# CSC 3201 Compiler Construction

Department of Computer Science SZABIST (Islamabad Campus)

Week 6 (Lecture 1)



- First(α) is a set of terminal symbols that begin in strings derived from α.
- Example:
  - A → abc | def | ghi
  - First(A) =  $\{a,d,g\}$

# First

#### Rules:

- For a production rule  $X \to \varepsilon$ , First(X) = {  $\varepsilon$  }
- For any terminal symbol 'a', First(a) = { a }
- For a production rule X → Y1Y2Y3, Calculating First(X)
  - If ε ∉ First(Y1), then First(X) = First(Y1)
  - If ε ε First(Y1), then First(X) = { First(Y1) − ε } ∪ First(Y2Y3), Calculating First(Y2Y3)
    - If ξ ∉ First(Y2), then First(Y2Y3) = First(Y2)
    - If ε ε First(Y2), then First(Y2Y3) = { First(Y2) − ε } ∪ First(Y3)

# Follow

- Follow(α) is a set of terminal symbols that appear immediately to the right of α.
- Rules:
  - For the start symbol S, place \$ in Follow(S)
  - For any production rule A → αB, Follow(B) = Follow(A)
  - For any production rule A → αBβ,
    - If ε ∉ First(β), then Follow(B) = First(β)
    - If ε ε First(β), then Follow(B) = { First(β) ε } ∪ Follow(A)



#### Note:

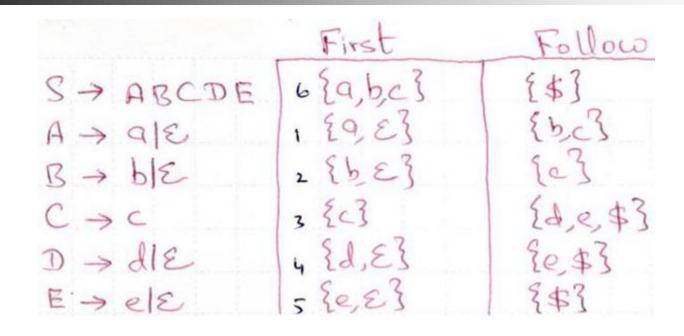
- ξ may appear in the first function of an NT.
- ξ will never appear in the follow function of an NT.
- Before calculating the first and follow functions, eliminate Left Recursion from the grammar.
- Calculate the follow function of an NT by looking where it is present on the RHS of a production rule.



	First	Follow
S -> ABCDE		
A → 9/E	1 29, 23	{bc}
B > b/E	2 { b E }	903
$C \rightarrow c$	3 { 6 }	{d,e,\$3
D > dle	4 {d, E}	Ee.\$3
E > e/E	5 Ee, E3	3\$3

- Follow(RHS) = Follow(LHS)
- Look at the occurrence on rHS in all productions.





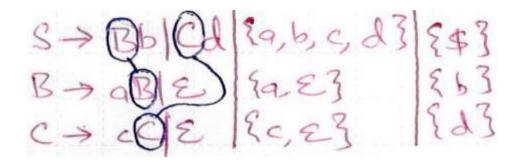
S is not in the RHS of any production



#### Example 2:

```
S \rightarrow Bb \mid Cd
B \rightarrow aB \mid \xi
C \rightarrow cC \mid \xi
```





# First and Follow

#### Example 3:

```
S \rightarrow aBDh

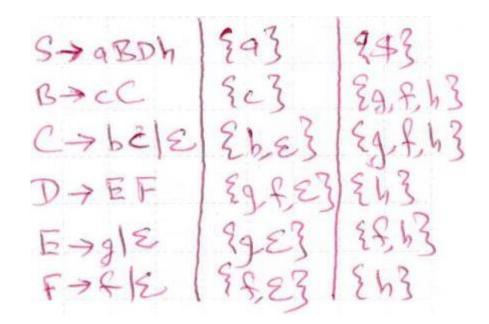
B \rightarrow cC

C \rightarrow bc \mid \epsilon

D \rightarrow EF

E \rightarrow g \mid \epsilon

F \rightarrow f \mid \epsilon
```



- First(S) = {a}
- First(B) = {c}
- First(C) =  $\{b, \epsilon\}$
- First(D) = { First(E) − ε } ∪ First(F) = {g,f,ε}
- First(E) = { g , ε }
- First(F) = { f , ε }

