

CSC 3201 Compiler Construction



Department of Computer Science
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Week 6 (Lecture 1)



First

- $\text{First}(\alpha)$ is a set of terminal symbols that begin in strings derived from α .
- Example:
 - $A \rightarrow abc \mid def \mid ghi$
 - $\text{First}(A) = \{a, d, g\}$



First

- Rules:

- For a production rule $X \rightarrow \epsilon$, $\text{First}(X) = \{ \epsilon \}$
- For any terminal symbol 'a', $\text{First}(a) = \{ a \}$
- For a production rule $X \rightarrow Y_1Y_2Y_3$, Calculating $\text{First}(X)$
 - If $\epsilon \notin \text{First}(Y_1)$, then $\text{First}(X) = \text{First}(Y_1)$
 - If $\epsilon \in \text{First}(Y_1)$, then $\text{First}(X) = \{ \text{First}(Y_1) - \epsilon \} \cup \text{First}(Y_2Y_3)$, Calculating $\text{First}(Y_2Y_3)$
 - If $\epsilon \notin \text{First}(Y_2)$, then $\text{First}(Y_2Y_3) = \text{First}(Y_2)$
 - If $\epsilon \in \text{First}(Y_2)$, then $\text{First}(Y_2Y_3) = \{ \text{First}(Y_2) - \epsilon \} \cup \text{First}(Y_3)$



Follow

- $\text{Follow}(\alpha)$ is a set of terminal symbols that appear immediately to the right of α .
- Rules:
 - For the start symbol S , place $\$$ in $\text{Follow}(S)$
 - For any production rule $A \rightarrow \alpha B$, $\text{Follow}(B) = \text{Follow}(A)$
 - For any production rule $A \rightarrow \alpha B \beta$,
 - If $\epsilon \notin \text{First}(\beta)$, then $\text{Follow}(B) = \text{First}(\beta)$
 - If $\epsilon \in \text{First}(\beta)$, then $\text{Follow}(B) = \{ \text{First}(\beta) - \epsilon \} \cup \text{Follow}(A)$



First and Follow

- Note:
 - ϵ may appear in the first function of an NT.
 - ϵ will never appear in the follow function of an NT.
 - Before calculating the first and follow functions, eliminate Left Recursion from the grammar.
 - Calculate the follow function of an NT by looking where it is present on the RHS of a production rule.

First and Follow

| | First | Follow |
|----------------------------|---------------------|----------------|
| $S \rightarrow ABCDE$ | | |
| $A \rightarrow a \epsilon$ | 1 $\{a, \epsilon\}$ | $\{b, c\}$ |
| $B \rightarrow b \epsilon$ | 2 $\{b, \epsilon\}$ | $\{c\}$ |
| $C \rightarrow c$ | 3 $\{c\}$ | $\{d, e, \$\}$ |
| $D \rightarrow d \epsilon$ | 4 $\{d, \epsilon\}$ | $\{e, \$\}$ |
| $E \rightarrow e \epsilon$ | 5 $\{e, \epsilon\}$ | $\{\$\}$ |

- $\text{Follow}(\text{RHS}) = \text{Follow}(\text{LHS})$
- Look at the occurrence on rHS in all productions.

First and Follow

| | First | Follow |
|------------------------------|---------------------|-----------------|
| $S \rightarrow ABCDE$ | 6 $\{a, b, c\}$ | $\{\$ \}$ |
| $A \rightarrow a \epsilon$ | 1 $\{a, \epsilon\}$ | $\{b, c\}$ |
| $B \rightarrow b \epsilon$ | 2 $\{b, \epsilon\}$ | $\{c\}$ |
| $C \rightarrow c$ | 3 $\{c\}$ | $\{d, e, \$ \}$ |
| $D \rightarrow d \epsilon$ | 4 $\{d, \epsilon\}$ | $\{e, \$ \}$ |
| $E \rightarrow e \epsilon$ | 5 $\{e, \epsilon\}$ | $\{\$ \}$ |

