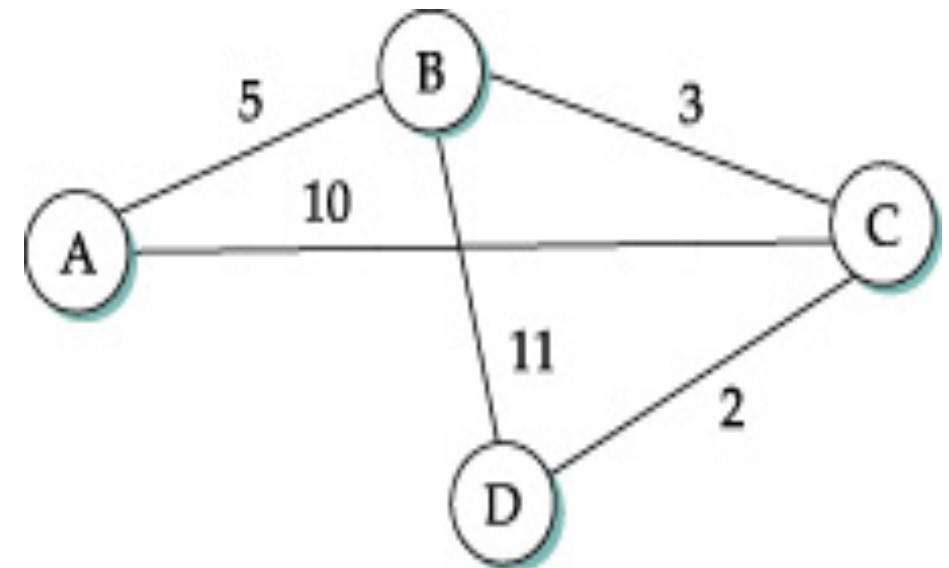
Dest	Cost	Hop
B	5	B
C	8	B
D	10	B



# Link State

Step	Confirmed	Tentative	Replaced
1	(A,0,-)		
2	(A,0,-)	(B,5,B) (C,10,C)	
3	(A,0,-) (B,5,B)	(C,8,B) (D,14,B)	(C,10,C)
4	(A,0,-) (B,5,B) (C,8,B)	(D,10,B)	(D,14,B)
5	+(D,10,B)		

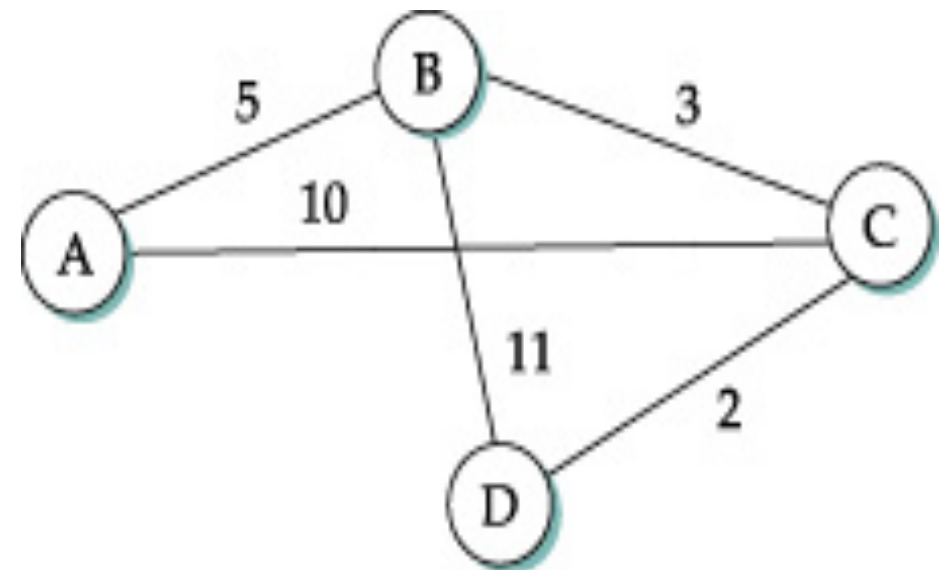
From: B cost 5  
 (A,5,A)  
 (C,3,C)  
 (D,11,D)





