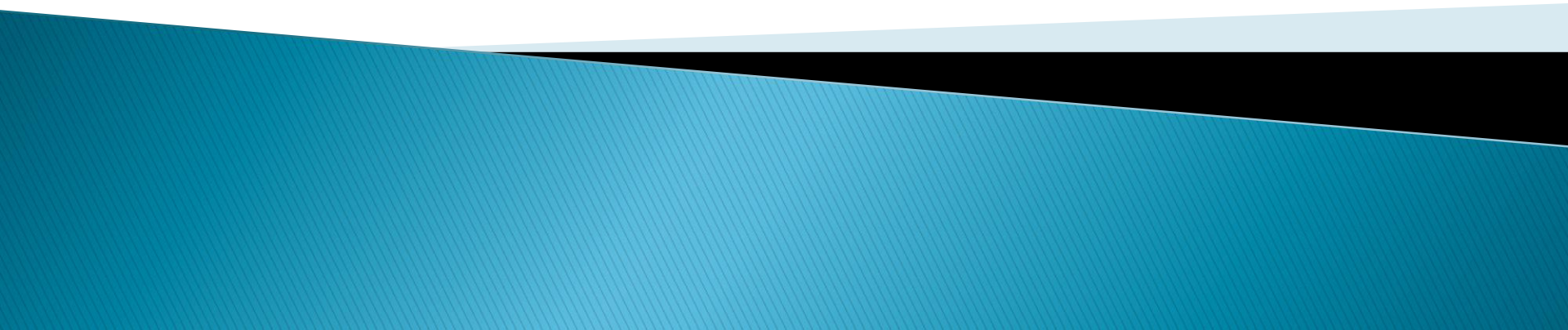


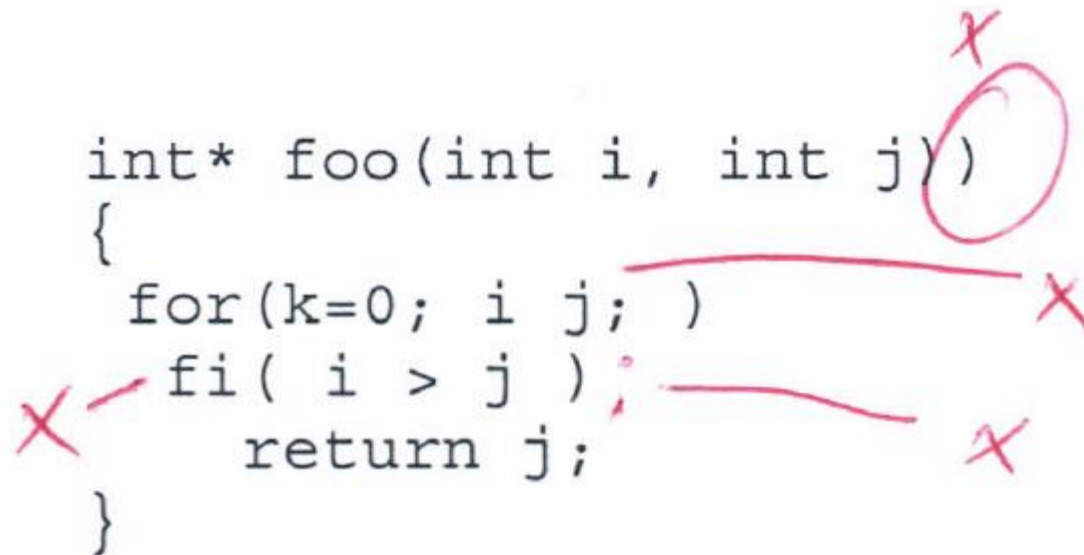
Compiler Construction

(Week 3, Lecture 2)



Syntactic Analysis

```
int* foo(int i, int j))  
{  
    for(k=0; i j; )  
    fi( i > j ) ;  
    return j;  
}
```



Semantic Analysis

```
int* foo(int i, int j)
{
    for(k=0; i < j; j++ )
        if( i < j-2 )
            sum = sum+i
    return sum;
}
```

Handwritten annotations in red:

- A large oval encircling the entire function definition.
- An oval around `int*` in the return type.
- An oval around `k=0` in the for loop.
- An oval around `sum` in the assignment statement.
- An oval around `sum` in the return statement.
- Two red question marks with underlines on the left margin, each with a line pointing to the `sum` variable in the assignment and return statements respectively.