- S is not in the RHS of any production



First and Follow

- Example 2:

$$S \rightarrow Bb \mid Cd$$
$$B \rightarrow aB \mid \varepsilon$$
$$C \rightarrow cC \mid \varepsilon$$

First and Follow

$S \rightarrow Bb | Cd \quad \{a, b, c, d\} \quad \{\$ \}$
 $B \rightarrow aB | \epsilon \quad \{a, \epsilon\} \quad \{b\}$
 $C \rightarrow cC | \epsilon \quad \{c, \epsilon\} \quad \{d\}$



First and Follow

- Example 3:

$S \rightarrow aBDh$

$B \rightarrow cC$

$C \rightarrow bc \mid \varepsilon$

$D \rightarrow EF$

$E \rightarrow g \mid \varepsilon$

$F \rightarrow f \mid \varepsilon$

First and Follow

| | | |
|-----------------------------|----------------------|---------------|
| $S \rightarrow aBDh$ | $\{a\}$ | $\{a\}$ |
| $B \rightarrow cC$ | $\{c\}$ | $\{a, f, h\}$ |
| $C \rightarrow bc \epsilon$ | $\{b, \epsilon\}$ | $\{a, f, h\}$ |
| $D \rightarrow EF$ | $\{g, f, \epsilon\}$ | $\{h\}$ |
| $E \rightarrow g \epsilon$ | $\{g, \epsilon\}$ | $\{f, h\}$ |
| $F \rightarrow f \epsilon$ | $\{f, \epsilon\}$ | $\{h\}$ |



First and Follow

- $\text{First}(S) = \{a\}$
- $\text{First}(B) = \{c\}$
- $\text{First}(C) = \{b, \epsilon\}$
- $\text{First}(D) = \{ \text{First}(E) - \epsilon \} \cup \text{First}(F) = \{g, f, \epsilon\}$
- $\text{First}(E) = \{ g , \epsilon \}$
- $\text{First}(F) = \{ f , \epsilon \}$



First and Follow

- $\text{Follow}(S) = \{ \$ \}$
- $\text{Follow}(B) = \{ \text{First}(D) - \epsilon \} \cup \text{First}(h) = \{ g, f, h \}$
- $\text{Follow}(C) = \text{Follow}(B) = \{ g, f, h \}$
- $\text{Follow}(D) = \text{First}(h) = \{ h \}$
- $\text{Follow}(E) = \{ \text{First}(F) - \epsilon \} \cup \text{Follow}(D) = \{ f, h \}$
- $\text{Follow}(F) = \text{Follow}(D) = \{ h \}$



First and Follow

- Example 4:

$$S \rightarrow ACB \mid CbB \mid Ba$$
$$A \rightarrow da \mid BC$$
$$B \rightarrow g \mid \varepsilon$$
$$C \rightarrow h \mid \varepsilon$$

First and Follow

| | | | |
|--------------------------------|---|-------------------------------|---------------------|
| $S \rightarrow ACB / CbB / Ba$ | 4 | $\{d, g, h, \epsilon, b, a\}$ | $\{\$ \}$ |
| $A \rightarrow da / BC$ | 3 | $\{d, g, h, \epsilon\}$ | $\{h, g, \$ \}$ |
| $B \rightarrow g / \epsilon$ | 1 | $\{g, \epsilon\}$ | $\{\$, a, h, g\}$ |
| $C \rightarrow h / \epsilon$ | 2 | $\{h, \epsilon\}$ | $\{g, \$, b, h, \}$ |



First and Follow

- Example 5:

$S \rightarrow aABb$

$A \rightarrow c \mid \varepsilon$

$B \rightarrow d \mid \varepsilon$



First and Follow

| | | |
|----------------------------|-------------------|------------|
| $S \rightarrow aABb$ | $\{a\}$ | $\{\$ \}$ |
| $A \rightarrow c \epsilon$ | $\{c, \epsilon\}$ | $\{d, b\}$ |
| $B \rightarrow d \epsilon$ | $\{d, \epsilon\}$ | $\{b\}$ |