# Link State

- Link State Packet (LSP)
- Distance/Cost of neighbors
- Flood to all routers



# Link State Game

1. Make “Routing Table” of network

Confirmed

Tentative

2. Flood messages to everyone with neighbor information only,  
keep your own

From: \_\_\_\_\_

Neighbor	Cost	NextHop
----------	------	---------

3. Use messages from others (Dijkstra’s Alg.)

4. Repeat from step 2, until tentative is empty

# Metrics

- Bandwidth
- Latency
- “hops”
- Cost (dollars)
- Utilization
- Geo-Political boundaries

# Inter-networking

✓ Internet Protocol (IP)

✓ Routing

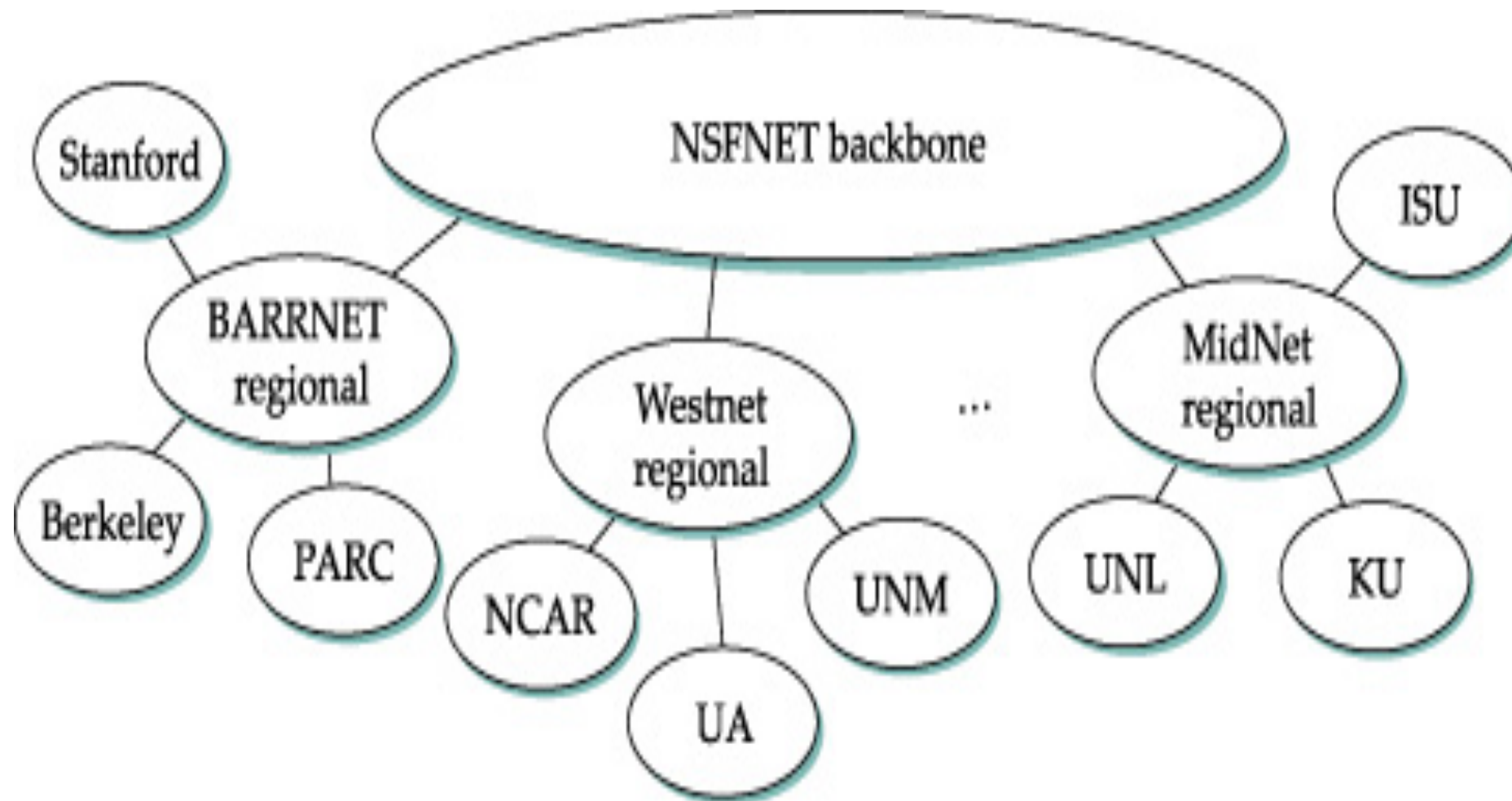
★ The Internet

- Multicast\*
- Multi-protocol Label Switching\*

