2025年春季学期《编译工程》



语法分析 自顶向下-递归下降

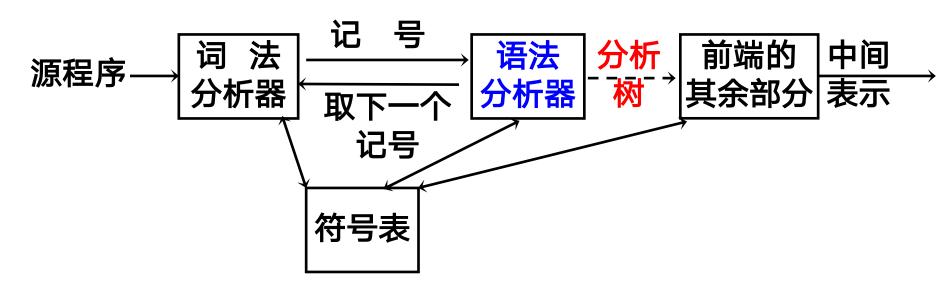
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2025年03月06日

❷ 本节提纲





- 自顶向下与自底向上方法的区别
- 自顶向下分析方法
 - 递归下降分析方法
 - 消除左递归、提取左公因子



语法分析的主要方法



・自顶向下(Top-down)

•针对输入串,从文法的开始符号出发,尝试根据产生式规则推导 (derive)出该输入串。

• 自底向上 (Bottom-up)

·针对输入串,尝试根据产生式规则归约 (reduce) 到文法的开始符号。



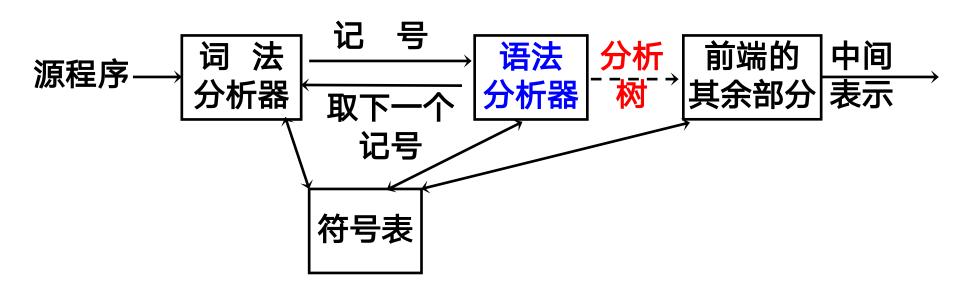
语法分析的主要方法



- · 自顶向下 (Top-down)
 - •针对输入串,从文法的开始符号出发,尝试根据产生式规则推导 (derive) 出该输入串。
 - 分析树的构造方法
 - 从根部开始
- 自底向上 (Bottom-up)
 - ·针对输入串,尝试根据产生式规则归约 (reduce) 到文法的开始符号。
 - 分析树的构造方法:
 - 从叶子开始

❷ 本节提纲





- 自顶向下与自底向上方法的区别
- 自顶向下分析方法
 - 递归下降分析方法
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递归下降语法分析



•数据结构

•一个输入缓冲区和向前看指针 lookahead

• 分析过程

- 自左向右扫描输入串
- •设计一个辅助过程match(),将lookahead指向的位置与产生式迭代生成的终结符进行匹配,如匹配,将lookahead挪到下一个位置
- 为每一个非终结符写一个分析过程
 - · 该过程可以调用其他非终结符的过程及match
 - 这些过程可能是递归的



递归下降语法分析——程序模拟推导



•考虑以下文法:

```
expr \rightarrow term
| term + expr
| term - expr
term \rightarrow num | (expr)
```

•分析过程:

- 从左到右扫描输入串
- 开始符号: *expr*
- 按顺序尝试产生式

```
void expr() {
   term();
   if (lookahead == '+/-') {
      match('+/-');
      expr();
   report("语法正确");
void term(){
  if (lookahead is num){
      match(lookahead);
  } else{ if (lookahead == '(') {
      match('(');
     expr();
      match(')');
  } else report("语法错误");}
```





```
expr \rightarrow term
| term + expr
| term - expr
term \rightarrow num | (expr)
```

```
( 2 - 3 )
```



```
void expr() {
   term();
   if (lookahead == '+/-') {
      match('+/-');
      expr();
   report("语法正确");
void term(){
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     match('(');
     expr();
     match(')');
  } else report("语法错误");}
```





