

Theory of Programming Languages

Example Problems

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Write a grammar for all unsigned numbers

- Unsigned numbers are strings such as 5280, 39.37, 86E9, 6.336E4, 1.894E-3, 2.3478E+11.
- The E is essentially scientific notation, stands for exponent
- You are not allowed to have a naked decimal point.
- For example, the number 5. or .5 is not allowed.
- It must be 0.5 or 5.0 .
- The fractional part, the exponential part, and the sign on the exponential part are optional.

Grammar for all unsigned numbers

<code><number></code>	<code>--></code>	<code><digit-list> <opt-fraction> <opt-exp></code>
<code><digit-list></code>	<code>--></code>	<code><digit> <digit> <digit-list></code>
<code><digit></code>	<code>--></code>	<code>0 1 2 ... 9</code>
<code><opt-fraction></code>	<code>--></code>	<code>. <digit-list> empty</code>
<code><opt-exp></code>	<code>--></code>	<code>E <opt-sign> <digit-list> empty</code>
<code><opt-sign></code>	<code>--></code>	<code>+ - empty</code>

C switch statement

C switch statement

```
e.g., switch (ch) {  
    case 'a': ...  
    case 'b': ...  
    default: ...    // optional  
}
```

Example: C switch statement

```
<switch>          --> switch ( <expr> ) { <case-list>
<opt-default> <case-list> }
<opt-default> --> empty | default: <statement>
<case-list>    --> empty | <one-case> <case-list>
<one-case>     --> case <const-value>: <statement>
<statement>    --> <block> | <assignment> | <if> |
...
<const-value> --> <int-const-value> | <char-const-
  value> | ...
<expr>        --> ...
```

Example: EBNF of a C switch statement

$\langle \text{switch_stmt} \rangle \rightarrow \textbf{switch} (\langle \text{expr} \rangle) \{ \textbf{case} \langle \text{literal} \rangle : \langle \text{stmt_list} \rangle$
 $\quad \{ \textbf{case} \langle \text{literal} \rangle : \langle \text{stmt_list} \rangle \} [\textbf{default} : \langle \text{stmt_list} \rangle] \}$

C- float literals (constants)

`<float-literal> --> <real> <suffix>`

`| <real> <exponent> <suffix>`

`| <integer> <exponent> <suffix>`

`<exponent> --> 'e' + <integer> (note: not empty, the letter 'e')`

`| 'e' - <integer>`

`| 'e' <integer>`

`| E + <integer>`

`| E - <integer>`

`| E <integer>`

`<integer> --> <digit> | <digit> <integer>`

`<real> --> .<integer>`

`| <integer> .`

`| <integer> . <integer>`

`<suffix> --> f | F | empty`

BNF of following example to give + precedence over * and force + to be right associative

Original

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\langle \text{id} \rangle \rightarrow A | B | C$

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle$
 | $\langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle * \langle \text{factor} \rangle$
 | $\langle \text{factor} \rangle$

$\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 | $\langle \text{id} \rangle$

$A = A * (B + (C * A))$

BNF of following example to give $+$ precedence over $*$ and force $+$ to be right associative

Modified

<assign> → <id> = <expr> **A=A*(B+(C*A))**

<id> → A | B | C

$$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle$$
$$| \langle \text{term} \rangle$$
$$\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle + \langle \text{term} \rangle$$
$$| \langle \text{factor} \rangle$$

```
<factor> → ( <expr> )
           | <id>
```

Left most derivation for $A=A*(B+(C*A))$
using following grammar

$$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$$

<id>-> A | B | C

```
<expr>-> <id>+<expr>
          | <id>*<expr>
          | (<expr>)
          | <id>
```