# The Internet

- Construction
  - Autonomous Systems
- Subnetting
- Classless Routing (CIDR)
- Border Gateway Routing (BGP)
- IPv6

# What does it look like



# Subnetting

- Network part of address uniquely identifies a physical network
- Class B network (16 & 16 bits) - unused
- Class C network (24 & 8 bits) - too small

# Subnetting

- Allocate one “network” to several physical networks

Network number	Host number
----------------	-------------

Class B address

- internal to network
- outside routes same

11111111111111111111	00000000
----------------------	----------

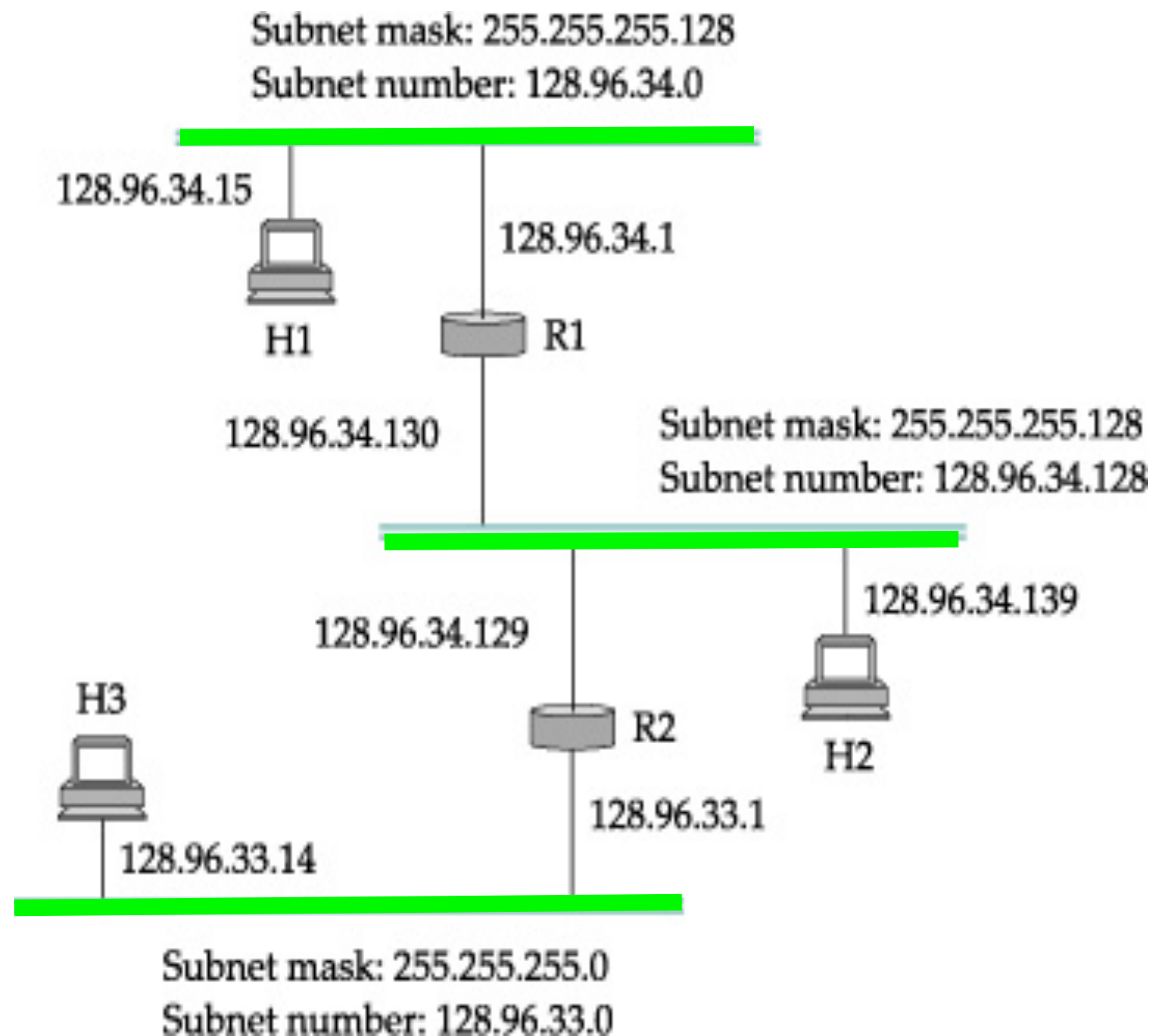
Subnet mask (255.255.255.0)