Parser

1. $\text{expr} \rightarrow \text{expr op expr}$
2. | num
3. | id
4. $\text{Op} \rightarrow +$
5. | $-$
6. | $*$
7. | $/$

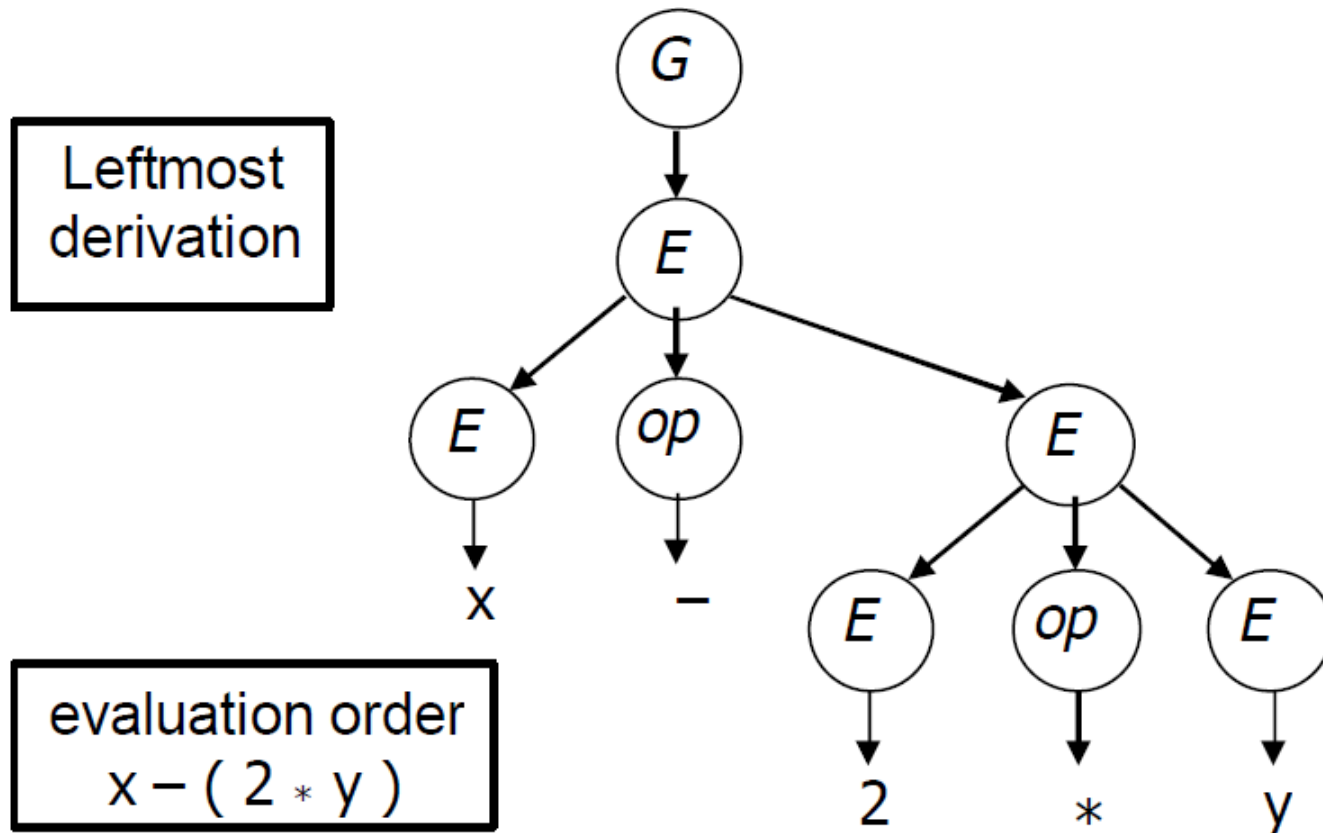
Parser

Rule	Sentential Form
–	expr
1	expr op expr
3	<id,x> op expr
5	<id,x> – expr
1	<id,x> – expr op expr
2	<id,x> – <num,2> op expr
6	<id,x> – <num,2> * expr
3	<id,x> – <num,2> * <id,y>

