
**CORRECTNESS OF:
LOOP TO COMPUTE A PRODUCT
THE DIVISION ALGORITHM
THE EUCLIDEAN ALGORITHM**

A LOOP TO COMPUTE A PRODUCT:

[pre-condition: m is a nonnegative integer,
 x is a real number, $i = 0$, and product = 0.]

while ($i \neq m$)
 1. product := product + x
 2. $i := i + 1$

end while

[post-condition: product = $m \cdot x$]

PROOF:

Let the loop invariant be

$I(n)$: $i = n$ and product = $n \cdot x$

The guard condition G of the while loop is

G : $i \neq m$

I. Basis Property:

[$I(0)$ is true before the first iteration of the loop.]

$I(0)$: $i = 0$ and product = $0 \cdot x = 0$

Which is true before the first iteration of the loop.

II. Inductive property:

[If the guard G and the loop invariant $I(k)$ are both true before a loop iteration (where $k \geq 0$), then $I(k + 1)$ is true after the loop iteration.]

Before execution of statement 1,

$$product_{old} = k \cdot x.$$

Thus the execution of statement 1 has the following effect:

$$product_{new} = product_{old} + x = k \cdot x + x = (k + 1) \cdot x$$

Similarly, before statement 2 is executed,

$$i_{old} = k,$$

So after execution of statement 2,

$$i_{new} = i_{old} + 1 = k + 1.$$

Hence after the loop iteration, the statement $I(k + 1)$ (i.e., $i = k + 1$ and product = $(k + 1) \cdot x$) is true. This is what we needed to show.

III. Eventual Falsity of Guard:

[After a finite number of iterations of the loop, the guard becomes false.]

IV. Correctness of the Post-Condition:

[If N is the least number of iterations after which G is false and $I(N)$ is true, then the values of the algorithm variables will be as specified in the post-condition of the loop.]

THE DIVISION ALGORITHM:

[pre-condition: a is a nonnegative integer and d is a positive integer, $r = a$, and $q = 0$]

while ($r \geq d$)

1. $r := r - d$
2. $q := q + 1$

end while

[post-condition: q and r are nonnegative integers with the property that $a = q \cdot d + r$ and $0 \leq r < d$.]

PROOF:

Let the loop invariant be

$$I(n): r = a - n \cdot d \text{ and } n = q.$$

The guard of the **while** loop is

$$G: r \geq d$$

I. Basis Property:

[$I(0)$ is true before the first iteration of the loop.]

$$I(0): r = a - 0 \cdot d = a \text{ and } 0 = q.$$

II. Inductive property:

[If the guard G and the loop invariant $I(k)$ are both true before a loop iteration (where $k \geq 0$), then $I(k + 1)$ is true after the loop iteration.]

$$I(k): r = a - k \cdot d \geq 0 \text{ and } k = q$$

$$I(k + 1): r = a - (k + 1) \cdot d \geq 0 \text{ and } k + 1 = q$$

$$\begin{aligned} r_{\text{new}} &= r - d \\ &= a - k \cdot d - d \\ &= a - (k + 1) \cdot d \\ q &= q + 1 \\ &= k + 1 \end{aligned}$$

also

$$\begin{aligned} r_{\text{new}} &= r - d \\ &\geq d - d = 0 \quad (\text{since } r \geq 0) \end{aligned}$$

Hence $I(k + 1)$ is true.

III. Eventual Falsity of Guard:

[After a finite number of iterations of the loop, the guard becomes false.]

IV. Correctness of the Post-Condition:

[If N is the least number of iterations after which G is false and $I(N)$ is true, then the values of the algorithm variables will be as specified in the post-condition of the loop.]

G is false and $I(N)$ is true.

That is, $r \geq d$ and $r = a - N \cdot d \geq 0$ and $N = q$.

or $r = a - q \cdot d$

or $a = q \cdot d + r$

Also combining the two inequalities involving r we get

$$0 \leq r < d$$

THE EUCLIDEAN ALGORITHM:

The greatest common divisor (gcd) of two integers a and b is the largest integer that divides both a and b . For example, the gcd of 12 and 30 is 6.