Network number	Subnet ID	Host ID
----------------	-----------	---------

Subnetted address



# Subnetting



# Subnetting

- Solves scalability problem
  - improve address assignment efficiency
  - aggregates information, from a distance

# Classless Routing (CIDR)

- Scalability problems in the backbone
  - Single location with multiple C nets
  - Lots of entries to maintain outside
  - Assignment of Class B wastes

# CIDR

- Solution: aggregate routes
  - $192.4.16/20 = 192.4.16 \text{ to } 192.4.31$
  - $192.4.16/24 = 192.4.16$

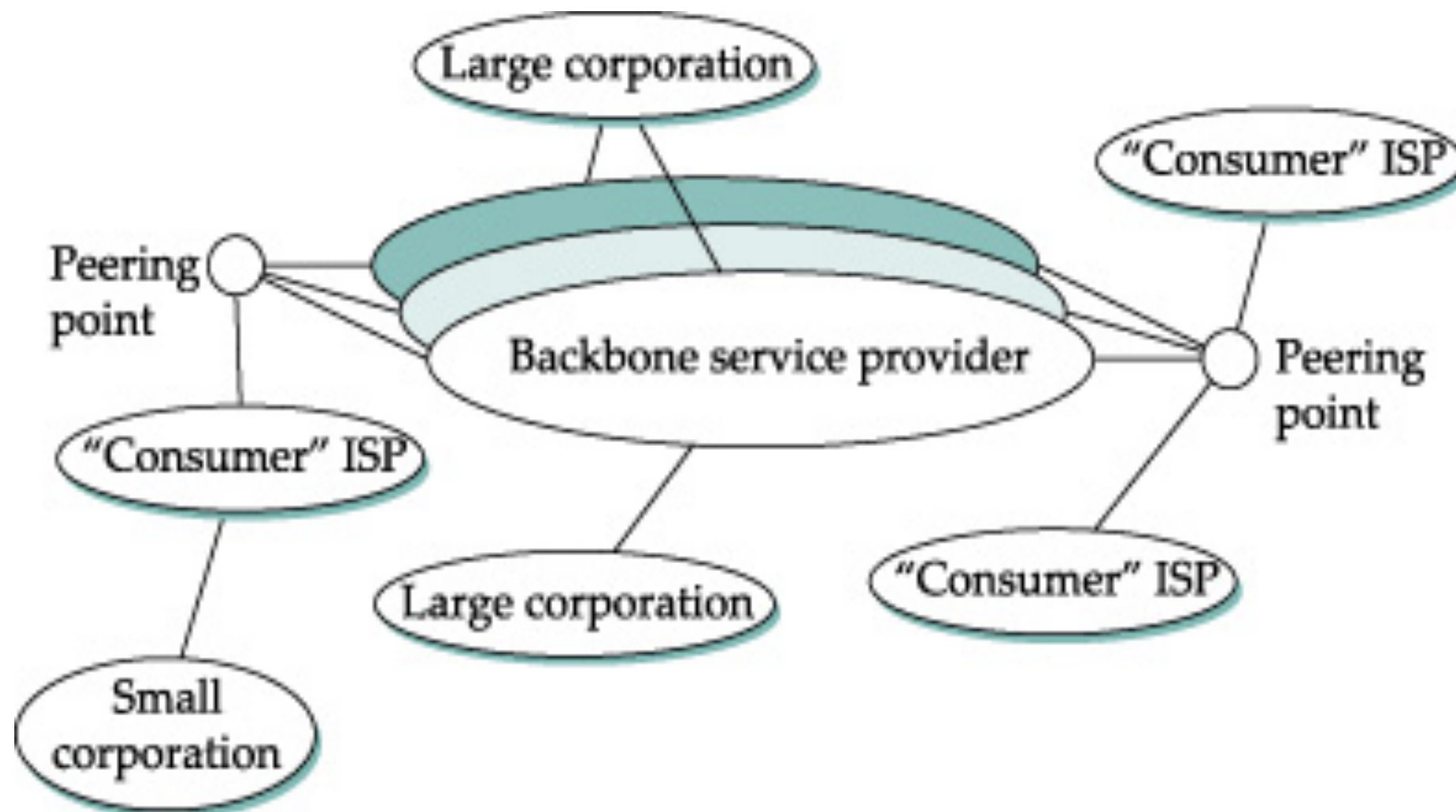
# CIDR

- Solution: Assign **blocks** of Class C addresses
  - creates aggregate routes to AS

# Interdomain Routing

- Each Autonomous System (AS)
  - Has a “border”
  - Handles its own internal routing
  - Has its own policies for in & out

# Interdomain Routing



# Interdomain Routing

- Find *some* path to destination
- Avoid *loops*
- Compliant with *policies* of ASs