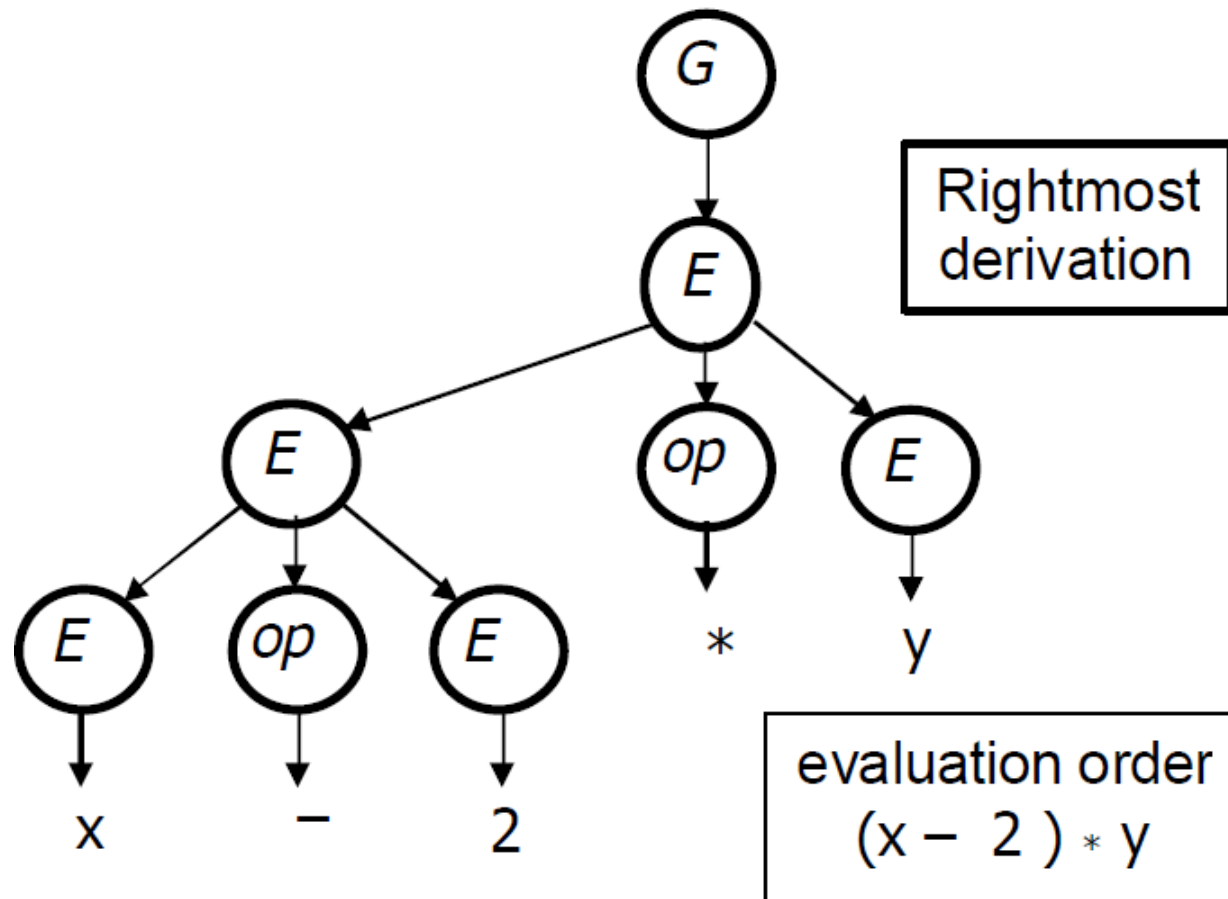
Parser

Rule	Sentential Form
–	expr
1	expr op expr
3	expr op <id,y>
6	expr * <id,y>
1	expr op expr * <id,y>
2	expr op <num,2> * <id,y>
5	expr – <num,2> * <id,y>
2	<id,x> – <num,2> * <id,y>



