

Design and Analysis of Algorithms

Prepared by Dr. Muhammad Imran

Modified by MN Khokhar

Pre-Conditions and Post-Conditions

- **Procedural Abstraction:**
 - Writing a complete operation, method or procedural declaration: Specify HOW.
 - Users do not need to know HOW.
 - Very powerful tool that allows programmers to consider operations one at a time rather than all together.

Pre-Conditions and Post-Conditions

- **Pre-Condition:**

- A statement or set of statements that outlines a condition that should be true when the operation is called.
- The operation is not guaranteed to perform as it should unless the pre-conditions have been met.

Pre-Conditions and Post-Conditions

- **Post-Condition:**
 - A statement or statements describing the condition that will be true when the operation has completed its task.
 - If the operation is correct and the pre-condition(s) met, then the post-condition is guaranteed to be true.

Correct Algorithms

- An algorithm is said to be correct if, for every input instance, it halts with the correct output. We say that a correct algorithm solves the given computational problem.
- An incorrect algorithm might not halt at all on some input instances, or it might halt with an answer other than the desired one.
- *Incorrect algorithms can sometimes be useful, if their error rate can be controlled. (An example of this when we study algorithms for finding large prime numbers.)*

What does an algorithm ?

- An algorithm is described by:
 - Input data
 - Output data
 - ***Preconditions***: specifies restrictions on input data
 - ***Postconditions***: specifies what is the result
- Example: Binary Search
 - Input data: `a:array of integer; x:integer;`
 - Output data: `found:boolean;`
 - Precondition: `a` is sorted in ascending order
 - Postcondition: `found` is true if `x` is in `a`, and `found` is false otherwise

Correct algorithms

- **An algorithm is correct if:**
 - for **any correct** input data:
 - it **stops** and
 - it produces **correct output**.
 - Correct input data: satisfies precondition
 - Correct output data: satisfies postcondition

Proving correctness

- An algorithm = a list of actions
- ***Proving that an algorithm is totally correct:***
 - 1. Proving that it will terminate***
 - 2. Proving that the list of actions applied to the precondition imply the postcondition***
- This is easy to prove for simple sequential algorithms
- This can be complicated to prove for repetitive algorithms (containing loops or recursivity)
 - use techniques based on ***loop invariants*** and ***induction***

Example – a sequential algorithm

Swap1 (x, y) :

 aux := x

 x := y

 y := aux

Precondition:

 x = a and y = b

Postcondition:

 x = b and y = a

Proof: *the list of actions applied to the precondition imply the postcondition*

1. Precondition:

 x = a and y = b

2. aux := x => aux = a

3. x := y => x = b

4. y := aux => y = a

5. x = b and y = a is
 the Postcondition

Example – a repetitive algorithm

Algorithm Sum_of_N_numbers

Input: a, an array of N numbers

Output: s, the sum of the N numbers
in a

s:=0;

k:=0;

While (k<N) do

 k:=k+1;

 s:=s+a[k];

end

***Proof: the list of actions
applied to the
precondition imply the
postcondition***

***BUT: we cannot enumerate
all the actions in case of
a repetitive algorithm !***

***We use techniques based on
loop invariants and
induction***