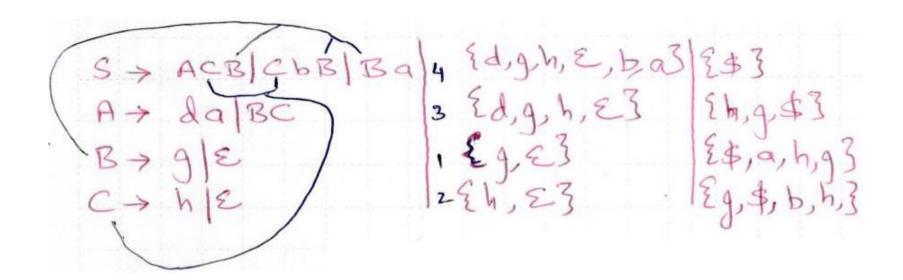
- Follow(S) = { \$ }
- Follow(C) = Follow(B) = { g , f , h }
- Follow(D) = First(h) = { h }
- Follow(F) = Follow(D) = { h }

# •

# First and Follow

#### Example 4:

```
S \rightarrow ACB \mid CbB \mid Ba
A \rightarrow da \mid BC
B \rightarrow g \mid \epsilon
C \rightarrow h \mid \epsilon
```





#### Example 5:

```
S \rightarrow aABb
A \rightarrow c \mid \epsilon
B \rightarrow d \mid \epsilon
```

# •



#### Example 6:

```
S \rightarrow A
A \rightarrow aB / Ad
B \rightarrow b
C \rightarrow g
```

# •

# First and Follow

#### Left recursion eliminated :

$$S \rightarrow A$$
 $A \rightarrow aBA'$ 
 $A' \rightarrow dA' / \epsilon$ 
 $B \rightarrow b$ 
 $C \rightarrow g$ 

```
First(S) = First(A) = { a }

First(A) = { a }

First(A') = { d , a }

First(B) = { b }

First(C) = { g }
```

```
Follow(S) = { $ }

Follow(A) = Follow(S) = { $ }

Follow(A') = Follow(A) = { $ }

Follow(B) = { First(A') - $ }

Follow(A) = { d , $ }

Follow(C) = NA
```



#### Example 7:

```
S \rightarrow (L) / a
L \rightarrow SL'
L' \rightarrow SL' / \epsilon
```



```
First(S) = \{ (, a) \}
First(L) = First(S) = \{ (, a) \}
First(L') = \{ , , \epsilon \}
Follow(S) = \{ \$ \} U
     {First(L') - \(\mathbb{E}\)} U Follow(L)
     U \text{ Follow}(L') = \{\$, , , \}
Follow(L) = \{ \}
Follow(L') = Follow(L) = { }
```



#### Example 8:

```
S \rightarrow AaAb \mid BbBa

A \rightarrow \epsilon

B \rightarrow \epsilon
```



```
First(S) = { First(A) - \( \epsilon \) \( \text{First(A)} \) \( \text{First(B)} - \( \epsilon \) \( \text{First(B)} - \( \epsilon \) \( \text{First(b)} = \{ a , b \} \) \( \text{First(A)} = \{ \( \epsilon \) \\ \( \epsilon \) \( \e
```



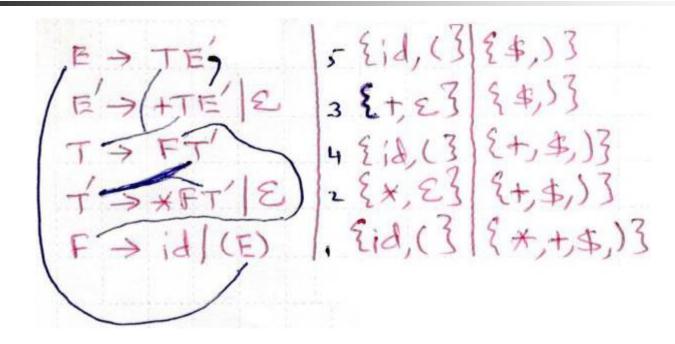
```
Follow(S) = { $ }
Follow(A) = First(a) U First(b) = {a,b}
Follow(B) = First(b) U First(a) = {a,b}
```

# First and Follow

#### Example 9:

```
E \rightarrow TE'
E' \rightarrow +TE' \mid \epsilon_{l}
T \rightarrow FT'
T' \rightarrow *FT' \mid \epsilon_{l}
F \rightarrow id \mid (E)
```





In P#1, Follow(E') contains Follow(E) too !!!