Parser



Parser



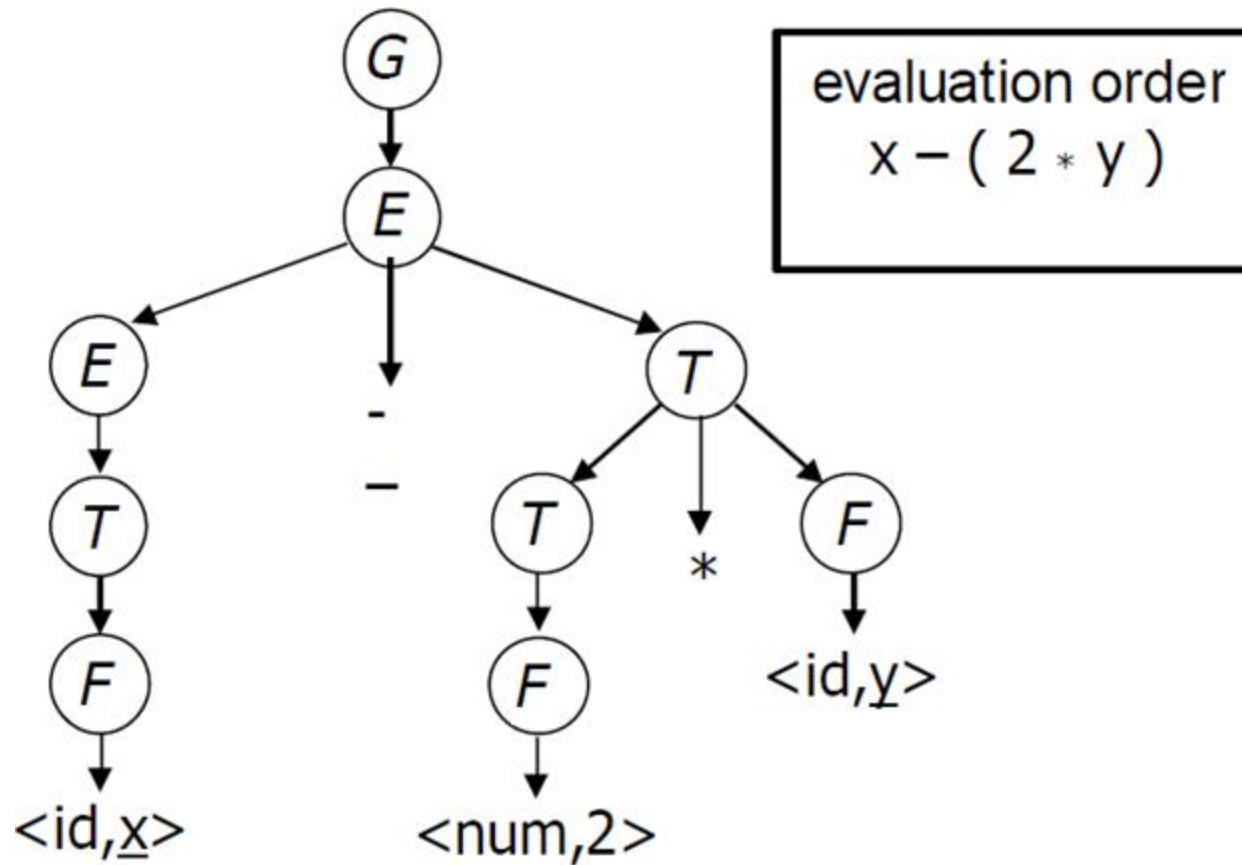
Precendence

1	goal	→	expr	
2	expr	→	expr + term	 Level 2
3			expr - term	
4			term	
5	term	→	term * factor	 Level 1
6			term / factor	
7			factor	
8	factor	→	number	
9			id	