•考虑以下文法:

```
\begin{array}{c} expr \rightarrow term \\ | term + expr \\ | term - expr \\ \\ term \rightarrow num \mid (expr) \end{array}
```

expr

```
( 2 - 3 )
```

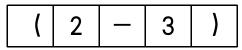


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```





```
expr \rightarrow term
| term + expr
| term - expr
term \rightarrow num | (expr)
```





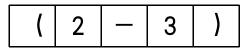
```
expr
|
term
```

```
void expr() {
   term();
   if (lookahead == '+/-') {
      match('+/-');
      expr();
   report("语法正确");
void term(){
  if (lookahead is num){
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```
expr \rightarrow term
| term + expr
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term \rightarrow num | (expr)
```





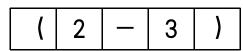
```
expr
|
term
```

```
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   if (lookahead == '+/-') {
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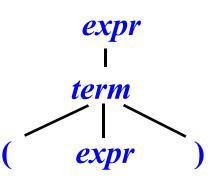




```
expr \rightarrow term
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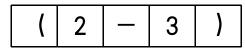




•考虑以下文法:

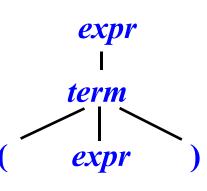
```
expr → term
| term + expr
| term - expr

term → num | (expr)
```





些配 箭头前进



```
void expr() {
   term();
   if (lookahead == '+/-') {
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```



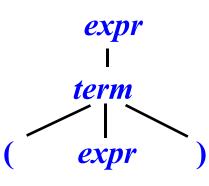


```
expr → term
| term + expr
| term - expr

term → num | (expr)
```





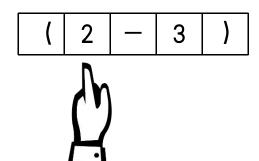


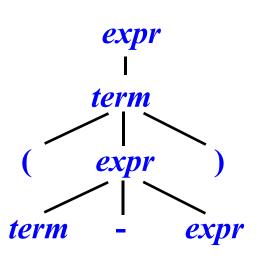
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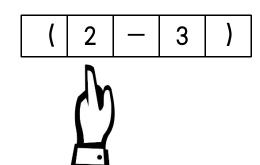


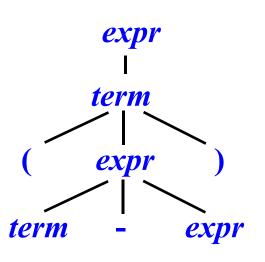
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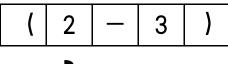
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   term();
   if (lookahead ==
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      expr();
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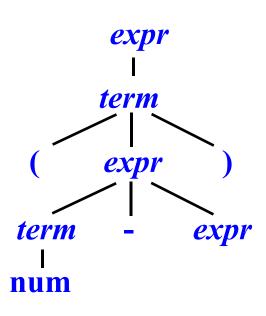


```
expr → term
| term + expr
| term - expr

term → num | (expr)
```







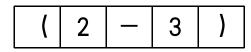
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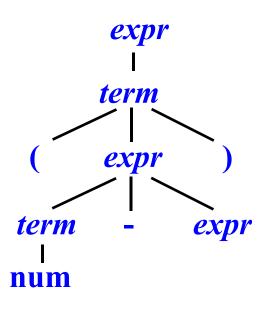




```
expr \rightarrow term
| term + expr
| term - expr
term \rightarrow num | (expr)
```







```
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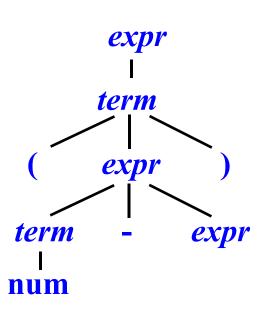


```
expr → term
| term + expr
| term - expr

term → num | (expr)
```







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```





```
expr \rightarrow term
         | term + expr
          term - expr
term \rightarrow num \mid (expr)
                                   expr
                                  term
                                  expr
               匹配
               箭头前进 term
                                          expr
                          num
```