$$A = A * (B + (C * A))$$

`<assign> => <id> = <expr>`

`=> A = <expr>`

`=> A = <id> * <expr>`

`=> A = A * <expr>`

`=> A = A * (<expr>)`

`=> A = A * (<id> + <expr>)`

`=> A = A * (B + <expr>)`

`=> A = A * (B + (<expr>))`

`=> A = A * (B + (<id> * <expr>))`

`=> A = A * (B + (C * <expr>))`

`=> A = A * (B + (C * <id>))`

`=> A = A * (B + (C * A))`

Parse tree for $A=A*(B+(C*A))$ using following grammar

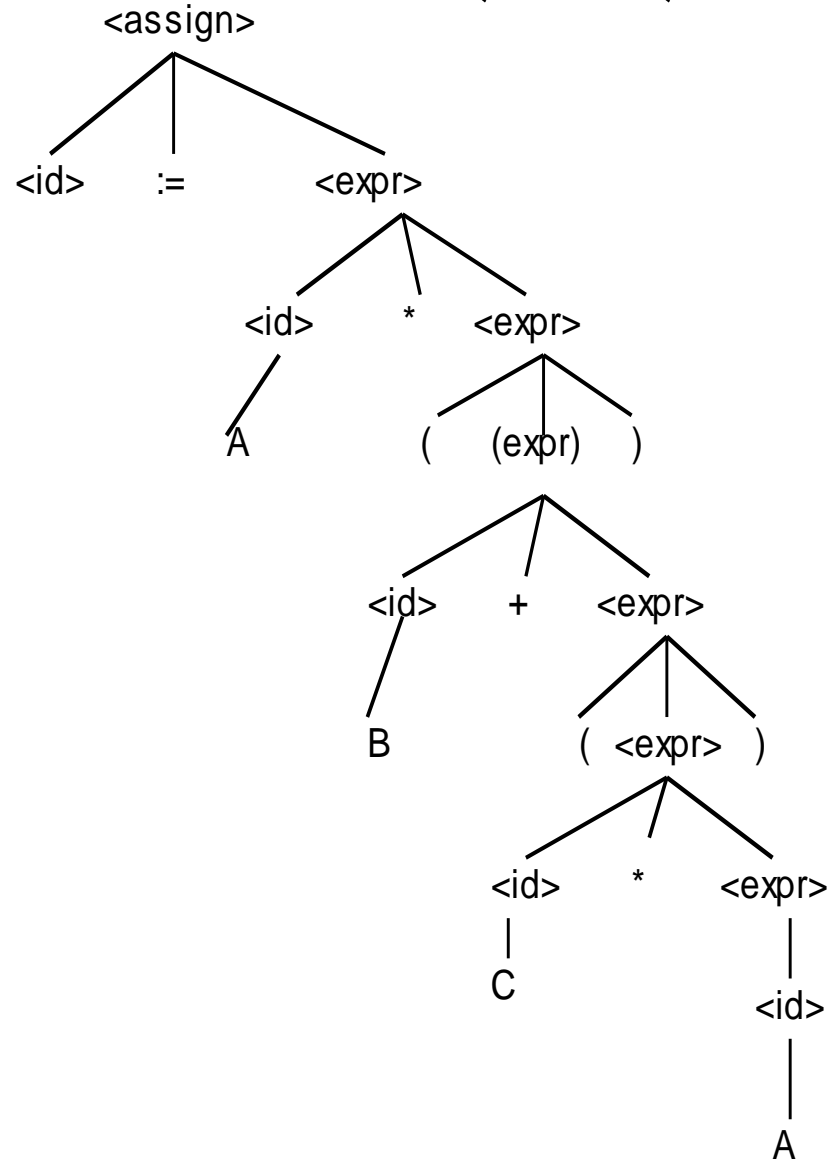
$$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$$
$$\langle id \rangle \rightarrow A | B | C$$

```

<expr>->  <id>+<expr>
           |  <id>*<expr>
           |  (<expr>)
           |  <id>

```

Parse tree for $A = A * (B + (C * A))$



Left most derivation for $A=(A+B)*C$ using following grammar

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\langle \text{id} \rangle \rightarrow A | B | C$

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle * \langle \text{factor} \rangle$
 $\quad \quad \quad | \langle \text{factor} \rangle$

$\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad \quad | \langle \text{id} \rangle$

Left most derivation for $A=(A+B)*C$

$\langle \text{assign} \rangle \Rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\Rightarrow A = \langle \text{expr} \rangle$

$\Rightarrow A = \langle \text{term} \rangle$

$\Rightarrow A = \langle \text{factor} \rangle * \langle \text{term} \rangle$

$\Rightarrow A = (\langle \text{expr} \rangle) * \langle \text{term} \rangle$

$\Rightarrow A = (\langle \text{expr} \rangle + \langle \text{term} \rangle) * \langle \text{term} \rangle$