Parser

Rule	Sentential Form
-	goal
1	expr
3	expr - term
5	expr - term * factor
9	expr - term * <id,y>
7	expr - factor * <id,y>
8	expr - <num,2> * <id,y>
4	term - <num,2> * <id,y>
7	factor - <num,2> * <id,y>
9	<id,x> - <num,2> * <id,y>

Parser



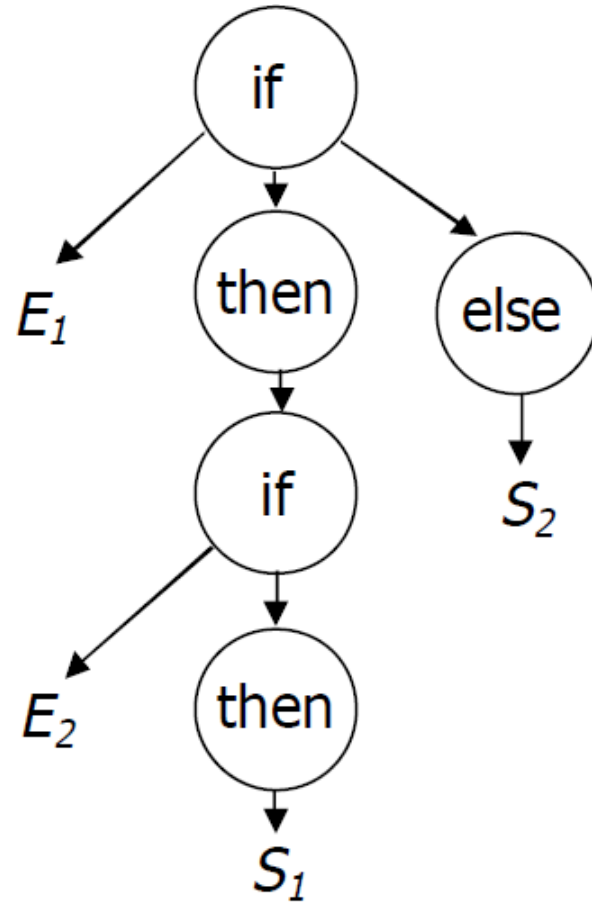
Ambiguity

```
stmt → if expr then stmt  
      | if expr then stmt else stmt  
      | ... other stmts ...
```

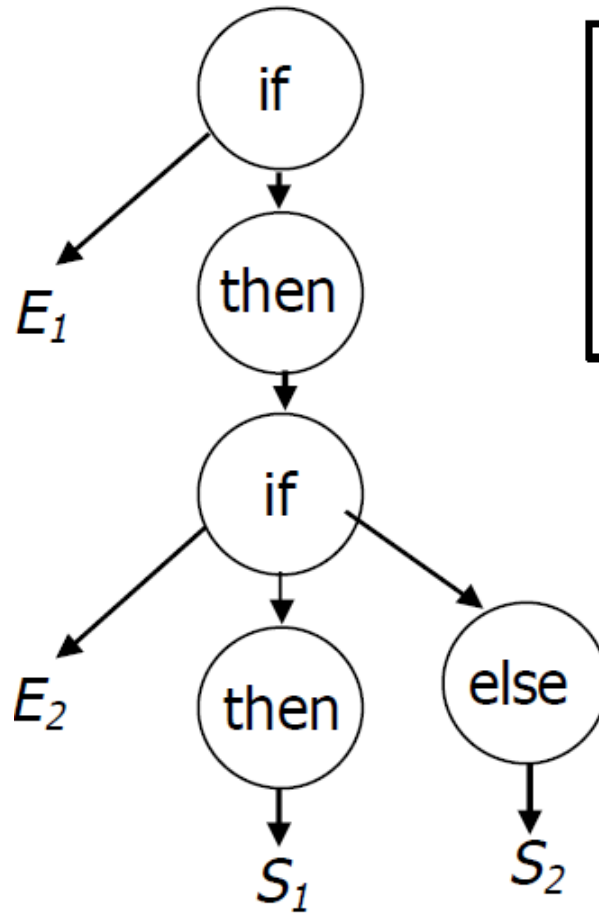
Two possible derivations for
if E1 then if E2 then S else S2