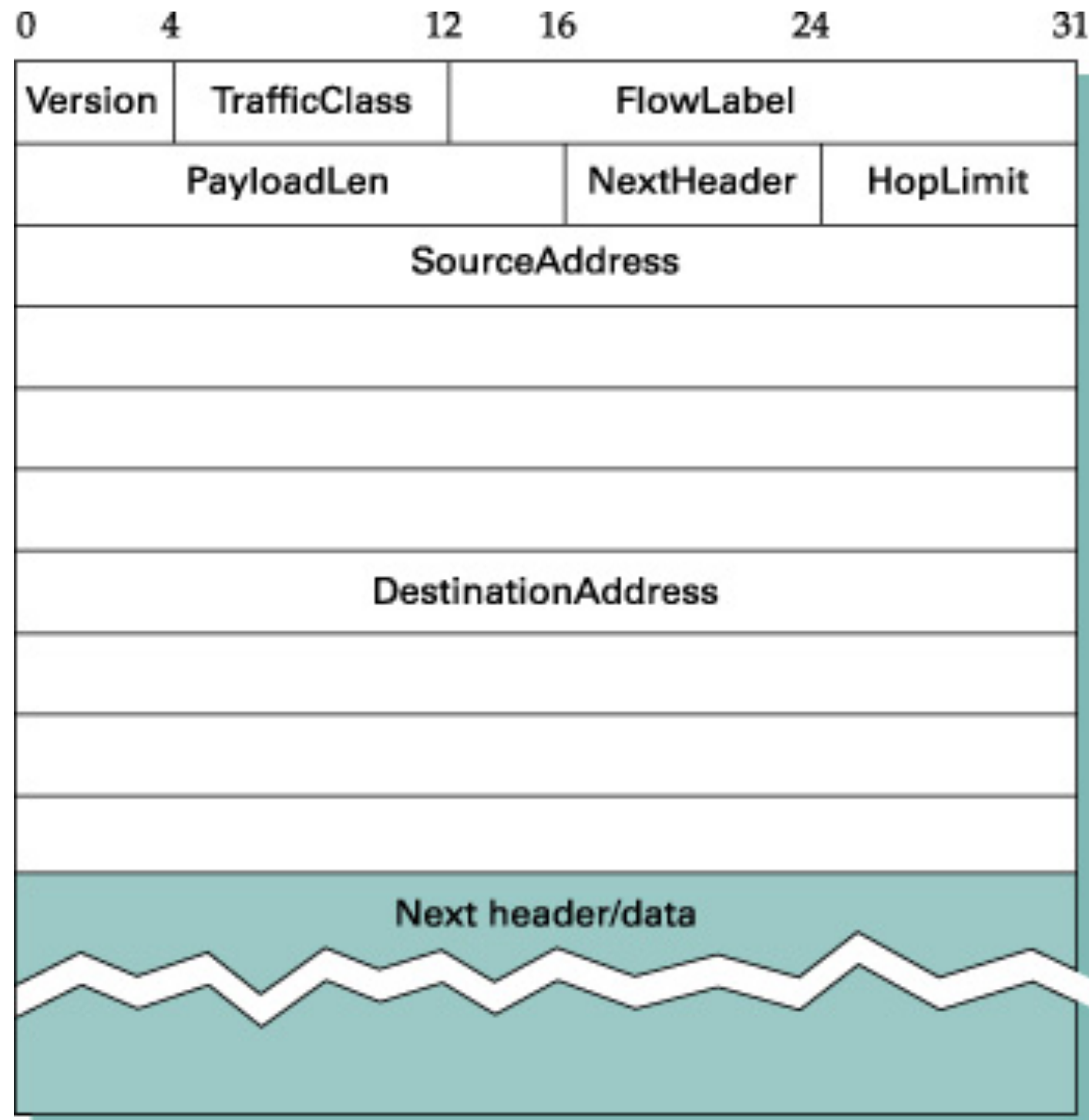
# Interdomain Routing

- Border Gateway Protocol (BGP)
  - AS numbers are *unique*
  - One node is the *speaker* for the network
  - Advertises *complete paths* to networks
  - Prevent *loops*
  - Path *withdrawn* messages

# Interdomain Routing

- Scalable (by hierarchy)
  - AS has own internal policies and routing
  - Only AS borders run BGP on backbone
  - Aggregate networks
  - Only need to find path to border