The Euclidean algorithm takes integers A and B with $A > B \geq 0$ and compute their greatest common divisor.

HAND CALCULATION OF gcd:

Use the Euclidean algorithm to find gcd(330, 156)

SOLUTION:

$$\begin{array}{r} 2 \\ 156 \overline{) 330} \\ \underline{312} \\ 18 \end{array} \qquad \begin{array}{r} 8 \\ 18 \overline{) 156} \\ \underline{144} \\ 12 \end{array}$$
$$\begin{array}{r} 1 \\ 12 \overline{) 18} \\ \underline{12} \\ 6 \end{array} \qquad \begin{array}{r} 2 \\ 6 \overline{) 12} \\ \underline{12} \\ 0 \end{array}$$

Hence gcd(330, 156) = 6

EXAMPLE:

Use the Euclidean algorithm to find gcd(330, 156)

Solution:

1. Divide 330 by 156:

$$\text{This gives } 330 = 156 \cdot 2 + 18$$

2. Divide 156 by 18:

$$\text{This gives } 156 = 18 \cdot 8 + 12$$

3. Divide 18 by 12:

$$\text{This gives } 18 = 12 \cdot 1 + 6$$

4. Divide 12 by 6:

$$\text{This gives } 12 = 6 \cdot 2 + 0$$

Hence $\gcd(330, 156) = 6$.

LEMMA:

If a and b are any integers with $b \neq 0$ and q and r are nonnegative integers such that

$$a = q \cdot b + r$$

then

$$\gcd(a, b) = \gcd(b, r)$$

[pre-condition: A and B are integers with

$A > B \geq 0, a = A, b = B, r = B.$]

while ($b \neq 0$)

1. $r := a \bmod b$

2. $a := b$

3. $b := r$

end while[post-condition: $a = \gcd(A, B)$]

PROOF:

Let the **loop invariant** be

$$I(n): \gcd(a, b) = \gcd(A, B) \text{ and } 0 \leq b < a.$$

The guard of the **while** loop is

$$G: b \neq 0$$

I. Basis Property:

[$I(0)$ is true before the first iteration of the loop.]

$$I(0): \gcd(a, b) = \gcd(A, B) \text{ and } 0 \leq b < a.$$

According to the precondition,

$$a = A, b = B, r = B, \text{ and } 0 \leq B < A.$$

Hence $I(0)$ is true before the first iteration of the loop.

II. Inductive property:

[If the guard G and the loop invariant $I(k)$ are both true before a loop iteration (where $k \geq 0$), then $I(k + 1)$ is true after the loop iteration.]

Since $I(k)$ is true before execution of the loop we have,

$$\gcd(a_{\text{old}}, b_{\text{old}}) = \gcd(A, B) \text{ and } 0 \leq b_{\text{old}} < a_{\text{old}}$$

After execution of statement 1,

$$r_{\text{new}} = a_{\text{old}} \bmod b_{\text{old}} \text{ Thus,}$$

$$a_{\text{old}} = b_{\text{old}} \cdot q + r_{\text{new}} \quad \text{for some integer } q$$

with,

$$0 \leq r_{\text{new}} < b_{\text{old}}.$$

But

$$\gcd(a_{\text{old}}, b_{\text{old}}) = \gcd(b_{\text{old}}, r_{\text{old}})$$

and we have,

$$\gcd(b_{\text{old}}, r_{\text{new}}) = \gcd(A, B)$$

When statements 2 and 3 are executed,

$$a_{\text{new}} = b_{\text{old}} \text{ and } b_{\text{new}} = r_{\text{new}}$$

It follows that

$$\gcd(a_{\text{new}}, b_{\text{new}}) = \gcd(A, B)$$

Also,

$$0 \leq r_{\text{new}} < b_{\text{old}}$$

becomes

$$0 \leq b_{\text{new}} < a_{\text{new}}$$

Hence $I(k + 1)$ is true.

III.Eventual Falsity of Guard:

[After a finite number of iterations of the loop, the guard becomes false.]

IV.Correctness of the Post-Condition:

[If N is the least number of iterations after which G is false and $I(N)$ is true, then the values of the algorithm variables will be as specified in the post-condition of the loop.]