Ambiguity

Production 1, then
Production 2:
if E_1 then
 if E_2 then S_1
else S_2



Ambiguity



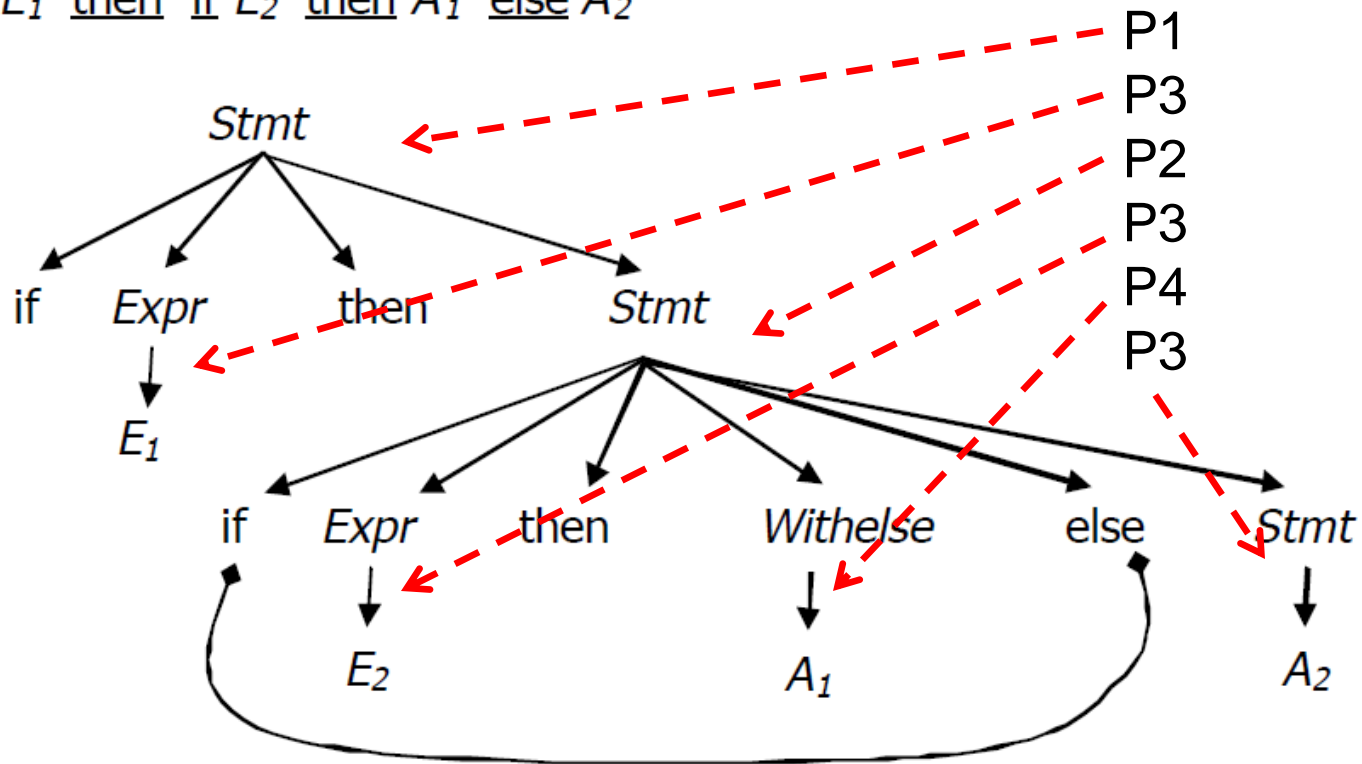
Production 2, then
Production 1:
if E_1 then
 if E_2 then S_1
 else S_2

Ambiguity

```
1 stmt      → if E then stmt
2           | if E then withelse else stmt
3           | ... other stmt ...
4 withelse  → if E then withelse else withelse
5           | ... other stmt ...
```

Ambiguity

if E_1 then if E_2 then A_1 else A_2



This binds the else controlling A_2 to inner if