$\Rightarrow A = (\langle \text{term} \rangle + \langle \text{term} \rangle) * \langle \text{term} \rangle$

$\Rightarrow A = (\langle \text{factor} \rangle + \langle \text{term} \rangle) * \langle \text{term} \rangle$

Left most derivation for $A=(A+B)*C$

$\Rightarrow A = (\text{<id>} + \text{<term>}) * \text{<term>}$

$\Rightarrow A = (A + \text{<term>}) * \text{<term>}$

$\Rightarrow A = (A + \text{<factor>}) * \text{<term>}$

$\Rightarrow A = (A + \text{<id>}) * \text{<term>}$

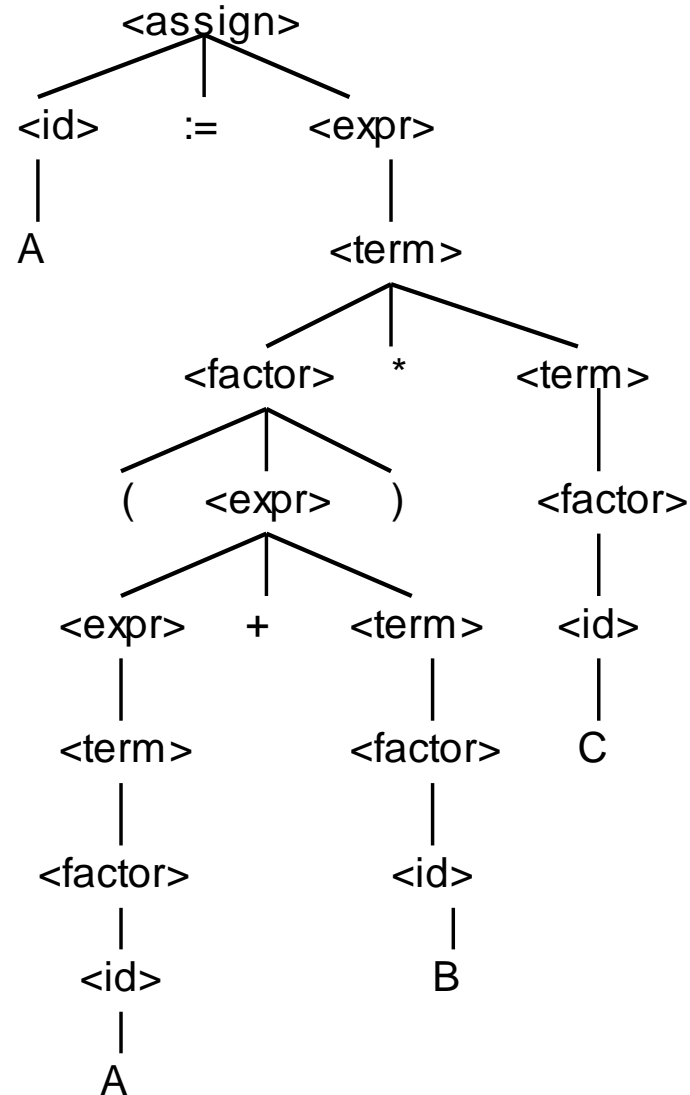
$\Rightarrow A = (A + B) * \text{<term>}$

$\Rightarrow A = (A + B) * \text{<factor>}$

$\Rightarrow A = (A + B) * \text{<id>}$

$\Rightarrow A = (A + B) * C$

Parse tree for $A = (A + B) * C$



Modify the grammar to add a unary minus operator that has higher precedence than + or *

Original

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\langle \text{id} \rangle \rightarrow A | B | C$

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle$
 $\quad \quad | \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle * \langle \text{factor} \rangle$
 $\quad \quad | \langle \text{factor} \rangle$

$\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad | \langle \text{id} \rangle$

Grammar having a unary minus operator that has higher precedence than + or *

Modified

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\langle \text{id} \rangle \rightarrow A | B | C$

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle$
 $| \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle * \langle \text{factor} \rangle$
 $| \langle \text{factor} \rangle$

$\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $| + \langle \text{id} \rangle$
 $| - \langle \text{id} \rangle$