# IPv6



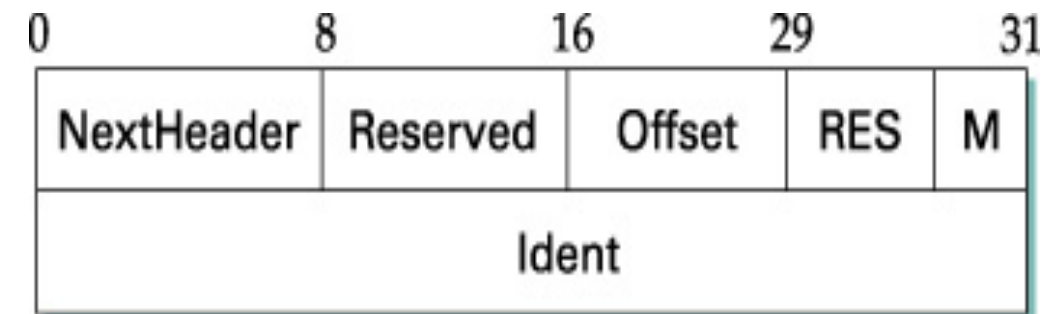
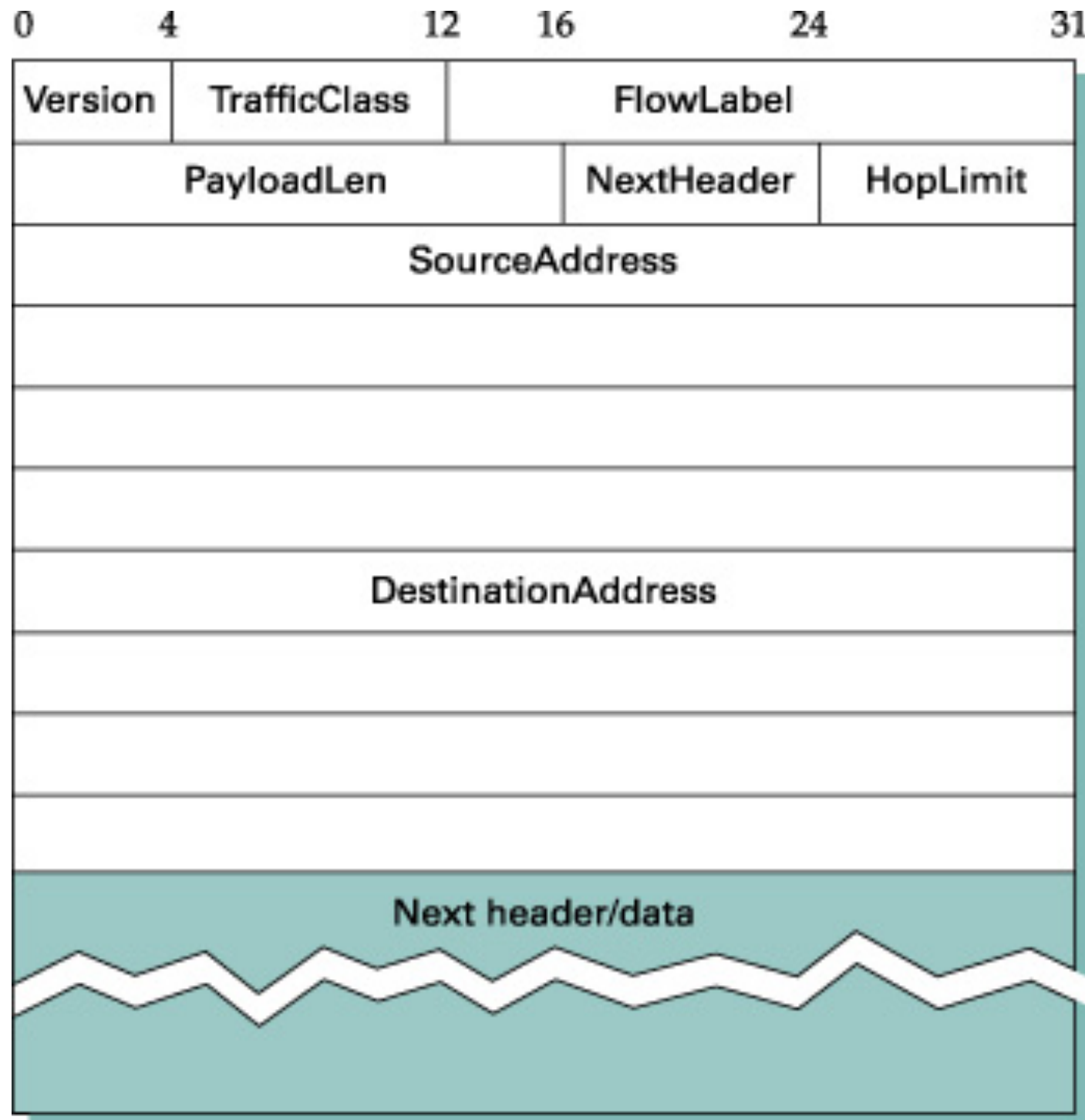
# IPv6 - Why?

- Address exhaustion (really Class B)
- Real-time services
- Security
- Configuration
- Routing - mobile hosts
- Jumbograms

# Addresses

- 128 bits =  $3.4 \times 10^{38}$  nodes
- Notation = x:x:x:x:x:x:x:x
  - :: = zeros; 45de:1230:FG::11:12
  - IPv4 (zero extended) ::FFFF:128.96.33.81

# Packet Format



# Configuration

- No more DHCP servers...
- Prefixes are assigned to networks (at AS)
  - Obtain interface ID
  - Obtain address prefix for subnet
    - local link only (1111 1110 10)
    - routers periodically advertise prefix
- Temporary addresses for mobile and privacy

# Inter-networking

- ✓ Internet Protocol (IP)
- ✓ Routing
- ✓ The Internet
  - Multicast\*
  - Multi-protocol Label Switching\*



# Inter-networking

fin