

# CSC 2204 Finite Automata Theory and Formal Languages



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Department of Computer Science  
SZABIST (Islamabad Campus)

Week 9 (Lecture 1)

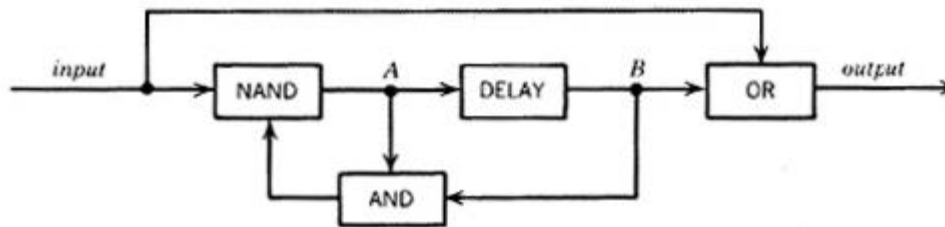


# Transducers as Models of Sequential Circuits

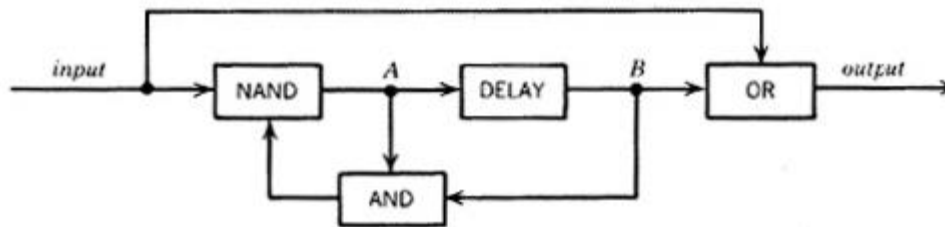
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- Automata with input and output are sometimes called “Transducers” because of their connection to electronics.

# Example 1



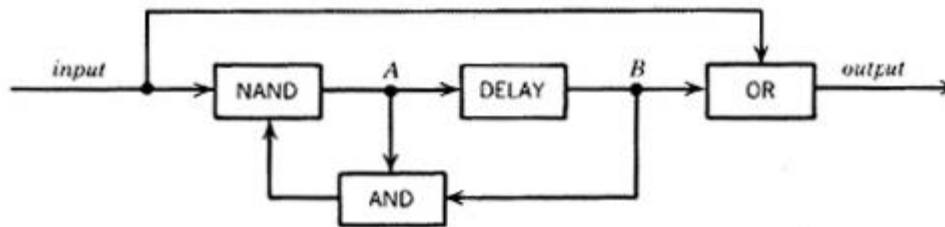
# Example 1



The following four types of boxes are used in this circuit

1. NAND box (NOT AND): For the given input, it provides the complement of Boolean AND output.
2. DELAY box (Flip Flop box): It delays the transmission of signal along the wire by one step (clock pulse).
3. OR box: For the given input, it provides the Boolean OR output.
4. AND box: For the given input, it provides the Boolean AND output.

# Example 1



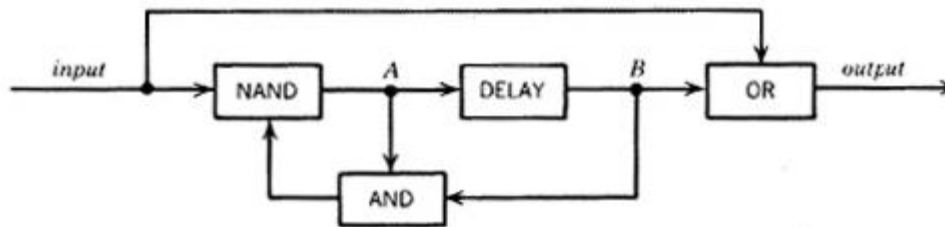
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There are two points A and B w.r.t. to which following four states of the machine are identified according to the presence and absence of current at these points *i.e.*

- 1)  $q_0(A=0, B=0) \equiv (0,0)$
- 2)  $q_1(A=0, B=1) \equiv (0,1)$
- 3)  $q_2(A=1, B=0) \equiv (1,0)$
- 4)  $q_3(A=1, B=1) \equiv (1,1)$

# Example 1



The operation of the circuit is such that the machine changes its state after reading 0 or 1. The transitions are determined using the following relations

new B = old A

new A = (input) NAND (old A AND old B)

output = (input) OR (old B)

It is to be noted that old A and old B indicate the presence or absence of current at A and B before inputting any letter. Similarly new A and new B indicate the presence or absence of current after reading certain letter. At various discrete pulses of a time clock, input is received by the machine and the corresponding output string is generated.



# Example 1

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The transition at the state  $q_0$  after reading the letter 0, can be determined, along with the corresponding output character as under

$$\begin{aligned}\text{new B} &= \text{old A} = 0 \\ \text{new A} &= (\text{input}) \text{ NAND } (\text{old A AND old B}) \\ &= 0 \text{ NAND } (0 \text{ AND } 0) = 0 \text{ NAND } 0 \\ &= 1 \\ \text{output} &= (\text{input}) \text{ OR } (\text{old B}) = 0 \text{ OR } 0 = 0\end{aligned}$$

Thus after reading 0 at  $q_0$  new B is 0 and new A is 1 i.e. machine will be at state  $(1,0) \equiv q_2$  and during this process its output character will be 0.



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Thus after reading 0 at  $q_0$  new B is 0 and new A is 1 i.e. machine will be at state  $(1,0) \equiv q_2$  and during this process it's output character will be 0.

Old state	Inputting 0		Inputting 1	
	State	Output	State	Output
$q_0 \equiv (0,0)$	$(1,0) \equiv q_2$	0	$(1,0) \equiv q_2$	1





# Example 1

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Old state	Inputting 0		Inputting 1	
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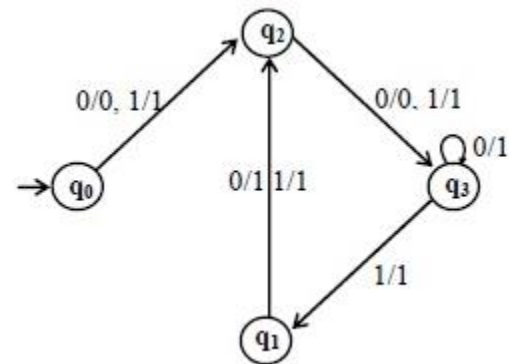
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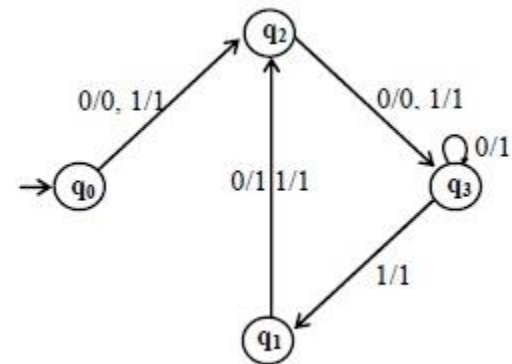
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Input		0	1	1	0	1	1	1	0
States	$q_0$	$q_2$	$q_3$	$q_1$	$q_2$	$q_3$	$q_1$	$q_2$	$q_3$
output		0	1	1	1	1	1	1	0