First and Follow

- Example 6:

$S \rightarrow A$

$A \rightarrow aB \ / \ Ad$

$B \rightarrow b$

$C \rightarrow g$



First and Follow

- Left recursion eliminated :

$$S \rightarrow A$$

$$A \rightarrow aBA'$$

$$A' \rightarrow dA' \quad / \quad \varepsilon$$

$$B \rightarrow b$$

$$C \rightarrow g$$



First and Follow

$\text{First}(S) = \text{First}(A) = \{ a \}$

$\text{First}(A) = \{ a \}$

$\text{First}(A') = \{ d, \epsilon \}$

$\text{First}(B) = \{ b \}$

$\text{First}(C) = \{ g \}$



First and Follow

$$\text{Follow}(S) = \{ \$ \}$$

$$\text{Follow}(A) = \text{Follow}(S) = \{ \$ \}$$

$$\text{Follow}(A') = \text{Follow}(A) = \{ \$ \}$$

$$\text{Follow}(B) = \{ \text{First}(A') - \epsilon \} \cup$$

$$\text{Follow}(A) = \{ d, \$ \}$$

$$\text{Follow}(C) = \text{NA}$$



First and Follow

- Example 7:

$$S \rightarrow (L) \quad / \quad a$$
$$L \rightarrow SL'$$
$$L' \rightarrow , SL' \quad / \quad \epsilon$$



First and Follow

$$\text{First}(S) = \{ (, , a \}$$
$$\text{First}(L) = \text{First}(S) = \{ (, , a \}$$
$$\text{First}(L') = \{ , , \epsilon \}$$
$$\text{Follow}(S) = \{ \$ \} \cup$$
$$\{ \text{First}(L') - \epsilon \} \cup \text{Follow}(L)$$
$$\cup \text{Follow}(L') = \{ \$, , ,) \}$$
$$\text{Follow}(L) = \{) \}$$
$$\text{Follow}(L') = \text{Follow}(L) = \{) \}$$



First and Follow

- Example 8:

$$S \rightarrow AaAb \mid BbBa$$
$$A \rightarrow \varepsilon$$
$$B \rightarrow \varepsilon$$



First and Follow

$$\text{First}(S) = \{ \text{First}(A) - \epsilon \} \cup$$

$$\text{First}(a) \cup$$

$$\{ \text{First}(B) - \epsilon \} \cup$$

$$\text{First}(b) = \{ a, b \}$$

$$\text{First}(A) = \{ \epsilon \}$$

$$\text{First}(B) = \{ \epsilon \}$$



First and Follow

$\text{Follow}(S) = \{ \$ \}$

$\text{Follow}(A) = \text{First}(a) \cup \text{First}(b) = \{a, b\}$

$\text{Follow}(B) = \text{First}(b) \cup \text{First}(a) = \{a, b\}$



First and Follow

- Example 9:

$$E \rightarrow TE'$$
$$E' \rightarrow +TE' \mid \varepsilon$$
$$T \rightarrow FT'$$
$$T' \rightarrow *FT' \mid \varepsilon$$
$$F \rightarrow id \mid (E)$$

First and Follow

| | | | |
|-------------------------------------|---|---------------------|---------------------|
| $E \rightarrow TE'$ | 5 | $\{id, (\}$ | $\{ \$,) \}$ |
| $E' \rightarrow +TE' \mid \epsilon$ | 3 | $\{ +, \epsilon \}$ | $\{ \$,) \}$ |
| $T \rightarrow FT'$ | 4 | $\{ id, (\}$ | $\{ +, \$,) \}$ |
| $T' \rightarrow *FT' \mid \epsilon$ | 2 | $\{ *, \epsilon \}$ | $\{ +, \$,) \}$ |
| $F \rightarrow id \mid (E)$ | 1 | $\{ id, (\}$ | $\{ *, +, \$,) \}$ |

- In P#1, Follow(E') contains Follow(E) too !!!