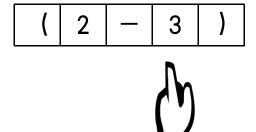
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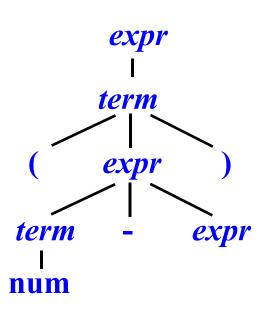




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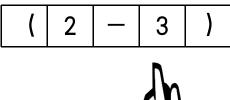


```
expr → term

| term + expr

| term - expr

term → num | (expr)
```





```
expr
term
( expr )
term - expr
num term
```

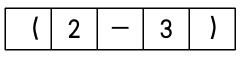
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```
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term
( expr )
term - expr
num term
num
```

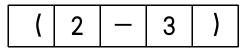
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•考虑以下文法:

```
expr \rightarrow term
| term + expr
| term - expr
term \rightarrow num | (expr)
```



匹配 箭头前进



```
expr
term
( expr )
term - expr
num term
num
```

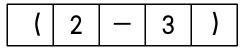
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expr \rightarrow term
| term + expr
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term \rightarrow num | (expr)
```



匹配 箭头前进



```
expr
term
expr
expr
expr
num
term
num
term
num
```

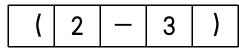
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•考虑以下文法:

```
expr \rightarrow term
| term + expr
| term - expr
term \rightarrow num | (expr)
```



