

## Recursive functions

$A(n) \{$   
     if (0) Nothing to count  
         return( $A(n/2) + A(n/2)$ )  
 $\}$

— Assume time  $T(n)$   
Some const. time  
Some const. time

So,  $T(n) = C + 2T(n/2)$  Recursive func.

$A(n) \{$   
     if ( $n > 1$ ) C or 1  
         return( $A(n-1)$ )  
 $\}$

$T(n) = 1 + T(n-1)$

slow

Back Substitution

$T(n) = 1 + T(n-1)$  — ①  
 $T(n-1) = 1 + T(n-2)$  — ②  
 $T(n-2) = 1 + T(n-3)$  — ③

Base/Anchor Condition

$T(1)$   
 $n-k=1$   
 $k=n-1$

Substitute ② in ①

$$T(n) = 1 + 1 + T(n-2) = 2 + T(n-2) \text{ — ④}$$

Substitute ③ in ④

$$T(n) = 2 + 1 + T(n-3) = 3 + T(n-3)$$

$$\dots k + T(n-k)$$

$$\dots (n-1) + T(n-(n-1))$$

$$= (n-1) + T(1) = n-1+1 = n$$

Stop condition: ~~n=0~~  $n=1$

$$O(n)$$

Ex:  $T(n) = n + T(n-1) ; n > 1$   
 $= 1 ; n = 1$

$$T(n) = n + T(n-1) \quad \text{--- ①}$$

$$T(n-1) = (n-1) + T(n-2) \quad \text{--- ②}$$

$$T(n-2) = (n-2) + T(n-3) \quad \text{--- ③}$$

...

Subs. ② in ①

$$T(n) = n + (n-1) + T(n-2) \quad \text{--- ④}$$

Subs. ③ in ④

$$T(n) = n + (n-1) + (n-2) + T(n-3)$$

...

$$T(n) = n + (n-1) + (n-2) + (n-3) + \dots + (n-k) + T(n-(k+1))$$

Base condition  $n-(k+1) = 1$

$$n-k-1 = 1$$

$$k = n-2$$

2

$$T(n) = n + (n-1) + (n-2) + \dots + (n-(n-2)) + T(n-(n-2+1))$$

$$= n + (n-1) + (n-2) + \dots + 2 + 1$$

$$= \frac{n(n+1)}{2}$$

$$O(n^2)$$

fact(int n) {

if (n=0)

return 1

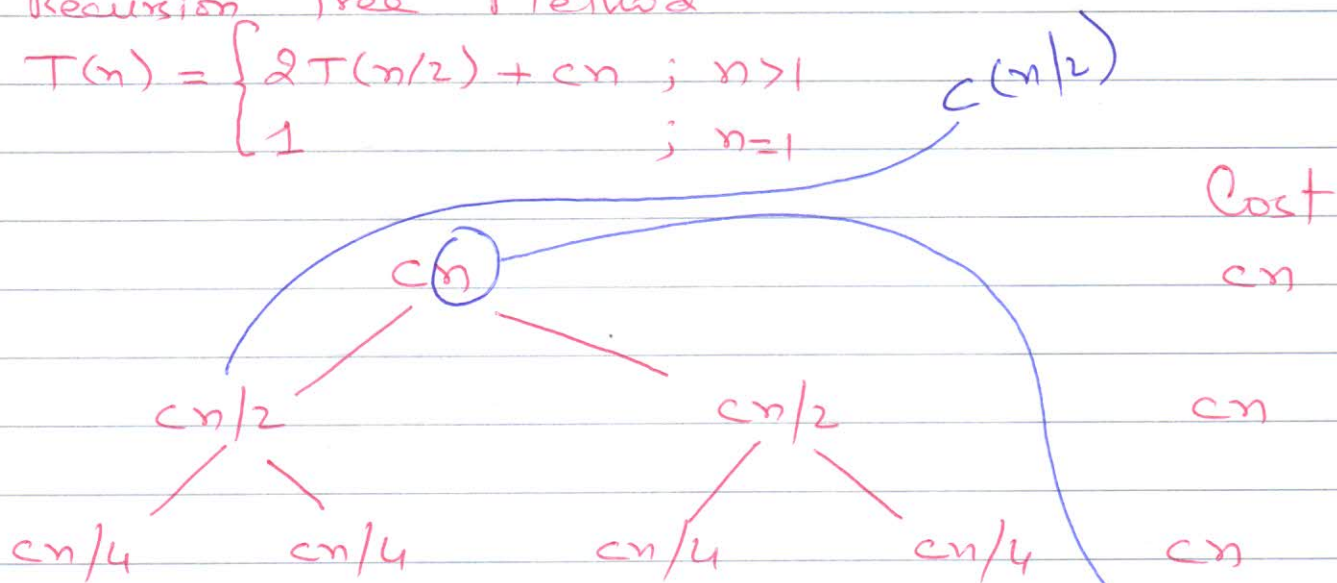
return n \* fact(n-1)

}

$$T(n) \begin{cases} n * T(n-1) & ; n > 0 \\ 1 & ; n = 1 \end{cases}$$

## Recursion Tree Method

$$T(n) = \begin{cases} 2T(n/2) + cn & ; n > 1 \\ 1 & ; n = 1 \end{cases}$$



Levels

0

1

2

3

⋮

L-1

Nodes

 $1 = 2^0$  $2 = 2^1$  $4 = 2^2$  $8 = 2^3$  $2^{L-1}$ 

$$2^{L-1} = n$$

$$L-1 = \log_2 n$$

$$L = \log_2 n + 1$$

Complexity = No of levels \* Cost per level

Cost

$$O(n \log_2 n)$$

$$= (\log_2 n + 1) * cn$$

$$= cn \log_2 n + cn$$

$$\left. \begin{aligned} T(n) &= 3T(n/3) + cn \\ L &= \log_3 n + 1 \end{aligned} \right\}$$