CSC 2204 Finite Automata Theory and Formal Languages

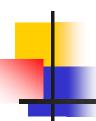
Department of Computer Science SZABIST (Islamabad Campus)

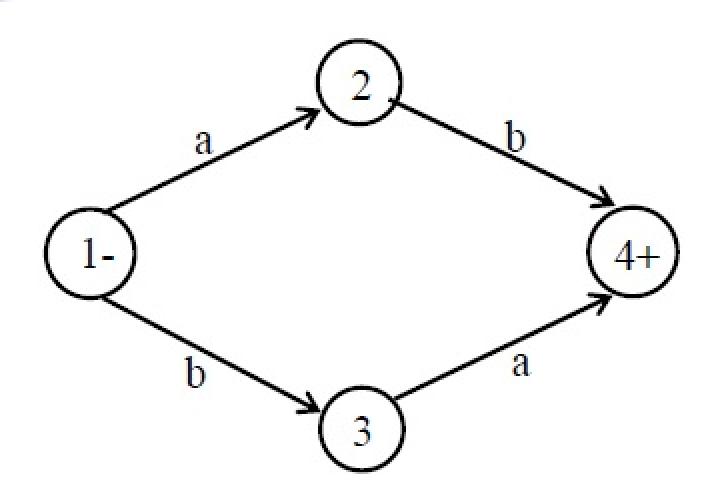
Week 4 (Lecture 1)



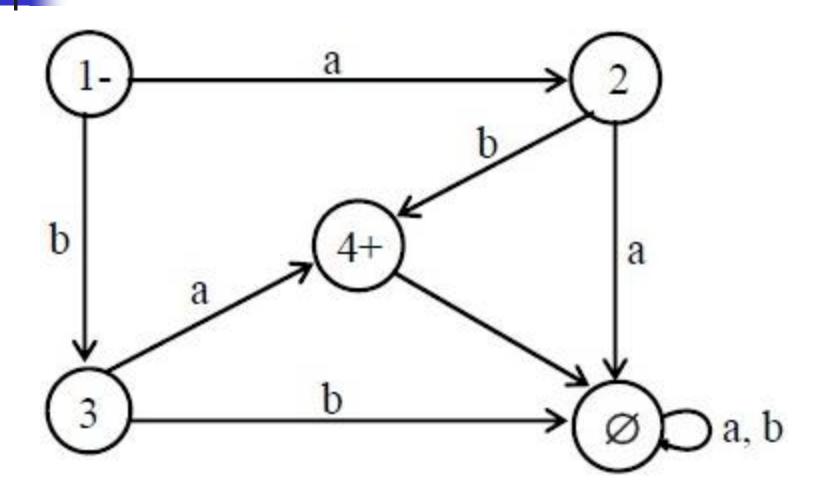
Method 1:

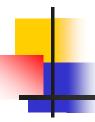
Since in an NFA, there may be more than one transition for a certain letter and there may not be any transition for certain letter, so starting from the initial state corresponding to the initial state of given NFA, the transition diagram of the corresponding FA, can be built introducing an empty state for a letter having no transition at certain state and a state corresponding to the combination of states, for a letter having more than one transitions.