分析完毕 接受该串

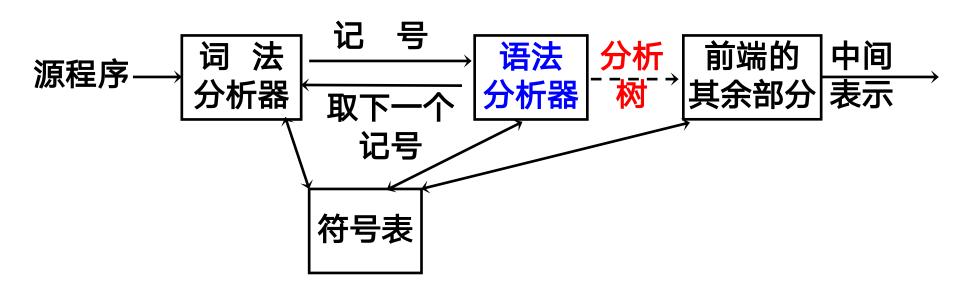


```
expr
term
( expr )
term - expr
num term
num
```

```
void expr() {
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   if (lookahead == '+/-') {
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      expr();
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- 自顶向下分析方法
 - 递归下降预测分析方法
 - 消除左递归、提取左公因子





- •可能进入无限循环
- 考虑以下文法

$$S \rightarrow Sa \mid b$$

· 该文法是左递归的(left-recursive)





- 可能进入无限循环
- 考虑以下文法

$$S \rightarrow Sa \mid b$$

- · 该文法是左递归的(left-recursive)
- 自顶向下分析方法无法处理左递归
 - Why?





- •可能进入无限循环
- 考虑以下文法

$$S \rightarrow Sa \mid b$$

- · 该文法是左递归的(left-recursive)
- 自顶向下分析方法无法处理左递归
 - 考虑输入文法符号串为baaaaa
 - 最左推导如下:
 - $S \Rightarrow Sa \Rightarrow Saaa \Rightarrow Saaaa \dots$
 - · 输入缓冲区lookahead指针纹丝未动





• 直接左递归

$$A \rightarrow Aa \mid \beta$$
, 其中 a , β 不以A开头

- 串的特点 $\beta a \dots a (A \Rightarrow ^+ A a)$
- •消除直接左递归

$$A \rightarrow \beta A'$$

$$A' \rightarrow \alpha A' \mid \varepsilon$$





・直接左递归

$$A \rightarrow Aa \mid \beta$$
, 其中 a , β 不以A开头

- ■串的特点 $\beta a \dots a (A \Rightarrow ^+ A a)$
- 消除直接左递归

$$A \to \beta A'$$

$$A' \to \alpha A' \mid \varepsilon$$

• 考虑之前的文法

$$\begin{array}{c|c}
S \to Sa \mid b \\
\hline
 & a & \beta
\end{array}$$



• 直接左递归

 $A \rightarrow Aa \mid \beta$, 其中a, β 不以A开头

- ■串的特点 $\beta a \dots a (A \Rightarrow +A a)$
- •消除直接左递归

$$A \to \beta A'$$

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• 考虑之前的文法

$$\begin{array}{c|c}
S \to Sa \mid b \\
\hline
 & a & \beta
\end{array}$$

$$S \rightarrow bS'$$

$$S' \rightarrow aS' \mid \varepsilon$$

baaaaa推导:

 $S \Rightarrow bS' \Rightarrow baS' \Rightarrow baaS' \Rightarrow baaaS' \Rightarrow baaaaS' \Rightarrow baaaaaS' \Rightarrow baaaaaS' 翰入缓冲区指针不停地移动$



• 直接左递归

 $A \rightarrow Aa \mid \beta$, 其中a , β 不以A开头

- ■串的特点 $\beta a \dots a (A \Rightarrow +A a)$
- 消除直接左递归

$$A \to \beta A'$$

$$A' \to \alpha A' \mid \varepsilon$$

• 考虑之前的文法

$$\begin{array}{c|c}
S \to Sa \mid b \\
\hline
 & a & \beta
\end{array}$$

$$S \to bS'$$

$$S' \to aS' \mid \varepsilon$$

baaaaa推导:

 $S \Rightarrow bS' \Rightarrow baS' \Rightarrow baaS' \Rightarrow baaaS' \Rightarrow baaaaS' \Rightarrow baaaaaS'$ 输入缓冲区指针不停地移动



• 例 算术表达文法

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

$$F \rightarrow (E) \mid id$$



• 例 算术表达文法

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

$$F \rightarrow (E) \mid id$$

$$(T+T...+T)$$

 $(F*F...*F)$





• 例 算术表达文法

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

$$F \rightarrow (E) \mid id$$

$$(T+T...+T)$$

 $(F*F...*F)$

• 消除左递归后文法

$$E \rightarrow TE'$$

$$E' \rightarrow + TE' \mid \varepsilon$$

$$T \rightarrow FT'$$

$$T' \rightarrow * F T' \mid \varepsilon$$

$$F \rightarrow (E) \mid id$$

注明:红色部分代表了 α ,蓝色部分代表了 β

消除左递归的推广



· 处理任意数量的A产生式

$$A \rightarrow A a_1 | A a_2 | \dots | A a_m | \beta_1 | \beta_2 | \dots | \beta_n$$

其中 β_i 都不以A开头

改为:

$$A \rightarrow \beta_1 A' \mid \beta_2 A' \mid \dots \mid \beta_n A'$$

$$A' \rightarrow \alpha_1 A' \mid \alpha_2 A' \mid \dots \mid \alpha_m A' \mid \varepsilon$$



消除间接左递归



• 非直接左递归

$$S \rightarrow Aa \mid b$$

$$A \rightarrow Sd \mid \varepsilon$$



消除间接左递归



・非直接左递归

$$S \rightarrow Aa \mid b$$

$$A \rightarrow Sd \mid \varepsilon$$

• 先变换成直接左递归

$$S \rightarrow Aa \mid b$$

$$A \rightarrow Aad \mid bd \mid \varepsilon$$

• 再消除左递归

$$S \rightarrow Aa \mid b$$

$$A \rightarrow bd \mid A' \mid A'$$

$$A' \rightarrow adA' \mid \varepsilon$$



- · 有左公因子的(left -factored)文法:
 - $A \rightarrow \alpha \beta_1 \mid \alpha \beta_2$

- ·提左公因子(left factoring)
 - 推后选择产生式的时机,以便获取更多信息 $A \rightarrow \alpha\beta_1 \mid \alpha\beta_2$ 等价于

$$A \to a A'$$

$$A' \to \beta_1 \mid \beta_2$$



提左公因子(left factoring)



· 例 悬空else的文法 $stmt \rightarrow if expr then stmt else stmt$ if expr then stmt other 提左因子 $stmt \rightarrow if expr then stmt optional else part$ other optional else part \rightarrow else stmt

算法仍然二义!!!





- 复杂的回溯→代价太高
 - 非终结符有可能有多个产生式
 - 由于信息缺失,无法准确预测选择哪一个
 - 考虑到往往需要对多个非终结符进行推导展开,因此尝试的路径可能 呈指数级爆炸
- 其分析过程类似于NFA
- ·问题:是否可以构造一个类似于DFA的分析方法?

2025年春季学期《编译工程》



一起努力 打造国产基础软硬件体系!

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