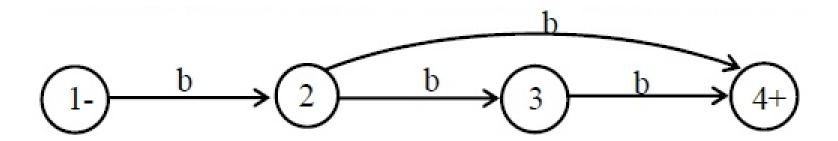




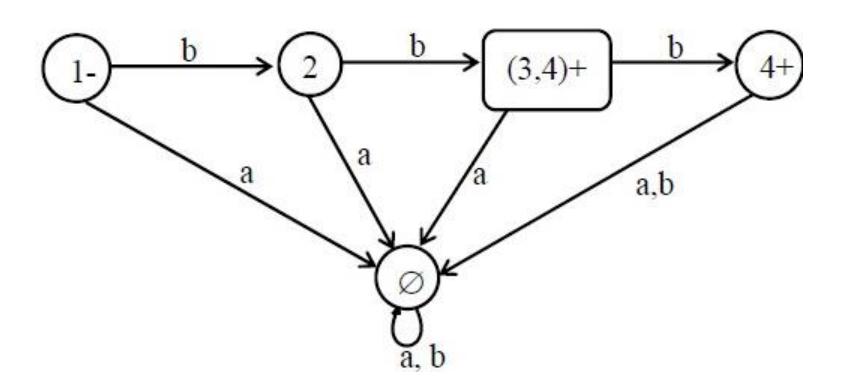
■ Language $L = \{bb,bbb\}$ defined over $\Sigma = \{a,b\}$.



Language L = {bb,bbb} defined over ∑ = {a,b}.







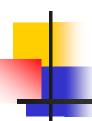


Method 2:

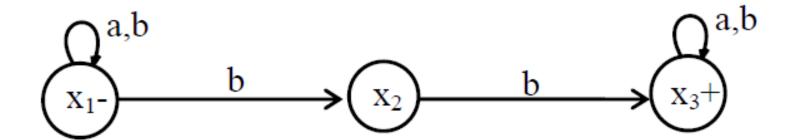
In an NFA, there may be more than one transition for a certain letter and there may not be any transition for certain letter, so starting from the initial state corresponding to the initial state of given NFA, the transition table along with new labels of states, of the corresponding FA, can be built introducing an empty state for a letter having no transition at certain state and a state corresponding to the combination of states, for a letter having more than one transitions



Language L defined over ∑ = {a,b} accepting words containing bb.



Language L defined over ∑ = {a,b} accepting words containing bb.



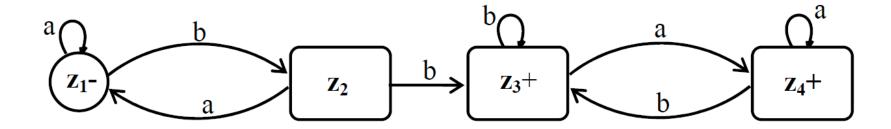
Old States	New States after reading	
	a	b
$z_1 = x_1$	$x_1 \equiv z_1$	$(x_1,x_2)\equiv z_2$

Old States	New States after reading	
	a	b
$z_1 = x_1$	$x_1 \equiv z_1$	$(x_1,x_2)\equiv z_2$
$z_2 \equiv (x_1, x_2)$	$(x_1, \emptyset) \equiv x_1 \equiv z_1$	$(x_1,x_2,x_3) \equiv z_3$

Old States	New States after reading	
	a	ь
$z_1 = x_1$	$x_1 \equiv z_1$	$(\mathbf{x}_1, \mathbf{x}_2) \equiv \mathbf{z}_2$
$z_2 \equiv (x_1, x_2)$	$(x_1, \emptyset) \equiv x_1 \equiv z_1$	$(x_1, x_2, x_3) \equiv z_3$
$z_3 + \equiv (x_1, x_2, x_3)$	$(x_1,x_3)\equiv z_4$	$(x_1, x_2, x_3) \equiv z_3$

Old States	New States after reading	
	a	b
$z_1 = x_1$	$x_1 = z_1$	$(x_1,x_2)\equiv z_2$
$z_2 \equiv (x_1, x_2)$	$(x_1, \emptyset) \equiv x_1 \equiv z_1$	$(x_1, x_2, x_3) \equiv z_3$
$z_3 + \equiv (x_1, x_2, x_3)$	$(x_1,x_3) \equiv z_4$	$(x_1, x_2, x_3) \equiv z_3$
$z_4 + \equiv (x_1, x_3)$	$(x_1,x_3) \equiv z_4$	$(x_1, x_2, x_3) \equiv z_3$





Exercises

- Language L defined over ∑ = {a,b} accepting words containing aa.
- Language L defined over ∑ = {a,b} accepting words ending in a.
- Language L defined over ∑ = {a,b} accepting words starting with a.
- Language L defined over ∑ = {a,b} accepting words starting and ending in same letters.